

UNIT - II

Set Theory (George Cantor)

* Set theory was developed by George Cantor.

Set: A well defined collection of objects (or) an idea is known as set.

* Well defined means that

- i. All the objects in the set should have a common feature in the property.
- ii. It should be possible to decide whether any given object belongs to the set or not.

Note:

Given a set 'S' and object 'P', exactly one of the statement following statements should be true
① $P \in S$
② $P \notin S$ (does not belongs to)

Representation of a set:

A set can be represented by 2 methods.

① Rule method (or) set builder method.

② Roster method (or) Tabulation.

① Rule method

In this method the set is represented by a statement (or) a rule (or) a common property which all the members of the set possess that is if the common property possessed by all the elements in 'P', the set is represented in the rule method by $S = \{x : x \text{ has the property } P\}$

P. this is read as "the set of all elements x such that x has property $P\}$ "

Ex:-

i. If P is the set of all prime numbers, then

$$P = \{x; x \text{ is a prime number}\}$$

ii. $A = \{1, 2, 5, 10, 25\}$ can be written as $A = \{x; x \text{ is a positive divisor of } 50 \text{ which is less than } 50\}$

iii. If A is the set of all natural numbers

b/w 10 and 100, then $A = \{x; x \in \mathbb{N} \text{ and } 10 < x < 100\}$

② Roster Method

In this method the elements of the given set are enlisted and enclosed in brackets (curly brackets).

Ex:-

i. If $A = \{a, b, c, 1, 2, 3\}$ then where

$$A = \{\text{set of letters } a, b, c \text{ and numbers } 1, 2, 3\}$$

ii. If A' is the set of letters of the word "MATHEMATICS", then in roster form

$$A' = \{M, A, T, H, E, M, A, T, I, C, S\}$$

Null set (or) Empty set

A set which does not contain any element is called an empty set (or) a null set. An empty set will be denoted by \emptyset (or) $\{\}$.

Ex:- A set of all even numbers b/w 2-4

Finite & Infinite Set:

A set is said to be finite if it has a finite number of elements otherwise it is said to be infinite.

Ex:- i) The set of vowels in English alphabet is a finite set.

ii) The set of natural numbers is an infinite set.

singleton set:

A set containing a single element is called
o singleton set
Ex: $\{1\}, \{a\}, \{\emptyset\}, \{\pi\}$

cardinality of a finite set

The number of elements in a finite set 'A' is called its cardinal number and is denoted by $n(A)$ (or) $|A|$

* If $A = \{1, 2, 3, 4\}$

$$n(A) = 4$$

complement of a set

Let 'U' be the universal set and $A \subset U$. Then the complement of A is denoted by A^c symbolically $A^c = U - A$

* $A^c = \{x : x \in U \text{ and } x \notin A\}$

$$A^c = \{x : x \notin A\}$$

Ex: If U is a set of natural numbers and A is the set of even numbers, then $A^c = U - A =$ the set of natural numbers which are not even.

* A^c = The set of odd natural numbers.

equal sets:

Two sets $A \& B$ is said to be equal if they have the same elements, then we write $A=B$.

Note: Two sets $A \& B$ are equal: $A \subseteq B$ and $B \subseteq A$.

we write $A=B$

Ex: $A = \{4, 5, 7\}$ $B = \{1, 4, 5, 7\}$

$$\boxed{A=B}$$

Sub Sets:

A set A is a subset of the set B, if and only if every element of A is also an element of B. we also say that A is contained B and use the notation $A \subseteq B$.

symbolically $x \in A \Rightarrow x \in B$ then $A \subseteq B$ otherwise we write $A \not\subseteq B$.

Proper set:

* A set A is called proper set of the set B.

If (i) A is subset of B

(ii) B is not a subset of A

i.e. A is said to be proper set of B if every element of A $\in B$ (belongs to set B) but there is at least one element of B which is not in A.

denoted $A \subset B$

Super Set:

If A is a subset of B, then B is called a super set of A.

Example:

Subset

$$A = \{1, 2, 3\}$$

$$B = \{0, 2, 4, 3, 7\}$$

proper set

$$A = \{1, 2, 3\}$$

$$B = \{1, 2, 3, 4, 5, 6\}$$

super set

$A = \{1, 2, 3\}$ $B = \{1, 2, 3, 4, 5, 6\}$, A is a proper subset of B, and B is a supersubset of A.

Power set

For a set A , a collection (or) family of all subsets of A is called power set of A .

* The power set of A is denoted by $P(A)$ (or) 2^A

$$* P(A) \text{ or } 2^A = \{x : x \subseteq A\}$$

Ex: For a set $S = \{a\}$; $P(S) = \{\emptyset, \{a\}\} = \{\emptyset, S\}$

$S = \{a, b\}$, $P(S) = \{\emptyset, \{a\}, \{b\}, \{a, b\}\} = \{\emptyset, S, \{a\}, \{b\}\}$

Universal Set,

$\{\{a\}, \{b\}, \{S\}\}$

A set 'U' which contains all the sets under consideration as subsets is called a universal set.

Ex: In plane geometry the universal set consists of all the points in the plane.

Disjoint set

Two sets are said to be disjoint, if they have no element in common.

Ex: $A = \{1, 2, 3\}$, $B = \{4, 5, 6\}$ then A & B are disjoint
 $A \cap B = \emptyset$

Ordered set

A set which an ordering imposed on it is said to be an ordered set.

Ex: $S = \{2, 3, 6, 8, \dots\}$

Same operations on set

Intersection

The intersection of any two sets A and B written as $A \cap B$, is the set consisting of all the elements which belong to both A and B .

* Symbolically $A \cap B = \{x : x \in A \cap x \in B\}$

* For any sets A and B , $A \cap B = B \cap A$

$$A \cap A = A$$

$$A \cap \emptyset = \emptyset, (A \cap B) \cap C = A \cap (B \cap C).$$

- ① If $A_1 = \{\{1, 2\}, \{3\}\}$, $A_2 = \{\{1\}, \{2, 3\}\}$ & $A_3 = \{\{1, 2, 3\}\}$
Then show that A_1 , A_2 and A_3 are mutually disjoint.

Sol $A_1 \cap A_2 = \emptyset$

$A_1 \cap A_3 = \emptyset$

$A_2 \cap A_3 = \emptyset$ the given sets are disjoint.

- ② Show that $A \subseteq B \iff A \cap B = A$

Sol For any x , $P \rightarrow Q \iff (P \wedge Q) \geq P$

$$x \in A \rightarrow x \in B \iff (x \in A \wedge x \in B) \geq x \in A$$

$$A \subseteq B \iff (\forall x)(x \in A \rightarrow x \in B) [:\exists x: x \in A \rightarrow x \in B]$$

$$A \cap B = A \iff (\forall x)(x \in A \wedge x \in B) \geq x \in A$$

Union:

* For any two sets A and B , the union of A and B is written as " $A \cup B$ " is the set of all elements which are members of the set A or the set B (or) both.

* symbolically $A \cup B = \{x: x \in A \vee x \in B\}$

$$A \cup B = B \cup A \text{ (Commutative)}$$

$$A \vee \emptyset = A$$

$$A \vee A = A$$

$$(A \cup B) \cup C = A \cup (B \cup C)$$

- ① Show that $A \cup A = A$

- Sol For any x

$$x \in (A \cup A) \iff x \in \{x: x \in A \vee x \in A\}$$

$$\iff x \in A \vee x \in A$$

$$\iff x \in A$$

$$\therefore A \cup A = A \iff x \in A$$

② If $A_1 = \{1, 2\}$, $A_2 = \{2, 3\}$ & $A_3 = \{1, 2, 3\}$ what
are $\bigcup_{i=1}^3 A_i$ & $\bigcap_{i=1}^3 A_i$

Sol Given sets $A_1 = \{1, 2\}$, $A_2 = \{2, 3\}$, $A_3 = \{1, 2, 3\}$

$$\begin{aligned}\bigcup_{i=1}^3 A_i &= A_1 \cup A_2 \cup A_3 \\ &= \{1, 2, 3, 6\}\end{aligned}$$

$$\bigcap_{i=1}^3 A_i = A_1 \cap A_2 \cap A_3$$

$$\bigcap_{i=1}^3 A_i = \{2\}$$

Relative complement (or) Difference

Let A and B be any two sets the relative complement of B in A is written as $A-B$, is the set consisting all elements of A which are not elements of B that is $A-B = \{x : x \in A \wedge x \notin B\}$
(or) $\{x : x \in A \wedge \neg(x \in B)\}$

If $A = \{4, 5, 6\}$, $B = \{3, 4, 2\}$, $C = \{1, 3, 4\}$ find
 $A-B$, $B-A$ show that $A-B \neq B-A$ and $A-C = A$

Sol Given $A = \{4, 5, 6\}$, $B = \{3, 4, 2\}$, $C = \{1, 3, 4\}$

$$\begin{aligned}A-B &= \{4, 5, 6\} - \{3, 4, 2\} \\ &= \{5, 6\}\end{aligned}$$

$$\begin{aligned}B-A &= \{3, 4, 2\} - \{4, 5, 6\} \\ &= \{3, 4\}\end{aligned}$$

$$A-B \neq B-A$$

$$\begin{aligned}A-C &= \{4, 5, 6\} - \{1, 3, 4\} \\ &= \{2, 5, 6\}\end{aligned}$$

$$\boxed{A-C \neq A}$$

② show that ① $A - B = A \cap B^c$ ② $A \subseteq B \Leftrightarrow B^c \subseteq A^c$

so i. for any x ,

$$x \in (A - B) \Leftrightarrow x \in \{x : x \in A \wedge x \notin B\}$$

$$\Leftrightarrow x \in A \wedge x \notin B$$

$$\Leftrightarrow x \in A \wedge x \in B^c$$

$$\Leftrightarrow x \in (A \cap B^c)$$

$$\therefore \text{i.e } A - B = A \cap B^c$$

ii. for any x ,

$$(A \subseteq B) \Leftrightarrow (\forall x)(x \in A \rightarrow x \in B)$$

$$(P \rightarrow Q \Leftrightarrow \neg Q \rightarrow \neg P)$$

$$\Leftrightarrow (\forall x)(\neg(x \in A) \rightarrow \neg(x \in B))$$

$$\Leftrightarrow (\forall x)(x \notin B \rightarrow x \notin A)$$

$$\Leftrightarrow (\forall x)(x \in B^c \rightarrow x \in A^c)$$

$$(A \subseteq B) \Leftrightarrow B^c \subseteq A^c$$

3. show that for any two sets A and B

$$A - (A \cap B) = A - B$$

so for any x

$$x \in (A - (A \cap B)) \Leftrightarrow x \in \{x : x \in A \wedge x \notin (A \cap B)\}$$

$$\Leftrightarrow x \in A \wedge x \notin (A \cap B)$$

$$\Leftrightarrow x \in A \wedge \neg(x \in (A \cap B))$$

$$\Leftrightarrow x \in A \wedge \neg(x \in A \wedge x \in B)$$

$$\Leftrightarrow (x \in A) \wedge (x \notin A \vee x \notin B)$$

using distributive law.

$$\Leftrightarrow (x \in A \wedge x \notin A) \vee (x \in A \wedge x \notin B)$$

$$\Leftrightarrow F \wedge (x \in A \wedge x \notin B)$$

$$(P \vee F \Leftrightarrow P)$$

$$\Leftrightarrow x \in A \wedge x \notin B$$

$$\Leftrightarrow x \in \{x : x \in A \wedge x \notin B\}$$

$$\Leftrightarrow x \in A \wedge x \notin B$$

$$\Leftrightarrow x \in A - B$$

$$A - (B \cap B) = A - B$$

Symmetric Difference (or) Boolean sum

Let A and B be any two sets. The symmetric difference of A and B is the set $A + B$ defined by

$$(or) A \cap B' \cup B \cap A'$$

$$A + B = (A - B) \cup (B - A)$$

* symbolically $x \in A + B \Leftrightarrow x \in \{x : x \in A \vee x \in B\}$
where \vee is the exclusive disjunction.

$$A + B = B + A$$

$$(A + B) + C = A + (B + C)$$

$$A + \emptyset = A, A + A = \emptyset$$

1. Prove that $A + \emptyset = A$

so) For any x

$$x \in (A + \emptyset) \Leftrightarrow x \in \{x : (x \in A \wedge x \notin \emptyset) \vee (x \in \emptyset \wedge x \notin A)\}$$

$$\Leftrightarrow (x \in A \wedge x \in \emptyset) \vee (x \in \emptyset \wedge x \notin A)$$

$$\Leftrightarrow (x \in A \wedge F) \vee (F \wedge x \notin A)$$

$$\Leftrightarrow (x \in A) \vee F$$

$$\Leftrightarrow x \in A$$

$$\Leftrightarrow x \in \{x : x \in A\}$$

$$\Leftrightarrow x \in A$$

$$\therefore A + \emptyset = A$$

2. Given $A = \{x : x \text{ is an integer and } 1 \leq x \leq 5\}$,
 $B = \{3, 4, 5, 17\}$ & $C = \{1, 2, 3, \dots\}$ find $A \cap B$,
 $A \cap C$, $A \cup B$ & $A \cup C$.

sol

Given $A = \{1, 2, 3, 4, 5\}$,

$B = \{3, 4, 5, 17\}$

$C = \{1, 2, 3, \dots\}$

$$A \cap B = \{3, 4, 5\}$$

$$A \cap C = \{1, 2, 3, 4, 5\} = A$$

$$A \cup B = \{1, 2, 3, 4, 5, 17\}$$

$$A \cup C = \{1, 2, 3, 4, 5, \dots\} = C$$

3. Given $A = \{2, 3, 4\}$, $B = \{1, 2\}$ & $C = \{4, 5, 6\}$, find
 $A+B$, $B+C$, $A+B+C$ & $(A+B)+B+C$.

Given $A = \{2, 3, 4\}$, $B = \{1, 2\}$; $C = \{4, 5, 6\}$

$$A+B = (A-B) \cup (B-A)$$

$$= (\{2, 3, 4\} - \{1, 2\}) \cup (\{1, 2\} - \{2, 3, 4\})$$

$$= \{3, 4\} \cup \{1\}$$

$$\boxed{A+B = \{1, 3, 4\}}$$

$$B+C = (B-C) \cup (C-B)$$

$$= (\{1, 2\} - \{4, 5, 6\}) \cup (\{4, 5, 6\} - \{1, 2\})$$

$$= \{1, 2\} \cup \{4, 5, 6\}$$

$$\boxed{B+C = \{1, 2, 4, 5, 6\}}$$

$$(A+B)+C = ((A+B)-C) \cup (C-(A+B))$$

$$= (\{1, 3, 4\} - \{4, 5, 6\}) \cup (\{4, 5, 6\} - \{1, 3, 4\})$$

$$= \{1, 3\} \cup \{5, 6\}$$

$$\boxed{A+B+C = \{1, 3, 5, 6\}}$$

$$(A+B)+(B+C) = \{(A+B)-(B+C)\} \cup \{(B+C)-(A+B)\}$$

$$= (\{1, 3, 4\} - \{1, 2, 4, 5, 6\}) \cup (\{1, 2, 4, 5, 6\} - \{1, 3, 4\})$$

$$= \{3\} \cup \{2, 5, 6\}$$

$$\boxed{(A+B)+(B+C) = \{3, 2, 5, 6\}}$$

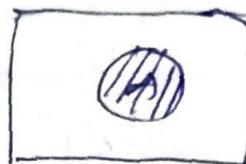
Venn Diagrams

The universal set permits the use of a pictorial device to study the connections b/w the sub-sets of a universal set and their intersection union, difference and other operations.

* A Venn diagram is a schematic representation of a set by a set of points. The universal set is represented by a set of points in a rectangle and a subset say 'A' of 'E' is represented by the interior of a circle (or) some other simple closed curve inside the rectangle.



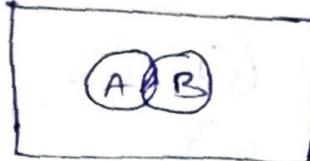
E



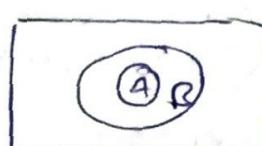
A



$A \cup B$



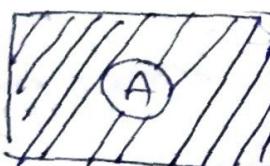
$A \cap B$



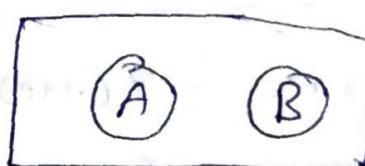
$A \subseteq B$



$A - B$



$A^c = E - A$



$A \cap B = \emptyset$

ordered pairs and n-tuples

An ordered pair consists of two objects in a given fixed order. Note that an ordered pair is not a set consisting two elements. The ordering of the two objects is important. The two objects need to be distinct. We shall denote an order pair by (x, y) .

* A familiar example of an ordered pair is the representation of a point in a two dimensional plane in Cartesian coordinates. Accordingly the ordered pairs $\langle 6,2 \rangle$, $\langle 2,4 \rangle$, $\langle 1,3 \rangle$, $\langle 4,5 \rangle$ represent different points in the plane.

* The equality of two ordered pairs $\langle x,y \rangle$ and $\langle u,v \rangle$ is defined by $\langle x,y \rangle = \langle u,v \rangle \Leftrightarrow (x=u) \wedge (y=v)$.

* so that $\langle 1,2 \rangle \neq \langle 2,1 \rangle$ and $\langle 1,1 \rangle \neq \langle 2,2 \rangle$.

