

Presented by

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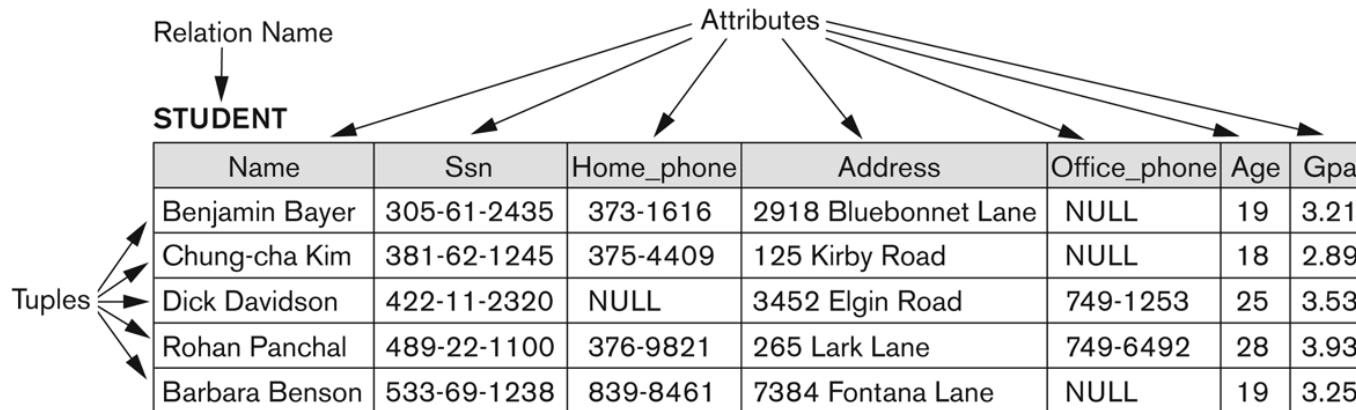
Advanced Databases UFCFU3-15-3

Relational databases, ER Modelling, and Normalization

Introduction - Database Management System

- A database management system
 - is a software system that enables users to define, create, maintain, and control access to the database
 - provides a systematic way to store, manage, and retrieve data
- In this module, we will look at both relational and no SQL databases

Relational Model



- All data is logically structured within relations (tables).
- Each relation has a name and is made up of named attributes (columns) of data.
- Each tuple (row) contains one value per attribute, representing a record
- Allows a high degree of data independence
- Provides substantial grounds for dealing with data semantics, consistency, and redundancy problems
- Enables the expansion of set-oriented data manipulation languages

Advantages of Relational Model

- Simplicity
- Flexibility
- Data Integrity
- Support for SQL
- Normalization

Database Design

- Conceptual

- high-level representation of the data, usually using an Entity-Relationship Diagram (ERD)

- Logical

- the process of transforming (or mapping) a conceptual schema of the application domain into a schema for the data model underlying a particular DBMS, e.g. relational

