

# STAT 321: Assignment 1

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## Problem 1

$$\mathbb{P}(\text{At least 2 students share a birthday}) = 1 - \mathbb{P}(\text{No students share a birthday}) \quad (1)$$

$$\text{Total number of students} = 20$$

Let us assume there are 365 days in a year and that all birthdays are equally likely. Then,

$$\begin{aligned} \mathbb{P}(\text{No students share a birthday}) &= \frac{365}{365} \cdot \frac{365-1}{365} \cdot \frac{365-2}{365} \cdots \frac{365-20}{365} \\ &= \frac{365}{365} \cdot \frac{364}{365} \cdot \frac{363}{365} \cdots \frac{345}{365} \\ &= \frac{365!}{345! \cdot 365^{20}} \\ &\approx 0.5886 \end{aligned}$$

Putting our result back in (1),

$$\mathbb{P}(\text{At least 2 students share a birthday}) = 1 - 0.5886 = 0.4114$$

So the probability of at least two students sharing a birthday is  $\approx 0.4114$

## Problem 2

### (a) FALSE

Let us assume  $A \subset B$  such that  $P(A) = 0.5$  and  $P(B) = 0.6$  and  $P(A \cap B) = 0.5$ .

Now  $P(A) \cdot P(B) = 0.3$ .

So  $P(A \cap B) = 0.5 \leq P(A) \cdot P(B) = 0.3$  does not hold.

Then  $P(A \cap B)$  is not less than or equal to  $P(A)P(B)$ .

### (b) FALSE

Let  $A$  be any the events in sample space  $\Omega$  and  $P_E(A) = P(A \cap E)$ . We know that  $0 < P(E) < 1$ .

$E \in \Omega$  so  $P_E(\Omega) = P(\Omega \cap E) = P(E)$

But  $P(E) < 1$ , which violates the Second Axiom of Probability, that  $P(\Omega) = 1$

Thus,  $P_E$  does not satisfy the axioms of probability.

**(c) TRUE**

$$\begin{aligned}\mathbb{P}(A \cap B^C) &= \mathbb{P}(A - B) \\ &= \mathbb{P}(A) - \mathbb{P}(A \cap B)\end{aligned}$$

And,

$$\begin{aligned}\mathbb{P}(A \cup B) - \mathbb{P}(B) &= \mathbb{P}(A) + \mathbb{P}(B) - \mathbb{P}(A \cap B) - \mathbb{P}(B) \\ &= \mathbb{P}(A) - \mathbb{P}(A \cap B)\end{aligned}$$

Thus,

$$\mathbb{P}(A \cap B^C) = \mathbb{P}(A \cup B) - \mathbb{P}(B)$$

**(d) TRUE**

We have to prove,

$$P(\cap_{i=1}^n A_i) \leq \min_i [P(A_i)] \quad (1)$$

Let  $n = 1$ . Then,

$$P(\cap_{i=1}^1 A_i) = P(A_1) \leq \min_i [P(A_i)]$$

Let  $n = 2$ . Then,

$$P(\cap_{i=1}^2 A_i) = P(A_1 \cap A_2)$$

If  $P(A) > P(B)$ , we know that  $P(A \cap B) \leq P(B)$ . So,

$$P(\cap_{i=1}^2 A_i) \leq \min_i [P(A_i)]$$

Now, let  $n = k$  be true for (1),

$$P(\cap_{i=1}^k A_i) \leq \min_i [P(A_i)] \quad (2)$$

For  $n = k + 1$ , we know that  $P(\cap_{i=1}^k A_i \cap A_{k+1}) \leq P(A_{k+1})$

So

$$P(\cap_{i=1}^{k+1} A_i) \leq \min_i [P(A_i)]$$

Thus, by induction Equation (1) holds true.

