

TASK 1

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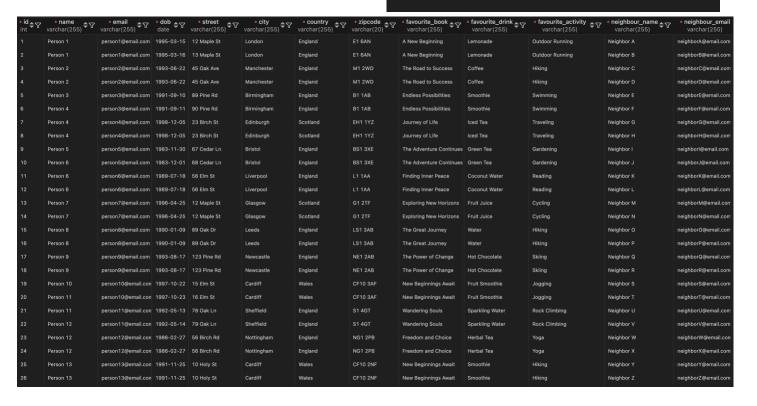
FIRST NORMAL FORM (1NF)

In First Normal Form, each column should contain an atomic value, and each record should be unique. Additionally, each row should have a unique identifier (primary key). For this purpose, a separate row was created for each neighbour of the person. The *id* was created to ensure each row in the table is uniquely identifiable, even if other columns contain duplicate values. This *id* was automatically incremented for each new entry, providing a unique identifier for each row.

1NF

- id : INTEGER NOT NULL «PK»
- name : VARCHAR(255) NOT NULLemail : VARCHAR(255) NOT NULL
- dob : DATE NOT NULL
- street : VARCHAR(255) NOT NULL
- city: VARCHAR(255) NOT NULL
- country: VARCHAR(255) NOT NULL
- zipcode : VARCHAR(255) NOT NULL
- favourite_books : VARCHAR(255) NOT NULL
- favourite_drink: VARCHAR(255) NOT NULL
- favourite_activity : VARCHAR(255) NOT NULL
- neighbour_name : VARCHAR(255) NOT NULL
- neighbour_email: VARCHAR(255) NOT NULL

- CREATE TABLE '1NF' (
- 'id' int NOT NULL AUTO_INCREMENT,
- 'name' varchar(255) NOT NULL,
- 'email' varchar(255) NOT NULL,
- 'dob' date NOT NULL,
- 'street' varchar(255) NOT NULL,
- 'city' varchar(255) NOT NULL,
- 'country' varchar(255) NOT NULL,
- 'zipcode' varchar(20) NOT NULL,
- `favourite_book` varchar(255) NOT NULL,
- `favourite_drink` varchar(255) NOT NULL,
- favourite_activity` varchar(255) NOT NULL,



SECOND NORMAL FORM (2NF)

In the Second Normal Form, all partial dependencies should be removed so that every non-key attribute should depend on the primary key. For this purpose, separate tables were created as *Person*, *Address*, *Favourite*, *Neighbour* and *Person-Neighbour*.

The Person table only contains information about the person:

```
CREATE TABLE 'Person' (
    'id' int NOT NULL AUTO_INCREMENT,
    'name' varchar(255) NOT NULL,
    'email' varchar(255) NOT NULL,
    'dob' date NOT NULL,
    PRIMARY KEY ('id')
)
```

*id int ≑⊽	* name varchar(255) ◆▽	* email varchar(255)	* dob date ◆▽
1	Person 1	person1@email.com	1995-03-15
2	Person 2	person2@email.com	1993-06-22
3	Person 3	person3@email.com	1991-09-10
4	Person 4	person4@email.com	1998-12-05
5	Person 5	person5@email.com	1983-11-30
6	Person 6	person6@email.com	1989-07-18
7	Person 7	person7@email.com	1996-04-25
8	Person 8	person8@email.com	1990-01-09
9	Person 9	person9@email.com	1993-08-17
10	Person 10	person10@email.com	1997-10-22
11	Person 11	person11@email.com	1992-05-13
12	Person 12	person12@email.com	1986-02-27
13	Person 13	person13@email.com	1991-11-25
14	Person 14	person14@email.com	1987-02-01
15	Person 15	person15@email.com	1984-08-12
16	Person 16	person16@email.com	1990-03-09
17	Person 17	person17@email.com	1995-11-17
18	Person 18	person18@email.com	1994-06-20
19	Person 19	person19@email.com	1992-12-11
20	Person 20	person20@email.com	1988-09-25