* Order triple $\langle \langle x,y,z \rangle, w \rangle \Leftrightarrow \langle \langle u,v,w \rangle, z \rangle$

$$((x=u) \wedge (y=v) \wedge (z=w))$$

* Order 'n' tuple $\langle \langle x_1, x_2, \dots, x_{n-1}, x_n \rangle, u_n \rangle = \langle \langle u_1, u_2, \dots, u_{n-1}, u_n \rangle, x_n \rangle \Leftrightarrow$

$$((x_1=u_1) \wedge (x_2=u_2) \wedge (x_3=u_3) \wedge \dots \wedge (x_n=u_n))$$

Cartesian products

Let 'A' and 'B' be any two sets. The set of all ordered pairs such that the first member of the ordered pair is an element of A and the second member is an element of B is called the Cartesian product of A & B is written as

$A \times B$

$$A \times B = \{ \langle x,y \rangle : (x \in A) \wedge (y \in B) \}$$

1. If $A = \{x, y\} \& B = \{1, 2, 3\}$ what are $A \times B$, $B \times A$, $A \times A$, $B \times B$, $(A \times B) \cap (B \times A)$.

Q1 Given $A = \{x, y\}$, $B = \{1, 2, 3\}$

$$A \times B = \{x, y\} \times \{1, 2, 3\}$$

$$= \{ \langle x, 1 \rangle, \langle x, 2 \rangle, \langle x, 3 \rangle, \langle y, 1 \rangle, \langle y, 2 \rangle, \langle y, 3 \rangle \}$$

$$B \times A = \{1, 2, 3\} \times \{x, y\}$$

$$= \{ \langle 1, \alpha \rangle, \langle 1, \beta \rangle, \langle 2, \alpha \rangle, \langle 2, \beta \rangle, \langle 3, \alpha \rangle, \langle 3, \beta \rangle \}$$

$$B \times D = \{1, 2, 3\} \times \{1, 2, 3\}$$

$$= \{ \langle 1, 1 \rangle, \langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 2, 1 \rangle, \langle 2, 2 \rangle, \langle 2, 3 \rangle, \langle 3, 1 \rangle, \langle 3, 2 \rangle, \langle 3, 3 \rangle \}$$

$$A \times A = \{ \alpha, \beta \} \times \{ \alpha, \beta \}$$

$$= \{ \langle \alpha, \alpha \rangle, \langle \alpha, \beta \rangle, \langle \beta, \alpha \rangle, \langle \beta, \beta \rangle \}$$

$$(A \times B) \cap (B \times A) = \emptyset$$

2 If $A = \emptyset$ & $B = \{1, 2, 3\}$ what are $A \times B$ & $B \times A$?

Sol Given $A = \emptyset$

$$B = \{1, 2, 3\}$$

$$A \times B = \{\emptyset\} \times \{1, 2, 3\}$$

$$A \times B = \emptyset$$

$$B \times A = \emptyset$$

$$A \times B = B \times A = \emptyset$$

Cartesian product more than two sets

$$\begin{aligned} (A \times B) \times C &= \{ \langle \langle a, b \rangle, c \rangle : \langle a, b \rangle \in A \times B \wedge c \in C \} \\ &= \{ \langle a, b, c \rangle : a \in A \wedge b \in B \wedge c \in C \} \end{aligned}$$

* write $A \times B$, $A \times B \times C$, B^2 , $A^3 \times B^2$, $B^2 \times A$ where

$$A = \{1\}, B = \{a, b\} \text{ & } C = \{2, 3\}$$

Sol Given $A = \{1\}$

$$B = \{a, b\}$$

$$C = \{2, 3\}$$

$$A \times B = \{1\} \times \{a, b\}$$

$$= \{ \langle 1, a \rangle, \langle 1, b \rangle \}$$

$$(A \times B) \times C = \{1\} \times \{a, b\} \times \{1, 2, 3\} \text{ (or)} \{<1, a>, <1, b>\} \times \{2, 3\}$$

$$= \{<1, a, 2>, <1, a, 3>, <1, b, 2>, <1, b, 3>\}$$

$$B^2 = B \times B$$

$$= \{a, b\} \times \{a, b\}$$

$$= \{<a, a>, <a, b>, <b, a>, <b, b>\}$$

$$A^3 = A \times A^2$$

$$= \{1\} \times \{1, 2\} \times \{1\}$$

$$= \{<1, 1, 1>\}$$

$$A^2 = \{1, 2\}$$

$$= \{1, 2\}$$

$$B^2 \times A = \{<a, a>, <a, b>, <b, a>, <b, b>\} \times \{1\}$$

$$= \{<a, a, 1>, <a, b, 1>, <b, a, 1>, <b, b, 1>\}$$

Relation and ordering

Relations

The word relation suggests some familiar examples of relations such as the relation of father to son, mother to son, brother to sister etc. Familiar examples in arithmetic are relations such as greater than, less than (or) that of equality b/w two numbers.

We also know the relation b/w the area of the circle and its radius and b/w the area of square and its sides. This examples suggest relations b/w two objects.

* The relation b/w parent & child the coincidence of three lines and that point lying b/w two points are examples of relations among three objects. Similar examples exist for relations among 4 (or) more objects.

Binary Relation

Any set of ordered pairs defines an binary relations.

* we shall call a binary relation simply a relation. It is sometimes convenient to express the fact that a particular order pair say $\langle x, y \rangle \in R$,

where ' R ' is a relation by writing xRy which may be read as " x is in relation R to y ".

* In mathematics relations are often denoted by special symbols rather than by capital letters. A familiar example is the relation greater than for real numbers. This relation is denoted by ' $>$ '.

* If a & b are two real numbers such that $a > b$ then we say that $\langle a, b \rangle \in (or)$ $a > b$.

* more precisely relation greater than is

$R = \{ \langle x, y \rangle ; x, y \text{ are real numbers and } x > y \}$,
similarly

$F = \{ \langle x, y \rangle ; x \text{ is the father of } y \}$

$S = \{ \langle x, y \rangle ; \langle x, z \rangle, \langle z, y \rangle \in S \}$
where F & S are relations,

Domain & Range

Let 'S' be a binary relation the set $D(S)$ of all objects 'x' such that for some 'y'; $\langle x, y \rangle \in S$ is called the domain of 'S' i.e $D(S) = \{ x : (\exists y) \langle x, y \rangle \in S \}$

Similarly the set $R(S)$ of all objects $\langle x, y \rangle \in S$ such that for some 'y'; $\langle x, y \rangle \in S$ is called the range of 'S' $R(S) = \{ y : (\exists x) \langle x, y \rangle \in S \}$

$$\text{Ex 1: } S = \{ \langle 2, 4 \rangle, \langle 5, 6 \rangle, \langle 7, 8 \rangle \}$$

$$D(S) = \{ 2, 5, 7 \}$$

$$R(S) = \{ 4, 6, 8 \}$$

1. Let $x = \{1, 2, 3, 4\}$ if $R = \{ \langle x, y \rangle; x \in X \cap y \in X \text{ and } (x-y) \text{ is an internal non-zero multiple of 2} \}$
 $S = \{ \langle x, y \rangle / x \in X \cap y \in X \cap (x-y) \text{ is an integral non-zero multiple of 3} \}$.

Find RUS & RNS.

sol Given $X = \{1, 2, 3, 4\}$

$$R = \{ \langle 1, 3 \rangle, \langle 3, 1 \rangle, \langle 2, 4 \rangle, \langle 4, 2 \rangle \}$$

$$S = \{ \langle 1, 4 \rangle, \langle 4, 1 \rangle \}$$

$$\text{RUS} = \{ \langle 1, 3 \rangle, \langle 3, 1 \rangle, \langle 2, 4 \rangle, \langle 4, 2 \rangle, \langle 1, 4 \rangle, \langle 4, 1 \rangle \}$$

$$R \cap S = \emptyset$$

2. Let ' L ' denote the relation "less than", ' D ' denote the relation "equal to" and ' D ' denote the relation "divides", $x D y$ means " x divides y ". Both ' L ' and ' D ' are defined in the set $\{1, 2, 3, 6\}$. Write " L " and " D " as sets and find " $L \cap D$ ".

sol Given set is $\{1, 2, 3, 6\}$

$$L = \{ \langle 1, 1 \rangle, \langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 1, 6 \rangle, \langle 2, 2 \rangle, \langle 2, 3 \rangle, \langle 2, 6 \rangle, \\ \langle 3, 3 \rangle, \langle 3, 6 \rangle \}$$

$$D = \{ \langle 1, 1 \rangle, \langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 1, 6 \rangle, \langle 2, 1 \rangle, \langle 2, 6 \rangle, \langle 3, 1 \rangle, \\ \langle 3, 2 \rangle, \langle 3, 3 \rangle, \langle 6, 1 \rangle \}$$

$$L \cap D = \{ \langle 1, 1 \rangle, \langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 1, 6 \rangle, \langle 2, 1 \rangle, \langle 2, 6 \rangle, \\ \langle 3, 1 \rangle, \langle 3, 2 \rangle, \langle 3, 3 \rangle, \langle 6, 1 \rangle \} = D$$

3. Let $P = \{\langle 4,2 \rangle, \langle 2,4 \rangle, \langle 3,3 \rangle\}$ and $Q = \{\langle 1,3 \rangle, \langle 2,4 \rangle, \langle 4,2 \rangle\}$ find $P \cup Q$, $D(P)$, $D(Q)$, $D(P \cup Q)$, $R(P)$, $R(Q)$ & $R(P \cap Q)$ and also show that $D(P \cup Q) = D(P) \cup D(Q)$, $R(P \cap Q) \subseteq R(P) \cap R(Q)$

Sol. Given $P = \{\langle 4,2 \rangle, \langle 2,4 \rangle, \langle 3,3 \rangle\}$

$$Q = \{\langle 1,3 \rangle, \langle 2,4 \rangle, \langle 4,2 \rangle\}$$

$$P \cup Q = \{\langle 1,2 \rangle, \langle 2,4 \rangle, \langle 3,3 \rangle, \langle 4,3 \rangle, \langle 4,2 \rangle\}$$

$$D(P) = \{1,2,3\}$$

$$D(Q) = \{1,2,4\}$$

$$D(P \cup Q) = \{1,2,3,4\}$$

$$R(P) = \{2,4,3\}$$

$$R(Q) = \{3,4,2\}$$

$$P \cap Q = \{\langle 2,4 \rangle\}$$

$$D(P \cap Q) = \{2\}$$

$$D(P \cup Q) = \{1,2,3,4\}$$

$$D(P) \cup D(Q) = \{1,2,3,4\}$$

$$\boxed{P(P \cup Q) = P(P) \cup P(Q)}$$

$$R(P \cap Q) = \{4\}$$

$$R(P) \cap R(Q) = \{4,2,3\}$$

$$R(P \cap Q) \subseteq R(P) \cap R(Q)$$

$$\{4\} \subseteq \{2,3,4\} \quad (\subseteq = \text{Subset})$$

Properties of Binary Relations in sets:

Reflexive: A binary relation R in a set X is reflexive if for every $x \in X$, $x R x$, $\langle x, x \rangle \in R$. (or) R is reflexive in X

$\Leftrightarrow \forall x (x \in X \rightarrow x R x)$

- * The relation \leq is reflexive in the set of real numbers.
- * Since for any ' x ' we have $x \leq x$.
- * Similarly the relation of inclusion is reflexive in the family of all subsets of universal sets. The relation of equality of sets is also reflexive.
- * The relation ' $<$ ' is not reflexive in the set of real numbers.

Symmetric:

A relation ' R ' in a set ' X ' is symmetric for every x and y in X whenever xRy then yRx i.e. R is symmetric in X $\Leftrightarrow (x)(y)$
 $(x \in X \wedge y \in X \wedge xRy \rightarrow yRx)$

- * The relations ' \leq ' and ' $<$ ' are not symmetric in the set of real numbers, while the relation of equality is the relation of similarity in the set of triangles in a plane is both reflexive and symmetric.

Transitive: The relation ' R ' in a set ' X ' is transitive if for every x, y and z in X : whenever xRy , yRz , then xRz .

i.e. R is transitive in ' X ', whenever $xRy \wedge yRz \rightarrow xRz$

- * The relations ' \leq ', ' $<$ ', ' $=$ ' are transitive in the set of real numbers. The relations ' \subseteq ', ' \subset ' and equality are also transitive in the family of subsets of a universal set.

Irreflexive: A relation ' R ' in a set ' X ' is irreflexive for every $x \in X$, $\langle x, x \rangle \notin R$.

* Note that any relation which is not reflexive is not necessarily irreflexive vice versa.

* The relation ' \neq ' in the set of real numbers is irreflexive.

Anti symmetric:

~~~~~

\* Relation 'R' in a set 'X' is anti symmetric if for every  $x \in X$ , whenever  $x R y$ , then  $x = y$ .

\* Symbolically 'R' is anti symmetric in 'X' if and only if ( $\forall x, y \in X$ ) ( $x R y \wedge x \neq y \rightarrow x = y$ )

Equivalence Relation:

~~~~~

A relation 'R' on a set 'A' is an equivalence relation if and only if ($\forall x, y, z \in A$)

- R is reflexive
- R is symmetric

for every $a, b \in A$, $\langle a, b \rangle \in R \Rightarrow \langle b, a \rangle \in R$

iii. R is transitive

$\langle a, b \rangle \in R, \langle b, c \rangle \in R \Rightarrow \langle a, c \rangle \in R$

Ex: $A = \{a, b, c\}$, $R = \{\langle a, a \rangle, \langle b, b \rangle, \langle c, c \rangle, \langle a, b \rangle, \langle b, c \rangle, \langle a, c \rangle, \langle a, b \rangle, \langle b, a \rangle, \langle a, c \rangle, \langle c, a \rangle\}$ is an equivalence relation.

Q1 Let $A = \{1, 2, 3, 4\}$ and $R = \{\langle 1, 1 \rangle, \langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 1, 2 \rangle, \langle 3, 1 \rangle, \langle 3, 2 \rangle, \langle 3, 3 \rangle, \langle 1, 3 \rangle, \langle 3, 1 \rangle, \langle 2, 1 \rangle, \langle 2, 3 \rangle, \langle 2, 2 \rangle, \langle 4, 1 \rangle, \langle 4, 4 \rangle\}$ is it an equivalence relation? why?

Sol Given $A = \{1, 2, 3, 4\}$

$R = \{\langle 1, 1 \rangle, \langle 1, 2 \rangle, \langle 2, 1 \rangle, \langle 2, 2 \rangle, \langle 3, 1 \rangle, \langle 3, 3 \rangle, \langle 1, 3 \rangle, \langle 3, 1 \rangle, \langle 4, 4 \rangle\}$

i. Reflexive: for every $a \in A$, $\langle a, a \rangle \in R$

1EA, $\langle 1,1 \rangle \in R$, $2 \in A, \langle 2,2 \rangle \in R$

$3 \in A, \langle 3,3 \rangle \in R$,

$4 \in A, \langle 4,4 \rangle \in R$.

Symmetric:

for every $a, b \in A, \langle a, b \rangle \in R$ then $\langle b, a \rangle \in R$.

If $\langle 1,2 \rangle \in R \Rightarrow \langle 2,1 \rangle \in R$.

If $\langle 2,1 \rangle \in R \Rightarrow \langle 1,2 \rangle \in R$

If $\langle 3,1 \rangle \in R \Rightarrow \langle 1,3 \rangle \in R$

If $\langle 1,3 \rangle \in R \Rightarrow \langle 3,1 \rangle \in R$

If $\langle 4,1 \rangle \in R \Rightarrow \langle 1,4 \rangle \notin R$

R is not symmetric.

$\therefore R$ is not an equivalence relation there is no need to check after transitive.

2. If $A = \{1, 2, 3, 4\}$, $R = \{\langle 1,1 \rangle, \langle 1,2 \rangle, \langle 2,1 \rangle, \langle 2,2 \rangle, \langle 3,3 \rangle, \langle 3,4 \rangle, \langle 4,3 \rangle, \langle 4,4 \rangle\}$. Check R is an equivalence or not?

Sol Given $A = \{1, 2, 3, 4\}$

$$R = \{\langle 1,1 \rangle, \langle 1,2 \rangle, \langle 2,1 \rangle, \langle 2,2 \rangle, \langle 3,3 \rangle, \langle 3,4 \rangle, \langle 4,3 \rangle, \langle 4,4 \rangle\}$$

Reflexive:

for every $a \in A, \langle a, a \rangle \in R$

1EA, $\langle 1,1 \rangle \in R$

2EA, $\langle 2,2 \rangle \in R$

3EA, $\langle 3,3 \rangle \in R$

4EA, $\langle 4,4 \rangle \in R$

Symmetric: for every $a, b \in A, \langle a, b \rangle \in R$, then $\langle b, a \rangle \in R$

If $\langle 1,2 \rangle \in R \Rightarrow \langle 2,1 \rangle \in R$

If $\langle 2,1 \rangle \in R \Rightarrow \langle 1,2 \rangle \in R$

If $\langle 3,4 \rangle \in R \Rightarrow \langle 4,3 \rangle \in R$

If $\langle 4,3 \rangle \in R \Rightarrow \langle 3,4 \rangle \in R$

$\therefore R$ is symmetric.

Transitive

For every $a, b, c \in A$

If $\langle a,b \rangle \in R \& \langle b,c \rangle \in R$, then $\langle a,c \rangle \in R$

If $\langle 1,2 \rangle \in R, \langle 2,1 \rangle \in R \Rightarrow \langle 1,1 \rangle \in R$

If $\langle 3,4 \rangle \in R, \langle 4,3 \rangle \in R \Rightarrow \langle 3,3 \rangle \in R$

If $\langle 2,1 \rangle \in R, \langle 1,2 \rangle \in R \Rightarrow \langle 2,2 \rangle \in R$

If $\langle 4,3 \rangle \in R, \langle 3,4 \rangle \in R \Rightarrow \langle 4,4 \rangle \in R$

$\therefore R$ is transitive.