## Problem 3

**(a)**

Computer 1 is connected to the network by 3 independent connections, each of which have a probability  $p$  of being broken at any time. For all these 3 connections to be broken,

$$\mathbb{P}(\text{All 3 connections are broken}) = \mathbb{P}(\text{Connector } n \text{ being broken})^3$$

where  $n$  is any of the connection between computer 1 and the others. So,

$$\mathbb{P}(\text{All 3 connections are broken}) = p^3$$

(b)

For  $j = 2$ ,

There are 2 connections shared between Computer 1 and Computer 2. Therefore, for both Computer 1 and Computer 2 to be isolated, 4 wires will have to be broken. So,

$$\mathbb{P}(4 \text{ connections are broken}) = p^4$$

For  $j = 3$ , There are no connections shared between Computer 1 and Computer 3. Therefore, for both Computer 1 and Computer 3 to be isolated, all 6 wires will have to be broken. So,

$$\mathbb{P}(\text{All connections are broken}) = p^6$$

For  $j = 4$ , There is only 1 connection shared between Computer 1 and Computer 4. Therefore, for both Computer 1 and Computer 4 to be isolated, 5 wires will have to be broken. So,

$$\mathbb{P}(5 \text{ connections are broken}) = p^5$$

(c)

$$\mathbb{P}(\text{No computer is isolated}) = 1 - \mathbb{P}(\text{At least 1 computer is isolated})$$

Now,

$$\begin{aligned} \mathbb{P}(\text{At least 1 computer is isolated}) &= \mathbb{P}(1 \text{ computer is isolated} \\ &\quad \cup 2 \text{ computers are isolated} \cup 3 \text{ computer are isolated}) \\ &\quad \cup \text{All computers are isolated}) \\ &= \sum_{1 \leq i \leq 4} \mathbb{P}(A_i) - \sum_{1 \leq i \leq j \leq 4} \mathbb{P}(A_i \cap A_j) + \sum_{1 \leq i \leq j \leq l \leq 4} \mathbb{P}(A_i \cap A_j \cap A_l) \\ &\quad - \sum_{1 \leq i \leq j \leq l \leq h \leq 4} \mathbb{P}(A_i \cap A_j \cap A_l \cap A_h) \end{aligned}$$

where each  $A_x (x = i, j, l, h)$  is the event that a computer is isolated.

$$\text{So, } \sum_{1 \leq i \leq 4} \mathbb{P}(A_i) = 4p^3$$

$$\text{And, } \sum_{1 \leq i \leq j \leq 4} \mathbb{P}(A_i \cap A_j) = p^4 + p^6 + p^5 + p^5 + p^6 + p^4 =$$

$$\text{And, } + \sum_{1 \leq i \leq j \leq l \leq 4} \mathbb{P}(A_i \cap A_j \cap A_l) = p^6 + p^6 + p^6 + p^6$$

$$\text{And, } \sum_{1 \leq i \leq j \leq l \leq h \leq 4} \mathbb{P}(A_i \cap A_j \cap A_l \cap A_h) = p^6$$

Now,

$$\begin{aligned} \mathbb{P}(\text{At least 1 computer is isolated}) &= 4p^3 - (2p^4 + 2p^5 + 2p^6) + (4p^6) - (p^6) \\ &= 4p^3 - 2p^4 - 2p^5 + p^6 \end{aligned}$$

$$\mathbb{P}(\text{No computer is isolated}) = 1 - (4p^3 - 2p^4 - 2p^5 + p^6)$$

## Problem 4

(a)

$$\begin{aligned}\mathbb{P}(\text{Signal from a fish will be received}) &= \mathbb{P}(A \cup B) \\ &= \mathbb{P}(A) + \mathbb{P}(B) - \mathbb{P}(A \cap B)\end{aligned}$$