- Physical

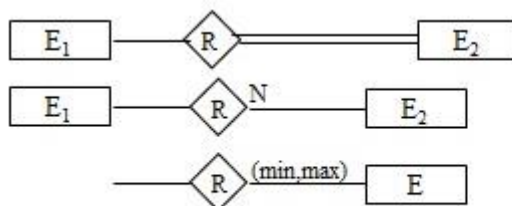
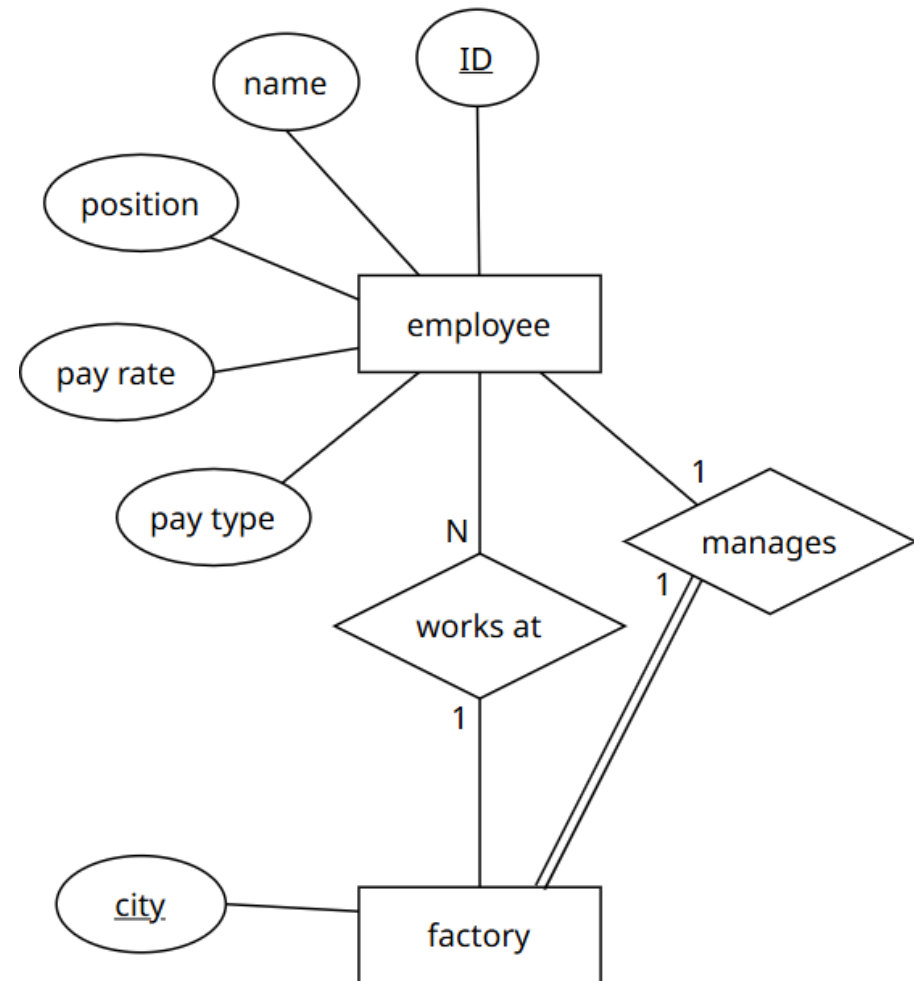
- the process of implementing a database's logical design on secondary storage using a specific DBMS

Entity Relationship Model (ER Model)

- Conceptual framework used for designing and structuring databases
- Key components of ERD:
 - Entity
 - Strong – typically has Primary key
 - Weak
 - Attributes
 - Simple
 - Composite
 - Derived
 - Multi-valued
 - Relationships and Cardinality
 - One to one (1:1)
 - One to Many (1:N)
 - Many to Many (M:N)
 - Participation constraints
 - Total
 - Partial

Entity Relationship Diagram Notations

Symbol	Meaning
	ENTITY TYPE
	WEAK ENTITY TYPE
	RELATIONSHIP TYPE
	IDENTIFYING RELATIONSHIP TYPE
	ATTRIBUTE
	KEY ATTRIBUTE
	MULTIVALUED ATTRIBUTE
	COMPOSITE ATTRIBUTE
	DERIVED ATTRIBUTE