$\therefore R$ is an equivalence relation.

IMP * composition of binary relation

Let R be a relation from X to Y and S be a relation written as "pos" is called a composite relation of $R \& S$ where $Ros = \{ \langle x, z \rangle : x \in X \wedge z \in Z \wedge (\exists y)(y \in Y \wedge \langle x, y \rangle \in R \wedge \langle y, z \rangle \in S) \}$. The operation of obtaining pos from $R \& S$ is called composition of relation.

$R \rightarrow X \leftarrow X$

$S \rightarrow Y \leftarrow Z$

$X \xrightarrow{R} Y \xrightarrow{S} Z \quad X \xrightarrow{Ros} Z$

Note

$$R \circ (S \circ P) = (R \circ S) \circ P$$

$$R \circ (S \circ P) = (R \circ S) \circ P = Ros \circ P$$

II Find $S \circ P$ and Ros , $R = \{ \langle 1,1 \rangle, \langle 1,3 \rangle, \langle 2,1 \rangle, \langle 2,2 \rangle, \langle 2,3 \rangle, \langle 3,2 \rangle \}$ and

$$S = \{ \langle 1,1 \rangle, \langle 2,2 \rangle, \langle 2,3 \rangle, \langle 3,1 \rangle, \langle 3,3 \rangle \}$$

SOL Given

$$R = \{ \langle 1,1 \rangle, \langle 1,3 \rangle, \langle 2,1 \rangle, \langle 2,2 \rangle, \langle 2,3 \rangle, \langle 3,2 \rangle \}$$

$$S = \{ \langle 1,1 \rangle, \langle 2,2 \rangle, \langle 2,3 \rangle, \langle 3,1 \rangle, \langle 3,3 \rangle \}$$

(\because repetition is neglect)

ROS

$\langle 1,1 \rangle \in R, \langle 1,1 \rangle \in S \Rightarrow \langle 1,1 \rangle \in ROS$
 $\langle 1,3 \rangle \in R, \langle 3,1 \rangle \in S \Rightarrow \langle 1,1 \rangle \in ROS$
 $\langle 3,3 \rangle \in S \Rightarrow \langle 1,3 \rangle \in ROS$

$\langle 2,1 \rangle \in R, \langle 1,1 \rangle \in S \Rightarrow \langle 2,1 \rangle \in ROS$
 $\langle 2,2 \rangle \in R, \langle 2,2 \rangle \in S \Rightarrow \langle 2,2 \rangle \in ROS$
 $\langle 2,3 \rangle \in S \Rightarrow \langle 2,3 \rangle \in ROS$
 $\langle 2,3 \rangle \in R, \langle 3,1 \rangle \in S \Rightarrow \langle 2,1 \rangle \in ROS$
 $\langle 3,3 \rangle \in S \Rightarrow \langle 2,3 \rangle \in ROS$
 $\langle 3,2 \rangle \in R, \langle 2,2 \rangle \in S \Rightarrow \langle 3,2 \rangle \in ROS$
 $\langle 2,3 \rangle \in S \Rightarrow \langle 3,3 \rangle \in ROS$

$ROS = \{\langle 1,1 \rangle, \langle 1,3 \rangle, \langle 2,1 \rangle, \langle 2,2 \rangle, \langle 2,3 \rangle, \langle 3,2 \rangle, \langle 3,3 \rangle\}$

Find

$SOR = \{\langle 1,1 \rangle \in S, \langle 1,1 \rangle \in R \Rightarrow \langle 1,1 \rangle \in SOR$
 $\langle 1,3 \rangle \in R \Rightarrow \langle 1,3 \rangle \in SOR$
 $\langle 2,2 \rangle \in S, \langle 2,1 \rangle \in R \Rightarrow \langle 2,1 \rangle \in SOR$
 $\langle 2,2 \rangle \in R \Rightarrow \langle 2,2 \rangle \in SOR$
 $\langle 2,3 \rangle \in R \Rightarrow \langle 2,3 \rangle \in SOR$
 $\langle 2,3 \rangle \in S, \langle 3,2 \rangle \in R \Rightarrow \langle 2,2 \rangle \in SOR$
 $\langle 3,1 \rangle \in S, \langle 1,1 \rangle \in R, \langle 1,3 \rangle \in R \Rightarrow \langle 3,1 \rangle \in SOR$
 $\langle 3,3 \rangle \in SOR$
 $\langle 3,3 \rangle \in S, \langle 3,2 \rangle \in R \Rightarrow \langle 3,2 \rangle \in SOR$

$SOR = \{\langle 1,1 \rangle, \langle 1,3 \rangle, \langle 2,1 \rangle, \langle 2,2 \rangle, \langle 2,3 \rangle, \langle 3,1 \rangle,$
 $\langle 3,3 \rangle, \langle 3,2 \rangle\}$

2. Let $R = \{\langle 4,2 \rangle, \langle 3,4 \rangle, \langle 2,2 \rangle\} \& S = \{\langle 4,1 \rangle, \langle 4,5 \rangle,$
 $\langle 3,1 \rangle, \langle 1,3 \rangle\}$ find ROS, SOR, $R_0(SOR)$, $(ROS)OR$,
 ROR , SOS , R_0ROR .

SOL Given $R = \{\langle 1,2 \rangle, \langle 3,4 \rangle, \langle 2,2 \rangle\}$

$$S = \{\langle 4,2 \rangle, \langle 2,5 \rangle, \langle 3,1 \rangle, \langle 1,3 \rangle\}$$

$$Ros = \{\langle 1,5 \rangle, \langle 3,2 \rangle, \langle 2,5 \rangle\}$$

$$SOR = \{\langle 4,2 \rangle, \langle 3,2 \rangle, \langle 1,4 \rangle\}$$

$$R_0(SOR) = \{\langle 3,2 \rangle\}$$

$$(Ros)OR = \{\langle 3,2 \rangle\}$$

$$ROR = \{\langle 1,2 \rangle, \langle 2,2 \rangle\}$$

$$SOS = \{\langle 4,5 \rangle, \langle 3,3 \rangle, \langle 1,1 \rangle\}$$

$$(ROR)OR = \{\langle 1,2 \rangle, \langle 2,2 \rangle\} \text{ or } \{\langle 1,2 \rangle, \langle 3,4 \rangle, \langle 4,1 \rangle\}$$
$$\Rightarrow \{\langle 1,2 \rangle, \langle 2,2 \rangle\}$$

3. Let R & S be two relations on a set of positive integers "I" are

$$R = \{\langle x, 2x \rangle : x \in I\}, S = \{\langle x, 7x \rangle : x \in I\}$$

Find Ros , SOR , ROR , $SOROR$.

SOL Given $R = \{\langle x, 2x \rangle : x \in I\}$

$$S = \{\langle x, 7x \rangle : x \in I\} \quad I = 1, 2, \dots$$

$$R = \{\langle 1,2 \rangle, \langle 2,4 \rangle, \langle 3,6 \rangle, \langle 4,8 \rangle, \langle 5,10 \rangle, \dots\}$$

$$S = \{\langle 1,7 \rangle, \langle 2,14 \rangle, \langle 3,21 \rangle, \langle 4,28 \rangle, \langle 5,35 \rangle, \langle 6,42 \rangle, \dots\}$$

Ros

If $\langle 1,2 \rangle \in R, \langle 2,14 \rangle \in S \Rightarrow \langle 1,14 \rangle \in Ros$

If $\langle 2,4 \rangle \in R, \langle 4,28 \rangle \in S \Rightarrow \langle 2,28 \rangle \in Ros$

If $\langle 3,6 \rangle \in R, \langle 6,42 \rangle \in S \Rightarrow \langle 3,42 \rangle \in Ros$

$$Ros = \{\langle 1,4 \rangle, \langle 2,8 \rangle, \langle 3,12 \rangle, \dots\}$$

$$Ros = \{\langle x, 14x \rangle / x \in I\}$$

SOR: $\langle 1,7 \rangle \in S, \langle 7,14 \rangle \in R \Rightarrow \langle 1,14 \rangle \in SOR$

If $\langle 2,14 \rangle \in S, \langle 14,28 \rangle \in R \Rightarrow \langle 2,28 \rangle \in SOR$

$$SOR = \{ \langle 1,14 \rangle, \langle 2,28 \rangle, \dots \}$$

$$SOR = \{ \langle x, 14x \rangle : x \in I \}$$

RoR

If $\langle 1,2 \rangle \in R, \langle 2,4 \rangle \in R \Rightarrow \langle 1,4 \rangle \in RoR$

If $\langle 2,4 \rangle \in R, \langle 4,8 \rangle \in R \Rightarrow \langle 2,8 \rangle \in RoR$

If $\langle 3,6 \rangle \in R, \langle 6,12 \rangle \in R \Rightarrow \langle 3,12 \rangle \in RoR$.

$$RoR = \{ \langle 1,4 \rangle, \langle 2,8 \rangle, \langle 3,12 \rangle, \dots \}$$

$$= \{ \langle x, 4x \rangle : x \in I \}$$

$$(RoR)OR = \{ \langle 1,4 \rangle, \langle 2,8 \rangle, \langle 3,12 \rangle, \dots \} \cup$$

$$\{ \langle 1,2 \rangle, \langle 2,4 \rangle, \langle 3,6 \rangle, \dots \}$$

$$= \{ \langle 1,8 \rangle, \langle 2,16 \rangle, \langle 3,24 \rangle, \dots \}$$

$$RoRoR = \{ \langle x, 8x \rangle : x \in I \}$$

$$(SOR)OR = \{ \langle 1,14 \rangle, \langle 2,28 \rangle, \dots \} \cup \{ \langle 1,2 \rangle, \langle 2,4 \rangle,$$

$$\langle 3,6 \rangle, \dots \}$$

$$= \{ \langle 1,28 \rangle, \langle 2,56 \rangle, \dots \}$$

$$= \{ \langle x, 28x \rangle : x \in I \}$$

* Representation of discrete structures!

i. Relation Matrix!

A relation 'R' from a finite set 'X' to a finite set 'Y' can also be represented by a matrix if 'R' is called the relation matrix of 'R'.

* $X = \{ x_1, x_2, x_3, \dots, x_m \} - m \text{ rows}$

$Y = \{ y_1, y_2, y_3, \dots, y_n \} - n \text{ columns}$

* R being a relation from 'X' to 'Y'

* The relation matrix 'R' can be obtained by first constructing a table whose columns are preceded by a column consisting of successive elements of 'x' and whose rows are headed by a row consisting of the successive elements of 'y'.

* If $x_i R y_j$, then we enter a '1' in the i^{th} row and j^{th} column. If $x_i \not R y_k$ then we enter a '0' in the k^{th} row and i^{th} column.

Ex:- consider $m=3$, $n=2$ and 'R' is given by
 $R = \{(x_1, y_1), (x_2, y_1), (x_3, y_3), (x_2, y_2)\}$.
the required columns is

	y_1	y_2
x_1	1	0
x_2	1	1
x_3	0	1

The relation matrix can be constructed in the following matrix.

$$Y_{ij} = \begin{cases} 1 & \text{if } x_i R y_j \\ 0 & \text{if } x_i \not R y_j \end{cases}$$

where ' Y_{ij} ' is the element in the i^{th} row and j^{th} column.

* If 'x' has 'm' elements and 'y' has 'n' elements then the relation matrix is an $m \times n$ matrix. For the relation R given in eq(1) the relation matrix is

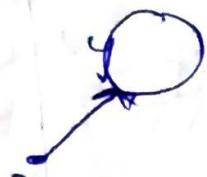
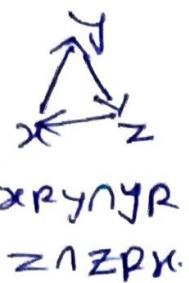
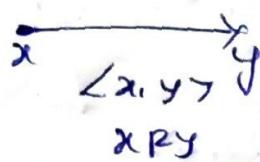
$$M_R = \begin{bmatrix} 1 & 0 \\ 1 & 1 \\ 0 & 1 \end{bmatrix}$$

Graph of a relation (Di Graph):

From the Graph of relation it is possible to observe some of its properties. The relation is reflexive, then their must be

a loop at each node (vertex).

* If the relation is irreflexive, then there is no loop at any node.



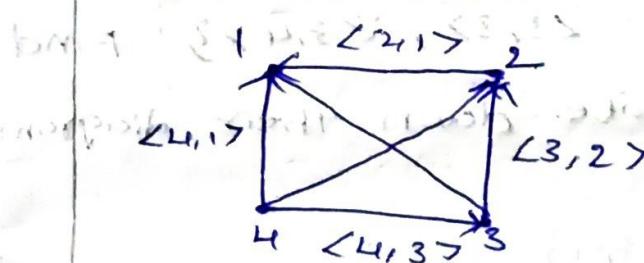
xRy and yRx

1. Let $X = \{1, 2, 3, 4\}$ and $R = \{\langle x, y \rangle ; x > y\}$
draw graph of 'R' and also give its matrix.

Sol Given $X = \{1, 2, 3, 4\}$

$$R = \{\langle 4, 1 \rangle, \langle 4, 2 \rangle, \langle 4, 3 \rangle, \langle 3, 1 \rangle, \langle 3, 2 \rangle, \langle 2, 1 \rangle\}$$

Graph of Relations



	1	2	3	4
1	0	0	0	0
2	1	0	0	0
3	1	1	0	0
4	1	1	1	0

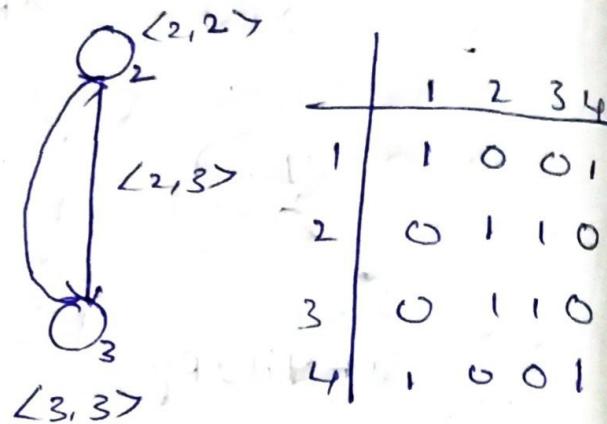
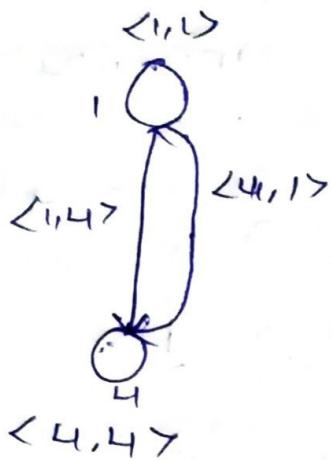
matrix $MR = \begin{bmatrix} 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 \end{bmatrix}$

2. Let $X = \{1, 2, 3, 4\}$ & $R = \{\langle 1, 1 \rangle, \langle 1, 4 \rangle, \langle 4, 1 \rangle, \langle 4, 4 \rangle, \langle 2, 2 \rangle, \langle 2, 3 \rangle, \langle 3, 2 \rangle, \langle 3, 3 \rangle\}$
write matrix of 'R' and sketch its graph.

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Given $X = \{1, 2, 3, 4\}$

$$R = \{\langle 1, 1 \rangle, \langle 1, 4 \rangle, \langle 4, 1 \rangle, \langle 4, 4 \rangle, \langle 2, 2 \rangle, \\ \langle 2, 3 \rangle, \langle 3, 2 \rangle, \langle 3, 3 \rangle\}$$



Matrix $MR = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 \end{bmatrix}$

3. If $A = \{1, 2, 3, 4\}$ and R is a relation on A defined by $R = \{\langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 2, 4 \rangle, \langle 3, 2 \rangle, \langle 3, 3 \rangle, \langle 3, 4 \rangle\}$ find

R^2 and R^3 . write down their diagrams.
OR

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Given $A = \{1, 2, 3, 4\}$

$$R = \{\langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 2, 4 \rangle, \langle 3, 2 \rangle, \langle 3, 3 \rangle, \\ \langle 3, 4 \rangle\}$$

$$R^2 = R \circ R$$

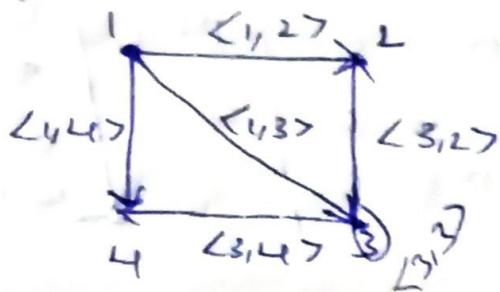
$$= \{\langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 2, 4 \rangle, \langle 3, 2 \rangle, \langle 3, 3 \rangle, \langle 3, 4 \rangle\}$$

$$\{ \langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 2, 4 \rangle, \langle 3, 2 \rangle, \langle 3, 3 \rangle, \langle 3, 4 \rangle \}$$

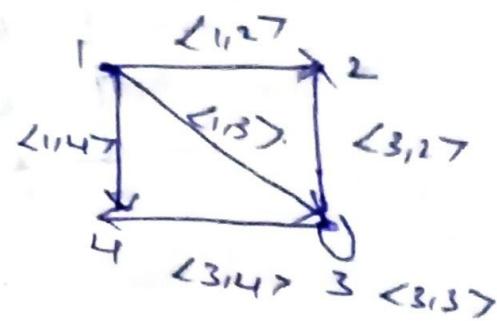
$$\geq \{\langle 1, 4 \rangle, \langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 1, 3 \rangle, \langle 3, 4 \rangle, \\ \langle 3, 2 \rangle\}$$

$$R^3 = \{\langle 1, 4 \rangle, \langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 3, 4 \rangle, \langle 3, 2 \rangle, \\ \langle 3, 3 \rangle\}$$

R^2 diagram



R^3



4. If $A = \{1, 2, 3, 4\}$ & R, S are relations on A defined by $R = \{\langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 2, 4 \rangle, \langle 4, 4 \rangle\}$ & $S = \{\langle 1, 1 \rangle, \langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 1, 4 \rangle, \langle 2, 3 \rangle, \langle 2, 4 \rangle\}$ find Ros , Sor , R^2 , S^2 , write down their matrices.