The address attributes were separated into a new table called Address. The *Person* table could have included the address information and still be in 2NF because the address attributes (*street*, city, *country* and *zipcode*) depend on the primary key of the *Person* table. However, separating the address information enhances the flexibility and scalability of the database design:

- 1. Even though it was not part of the user, this design avoids repeating the same address information for multiple people who might share the same address in the future and, therefore, reduces data redundancy.
- 2. This design allows for easier management and expansion as the system grows. This way, if any address information needs to be updated, the changes need to be made only once in the *Address* table rather than updating multiple records of the Person table when multiple users share the same address.
- 3. This separation also provides a cleaner and more modular database design.

```
CREATE TABLE `Address` (
    `id` int NOT NULL AUTO_INCREMENT,
    `person_id` int DEFAULT NULL,
    `street` varchar(255) NOT NULL,
    `city` varchar(255) NOT NULL,
    `country` varchar(255) NOT NULL,
    `zipcode` varchar(20) NOT NULL,
    PRIMARY KEY (`id`),
    KEY `person_id` (`person_id`),
    CONSTRAINT `Address_ibfk_1` FOREIGN KEY (`person_id`) REFERENCES `Person` (`id`)
)
```

* id int ◆♡	person_id ♦Ţ	* street varchar(255) ◆▽	* city varchar(255) ◆♡	* country varchar(255) ◆▽	* zipcode varchar(20) ◆ ▽
1		12 Maple St	London	England	E1 6AN
2	2	45 Oak Ave	Manchester	England	M1 2WD
3	3	89 Pine Rd	Birmingham	England	B1 1AB
4	4	23 Birch St	Edinburgh	Scotland	EH1 1YZ
5	5	67 Cedar Ln	Bristol	England	BS1 3XE
6	6	56 Elm St	Liverpool	England	L1 1AA
7	7	12 Maple St	Glasgow	Scotland	G1 2TF
8	8	89 Oak Dr	Leeds	England	LS1 3AB
9	9	123 Pine Rd	Newcastle	England	NE1 2AB
10	10	15 Elm St	Cardiff	Wales	CF10 3AF
11	11	78 Oak Ln	Sheffield	England	S1 4GT
12	12	56 Birch Rd	Nottingham	England	NG1 2PB
13	13	10 Holy St	Cardiff	Wales	CF10 2NF
14	14	34 Willow Rd	Edinburgh	Scotland	EH1 1AB
15	15	78 Cedar Ave	Cambridge	England	CB1 2SE
16	16	45 Maple Rd	Oxford	England	OX2 6TP
17	17	23 Birch Ave	Southampton	England	SO14 3HL
18	18	12 Elm Blvd	Leicester	England	LE1 3PL
19	19	56 Oak Rd	Norwich	England	NR1 4BE

The *Favourite* table only consists of the user's favourite information. This table was created with attributes as *id* for primary key, *type*, value, and *person_id* as a foreign key to the Person table. The Favourite attributes (favourite_books, *favoruite_drink*, *favourite_activity*) could have been placed in a *Favourite* table with separate columns for each type of favourite and still be in 2NF. However, creating a *Favourite* table with *type* and value columns enhances the flexibility and scalability of the database design:

- 1. By using a type-value structure, uniforms the data model. This way, the system can easily accommodate new types of favourites without altering the table schema, providing more dynamic expansion. For instance, adding a new favourite type called favourite destination only requires a new row in the table. This avoids the need to add a new column called *favourite_destination* to the table and update existing records with null values for this attribute.
- 2. This design simplifies queries to retrieve favourite information. Instead of querying multiple columns for different favourite types, the system can query a single table and filter by type.

```
CREATE TABLE `Favourite` (

`id` int NOT NULL AUTO_INCREMENT,

`type` varchar(255) NOT NULL,

`value` varchar(255) NOT NULL,

`person_id` int DEFAULT NULL,

PRIMARY KEY (`id`),

KEY `person_id` (`person_id`),

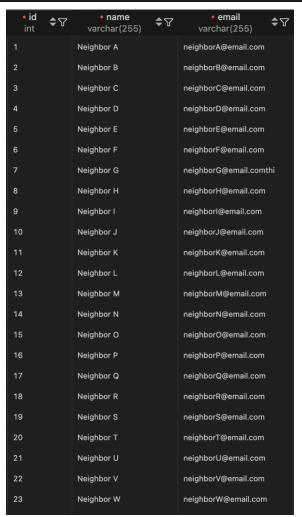
CONSTRAINT `Favourite_ibfk_1` FOREIGN KEY (`person_id`) REFERENCES `Person` (`id`)

)
```

* id int ◆∀	* type varchar(255) ♦ 🕆	* value varchar(255)	person_id int
1	Book	A New Beginning	1
2	Drink	Lemonade	1
3	Activity	Outdoor Running	1
4	Book	The Road to Success	2
5	Drink	Coffee	2
6	Activity	Hiking	2
7	Book	Endless Possibilities	3
8	Drink	Smoothie	3
9	Activity	Swimming	3
10	Book	Journey of Life	4
11	Drink	Iced Tea	4
12	Activity	Traveling	4
13	Book	The Adventure Continues	5
14	Drink	Green Tea	5
15	Activity	Gardening	5
16	Book	Finding Inner Peace	6
17	Drink	Coconut Water	6
18	Activity	Reading	6
19	Book	Exploring New Horizons	7

The Neighbour table contains the neighbour information:

```
CREATE TABLE 'Neighbour' (
  'id' int NOT NULL AUTO_INCREMENT,
  'name' varchar(255) NOT NULL,
  'email' varchar(255) NOT NULL,
  PRIMARY KEY ('id')
)
```



The *PersonNeighbour* table is a junction table that establishes a many-to-many relationship between the *Person* and *Neighbour* entities to eliminate partial dependency. Each row in the *PersonNeighbour*

table represents a unique relationship between a person and their neighbour. The *Neighbour* table could have included person_id as a foreign key to the *Person* table, but this would imply that each neighbour is associated with only one person, or their information would need to be repeated for each person, leading to data duplication. This does not accurately represent the real-world scenario where a neighbour can also be another person's neighbour even though the use case did not include any example. This design minimises data redundancy and maintains data integrity properly.

```
CREATE TABLE `PersonNeighbour`

`id` int NOT NULL AUTO_INCREMENT,

`person_id` int DEFAULT NULL,

`neighbour_id` int DEFAULT NULL,

PRIMARY KEY (`id`),

KEY `person_id` (`person_id`),

KEY `neighbour_id` (`neighbour_id`),

CONSTRAINT `PersonNeighbour_ibfk_1` FOREIGN KEY (`person_id`) REFERENCES `Person` (`id`),

CONSTRAINT `PersonNeighbour_ibfk_2` FOREIGN KEY (`neighbour_id`) REFERENCES `Neighbour` (`id`)

)
```

* id int ◆∀	person_id ♣√	neighbour_id
1	1	1
2	1	2
3	2	3
4	2	4
5	3	5
6	3	6
7	4	7
8	4	8
9	5	9
10	5	10
11	6	11
12	6	12
13	7	13
14	7	14
15	8	15
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17	9	17
18	9	18
19	10	19
20	10	20
21	11	21
22	11	22
23	12	23

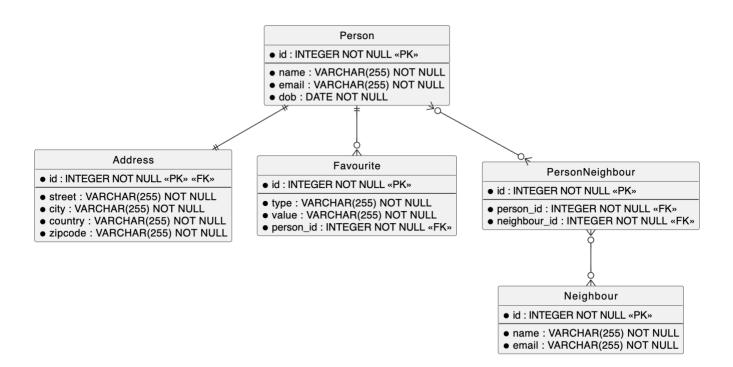
THIRD NORMAL FORM (3NF)

In Third Normal Form, all the attributes must be functionally dependent on the primary key, and there should be no transitive dependency, which means non-key attributes should not be dependent on other non-key attributes. By fulfilling the 2NF for each table, the design already satisfies the requirements of 3NF.