TOTAL PARTICIPATION OF E_2 IN R

CARDINALITY RATIO 1:N FOR $E_1:E_2$ IN R

STRUCTURAL CONSTRAINT (min,max) ON PARTICIPATION OF E IN R

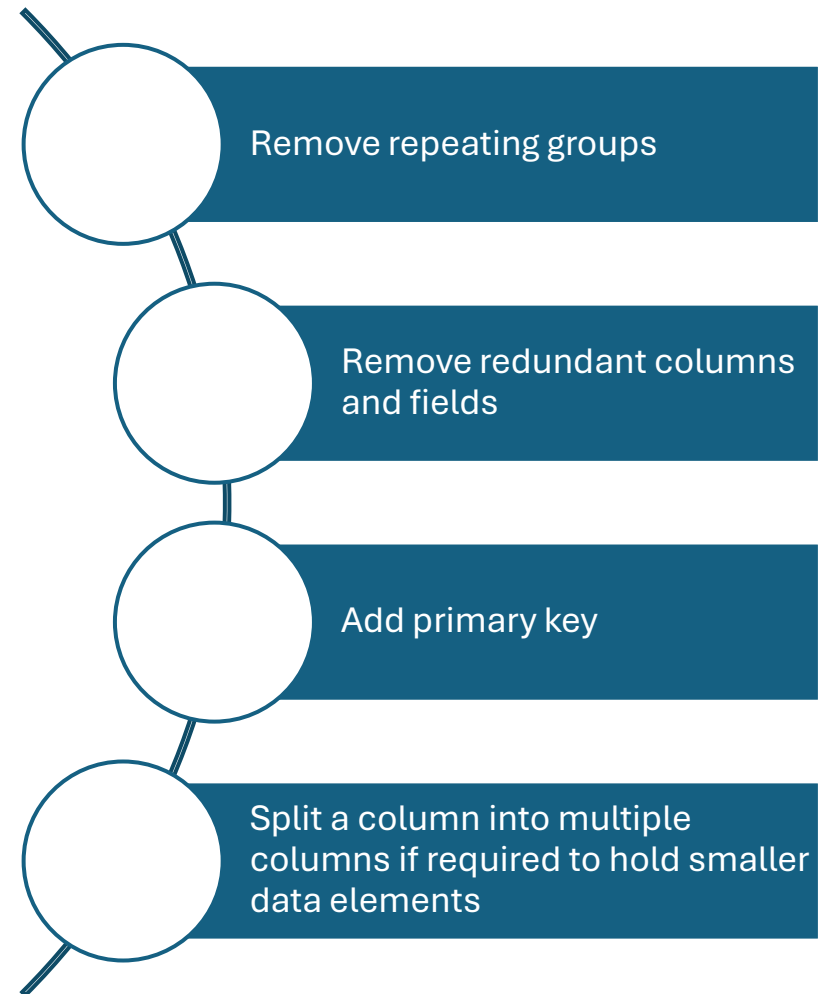
Normalization

- A technique for producing a set of suitable relations that support the data requirements of an enterprise.
 - Reduces redundancy
 - Minimizes anomalies (Data inconsistencies)
 - Improves data integrity

First Normal Form (1NF)

How to convert from UNF to 1NF?

- No repeating group
- A repeating group is an attribute, or group of attributes, within a table that occurs with multiple values for a single occurrence of the nominated key attribute(s) for that table.
- Atomic data
- A relation in which the intersection of each row and column contains one and only one value.
- Primary key
 - Each record must have a unique identifier



Example: UNF to 1 NF

- The repeating group in the unnormalized table as the property rented details, which repeats for each client.
- The structure of the repeating group is: Repeating Group = (propertyNo, pAddress, rentStart, rentFinish, rent, ownerNo, oName)
- Multiple values at the intersection of certain rows and columns. For example, there are two values for propertyNo(PG4 and PG16) for the client named John Kay.
- To transform an unnormalized table into 1NF, we ensure that there is a single value at the intersection of each row and column.

ClientRental

clientNo	cName	propertyNo	pAddress	rentStart	rentFinish	rent	ownerNo	oName
CR76	John Kay	PG4	6 Lawrence St, Glasgow	1-Jul-07	31-Aug-08	350	CO40	Tina Murphy
		PG16	5 Novar Dr, Glasgow	1-Sep-08	1-Sep-09	450	CO93	Tony Shaw
CR56	Aline Stewart	PG4	6 Lawrence St, Glasgow	1-Sep-06	10-June-07	350	CO40	Tina Murphy
		PG36	2 Manor Rd, Glasgow	10-Oct-07	1-Dec-08	375	CO93	Tony Shaw
		PG16	5 Novar Dr, Glasgow	1-Nov-09	10-Aug-10	450	CO93	Tony Shaw