Sol

Given $A = \{1, 2, 3, 4\}$

$$R = \{\langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 2, 4 \rangle, \langle 4, 4 \rangle\}$$

$$S = \{\langle 1, 1 \rangle, \langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 1, 4 \rangle, \langle 2, 3 \rangle, \langle 2, 4 \rangle\}$$

$$Ros = \{\langle 1, 3 \rangle, \langle 1, 4 \rangle\}$$

$$Sor = \{\langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 1, 4 \rangle, \langle 2, 4 \rangle\}$$

$$R^2 = ROR = \{\langle 1, 4 \rangle, \langle 2, 4 \rangle, \langle 4, 4 \rangle\}$$

$$S^2 = SOS = \{\langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 1, 4 \rangle\}$$

$$R^2 = \begin{array}{c|cccc} & 1 & 2 & 3 & 4 \\ \hline 1 & 0 & 1 & 0 & 1 \\ 2 & 0 & 0 & 0 & 1 \\ 3 & 0 & 0 & 0 & 0 \\ 4 & 0 & 0 & 0 & 1 \end{array}$$

$$S^2 = \begin{array}{c|cccc} & 1 & 2 & 3 & 4 \\ \hline 1 & 0 & 1 & 1 & 1 \\ 2 & 0 & 0 & 0 & 0 \\ 3 & 0 & 0 & 0 & 0 \\ 4 & 0 & 0 & 0 & 0 \end{array}$$

matrices of above composite relations.

$$Ros = \begin{array}{c|cccc} & 1 & 2 & 3 & 4 \\ \hline 1 & 0 & 0 & 1 & 1 \\ 2 & 0 & 0 & 0 & 0 \\ 3 & 0 & 0 & 0 & 0 \\ 4 & 0 & 0 & 0 & 0 \end{array}$$

$$Sor = \begin{array}{c|cccc} & 1 & 2 & 3 & 4 \\ \hline 1 & 0 & 1 & 1 & 1 \\ 2 & 0 & 0 & 0 & 1 \\ 3 & 0 & 0 & 0 & 0 \\ 4 & 0 & 0 & 0 & 0 \end{array}$$

$$M_{ROS} = \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} \quad M_{SOF} = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$M_{P^2} = \begin{bmatrix} 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix} \quad M_{S^2} = \begin{bmatrix} 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

complementation of a Relation

Given a relation R from a set A to set B , the complement of R is denoted by \bar{R} is defined as a relation from A to B with the property ordered pair $\langle a, b \rangle \in \bar{R}$ iff $\langle a, b \rangle \notin R$.

* In other words \bar{R} is complement of the set R in the universal set $A \times B$.

$$\bar{R} = (A \times B) - R$$

Converse of a Relation

Given a relation R from a set A to set B , the converse of R is denoted by R^C is defined as a relation from B to A with the property ordered pair $\langle a, b \rangle \in R^C$ iff $\langle a, b \rangle \in R$.

* If "matrix R is then (MR) ", The Transp

MR is the matrix of R^C and

$$*(R^C)^C = R$$

1. consider set $A = \{a, b, c\}$ and $B = \{1, 2, 3\}$ and the relations $R = \{\langle a, 1 \rangle, \langle a, 2 \rangle, \langle b, 1 \rangle, \langle b, 2 \rangle, \langle c, 3 \rangle\}$ from A to B . Determine \bar{R} , \bar{S} , R_{US} , R_{NS} , R^C , S^C .

Sol

$$\text{Given } A = \{a, b, c\}$$

$$B = \{1, 2, 3\}$$

$$R = \{\langle a, 1 \rangle, \langle a, 2 \rangle, \langle b, 1 \rangle, \langle b, 2 \rangle, \langle c, 3 \rangle\}$$

$$S = \{\langle a, 1 \rangle, \langle a, 2 \rangle, \langle b, 1 \rangle, \langle b, 2 \rangle, \langle c, 3 \rangle\}$$

$$A \times B = \{a, b, c\} \times \{1, 2, 3\}$$

$$= \{\langle a, 1 \rangle, \langle a, 2 \rangle, \langle a, 3 \rangle, \langle b, 1 \rangle, \langle b, 2 \rangle, \langle b, 3 \rangle, \langle c, 1 \rangle, \langle c, 2 \rangle, \langle c, 3 \rangle\}$$

\bar{R} = complement of R in $A \times B$

$$= (A \times B) - R$$

$$= \{\langle a, 1 \rangle, \langle a, 2 \rangle, \langle a, 3 \rangle, \langle b, 1 \rangle, \langle b, 2 \rangle, \langle b, 3 \rangle, \langle c, 1 \rangle, \langle c, 2 \rangle, \langle c, 3 \rangle\} - \{\langle a, 1 \rangle, \langle a, 2 \rangle, \langle b, 1 \rangle, \langle b, 2 \rangle, \langle c, 3 \rangle\}$$

$$= \{\langle a, 3 \rangle, \langle b, 3 \rangle, \langle c, 1 \rangle, \langle c, 2 \rangle, \langle c, 3 \rangle\}$$

\bar{S} = complement of S in $A \times B$

$$= (A \times B) - S$$

$$= \{\langle a, 1 \rangle, \langle a, 2 \rangle, \langle a, 3 \rangle, \langle b, 1 \rangle, \langle b, 2 \rangle, \langle b, 3 \rangle, \langle c, 1 \rangle, \langle c, 2 \rangle, \langle c, 3 \rangle\} - \{\langle a, 1 \rangle, \langle a, 2 \rangle, \langle b, 1 \rangle, \langle b, 2 \rangle, \langle c, 3 \rangle\}$$

$$= \{\langle a, 3 \rangle, \langle b, 3 \rangle, \langle c, 1 \rangle, \langle c, 2 \rangle, \langle c, 3 \rangle\}$$

$$= \{\langle a, 3 \rangle, \langle b, 3 \rangle, \langle c, 1 \rangle, \langle c, 2 \rangle, \langle c, 3 \rangle\}$$

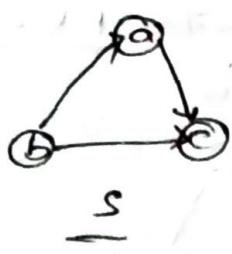
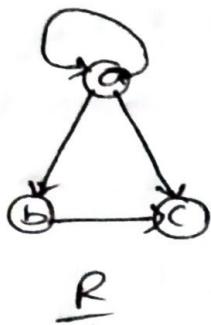
$$RUS = \{ \langle a, 1 \rangle, \langle a, 2 \rangle, \langle b, 1 \rangle, \langle b, 2 \rangle, \langle c, 2 \rangle, \langle c, 3 \rangle \}$$

$$RNS = \{ \langle a, 1 \rangle, \langle b, 1 \rangle \}$$

$$R^C = \{ \langle 1, a \rangle, \langle 1, b \rangle, \langle 2, c \rangle, \langle 3, c \rangle \}$$

$$S^C = \{ \langle 1, a \rangle, \langle 2, a \rangle, \langle 1, b \rangle, \langle 2, b \rangle \}$$

2. The digraphs of two relations R and S on the set $A = \{a, b, c\}$ are given below. Draw the digraphs of \bar{R} , RUS , RNS and R^C .



Sol

$$R = \{ \langle a, a \rangle, \langle a, b \rangle, \langle a, c \rangle, \langle b, c \rangle \}$$

$$S = \{ \langle a, c \rangle, \langle b, a \rangle, \langle b, c \rangle \}$$

$$A \times A = \{a, b, c\} \times \{a, b, c\}$$

$$= \{ \langle a, a \rangle, \langle a, b \rangle, \langle a, c \rangle, \langle b, a \rangle, \langle b, b \rangle, \langle b, c \rangle, \langle c, a \rangle, \langle c, b \rangle, \langle c, c \rangle \}$$

\bar{R} = complement of R in $A \times A$

$$= A \times A - R$$

$$= \{ \langle a, a \rangle, \langle a, b \rangle, \langle a, c \rangle, \langle b, a \rangle, \langle b, b \rangle, \langle b, c \rangle, \langle c, a \rangle, \langle c, b \rangle, \langle c, c \rangle \} - \{ \langle a, a \rangle, \langle a, b \rangle, \langle a, c \rangle, \langle b, c \rangle \}$$

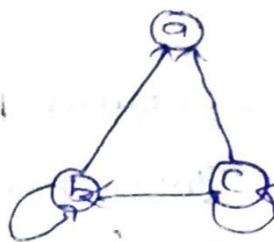
$$= \{ \langle b, a \rangle, \langle b, b \rangle, \langle c, a \rangle, \langle c, b \rangle, \langle c, c \rangle \}$$

$$RUS = \{ \langle a, a \rangle, \langle a, b \rangle, \langle a, c \rangle, \langle b, a \rangle, \langle b, c \rangle \}$$

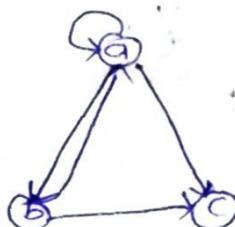
$$RNS = \{ \langle a, c \rangle, \langle b, c \rangle \}$$

$$R^c = \{ \langle a, a \rangle, \langle b, a \rangle, \langle c, a \rangle, \langle c, b \rangle \}$$

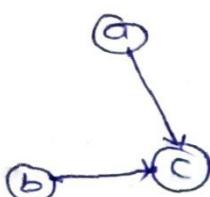
The diagraph of \bar{R}



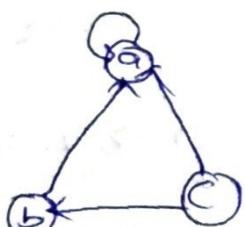
The diagraph of $R \cup S$



The diagraph of $R \cap S$



The diagraph of R^c



covering of a set

Let 'S' be a given set and $A = \{A_1, A_2, A_3, \dots, A_m\}$ where each A_i is subset of S and union $\bigcup A_i = S$.

Then set A is called a covering of set S , and the subsets $A_1, A_2, A_3, \dots, A_m$ are said to cover set S .

Ex:-

Let $S = \{a, b, c\}$ and consider the following collection of subsets of 'S' $A = \{\{a, b\}, \{b, c\}\}$,

$B = \{\{a\}, \{a, c\}, \{a\} \cup \{c\}, \{b, c\}\}$, $D = \{\{a, b, c\} \in \{\{a\}, \{b\}, \{c\}\}\}$, $F = \{\{a\}, \{a, b\}, \{a, c\}\}$.

Among them which which subsets are covering sets?

Sol

Except 'B' all subsets are covering sets of 'S'.

Partition of a set

Let 'S' be a non-empty set, $S_1, S_2, S_3, \dots, S_m$ are the subsets of 'S', the collection of subsets S_i is a partition of 'S' iff and only if

- i. $S_i \neq \emptyset$, for each i
- ii. $S_i \cap S_j = \emptyset$ for $i \neq j$
- iii. $\bigcup_{i=1}^m S_i = S$.

Where $\bigcup_{i=1}^m$ represents the union of the subsets

S_i , for all i and $S_1, S_2, S_3, \dots, S_m$ are called blocks or cells of partitions.

Ex

* Let $S = \{1, 2, 3, 4, 5, 6\}$ and the subsets $S_1 = \{1, 3, 5\}$, $S_2 = \{2, 4\}$, $S_3 = \{1, 3\}$, $S_4 = \{3, 5\}$, $S_5 = \{2, 5, 6\}$, $S_6 = \{1, 2, 3\}$, $S_7 = \{4, 5\}$, $S_8 = \{6\}$ check whether the following partition are valid or not?

$$\Pi_1 = \{S_6, S_7, S_8\}.$$

Sol

$$S_6 \neq \emptyset, S_7 \neq \emptyset, S_8 \neq \emptyset$$

The first condition is satisfied.

$$\begin{aligned} \{S_6, S_7\} &= \{S_6 \cap S_7\} \\ &= \{1, 2, 3\} \cap \{4, 5\} \end{aligned}$$

$$= \emptyset$$

$$\{S_7, S_8\} = \{S_7\} \cap \{S_8\}$$

$$= \emptyset$$

$$\{S_6, S_8\} = \{S_6\} \cap \{S_8\}$$

$$= \emptyset$$

The second condition is satisfied.

$$\text{iii. } \{S_6, S_7, S_8\} = \{S_6\} \cup \{S_7\} \cup \{S_8\}$$

$$= \{1, 2, 3\} \cup \{4, 5\} \cup \{6\}$$

$$= \{1, 2, 3, 4, 5, 6\}$$

$$= S$$

\therefore The third condition is satisfied.

$\Pi_1 = \{S_6, S_7, S_8\}$ is a valid Partition.

Partial ordering

* A binary relation R in a set P is called
a partial order relation. (or) A partial ordering
in $P \Leftrightarrow R$ is reflexive, Anti symmetric and
transitive.

* We denote a Partial ordering by the
symbol " \leq ".

* If " \leq " is a partial ordering on P , then the
order pair $\langle P, \leq \rangle$ is called Partial ordering
Set or "Poset".

Totally ordered set

Let $\langle P, \leq \rangle$ be a partially ordered set. If
for every $x, y \in P$ we have either $x \leq y$ or $y \leq x$

Then \leq is called simple ordering or linear ordering on P , and $\langle P, \leq \rangle$ is called a totally ordered set or simply ordered or a chain.

Compatibility Relation

* A relation R in X is said to be a compatibility relation if it is reflexive and symmetric.

* The relation R is given by

$$R = \{ \langle x, y \rangle \mid x, y \in X \wedge R(x, y) \text{, if } x \text{ & } y$$

contain some common letter }

Hasse Diagram

* A partial ordering \leq on a set P can be represented by means of a diagram is known as Hasse diagram of $\langle P, \leq \rangle$.

* In Hasse diagram each element is represented by small circle or dots.

* In Hasse diagram we represent their vertices by dots or circles, we does not put arrow on edges and we does not draw selfloops at vertices.

* In diagram of partial order, there is an edge from vertex A to vertex B and there is an edge from vertex B to vertex C, There should be an edge from vertex A to vertex C as such we need not exhibit an edge from A to C explicitly.

b) Let $X = \{2, 3, 6, 12, 24, 36\}$ and the relation \leq be such that $x \leq y$. If x divides y . Draw the Hasse diagram of $\langle X, \leq \rangle$.

Sol Given $X = \{2, 3, 6, 12, 24, 36\}$

$2 \rightarrow 2, 6, 12, 24, 36$

ordered pairs: $\langle 2, 2 \rangle, \langle 2, 6 \rangle, \langle 2, 12 \rangle, \langle 2, 24 \rangle, \langle 2, 36 \rangle$

$3 \rightarrow 3, 6, 12, 24, 36$

ordered pairs: $\langle 3, 3 \rangle, \langle 3, 6 \rangle, \langle 3, 12 \rangle, \langle 3, 24 \rangle, \langle 3, 36 \rangle$

$6 \rightarrow 6, 12, 24, 36$

ordered pairs: $\langle 6, 6 \rangle, \langle 6, 12 \rangle, \langle 6, 24 \rangle, \langle 6, 36 \rangle$

$12 \rightarrow 12, 24, 36$

ordered pairs: $\langle 12, 12 \rangle, \langle 12, 24 \rangle, \langle 12, 36 \rangle$

$24 \rightarrow \langle 24, 24 \rangle$

$36 \rightarrow \langle 36, 36 \rangle$

$$R = \{\langle 2, 2 \rangle, \langle 2, 6 \rangle, \langle 2, 12 \rangle, \langle 2, 24 \rangle, \langle 2, 36 \rangle, \langle 3, 3 \rangle, \langle 3, 6 \rangle, \langle 3, 12 \rangle, \langle 3, 24 \rangle, \langle 3, 36 \rangle, \langle 6, 6 \rangle, \langle 6, 12 \rangle, \langle 6, 24 \rangle, \langle 6, 36 \rangle, \langle 12, 12 \rangle, \langle 12, 24 \rangle, \langle 12, 36 \rangle, \langle 24, 24 \rangle, \langle 36, 36 \rangle\}$$

Reflexive: For all $a \in X$, $\langle a, a \rangle \in R$

R is a Reflexive

Antisymmetric: For all $a, b \in X$, if $\langle a, b \rangle \in R$ & $\langle b, a \rangle \in R$ then $a = b$.

R is anti symmetric

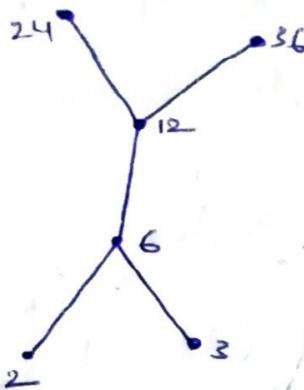
Transitive:

For all $a, b, c \in X$, if $\langle a, b \rangle \in R$, $\langle b, c \rangle \in R$ and $\langle a, c \rangle \in R$

$\therefore R$ is a Partially ordered relation.

Hasse Diagram

$$X = \{2, 3, 6, 12, 24, 36\}$$



2. Let $A = \{1, 2, 3, 4, 6, 12\}$ on A, defined a relation R by aRb if and only if a divides b. prove that R is a partial order on A. Draw the hasse diagram for this relation.