Since A and B are independent receivers,

$$\mathbb{P}(A \cap B) = \mathbb{P}(A) \cdot \mathbb{P}(B)$$

$$\mathbb{P}(A \cap B) = 0.8 \cdot 0.9 = 0.72$$

$$\text{And, } \mathbb{P}(A \cup B) = 0.8 + 0.9 - 0.72 = 0.98$$

So,

$$\mathbb{P}(\text{Signal from a fish will be received}) = 0.98$$

(b)

$$\begin{aligned}\mathbb{P}(\text{Signal from a fish will be received by only 1 receiver}) &= \mathbb{P}(A \Delta B) \\ &= \mathbb{P}[(A \cap B^C) \cup (A^C \cap B)] \\ &= \mathbb{P}(A \cap B^C) + \mathbb{P}(A^C \cap B) - \mathbb{P}(A \cap A^C \cap B \cap B^C)\end{aligned}$$

We know that  $\mathbb{P}(A \cap A^C \cap B \cap B^C) = 0$ .

So,

$$\begin{aligned}\mathbb{P}(\text{Signal from a fish will be received by only 1 receiver}) &= \mathbb{P}(A \cap B^C) + \mathbb{P}(A^C \cap B) \\ &= \mathbb{P}(A) - \mathbb{P}(A \cap B) + \mathbb{P}(B) - \mathbb{P}(A \cap B) \\ &= 0.8 - 0.72 + 0.9 - 0.72\end{aligned}$$

So,

$$\mathbb{P}(\text{Signal from a fish will be received by only 1 receiver}) = 0.26$$

## Problem 5

(a)

For a Monte Carlo Approximation, we can simulate a sample of 20 8-sided die and counting the number of times we get  $i$  rolls greater than or equal to 6 and  $j$  rolls greater than or equal to 4. By counting these we can estimate the probability.

```
# Run number of simulations to estimate probability
number_sim <- 10000

# Add conditions to satisfy calculating probability for events
condition_1 = condition_2 <- c(0)

# Simulating n= number_sim rolls
for(k in 1:number_sim){
  sample_rolls <- sample(1:8, size = 20, replace = TRUE)

  # Counting the number of dice rolls greater than or equal to 6 in the sample
  i_check <- sum(sample_rolls[] >= 6)
  # Counting the number of dice rolls greater than or equal to 4 in the sample
  j_check <- sum(sample_rolls[] >= 4)

  # Counting the samples which satisfy conditions 1, that is there are at least i rolls greater than
  if(i_check >= i){
    condition_1 = condition_1 + 1
  }
  # Counting the samples which satisfy conditions 1, that is there are at least j rolls greater than
  if(j_check >= j){
    condition_2 = condition_2 + 1
  }
}
# Estimating conditional probability
prob <- condition_1/condition_2
```

(b) and (c)

We can create a function with inputs  $i$  and  $j$  to output the probabilities for the given  $i$  and  $j$ . Then we can run a loop to create a matrix for the probabilities for each  $i \in [12, 14]$  and  $j \in [8, 12]$ .

```
sim_dice <- function(i,j){
  number_sim <- 10000
  condition_1 = condition_2 <- c(0)
  for(k in 1:number_sim){
    sample_rolls <- sample(1:8, size = 20, replace = TRUE)
    if(sum(sample_rolls[] >= 6) >= i){
      condition_1 = condition_1 + 1
    }
    if(sum(sample_rolls[] >= 4) >= j){
      condition_2 = condition_2 + 1
    }
  }
}
```

```

prob <- condition_1/condition_2

return(prob)
}

output <- data.frame(row.names = c("j=8","j=9","j=10","j=11","j=12"))
for(m in 1:3){
  for(o in 1:5){
    output[o,m] <- sim_dice(m+11,o+7)
  }
}
colnames(output) <- c("i=12","i=13","i=14")

output

```

```

##           i=12           i=13           i=14
## j=8  0.03395654 0.01354630 0.003644094
## j=9  0.03656894 0.01108923 0.003408739
## j=10 0.03432171 0.01144788 0.004779492
## j=11 0.04113267 0.01381249 0.004506150
## j=12 0.05020313 0.01763928 0.004199855

```

The above are the estimated probabilities for every combination of *i*s and *j*s.