• The *Person* table has a one-to-many relationship with the *Favourite* table, represented by the foreign key *person_id* in the *Favourite* table, which means the one person can have 0 or more favourites.

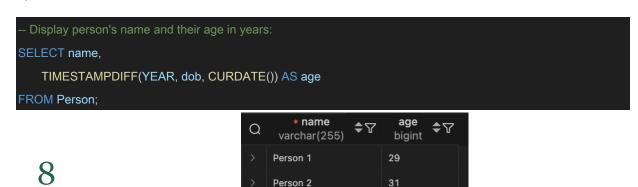
- The *Person* table has a one-to-one relationship with the *Address* table, represented by the foreign key *id* in the *Address* table, which means the one person can have only one address.
- The *Person* table has a many-to-many relationship with the *Neighbour* table, represented by *PersonNeighbour* junction table, which means one person can have 0 or more neighbours.

ENTITY RELATIONSHIP DIAGRAM



SQL QUERIES

1)



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Person 3

2)

```
-- Group Persons by their favourite drink and return average age of each group

SELECT f.value AS favourite_drink,

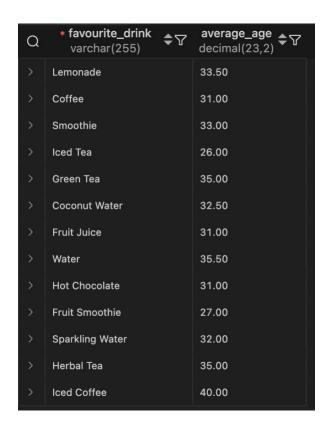
ROUND(AVG(TIMESTAMPDIFF(YEAR, p.dob, CURDATE())), 2) AS average_age

FROM Person p

JOIN Favourite f ON p.id = f.person_id

WHERE f.type = 'Drink'

GROUP BY f.value;
```



3)

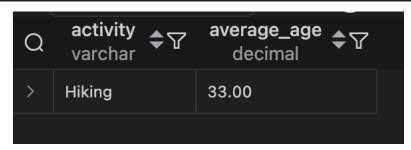
-- Display average age of people who likes Hiking

SELECT f.value as activity, ROUND(AVG(TIMESTAMPDIFF(YEAR, dob, CURDATE())), 2) AS average_age

FROM Person p

JOIN Favourite f ON p.id = f.person_id

WHERE f.type = 'Activity' AND f.value = 'Hiking';



4)

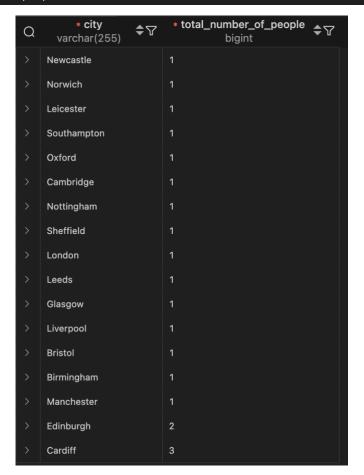
-- Display the total number of people from each City and sort it in ascending order by total number of people

SELECT a.city, COUNT(p.id) AS total_number_of_people

FROM Person p

JOIN Address a ON p.id = a.person_id

ORDER BY total_number_of_people ASC;



5)

-- Display name of person(s) whose neighbour is neighbour C

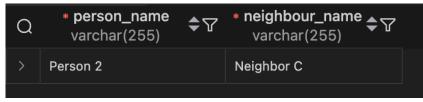
SELECT p.name as person_name, n.name as neighbour_name

FROM Person p

JOIN PersonNeighbour pn ON p.id = pn.person_id

JOIN Neighbour n ON pn.neighbour_id = n.id

WHERE n.name = 'Neighbor C';



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