Sol Let $A = \{1, 2, 3, 4, 6, 12\}$

$$R = \{(a, b) / a, b \in A \text{ and } a \text{ divides } b\}$$

$$R = \{(1, 1), (1, 2), (1, 3), (1, 4), (1, 6), (1, 12), (2, 2), (2, 4), (2, 6), (2, 12), (3, 3), (3, 6), (3, 12), (4, 4), (4, 12), (6, 6), (6, 12), (12, 12)\}$$

Reflexive: For all $a \in A$, $(a, a) \in R$

$\therefore R$ is reflexive.

Anti symmetric:

For all $a, b \in A$, if a divides b and b divides a then $a = b$

$\therefore R$ is antisymmetric.

Transitive:

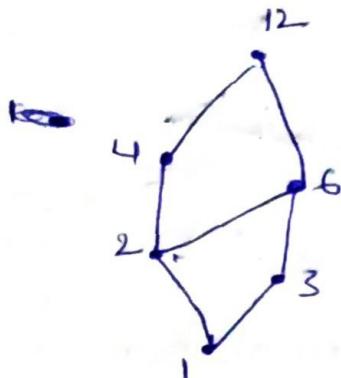
For all $a, b, c \in R$ if $(a, b) \in R, (b, c) \in R$ then $(a, c) \in R$

$\therefore R$ is transitive

$\therefore R$ is poset.

Hasse Diagram

$$A = \{1, 2, 3, 4, 6, 12\}$$



3. Draw the Hasse diagram for $\{1, 2, 3, 5, 6, 10, 15, 30\}$
- []

Sol Given $A = \{1, 2, 3, 5, 6, 10, 15, 30\}$

$R = \{(a, b) / a, b \in A \text{ and } a \text{ divides } b\}$

$$\begin{aligned} R = & \{(1, 1), (1, 2), (1, 3), (1, 5), (1, 6), (1, 10), (1, 15), (1, 30), \\ & (2, 2), (2, 6), (2, 10), (2, 30), (3, 3), (3, 6), (3, 15), \\ & (3, 30), (5, 5), (5, 10), (5, 15), (5, 30), (6, 6), (6, 30), \\ & (10, 10), (10, 30), (15, 15), (15, 30), (30, 30)\} \end{aligned}$$

Reflexive : For all $a \in A$, $(a, a) \in R$.

$\therefore R$ is reflexive.

Anti Symmetric

For all $a, b \in A$, if a divides b and b divides a then $a = b$.

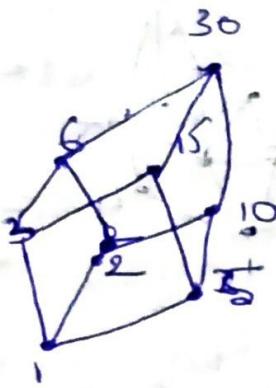
$\therefore R$ is antisymmetric

Transitive | For all $a, b, c \in A$ if $(a, b) \in R$ and $(b, c) \in R$ then $(a, c) \in R$. $\rightarrow R$ is transitive

R is a Poset.

Hasse Diagram

$$A = \{1, 2, 3, 5, 6, 10, 15, 30\}$$



4. Draw the Hasse diagram representing the positive divisors of 36

Sol The set of all positive divisors of 36 is

$$D_{36} = \{1, 2, 3, 4, 6, 9, 12, 18, 36\}$$

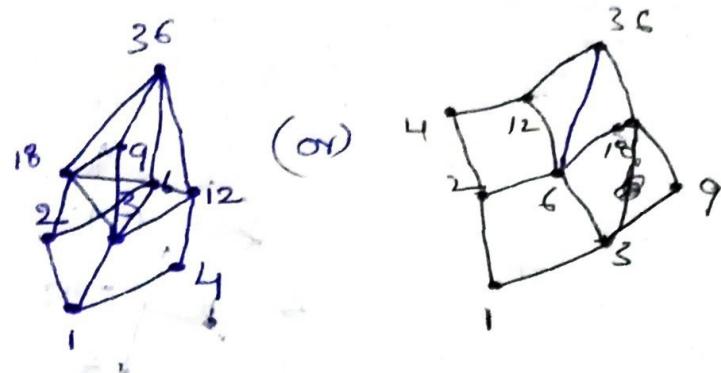
$$R = \{(a, b) / a, b \in A \text{ and } a \text{ divides } b\}$$

$$R = \{\langle 1, 1 \rangle, \langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 1, 4 \rangle, \langle 1, 6 \rangle, \langle 1, 9 \rangle, \langle 1, 12 \rangle, \langle 1, 18 \rangle, \langle 1, 36 \rangle, \langle 2, 2 \rangle, \langle 2, 4 \rangle, \langle 2, 6 \rangle, \langle 2, 12 \rangle, \langle 2, 18 \rangle, \langle 2, 36 \rangle, \langle 3, 3 \rangle, \langle 3, 6 \rangle, \langle 3, 9 \rangle, \langle 3, 12 \rangle, \langle 3, 18 \rangle, \langle 3, 36 \rangle, \langle 4, 4 \rangle, \langle 4, 12 \rangle, \langle 4, 36 \rangle, \langle 6, 6 \rangle, \langle 6, 12 \rangle, \langle 6, 18 \rangle, \langle 6, 36 \rangle, \langle 9, 9 \rangle, \langle 9, 18 \rangle, \langle 9, 36 \rangle, \langle 12, 12 \rangle, \langle 12, 36 \rangle, \langle 18, 18 \rangle, \langle 18, 36 \rangle, \langle 36, 36 \rangle\}$$

$\because R$ is reflexive, antisymmetric, transitive then
 R is a Poset.

Hasse Diagram

$$D_{36} = \{1, 2, 3, 4, 6, 9, 12, 18, 36\}$$



5. Let A be a given finite set and $P(A)$ its power set. Let \subseteq be the inclusion relation on the elements of $P(A)$. Draw the Hasse diagram of $\langle P(A), \subseteq \rangle$ for i. $A = \{a\}$ ii. $A = \{a, b\}$, iii. $A = \{a, b, c\}$

SOL: Given $A = \{a\}$

$$P(A) = \{\emptyset, \{a\}\}$$

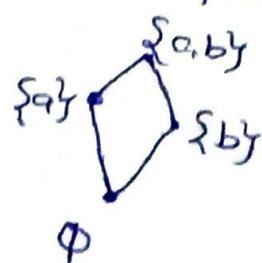
Hasse Diagram for $\langle P(A), \subseteq \rangle$



iii. Given: $A = \{a, b\}$

$$P(A) = \{\emptyset, \{a\}, \{b\}, \{a, b\}\}$$

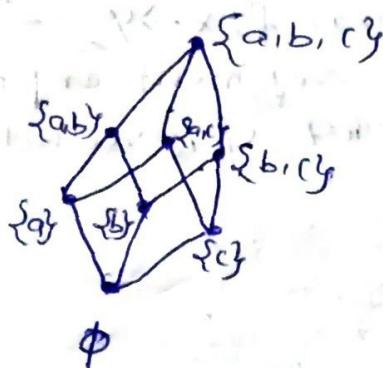
Hasse Diagram for $\langle P(A), \subseteq \rangle$



iii. Given $A = \{a, b, c\}$

$$P(A) = \{\emptyset, \{a\}, \{b\}, \{c\}, \{a, b\}, \{a, c\}, \{b, c\}, \{a, b, c\}\}$$

Hasse diagram for $\langle P(A), \subseteq \rangle$



6. Draw the Hasse diagram for $\{1, 2, 3, 4, 6, 9\}$.

So Let $A = \{1, 2, 3, 4, 6, 9\}$

$R = \{(a, b) / a, b \in A \text{ and } a \text{ divides } b\}$

$$R = \{(1, 1), (1, 2), (1, 3), (1, 4), (1, 6), (1, 9),$$

Functions

* Let x and y be any two sets. A Relation f from $x \rightarrow y$ ($f: x \rightarrow y$) is called function.

If for every $x \in x$, there is a unique $y \in y$ such that order pair $(x, y) \in f$. It is denoted as $f: x \rightarrow y$.

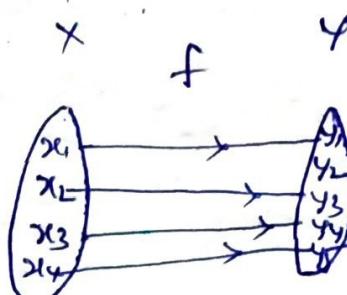
* For a function $f: x \rightarrow y$, $(x, y) \in f$, then x is called an argument or preimage of y and the corresponding y is called image of x under f .

Types of functions

1. One to One function (or) Injective or 1-1

A function $f: x \rightarrow y$ is said to be a one to one function. If different elements of x have different images in y under f .

* That is if whenever $x_1, x_2 \in x$ with $x_1 \neq x_2$ then $f(x_1) \neq f(x_2)$ (or) equivalently if whenever $f(x_1) = f(x_2)$ for $x_1, x_2 \in x$ then $x_1 = x_2$



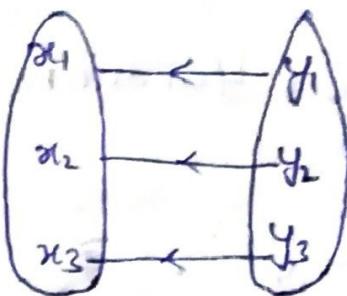
* If $f: x \rightarrow y$ is a one to one function, then Every element of x has a unique image in y and Every element of $f(x)=y$ has a unique preimage in x .

2. on to function (or) surjection

A function $f: x \rightarrow y$ is said to be onto function if every element of y has a preimage in x under f .

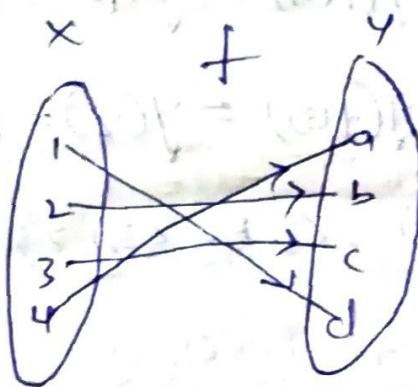
* In other words f is an onto function from X to Y if the range of f is equal to Y ($f(X) = Y$).
i.e. ($y = f(x)$).

$X \rightarrow Y$



3. one to one correspondence (or) Bijection (or) one to one onto.

* A function which is both one to one and onto is called a one to one correspondence. If $f: X \rightarrow Y$ is such a function, then every element of X has a unique image in Y and every element of Y has a unique preimage in X .

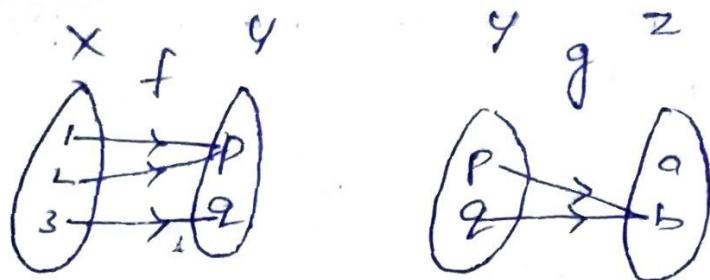


1. Let $X = \{1, 2, 3\}$, $Y = \{P, Q\}$, $Z = \{a, b\}$. Also let $f: X \rightarrow Y$ be $f = \{\langle 1, P \rangle, \langle 2, P \rangle, \langle 3, Q \rangle\}$ and $g: Y \rightarrow Z$ be given by $g = \{\langle P, a \rangle, \langle Q, b \rangle\}$. Find $g \circ f$.

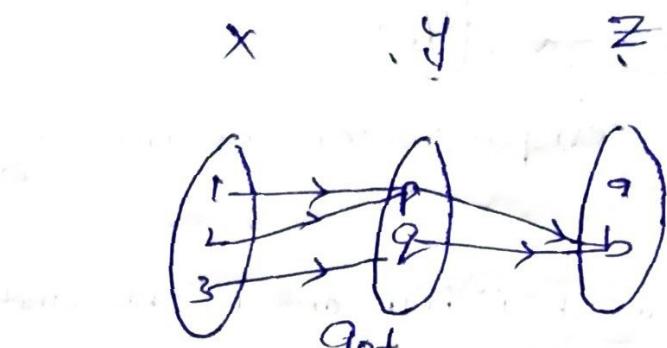
Sol Given $X = \{1, 2, 3\}$, $Y = \{P, Q\}$, $Z = \{a, b\}$

$f: X \rightarrow Y$, $f = \{\langle 1, P \rangle, \langle 2, P \rangle, \langle 3, Q \rangle\}$

$g: Y \rightarrow Z$, $g = \{\langle P, a \rangle, \langle Q, b \rangle\}$



The composition function of $g \circ f$.



Finding $g \circ f$

$$g \circ f(1) = g(f(1)) = g(P) = b$$

$$g \circ f(2) = g(f(2)) = g(P) = b$$

$$g \circ f(3) = g(f(3)) = g(Q) = b$$

$$\therefore g \circ f = \{(1, b), (2, b), (3, b)\}$$

composition of function :-

* Let $f: X \rightarrow Y$ and $g: Y \rightarrow Z$ be two functions
the composite relation $g \circ f$ such that.

$$g \circ f = \{(x, z) / (x \in X) \wedge (z \in Z) \wedge (\exists y) (y \in Y)$$

$y = f(x) \wedge z = g(y)\}$ is called
composition of function (or) relative product of
functions f and g .

2. Let $X = \{1, 2, 3\}$ and f, g, h, s be functions from $X \times X$ given by $f = \{\langle 1, 2 \rangle, \langle 2, 3 \rangle, \langle 3, 1 \rangle\}$,
 $g = \{\langle 1, 2 \rangle, \langle 2, 1 \rangle, \langle 3, 3 \rangle\}$, $h = \{\langle 1, 1 \rangle, \langle 2, 2 \rangle, \langle 3, 2 \rangle\}$
 $s = \{\langle 1, 1 \rangle, \langle 2, 2 \rangle, \langle 3, 3 \rangle\}$. Find fog , gof , $fahog$,
 sog , gas , sas & fos .

Sol

Given $X = \{1, 2, 3\}$

$$f = \{\langle 1, 2 \rangle, \langle 2, 3 \rangle, \langle 3, 1 \rangle\}$$

$$g = \{\langle 1, 2 \rangle, \langle 2, 1 \rangle, \langle 3, 3 \rangle\}$$

$$h = \{\langle 1, 1 \rangle, \langle 2, 2 \rangle, \langle 3, 2 \rangle\}$$

$$s = \{\langle 1, 1 \rangle, \langle 2, 2 \rangle, \langle 3, 3 \rangle\}$$

fog

$$fog(1) = f(g(1)) = f(2) = 3$$

$$fog(2) = f(g(2)) = f(1) = 1$$

$$fog(3) = f(g(3)) = f(3) = 3$$

$$\therefore fog = \{\langle 1, 3 \rangle, \langle 2, 1 \rangle, \langle 3, 3 \rangle\}$$

gof

$$gof(1) = g(f(1)) = g(2) = 1$$

$$gof(2) = g(f(2)) = g(3) = 3$$

$$gof(3) = g(f(3)) = g(1) = 2$$

$$\therefore gof = \{\langle 1, 1 \rangle, \langle 2, 3 \rangle, \langle 3, 2 \rangle\}$$

fahog

$$fahog(1) = fah(g(1)) = fah(h(2)) = f(2) = 3$$

$$fahog(2) = fah(g(2)) = fah(h(1)) = f(1) = 2$$

$$fahog(3) = fah(g(3)) = fah(h(3)) = f(1) = 2$$

$$\therefore f \circ h \circ g = \{ \langle 1, 3 \rangle, \langle 2, 2 \rangle, \langle 3, 1 \rangle \}$$

sog

$$s(g(1))$$

$$sog(1) = s(\underline{g(1)}) = 2$$

$$sog(2) = s(g(2)) = s(1) = 1$$

$$sog(3) = s(g(3)) = s(3) = 3$$

$$\therefore sog = \{ \langle 1, 2 \rangle, \langle 2, 1 \rangle, \langle 3, 3 \rangle \}$$

gos

$$g \circ s(1) = g(s(1)) = g(1) = 2$$

$$g \circ s(2) = g(s(2)) = g(2) = 1$$

$$g \circ s(3) = g(s(3)) = g(3) = 3$$

$$\therefore gos = \{ \langle 1, 2 \rangle, \langle 2, 1 \rangle, \langle 3, 3 \rangle \}$$

sos

$$s \circ g(1) = s(g(1)) = s(1) = 1$$

$$s \circ g(2) = s(g(2)) = s(2) = 2$$

$$s \circ g(3) = s(g(3)) = s(3) = 3$$

$$\therefore sos = \{ \langle 1, 1 \rangle, \langle 2, 2 \rangle, \langle 3, 3 \rangle \}$$

fos

$$f \circ s(1) = f(s(1)) = f(1) = 2$$

$$f \circ s(2) = f(s(2)) = f(2) = 3$$

$$f \circ s(3) = f(s(3)) = f(3) = 1$$

$$\therefore fos = \{ \langle 1, 2 \rangle, \langle 2, 3 \rangle, \langle 3, 1 \rangle \} = f$$

3. Let $A = \{1, 2, 3, 4\}$, $B = \{a, b, c\}$, $C = \{w, x, y, z\}$
 with $f: A \rightarrow B$, $g: B \rightarrow C$ given by $f = \{1, a\}, \{2, a\},$
 $\{3, b\}, \{4, c\}$ and $g = \{a, x\}, \{b, y\}, \{c, z\}$,
 find gof

Sol

$$\text{Given } A = \{1, 2, 3, 4\}$$

$$B = \{a, b, c\}$$

$$C = \{w, x, y, z\}$$

$$f: A \rightarrow B, g: B \rightarrow C$$

$$f = \{1, a\}, \{2, a\}, \{3, b\}, \{4, c\}$$

$$g = \{a, x\}, \{b, y\}, \{c, z\}$$

gof

$$gof(1) = g(f(1)) = g(a) = x$$

$$gof(2) = g(f(2)) = g(a) = x$$

$$gof(3) = g(f(3)) = g(b) = y$$

$$gof(4) = g(f(4)) = g(c) = z$$

$$gof = \{1, x\}, \{2, x\}, \{3, y\}, \{4, z\}$$

4. Let $f(x) = x+2$, $g(x) = x-2$ and $h(x) = 3x$ for $x \in R$
 where R is the set of real numbers find gof , fog ,
 fob , gog , fob , hog , $\text{if } fohog$.