ClientRental

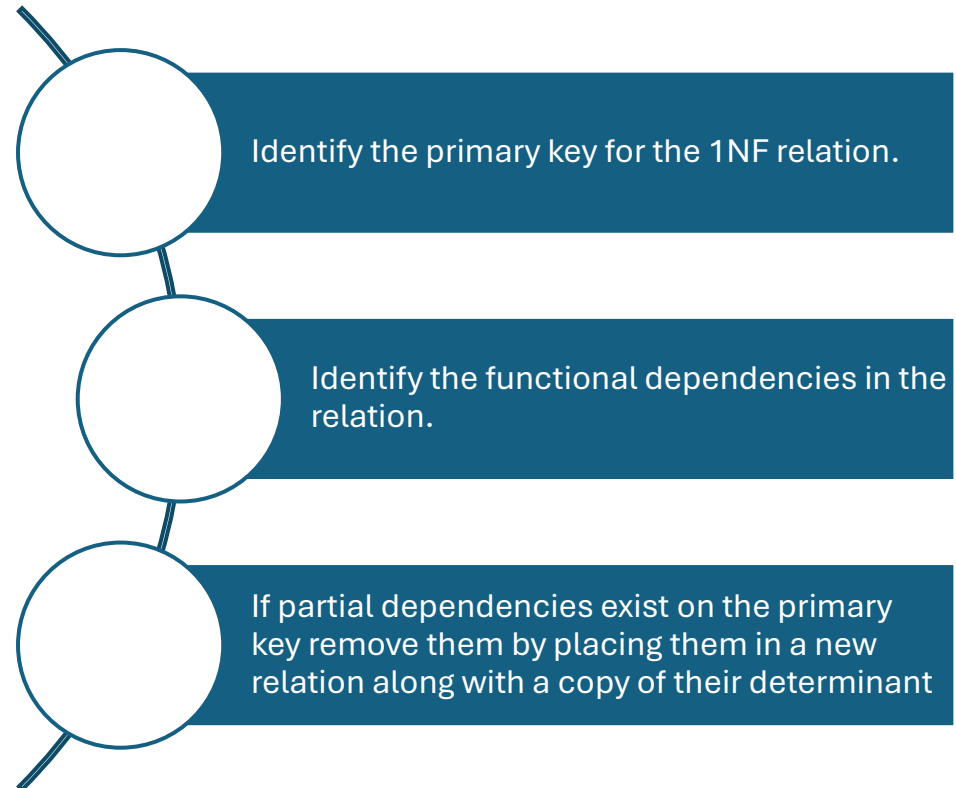
clientNo	propertyNo	cName	pAddress	rentStart	rentFinish	rent	ownerNo	oName
CR76	PG4	John Kay	6 Lawrence St, Glasgow	1-Jul-07	31-Aug-08	350	CO40	Tina Murphy
CR76	PG16	John Kay	5 Novar Dr, Glasgow	1-Sep-08	1-Sep-09	450	CO93	Tony Shaw
CR56	PG4	Aline Stewart	6 Lawrence St, Glasgow	1-Sep-06	10-Jun-07	350	CO40	Tina Murphy
CR56	PG36	Aline Stewart	2 Manor Rd, Glasgow	10-Oct-07	1-Dec-08	375	CO93	Tony Shaw
CR56	PG16	Aline Stewart	5 Novar Dr, Glasgow	1-Nov-09	10-Aug-10	450	CO93	Tony Shaw

Second Normal Form (2NF)

A relation that is in first normal form and every non-primary-key attribute is fully functionally dependent on any candidate key.

- A relation in 2NF **must be**
 - in 1NF and
 - every non-primary-key attribute is fully functionally dependent on the primary key.

How to convert from 1NF to 2NF?



1NF to 2NF

- This results in the creation of three new relations called Client, Rental, and PropertyOwner,
- These three relations are in second normal form, as every non-primary-key attribute is fully functionally dependent on the primary key of the relation.

Client

clientNo	cName
CR76	John Kay
CR56	Aline Stewart

Rental

clientNo	propertyNo	rentStart	rentFinish
CR76	PG4	1-Jul-07	31-Aug-08
CR76	PG16	1-Sep-08	1-Sep-07
CR56	PG4	1-Sep-06	10-Jun-07
CR56	PG36	10-Oct-07	1-Dec-08
CR56	PG16	1-Nov-09	10-Aug-10

PropertyOwner

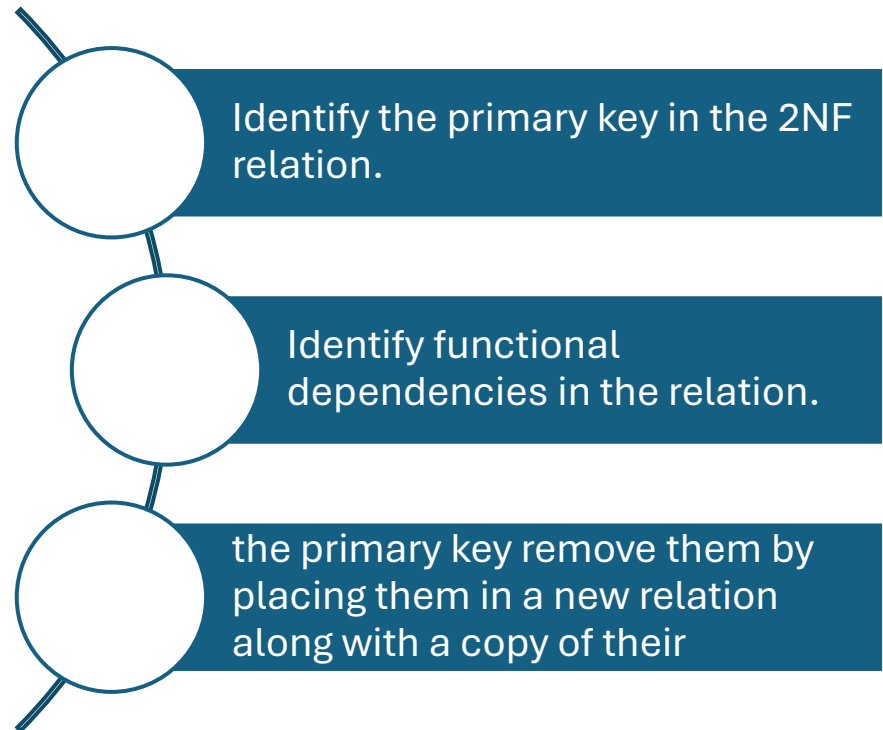
propertyNo	pAddress	rent	ownerNo	oName
PG4	6 Lawrence St, Glasgow	350	CO40	Tina Murphy
PG16	5 Novar Dr, Glasgow	450	CO93	Tony Shaw
PG36	2 Manor Rd, Glasgow	375	CO93	Tony Shaw