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$$\text{Given } f(x) = x+2, g(x) = x-2, h(x) = 3x$$

gof

$$gof(x) = g(f(x)) = g(x+2) = x+2-2 = x$$

$$gof = \{x, x\} \quad x \in R$$

$$f \circ g(x) = f(g(x)) = f(x-2) = x-2+2 = x$$

$$\therefore f \circ g = \{ \langle x, x \rangle / x \in \mathbb{R} \}$$

$$f \circ f(x) = f(f(x)) = f(x+2) = x+2+2 = x+4$$

$$\therefore f \circ f = \{ \langle x, x+4 \rangle / x \in \mathbb{R} \}$$

$$g \circ g(x) = g(g(x)) = g(x-2) = x-2-2 = x-4$$

$$\therefore g \circ g = \{ \langle x, x-4 \rangle / x \in \mathbb{R} \}$$

$$f \circ h(x) = f(h(x)) = f(3x) = 3x+2$$

$$f \circ h = \{ \langle x, 3x+2 \rangle / x \in \mathbb{R} \}$$

$$h \circ g(x) = h(g(x)) = h(x-2) = 3(x-2)$$

$$h \circ g = \{ \langle x, 3x-6 \rangle / x \in \mathbb{R} \}$$

$$f \circ h \circ g(x) = f(h(g(x))) = f(h(x-2)) = f(3x-6)$$

$$= 3x-6+2$$

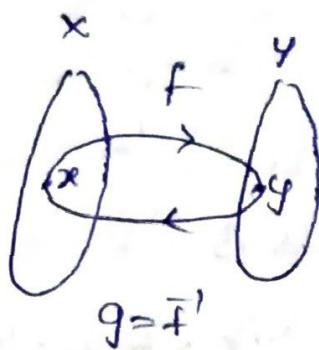
$$= 3x-4$$

$$f \circ h \circ g = \{ \langle x, 3x-4 \rangle / x \in \mathbb{R} \}$$

Inverse Functions (Or) Invertible Functions

* A function $f: x \rightarrow y$ is said to be invertible if there exist a function $g: y \rightarrow x$ such that $g \circ f = I_x$ and $f \circ g = I_y$, where I_x is the identity function on x and I_y is the identity function on y .

* Then g is called an inverse of f and we write $g = f^{-1}$



NOTE:

Let $X = \{1, 2, 3\}$, $Y = \{P, Q, R\}$ and $f: X \rightarrow Y$ be given by $f = \{(1, P), (2, Q), (3, R)\}$ then $f^{-1} = \{(P, 1), (Q, 2), (R, 3)\}$ and f^{-1} is not a function

Identity Mapping:-

A mapping $I_x: X \rightarrow X$ is called an Identity mapping if $I_x = \{(x, x) / x \in X\}$

1. Show that the functions $f(x) = x^3$ and $g(x) = x^{1/3}$ for $x \in \mathbb{R}$ are inverse of one another.

Sol. Given $f(x) = x^3$
 $g(x) = x^{1/3}$

fog

$$fog(x) = f(g(x)) = f(x^{1/3}) = (x^{1/3})^3 = x = I_x(x) I_{\mathbb{R}}(x)$$

gof $\rightarrow fog = \{(x, x) / x \in \mathbb{R}\}$

$$gof(x) = g(f(x)) = g(x^3) = (x^3)^{1/3} = x = I_x(x) I_{\mathbb{R}}(x)$$

$$gof = \{(x, x) / x \in \mathbb{R}\}$$

$$\therefore fog = gof$$

$$\therefore f = g^{-1} \text{ and } g = f^{-1}$$

$\therefore f$ and g are inverse of each other.

2. Let $A = \{1, 2, 3, 4\}$ and f and g be the functions from A to A given by $f = \{\langle 1, 4 \rangle, \langle 2, 1 \rangle, \langle 3, 2 \rangle, \langle 4, 3 \rangle\}$ and $g = \{\langle 1, 2 \rangle, \langle 2, 3 \rangle, \langle 3, 4 \rangle, \langle 4, 1 \rangle\}$. Prove that f and g are inverse of each other.

Sol.

Given $A = \{1, 2, 3, 4\}$

$$f = \{\langle 1, 4 \rangle, \langle 2, 1 \rangle, \langle 3, 2 \rangle, \langle 4, 3 \rangle\}$$

$$g = \{\langle 1, 2 \rangle, \langle 2, 3 \rangle, \langle 3, 4 \rangle, \langle 4, 1 \rangle\}$$

fog

$$fog(1) = f(g(1)) = f(2) = 1 = IA(1)$$

$$fog(2) = f(g(2)) = f(3) = 2 = IA(2)$$

$$fog(3) = f(g(3)) = f(4) = 3 = IA(3)$$

$$fog(4) = f(g(4)) = f(1) = 4 = IA(4)$$

$$fog = \{\langle 1, 1 \rangle, \langle 2, 2 \rangle, \langle 3, 3 \rangle, \langle 4, 4 \rangle\}$$

gof

$$gof(1) = g(f(1)) = g(4) = 1 = IA(1)$$

$$gof(2) = g(f(2)) = g(1) = 2 = IA(2)$$

$$gof(3) = g(f(3)) = g(2) = 3 = IA(3)$$

$$gof(4) = g(f(4)) = g(3) = 4 = IA(4)$$

$$gof = \{\langle 1, 1 \rangle, \langle 2, 2 \rangle, \langle 3, 3 \rangle, \langle 4, 4 \rangle\}$$

$$\therefore fog \Rightarrow gof$$

$$\therefore f^{-1} = g \text{ and } g^{-1} = f$$

$\therefore f$ and g are inverse of each other.

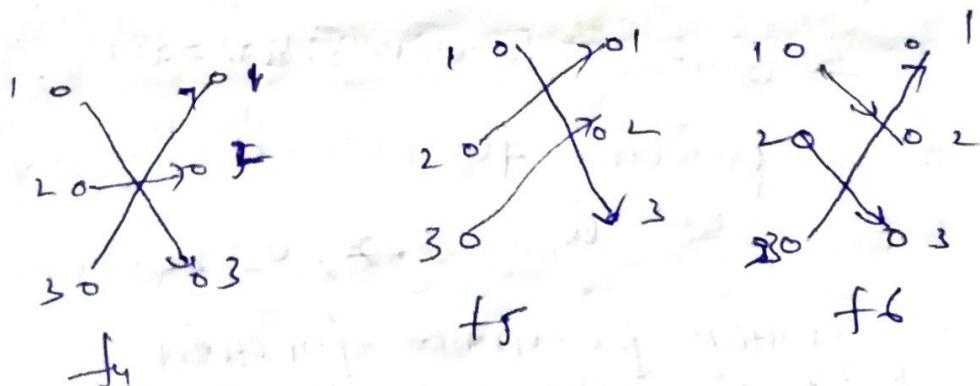
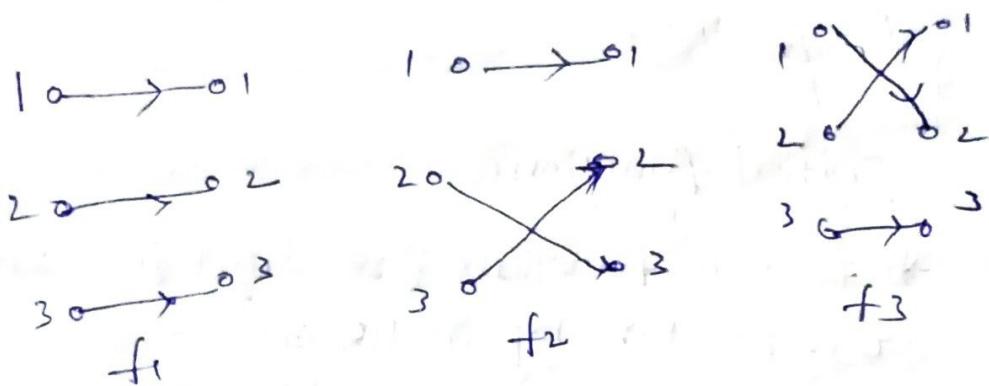
Let F_x be the set of all one to one mapping from X onto X where $X = \{1, 2, 3\}$. Find all the elements of F_x and find the inverse of each element.

NOTE: If a set X has n elements, then there are $n!$ functions from $X \rightarrow X$ which are bijective.

Given $X = \{1, 2, 3\}$

Since X consists of 3 elements then a set of bijective function from $X \rightarrow X$ consists of $3! = 6$ elements.

$$F_x = \{f_1, f_2, f_3, f_4, f_5, f_6\}$$



$$f_1^{-1} = f_1, f_2^{-1} = f_3 = f_5, f_4^{-1} = f_4, \text{ and } f_6^{-1} = f_6$$

Recursive function

* Any function $f: \mathbb{N}^n \rightarrow \mathbb{N}$ is called total because it is defined for every n -tuple in \mathbb{N}^n . On the other hand $f: D \rightarrow \mathbb{N}$ where $D \subseteq \mathbb{N}^n$, f is called partial.

Ex:

1. $f(x, y) = x + y$ which is defined for all $x, y \in \mathbb{N}$ and hence is a total function.
2. $g(x, y) = x - y$ which is defined for only those $x, y \in \mathbb{N}$ which satisfies $x \geq y$. Hence $g(x, y)$ is partial.

Initial functions

* Initial functions are used in defining other functions by induction

1. zero function $z: z(x) = 0$
2. successor function $s: s(x) = x + 1$
3. projection function $u_i^n: u_i^n(x_1, x_2, \dots, x_n) = x_i$

$$\text{Ex: } u_1^1(x, y) = x, u_2^3(2, 4, 6) = 4$$

Primitive Recursive function

* A function f is called primitive recursive iff it can be obtain from the initial functions by a finite number of operations of composition and recursion.

1. Show that the function $f(x, y) = x+y$ is primitive recursive.

Sol

Given $f(x, y) = x+y$

zero function:

$$f(x, 0) = x+0 = x$$

successor function

$$\begin{aligned} f(x, y+1) &= x+(y+1) = (x+y)+1 && \because S: S(x) = x+1 \\ &= f(x, y)+1 \\ &= S(f(x, y)) \end{aligned}$$

projection function

$$f(x, 0) = x = u_1(x) \rightarrow \text{base function.}$$

$$f(x, y+1) = S(f(x, y)) = S(u_3(x, y) + f(x, y))$$

Induction
Step function.

2. Using recursion, define the multiplication function * given by $g(x, y) = x*y$.

Sol

Given $g(x, y) = x*y$

zero function

$$g(x, 0) = x*0 = 0$$

successor function

$$\begin{aligned} g(x, y+1) &= x*(y+1) \\ &= (x*y) + x \end{aligned}$$

$$= g(x, y) + x$$

$$= S(g(x, y))$$

Projection function

$$g(x,y+1) = g(x,y) + x = f(u^3(x,y, g(x,y))),$$

$$u^3(x,y, g(x,y))$$

Primitive recursive function:

* Sign function or Non-test function (sg)

$$sg(0) = 0, sg(y+1) = 1$$

* zero test function (\bar{sg})

$$\bar{sg}(0) = 1, \bar{sg}(y+1) = 0$$

* Predecessor function

Predecessor

$$P(0) = 0, P(y+1) = y = u^2(y, P(y))$$

$$\text{Ex: } P(0) = 7, P(7) = 8.$$

* Proper subtraction ($\bar{-}$)

$$x \bar{-} 0 = x, x \bar{-} (y+1) = P(x \bar{-} y)$$

NOTE

* $x \bar{-} y = 0$ for $x \leq y$

* $x \bar{-} y = x - y$ for $x \geq y$

3. Show that $f(x,y) = xy$ is a primitive recursive function.

Sol

Given $f(x,y) = xy$

zero function

$$f(x,0) = x^0 = 1$$

Successor function

$$f(x, y+1) = x^{y+1}$$

$$= x^y \cdot x$$

$$= f(x, y) * x$$

Projection function $\therefore x^0 = 1$, for $(x \neq 0)$

$$f(x, 0) = sg(x) \quad \text{Put } x^0 = 0, \text{ for } x = 0.$$

$$f(x, y+1) = f(x, y) * x$$

$$= u_3^3 \langle x, y, f(x, y) \rangle * u_1^3 \langle x, y, f(x, y) \rangle.$$

4. Let $\lfloor \sqrt{x} \rfloor$ be the greatest integer $\leq \sqrt{x}$.
 Show that $\lfloor \sqrt{x} \rfloor$ is primitive recursive.

sol Let $\lfloor \sqrt{x} \rfloor$ be the greatest integer $\leq \sqrt{x}$.
 we observe that $(y+1)^2 - x$ is zero for $(y+1)^2 \leq x$
 and non-zero for $(y+1)^2 > x$.

* Therefore $\overline{sg}((y+1)^2 - x)$ is 1 for $(y+1)^2 \leq x$
 and can't be equal to zero.

* The smallest value of y for which
 $(y+1)^2 \geq x$ is the required number $\lfloor \sqrt{x} \rfloor$

* Hence $\lfloor \sqrt{x} \rfloor = u_y \langle \overline{sg}((y+1)^2 - x) \rangle = 0$

Algebraic Structures

Definitions and Examples

* Let $f: X^n \rightarrow X$, for $n=1$, such an operation is called unary operation.

* Similarly $n=2$ is called a Binary operation. any

* distinguished element of X such as an identity

element or zero element with respect to a
Binary operation is considered as a Zero operation

* A Binary operation will be denoted by means of
symbols as $\ast, \Delta, +, \oplus$ and the result of
Binary operation on the elements say $x_1, x_2 \in X$
is expressed by writing $x_1 \ast x_2$.

* If f denotes an n -ary operation when
 $f: \langle x_1, x_2, \dots, x_n \rangle \rightarrow \text{operation}$ is the image in
 X of the n -tuple $\langle x_1, x_2, \dots, x_n \rangle \in X^n$.

* We shall denote an algebraic system $\langle S, f_1, f_2, \dots \rangle$ where f_i is a non-empty set f_1, f_2, \dots
are the operations on S .

* Since the operations and relations on the set
 S define a structure on the elements of S an algebraic
system is called a algebraic structure.

Exemples:

* Let I be set of integers consider the
algebraic system $\langle I, +, \ast \rangle$ where $+$ and \ast are
the operations of addition and multiplication on I .

* (A-1) for any $a, b \in I$.

* $a + (b + c) = (a + b) + c \rightarrow$ associative.

* (A-2) for any $a, b \in I$

* $a + b = b + a \rightarrow$ commutative.

* (A-3) There exist a distinguished element
 $0 \in I$ such that for any $a \in I$

* $a + 0 = 0 + a = a$ (\rightarrow Identity law)

* $0 \in I$ is the Identity element under
addition

* (A-4) for each $a \in \mathbb{P}$ there exist an element in \mathbb{P} denoted by $-a$ and called the negative off a such that $a + (-a) = a - a = 0$ (Inverse law).

* (M-1) for any $a, b, c \in \mathbb{P}$

$$a * (b * c) = (a * b) * c \rightarrow \text{Associativity}$$

* (M-2) for any $a, b, c \in \mathbb{P}$

$$a * b = b * a \text{ (commutative)}$$

* (M-3) there exists a distinguished element $1 \in \mathbb{P}$ such that for any $a \in \mathbb{P}$.

$$a * 1 = 1 * a = a \text{ (Identity law)}$$

* (D), for any $a, b, c \in \mathbb{P}$, the

$$a * (b + c) = (a * b) + (a * c) \text{ (Distributive)}$$

* (C) for any $a, b, c \in \mathbb{P}$ and $a \neq 0$

$$a * b = a * c \Rightarrow b = c \text{ (left cancellation law)}$$

* consider a set $B = \{0, 1\}$ and the operations $+_2$ & $*_2$ on B given by the following tables

$+_2$	0	1
0	0	1
1	1	0

$*_2$	0	1
0	0	0
1	0	1

- which of the following system satisfies the properties of $\langle I, +, * \rangle$. i.e. $\langle \mathbb{Z}_6, +_6, *_6 \rangle$

Qd Given $\mathbb{Z}_6 = \{0, 1, 2, 3, 4, 5\}$

+_6	0	1	2	3	4	5
0	0	1	2	3	4	5
1	0	2	3	4	5	0
2	2	3	4	5	0	1
3	3	4	5	0	1	2
4	4	5	0	1	2	3
5	5	0	1	2	3	4

\times_6	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0	1	2	3	4	5
2	0	2	4	0	2	4
3	0	3	0	3	0	3
4	0	4	2	0	4	2
5	0	5	4	3	2	1

* $\langle z_6, +_6 \rangle$ satisfies closure and associative laws.

* $\langle z_6, +_6, \times_6 \rangle$ satisfies Identity Law.

* $\langle z_6, +_6 \rangle$ satisfies Inverse law.

$$\text{* i.e } \bar{0} = 0$$

$$\bar{1} = 5$$

$$\bar{2} = 4$$

$$\bar{3} = 3$$

$$\bar{4} = 2$$

$$\bar{5} = 1$$

* $\langle z_6, \times_6 \rangle$ does not satisfies the Inverse Law.

Sub-algebra:-

* Let $\langle X, \circ \rangle$ be an algebraic system and $S \subseteq X$ which is closed under the operation \circ then $\langle S, \circ \rangle$ is called a sub-algebra of $\langle X, \circ \rangle$.

Ph.

Homomorphism :-

* Let $\langle X, \circ \rangle$ and $\langle Y, * \rangle$ be two algebraic systems of the same type in the sense that both \circ and $*$ are binary operations (\oplus -ary) (n -ary).

* A mapping $g: X \rightarrow Y$ is called homomorphism (or) simply morphism from $\langle X, \circ \rangle$ to $\langle Y, * \rangle$ if x_1, x_2

Ex

$$g(x_1 \circ x_2) = g(x_1) * g(x_2)$$

Definition

Let g be a homomorphism from $\langle X, \circ \rangle$ to $\langle Y, * \rangle$.

If $g: X \rightarrow Y$ is onto then g is called epimorphism.

If $g: X \rightarrow Y$ is one to one, then g is called monomorphism.

If $g: X \rightarrow Y$ is one to one onto or bijective then g is called an isomorphism.

Isomorphic

Let $\langle X, \circ \rangle$ and $\langle Y, * \rangle$ be two algebraic systems

of the same type if there exist an isomorphic

mapping $g: X \rightarrow Y$, then $\langle X, \circ \rangle$ and $\langle Y, * \rangle$ are

said to be Isomorphic.

Q. Prove that symmetric difference is commutative operation on sets.

So Consider any two sets A and B .

$$A + B = (A - B) \cup (B - A)$$

$$A \oplus = (A - B) \cup (B - A)$$

$$= \{x/x \in A \cup B \text{ and } x \notin A \cap B\}$$

$$= \{x/x \in B - A \text{ and } x \notin A - B\}$$

$$= B \oplus A$$

2. If $f: S \rightarrow T$ is a homomorphism $\langle S, * \rangle$ to $\langle T, \Delta \rangle$ and $g: T \rightarrow P$ is also a homomorphism $\langle T, \Delta \rangle$ to $\langle T, \nabla \rangle$, then $g \circ f: S \rightarrow P$ is a homomorphism $\langle S, * \rangle$ to $\langle P, \nabla \rangle$.

Sol

Given that

$f: S \rightarrow T$ and $g: T \rightarrow P$ is a homomorphism.

We have to show that $g \circ f: S \rightarrow P$ is a homomorphism.

If f is a homomorphism, then for any $s_1, s_2 \in S$

we have $f(s_1 * s_2) = f(s_1) \Delta f(s_2)$. But for $f(s_1)$ and $f(s_2) \in T$.

We have $g(f(s_1) \Delta f(s_2)) = g(f(s_1)) \nabla g(f(s_2))$.

For any $s_1, s_2 \in S$ the composition function $g \circ f$ we have

$$g \circ f(s_1 * s_2) = g(f(s_1 * s_2))$$

$$= g(f(s_1) \Delta f(s_2)) \quad [f \text{ is a}$$

homomorphism]

$$= g(f(s_1)) \nabla g(f(s_2)) \quad [g \text{ is a}$$

homomorphism]

Hence

$g \circ f: S \rightarrow P$ is a homomorphism.

Semi groups and Monoids

Semigroups:

- * Let 'S' be a non-empty set, 'o' be a binary operation on S. The algebraic system $\langle S, o \rangle$ is called a semigroup. If the operation circle is associative.
- * In other words $\langle S, o \rangle$ is a semi group if for any $x, y, z \in S$, $x o (y o z) = (x o y) o z$

Monoids

- * A semi group $\langle M, o \rangle$ with an identity element with respect to the operation 'o' is called a monoid.
 - * In other words algebraic system $\langle M, o \rangle$ is a called a monoid if for any $x, y, z \in M$,
 $x o (y o z) = (x o y) o z$ and there exist an element $e \in M$ such that for any $x \in M$
 $x o e = e o x = x$
- NOTE: A monoid is always a semi group

Examples:

- * Let 'N' be the set of natural numbers, then the algebraic system $\langle N, * \rangle$ is a monoid with identity element "1", $\langle N, + \rangle$ is not a monoid.
- * If 'E' be the set of all the even numbers then $\langle E, + \rangle$ and $\langle E, * \rangle$ are semi groups but not monoids.
- * Let 'S' be a non-empty set and $P(S)$ be its powerset. The algebraic systems

$\langle P(S), \cup \rangle$

$\langle P(S), \cap \rangle$ are monoids with the identities \emptyset and S respectively.

Semigroup homomorphism

- * Let $\langle S, * \rangle$ and $\langle T, \cdot \rangle$ be any two semigroups. A mapping $g: S \rightarrow T$ such that for any two elements $a, b \in S$.
- * $g(a * b) = g(a) \cdot g(b)$ is called a semigroup homomorphism.

Monoid homomorphism :-

Let $\langle M, *, e_M \rangle$ and $\langle T, \cdot, e_T \rangle$ be any two monoids. A mapping $g: M \rightarrow T$ such that for any two elements $a, b \in M$.

$$g(a * b) = g(a) \cdot g(b)$$

$g(e_M) = e_T$ is called a monoid

homomorphism.

Theorem

Let $\langle S, * \rangle$, $\langle T, \cdot \rangle$ and $\langle V, \oplus \rangle$ be semigroups and $g: S \rightarrow V$ and $h: T \rightarrow V$ be semigroup homomorphisms. Then $(h \circ g): S \rightarrow V$ is a semigroup homomorphism from $\langle S, * \rangle$ to $\langle V, \oplus \rangle$.

Proof

Given that

$\langle S, * \rangle$, $\langle T, \cdot \rangle$, $\langle V, \oplus \rangle$ are semigroups.

also given that $g: S \rightarrow V$ and $h: T \rightarrow V$ be semigroup homomorphism.

We have to show that $(h \circ g): S \rightarrow V$ is a semi group homomorphism.

Let $a, b \in S$

$$\begin{aligned} \text{Then } h \circ g(a * b) &= h(g(a * b)) \\ &= h(g(a) \Delta g(b)) \\ &= h(g(a)) \oplus h(g(b)) \\ &= h \circ g(a) \oplus h \circ g(b) \end{aligned}$$

Therefore $\langle S, *\rangle$ to $\langle V, \oplus \rangle$ is a semi group homomorphism.

Theorem - 2

Let $\langle S, *\rangle$ and $\langle T, \Delta \rangle$ be two semi groups and g be a semigroup homomorphism from $\langle S, *\rangle$ to $\langle T, \Delta \rangle$ corresponding to the homomorphism g , there exist a congruence relation R on $\langle S, *\rangle$ defined by $x R y \iff g(x) = g(y)$ for

$x, y \in S$.

Proof:-

Given $\langle S, *\rangle$ and $\langle T, \Delta \rangle$ be two semi groups.

$g: S \rightarrow T$ is a semi group homomorphism and

R is equivalence relation on $\langle S, *\rangle$.

Let x_1, x_2, x'_1, x'_2 such that $x_1 R x'_1$ and $x_2 R x'_2$

$$\begin{aligned} \text{consider } g(x_1 * x_2) &= g(x_1) \Delta g(x_2) \\ &= g(x'_1) \Delta g(x'_2) \\ &= g(x'_1 * x'_2) \end{aligned}$$

* It follows that R is a congruence relation on $\langle S, *\rangle$

Sub-semigroups and Sub monoids

* Let $\langle S, * \rangle$ be a semi group and $T \subseteq S$.
If the set T is closed under the operation $*$, then $T, *$ is said to be a sub-semi group of $\langle S, * \rangle$.

* Let $\langle M, * \rangle$ be a monoid and $T \subseteq M$.
If the set T is closed under the operation $*$ and $e \in T$, Then $T, *, e$ is said to be a Sub-monoid of $\langle M, *, e \rangle$.

* For any $a \in S$, the set consisting of all powers of a under the operation $*$ is a sub-semi group i.e $\overline{B} = \{a, a^2, a^3, \dots\}$
where $a^k = a * a^{k-1}$. Then $\langle \overline{B}, * \rangle$ is a cyclic semi group which is a sub-semi group of $\langle S, * \rangle$ generated by the element a .

Theorem-3

For any commutative monoid $\langle M, * \rangle$, the set Idempotent element of M forms a sub-monoid.

Proof: Given that $\langle M, * \rangle$ is a commutative monoid i.e $a * b = b * a \quad a, b \in M$

Since the identity $e \in M$ is Idempotent.

Let S be set of Idempotent of M (SCM)

i.e let $a, b \in S$

Then $a * a = a$ and

$$b * b = b$$

$$\begin{aligned} \text{Consider } (a * b) * (a * b) &= a * (b * a) * b \\ &= a * (a * b) * b \end{aligned}$$

$$= (a * a) * (b * b)$$

$$= a * b$$

Hence $a * b \in S$ and $\langle S, * \rangle$ is a sub-Monoid.

Groups:-

Let G' be a non-empty set in which an operation $*$ is defined then the algebraic structure $\langle G', * \rangle$ is said to be a group. If G' satisfies the following axioms.

1. $\forall a, b \in G', a * b \in G' \rightarrow$ closure axiom
2. $\forall a, b, c \in G', a * (b * c) = (a * b) * c \rightarrow$ associative axiom
3. $\forall a \in G', \exists e \in G' \Rightarrow a * e = e * a = a$
where e is an Identity element
4. $\forall a \in G', \exists a' \in G' \Rightarrow a * a' = a' * a = e$
 a' is called inverse of a .

Abelian group

A Group $\langle G, * \rangle$ is said to an abelian group if $a * b = b * a, \forall a, b \in G$ (Commutative property)

Non-abelian group:-

If the group does not satisfy commutative property, then the group G is called non-abelian.

Finite and Infinite group:-

If the no: of elements in a group G' is countable then G' is finite group otherwise infinite group.

order of a group

The no: of elements in a finite a group is called order of a group. It is denoted by $O(G)$ or $|G|$

* show that the set of integers \mathbb{Z} with respect to addition is an infinite abelian group.

Given that $\langle \mathbb{Z}, + \rangle$

$$\text{where } \mathbb{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}$$

and the binary operation '+' is the binary composition in \mathbb{Z} .

i. Closure Property

Let $a, b \in \mathbb{Z}$, then $a+b \in \mathbb{Z}$

$\langle \mathbb{Z}, + \rangle$ satisfies closure property.

ii. Associative Property:

Let $a, b, c \in \mathbb{Z}$, then $a+(b+c) = (a+b)+c$

$\langle \mathbb{Z}, + \rangle$ satisfies associative property.

iii. Identity Property

Let $a \in \mathbb{Z}$, then there exist an element $0 \in \mathbb{Z}$ such that

$$a+0=0+a=a$$

'0' is the identity element under addition,
 $\langle \mathbb{Z}, + \rangle$ satisfies identity property.

iv. Inverse Property:

Let $a \in \mathbb{Z}$, there exist an element $-a \in \mathbb{Z}$, such that $a+(-a) = a-a=0$

• if $-a$ is the additive inverse of a .

∴ $\langle \mathbb{Z}, + \rangle$ is a group.

v. Commutative Property:

Let $a, b \in \mathbb{Z}$, then $a+b=b+a$.

$\langle \mathbb{Z}, + \rangle$ is a abelian group.

The number of integers is infinite.
 $\langle \mathbb{Z}, + \rangle$ is an infinite abelian group.

Sub-group

Let $\langle G, * \rangle$ be a group and $S \subseteq G$ be such that it satisfies the following conditions,

1. $e \in S$, where e is the identity $\langle G, * \rangle$.
2. For any $a \in S$, $\bar{a} \in S$.
3. For $a, b \in S$, $a * b \in S$

Then $\langle S, * \rangle$ is called a sub-group of $\langle G, * \rangle$

* A subset $S \neq \emptyset$ of G is a subgroup of $\langle G, * \rangle$ iff for any pair of elements $a, b \in S$, $a * b^{-1} \in S$

Sol Given that $\langle G, * \rangle$ be a group.

that
assume S is a sub group of G

Let we know that $a \in S$, then $\bar{a} \in S$

$$\therefore a * \bar{a} \in S$$

Converse

Let us assume that $a \in S$ and $a * \bar{a} \in S$
we have to prove that S is a subgroup of G .

$$\text{Taking } b \in S, \text{ then } a * \bar{a} = e \in S$$

From $a, b \in S$, then

$$a * e = a * \bar{a} = \bar{a} \in S$$

$$\text{Similarly } b * e = e * b = b \in S$$

Finally $a, b \in S$ we have

$$a * b \in S$$

$\therefore \langle S, * \rangle$ is a sub group of $\langle G, * \rangle$

Group Homomorphism:

- * Let $\langle G, * \rangle$ and $\langle H, \Delta \rangle$ be two groups
a mapping $g: G \rightarrow H$ is called a group homomorphism from $\langle G, * \rangle$ to $\langle H, \Delta \rangle$.
- * If for any $a, b \in G$

$$g(a * b) = g(a) \Delta g(b)$$

$$g(e_G) = e_H$$

$$g(a^{-1}) = [g(a)]^{-1}$$

Kernel

- * Let g be a group homomorphism from $\langle G, * \rangle$ to $\langle H, \Delta \rangle$. The set of elements of G which are mapped into e_H , the identity of H is called the kernel of the homomorphism g and denoted by $\text{Ker}(g)$

or

kernel of g

Theorem:

The kernel of a homomorphism g from a group $\langle G, * \rangle$ to $\langle H, \Delta \rangle$ is a subgroup of $\langle G, * \rangle$

Sol Given that $\langle G, * \rangle$ and $\langle H, \Delta \rangle$ be two groups. We have to show that kernel of g is a subgroup of $\langle G, * \rangle$.

$g: G \rightarrow H$ is a homomorphism from $\langle G, * \rangle$ to $\langle H, \Delta \rangle$.

Since $g(e_G) = e_H$, $e_G \in \text{Ker}(g)$

FOR $a, b \in \text{kernel}(g)$ that is

$$g(a) = g(b) = e_H$$

$$\text{Then } g(a * b) = g(a) \Delta g(b)$$

$$= e_H \Delta e_H$$

$$= e_H$$

$$\therefore a * b \in \text{ker}(g)$$

$$\text{If } a \in \text{ker}(g) \text{ then } g(\bar{a})! = (\bar{g}(a))! = e_H^{-1}$$

$$\therefore \bar{a} \in \text{ker}(g) \text{ and kernel of } g \text{ is } a = C_H$$

sub group of $\langle G, * \rangle$.

Lattices:-

A lattice is a partially ordered set or poset $\langle L, \leq \rangle$ in which every pair of elements $a, b \in L$ has a greatest lower bound (GLB) of a subset $\{a, b\} \subseteq L$ and a least upper bound.

* The greatest lower bound of a subset $\{a, b\} \subseteq L$ will be denoted by $a \wedge b$, and the least upper bound by $a \vee b$.

* It is customary to call the $GLB\{a, b\} = a \wedge b$ the meet or product of a and b and $LUB\{a, b\} = a \vee b$ the join or sum of a and b .

* Other symbols such as \wedge and \vee (or \cdot and $+$) are also used to denote the meet and join of two elements respectively.

Some Properties Of Lattices

* We shall first list sum of the properties of the two binary operations Meet and Join denoted by $*$ and \oplus and lattice $\langle L, \leq \rangle$ for any $a, b \in L$ we have

1. Idempotent property

$$a * a = a$$

$$a \oplus a = a$$

2. $a * b = b * a$ } Commutative Property
 $a \oplus b = b \oplus a$

3. $a * (b * c) = (a * b) * c$ } Associative
 $a \oplus (b \oplus c) = (a \oplus b) \oplus c$ } Property

4. $a * (a \oplus b) = a$ } Absorption law
 $a \oplus (a * b) = a$

Theorem:-

Let $\langle L, \leq \rangle$ be a lattice in which $*$ and \oplus denote the operations of meet and join respectively. For any $a, b, c \in L$, $a \leq b \Leftrightarrow a * b = a \leq a \oplus b = b$

Proof:-

Let $\langle L, \leq \rangle$ be a lattice in which $*$ and \oplus denote the operations of meet and join respectively

We have show that

$$a, b, c \in L, a \leq b \Leftrightarrow a * b = a \Leftrightarrow a \oplus b = b$$

for this it is enough to show that

i. $a * b \leq a \oplus b = a$

ii. $a \leq b \Leftrightarrow a \oplus b = b$

i. First suppose that $a \leq b$

To prove that $a \leq a \ast b$

By additive property $a \leq a \oplus b \leq b$

For the set $\{a, b\}$, 'a' is the lower bound.

Greatest lower bound (g.l.b) of $\{a, b\}$ is $a \ast b$

$$\Rightarrow a \leq a \ast b \quad \text{--- (1)}$$

Again $a \ast b$ is the g.l.b of $\{a, b\}$

$\Rightarrow a \ast b$ is the lower bound

$$\Rightarrow a \ast b \leq a, b$$

$$\Rightarrow a \ast b \leq a \quad \text{--- (2)}$$

From (1) & (2)

$$\therefore a = a \ast b$$

$$a \leq b \Leftrightarrow a = a \ast b.$$

Converse

Suppose that $a \ast b = a = \text{g.l.b of } \{a, b\}$

$$\Rightarrow a \leq a, b$$

$$\Rightarrow a \leq b$$

$$\therefore a \leq b \Leftrightarrow a \ast b = a.$$

iii. Similarly we can prove that

$$a \leq b \Leftrightarrow a \oplus b = b$$

$$\therefore a \leq b \Leftrightarrow a \ast b = a \Leftrightarrow a \oplus b = b.$$

Theorem - 2

Let $\langle L, \leq \rangle$ be a lattice. For any $a, b, c \in L$,
the following properties called isotonicity hold

$$b \leq c \Rightarrow \begin{cases} a \ast b \leq a \ast c \\ a \oplus b \leq a \oplus c \end{cases}$$

Proof

Let (L, \leq) be a lattice and $a, b, c \in L$.