Third Normal Form (3NF)

A relation that is in first and second normal form and in which no non-primary-key attribute is transitively dependent on any candidate key

- Based on the concept of transitive dependency.
- Transitive Dependency is a condition where
 - A, B and C are attributes of a relation such that if $A \rightarrow B$ and $B \rightarrow C$,
 - then C is transitively dependent on A through B. (Provided that A is not functionally dependent on B or C).

How to convert from 2NF to 3NF?



2NF to 3NF

- To transform the PropertyOwnerrelation into 3NF, we must first remove this transitive dependency by creating two new relations called **PropertyForRent** and **Owner**

PropertyForRent

propertyNo	pAddress	rent	ownerNo
PG4	6 Lawrence St, Glasgow	350	CO40
PG16	5 Novar Dr, Glasgow	450	CO93
PG36	2 Manor Rd, Glasgow	375	CO93

Owner

ownerNo	oName
CO40	Tina Murphy
CO93	Tony Shaw

Passive DBMSs vs Active DBMSs

- Conventional DBMSs are **passive**: they execute operations only upon explicit requests
- Often, however, there is the need of **reactive** capabilities: the DBMS autonomously reacts to some events and executes specified operations
- We refer to them as **activeDBMS (ADBMS)** for which we can specify *active rules*, also called **triggers**

Applications for Active Databases

- Logging/tracking events
- Security
- Notification
 - Automatic notification when certain condition occurs
- Enforcing business constraints/rules
- Maintenance of derived data
 - Automatically update derived data and avoid anomalies due to redundancy e.g.trigger to update the Total salary in a department
- Maintain replicated tables

Triggers

- A trigger is a statement that is executed automatically by the DBMS as a side effect of a modification to the database
- A trigger describes an action the database should take under certain conditions when some database-related event (such as inserts, updates, deletes) occurs.

Event-Condition-Action (ECA) Model

- The model that has been used to specify active database rules is referred to as Event-Condition-Action (ECA) model
- **ON** *event* **IF** *condition* **THEN** *action*
- **Event(s)** that trigger the rule to be checked: e.g. insert, delete, update
- **Condition** that determines whether the action should be executed: A logical expression (true/false)
 - Optional: If no condition is specified then condition is always true
- **Action:** Sequence of SQL statements and code to be executed

Triggers –An Example

```

create trigger overdraft_trigger
after update on account
referencing new row as nrow
for each row
    when: nrow.balance < 0
    begin atomic
        insert into borrower
        (select customer_name, account_number From depositor
        where: nrow.account_number= depositor.account_number);

        insert into loan values
        (:nrow.account_number, :nrow.branch_name, :nrow.balance);
        update account set balance = 0
        where account.account_number= :nrow.account_number
    end

```

Resources – some helpful links for further learning

- Reading list: [UFCFU3-15-3 Advanced Databases | UWE Bristol](#)
- https://www.w3schools.com/mysql/mysql_rdbms.asp
- <https://cloud.google.com/learn/what-is-a-relational-database>
- <https://learn.microsoft.com/en-us/office/troubleshoot/access/database-normalization-description>
- Normalization: <https://www.youtube.com/watch?v=Yp82NgeQZ9o>