Now we have to prove that

$$b \leq c \Rightarrow \begin{cases} a * b \leq a * c \\ a \oplus b \leq a \oplus c \end{cases}$$

we know that

$$a \leq b \Rightarrow a * b = a \Leftrightarrow a \oplus b = b$$

for $b, c \in L$

we have

$$b \leq c \Rightarrow b * c = b \Leftrightarrow b \oplus c = c \quad (1)$$

First suppose that $b \leq c$.

i. To prove that

$$a * b \leq a * c$$

For this it is enough to show that $(a * b) * (a * c) = a * b$

$$\begin{aligned} \text{consider } (a * b) * (a * c) &= a * (b * a) * c \\ &= a * (a * b) * c \\ &= (a * a) * (b * c) \\ &= a * (b * c) \text{ (from (1))} \\ &= a * b \end{aligned}$$

$$\therefore (a * b) * (a * c) = a * b$$

$$\therefore (a * b) \leq (a * c).$$

$$\therefore b \leq c \Rightarrow a * b \leq a * c$$

ii. To prove that

$$a \oplus b \leq a \oplus c$$

For this it is enough to show that

$$(a \oplus b) \oplus (a \oplus c) = a \oplus c$$

$$\begin{aligned}
 \text{consider } (a \oplus b) \oplus (a \oplus c) &= a \oplus (b \oplus a) \oplus c \\
 &= a \oplus (a \oplus b) \oplus c \\
 &= (a \oplus a) \oplus (b \oplus c) \\
 &= a \oplus (b \oplus c) \quad (\text{from i}) \\
 &= a \oplus \mathbb{1}
 \end{aligned}$$

$$\therefore (a \oplus b) \oplus (a \oplus c) = a \oplus \mathbb{1}$$

$$\begin{aligned}
 \text{since } (a \oplus b) \leq (a \oplus c) \\
 \therefore b \leq c \Rightarrow a \oplus b \leq a \oplus c.
 \end{aligned}$$

$$\begin{aligned}
 \therefore b \leq c \Rightarrow \begin{cases} a \oplus b \leq a \oplus c & \text{hence proved} \\ a \oplus b \leq a \oplus c \end{cases}
 \end{aligned}$$

NOTE:

1. $a \leq b \wedge a \leq c \Rightarrow a \leq (b \oplus c)$
2. $a \leq b \wedge a \leq c \Rightarrow a \leq (b * c)$
3. $a \geq b \wedge a \geq c \Rightarrow a \geq (b * c)$
4. $a \geq b \wedge a \geq c \Rightarrow a \geq (b \oplus c)$

Theorem - 3

Let $\langle L, \leq \rangle$ be a lattice for any $a, b, c \in L$.
 the following inequalities called the distributive
 inequalities hold.

- i. $a * (b \oplus c) \geq (a * b) \oplus (a * c)$
- ii. $a \oplus (b * c) \leq (a \oplus b) * (a \oplus c)$

proof:

Let $\langle L, \leq \rangle$ be a lattice for any $a, b, c \in L$.

Now we have to prove that

- i. $a * (b \oplus c) \geq (a * b) \oplus (a * c)$
- ii. $a \oplus (b * c) \leq (a \oplus b) * (a \oplus c)$

i. To prove that

$$a * (b \oplus c) \geq (a * b) \oplus (a * c)$$

By the definition of g.l.b

$$a * b \leq a \oplus b \quad a * c \leq a \oplus c$$

$$a * b \leq a \quad a * c \leq a$$

'a' is the upper bound

- the least upper bound of $\{a * b, a * c\}$

$$= a * b \oplus a * c$$

$$\leq a \oplus a$$

$$\leq a \quad \text{---(1)}$$

we have $a * b \leq b, a * c \leq c$

- the l.u.b. $\{a * b, a * c\} = (a * b) \oplus (a * c)$

$$\leq b \oplus c$$

$$\therefore (a * b) \oplus (a * c) \leq b \oplus c \quad \text{---(2)}$$

From (1) & (2)

$$(a * b) \oplus (a * c) \leq \text{g.l.b.} \{ a * b \oplus a * c \}$$

$$(a * b) \oplus (a * c) \leq a * (b \oplus c)$$

$$\Rightarrow a * (b \oplus c) \geq (a * b) \oplus (a * c) \quad \text{---(3)}$$

Dualizing eq (3) [changing $*$ by \oplus and \oplus by $*$ and \leq by \geq]
we get.

$$\Rightarrow a \oplus (b * c) \leq (a \oplus b) * (a \oplus c)$$

\therefore Hence proved.

Theorem-4

Let $\langle L, \leq \rangle$ be a lattice for any $a, b, c \in L$,
following inequalities called the modulus inequalities hold as $c \Leftarrow a \oplus (b * c) \leq (a \oplus b) * c$.

Proof:

Let $\langle L, \leq \rangle$ be a lattice for any $a, b, c \in L$.

Now we have to prove that

$$a \leq c \Leftrightarrow a \oplus (b * c) \leq (a \oplus b) * c$$

for $a, c \in L$.

We have $a \leq c \Leftrightarrow a * c = a \Leftrightarrow a \oplus c = c$ — (1)
from distributive inequality

$$a \oplus (b * c) \leq (a \oplus b) * (a \oplus c) — (2)$$

Suppose that $a \leq c$

we have to prove that $a \oplus (b * c) \leq (a \oplus b) * c$

from (2) we have

$$\begin{aligned} a \oplus (b * c) &\leq (a \oplus b) * (a \oplus c) \\ &\leq (a \oplus b) * c \quad (\text{from (1)}) \end{aligned}$$

converse

Suppose that $a \oplus (b * c) \leq (a \oplus b) * c$ — (3)

from (2) and (3), we have

$$(a \oplus b) * (a \oplus c) = (a \oplus b) * c$$

By using left cancellation law,

$$\begin{aligned} a \oplus c &= c \\ \Rightarrow a &\leq c \quad (\text{from (1)}) \end{aligned}$$

$$\therefore a \leq c \Leftrightarrow a \oplus (b * c) \leq (a \oplus b) * c$$

Theorem - 5

Show that in a lattice if $a \leq b \leq c$,
then i. $a \oplus b = b \oplus c$

$$\text{ii. } (a * b) \oplus (b * c) = b = (a \oplus b) * (a \oplus c)$$

Proof:-

Let $L_{L, \leq}$ be a Lattice.

for any $a, b, c \in L$

such that $a \leq b \leq c$

i. we know that

$$a \leq b \Leftrightarrow a * b = a \quad \textcircled{1}$$

$$\Leftrightarrow a \oplus b = b \quad \textcircled{2}$$

$$b \leq c \Leftrightarrow b * c = b \quad \textcircled{3}$$

$$\Leftrightarrow b \oplus c = c \quad \textcircled{4}$$

$$a \leq c \Leftrightarrow a * c = a \quad \textcircled{5}$$

$$\Leftrightarrow a \oplus c = c \quad \textcircled{6}$$

from eq $\textcircled{2}$ and $\textcircled{3}$ we get

$$a \oplus b = b * c$$

ii. L.H.S

$$(a * b) \oplus (b * c) = a \oplus b$$

$$= b \quad \textcircled{7}$$

R.H.S

$$(a \oplus b) * (a \oplus c) = b * c$$

$$\text{from } \textcircled{6} \text{ and } \textcircled{7} \Rightarrow b \quad \textcircled{8}$$

$$\text{LHS} = \text{RHS}$$

$$\therefore (a * b) \oplus (b * c) = b = (a \oplus b) * (a \oplus c)$$

Sub-lattice

* Let $\langle L, *, \oplus \rangle$ be a lattice and let $S \subseteq L$ be a subset of L . The algebra $\langle S, *, \oplus \rangle$ is a sub-lattice of $\langle L, *, \oplus \rangle \Leftrightarrow S$ is closed under both operations $*$ and \oplus .

Some special lattices:-

1. complete lattice:

A lattice is called complete if each of its non-empty sub-sets has a least upper bound and greatest lower bound.

* Obviously every finite lattice must be complete. Also every complete lattice must have at least one element and a greatest element. The least and greatest element of a lattice, if they exist are called the bounds of lattice and are denoted by $\textcircled{0}$ and $\textcircled{1}$ respectively.

* $\langle L, *, \oplus \rangle$ satisfies the following identities for any $a \in L$.

Properties:-

$$a \oplus 0 = a, a * 1 = a \rightarrow \textcircled{1}$$

$$a \oplus 1 = 1, a * 0 = 0 \rightarrow \textcircled{2}$$

complement of an element

* In a bounded lattice $\langle L, *, \oplus, 0, 1 \rangle$ and element $b \in L$ is called a complement of an element $a \in L$. If $a * b = 0$ and $a \oplus b = 1$.

* From the identities $\textcircled{1}$ and $\textcircled{2}$ we have $0 * 1 = 0$ and $0 \oplus 1 = 1$.

which shows that D and I are complements of each other.

Complemented Lattice:

* A lattice $\langle L, \wedge, \vee, 0, 1 \rangle$ is said to be a Complemented Lattice, if every element of L has atleast one complement.

Distributive Lattice

* A lattice $\langle L, \wedge, \vee, \oplus, \ominus \rangle$ is called a distributive lattice. If for any $a, b, c \in L$.

Properties:

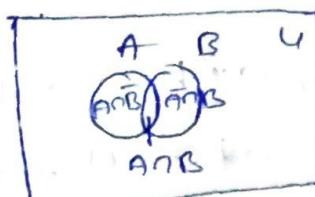
$$a * (b \oplus c) = (a * b) \oplus a * c$$

$$a \oplus (b * c) = (a \oplus b) * (a \oplus c)$$

The principle of inclusion - Exclusion:

The principle of inclusion - Exclusion for two sets.

If A and B are subsets of some universal set U , then $|A \cup B| = |A| + |B| - |A \cap B|$



For Three sets

$$\begin{aligned} |A \cup B \cup C| &= |A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C| \\ &\quad + |A \cap B \cap C| \end{aligned}$$

DeMorgan's Law

$$*(A \cap B \cap C) = \overline{A \cup B \cup C} = 1 - |A \cup B \cup C|$$

For n-sets

If A_i are finite subsets of a universal set, Then

$$\begin{aligned}|A_1 \cup A_2 \cup A_3 \cup \dots \cup A_n| &= \sum_i |A_i| \\&= \sum_{i,j} |A_i \cap A_j| + \sum_{i,j,k} |A_i \cap A_j \cap A_k| \\&\quad + \dots + (-1)^{n-1} |A_1 \cap A_2 \cap \dots \cap A_n|\end{aligned}$$

- A computer company requires 120 programmers to handle systems programming jobs and 40 programmers for applications programming. If the company appoints 55 programmers to carry out these jobs, how many of these perform jobs of both types? How many handle only system programming jobs? How many handle only applications programming jobs?

Sol Let $|A|$ denote the set of programmers who handle system programming jobs and $|B|$ denotes the set of programmers who handle application programming jobs.

$|A \cup B|$ is the set of programmers appointed to carry out these jobs.

we have $|A|=30$

$$|B|=40$$

$$|A \cup B|=55$$

$$|A \cap B|=?$$

By addition rule:

$$|A \cup B| = |A| + |B| - |A \cap B|$$

$$|A \cap B| = (|A| + |B|) - |A \cup B|$$

$$= 30 + 40 - 55$$

$$= 15$$

∴ 15 programmers to perform jobs.

ii. The number of programmers who handle only systems programming jobs is

$$|A - B| = |A| - |A \cap B|$$

$$= 30 - 15$$

$$= 15$$

∴ 15 system programmers to perform jobs.

iii. The number of programmers who handle only application programming jobs is

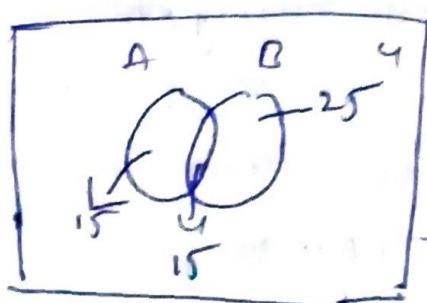
$$|B - A| = |B| - |A \cap B|$$

$$= 40 - 15$$

$$\underline{=} 25$$

∴ 25 Application programmers to Perform jobs.

Venn diagrams



2. In a class of 52 students, 30 are studying C++; 28 are studying pascal; and 13 are studying both languages how many in this class are studying atleast one of these languages? How many are studying neither of these languages?

Sol.

Let U denotes the set of all students in the class.

$|A|$ denote the set of students in the class who are studying C++

$|B|$ denotes the set of students in the class who are studying pascal.

$|A \cap B| \rightarrow$ Then the set of students in the class who are studying both languages.

The set of students who are studying atleast one ^{two} ~~these~~ languages. $|A \cup B|$

The set of students who are studying neither of these languages is $|\overline{A \cup B}|$

We have,

$$|U| = 52$$

$$|A| = 30$$

$$|B| = 28$$

$$|A \cap B| = 13$$

$$|A \cup B| = ?$$

$$|\overline{A \cup B}| = ?$$

By addition Rule,

$$|A \cup B| = |A| + |B| - |A \cap B|$$

$$= 30 + 28 - 13$$

$$= 45$$

$$\text{ii. } |A \cup B| = 44 - |A \cap B|$$

$$= 52 - 45$$

$$= 7.$$

3. In a sample of 100 logic chips, 23 have a defect D_1 , 26 have a defect D_2 , 30 have a defect D_3 , 7 have defects D_1 and D_2 , 8 have defects D_1 and D_3 , 10 have defects D_2 and D_3 and 3 have all the 3 defects. Find the no. of chips having i. atleast one defect.
ii. no defect.

Sol

Let 'u' denotes all logic chips and A, B, C denotes the sets of chips having defects D_1, D_2, D_3 respectively

$$|u| = 100$$

$$|A| = 23$$

$$|B| = 26$$

$$|C| = 30$$

$$|A \cap B| = 7$$

$$|A \cap C| = 8$$

$$|B \cap C| = 10$$

$$|A \cap B \cap C| = 3$$

- i. The set of logic chips having atleast one defect is $(A \cup B \cup C)$

$$\begin{aligned} |A \cup B \cup C| &= |A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C| + |A \cap B \cap C| \\ &= 23 + 26 + 30 - 7 - 8 - 10 + 3 \\ &= 57 \end{aligned}$$

ii. The set of logic chips having no chip defect (AUBUC)

$$\begin{aligned}|AUBUC| &= u - |AUBUC| \\&= 100 - 57 \\&= 43.\end{aligned}$$

- A. Suppose that 200 faculty members can speak French and 50 can speak Russian, while only 20 can speak both French and Russian. How many faculty members can speak either French or Russian.

Sol

Given that

$$\text{French} = 200 \Rightarrow |F| = 200$$

$$\text{Russian} = 50 \Rightarrow |R| = 50$$

$$F \cap R = 20$$

Faculty members can speak either French or

$$\text{Russian } FUR = |F| + |R| - |F \cap R|$$

$$= 200 + 50 - 20$$

$$= 230$$

Pigeon Hole Principle

If m pigeons occupy n pigeon holes, and if $m > n$, then 2 or more pigeons occupy the same pigeon hole. This statement is known as the Pigeon Hole Principle.

Generalization of the Pigeon Hole Principle:-

If m pigeons occupy n pigeon holes, then at least one pigeon hole must contain $(P+1)$ or more pigeons, where $P = \frac{m-1}{n}$

1. Prove that if 30 dictionaries containing a total of 61,327 pages then atleast one of the dictionaries must have atleast 2045 pages.

Sol Treating the pages as Pigeons and dictionaries as pigeon holes.

we find by using the Generalized Pigeon hole principle that atleast one of the dictionaries must contain $(p+1)$ or more pages.

$$\begin{aligned} p+1 &= \frac{m-1}{n} + 1 \\ &= \frac{61,327 - 1}{30} + 1 \\ &= 2045.2 \\ &\cong 2045 \end{aligned}$$

2. prove that in any set of 29 persons atleast 5 persons must have been born on the same day of the week.

Sol Treating the persons as Pigeons and Treating the 7 days of week as 7 pigeon holes.

we find by using Generalized Pigeon hole principle that atleast one day of the week is assigned to $(p+1)$.

$$\begin{aligned} p+1 &= \frac{m-1}{n} + 1 \\ &= \frac{29 - 1}{7} + 1 = \frac{28}{7} + 1 = 5 \end{aligned}$$

3. How many persons must be chosen in order that at least 5 of them will have birthdays in the same calendar month.

Let 'm' be required number of persons.

Since the no. of months over which the birthdays are distributed is 12.

The least no. of who have their birthdays in the same month is, by the generalized pigeon hole principle

$$\frac{m-1}{n} + 1 = 5$$

$$\frac{m-1}{12} + 1 = 5$$

$$m-1 = 4 \times 12$$

$$\boxed{m=49}$$

* The required no. of persons is 49.

Cyclic Group:-

A Group ' G_i ' is said to be cyclic if there exist an element $a \in G_i$, such that every element of ' G_i ' is an integral power of a .

Here 'a' is called the Generator of ' G_i ' and is denoted by $\langle a \rangle$.

$$\text{i.e } G_i = \{a^n | n \in \mathbb{Z}\}$$

Theorem:

Show that Every cyclic Group of order n is isomorphic to the group $\langle z_n, +_n \rangle$

Proof:

Let the cyclic group $\langle G_i, * \rangle$ of order n be generated by an element $a \in G_i$.

so that the elements of g are a ,
 $a^2, a^3, \dots, a^n = e$

* Define $g: \mathbb{Z}_n \rightarrow G$ such that $g(0) = e$

* so that c_{ij} is the generator of $\langle z_n \rangle$

Then $(gc_{ij}) = a^j$ for $j=0, 1, 2, \dots, n-1$

∴ Hence the isomorphism is established,