

## ADDRESSING MODES

- The different ways in which the location of an operand is specified in an instruction are referred to as **Addressing Modes** (Table 2.1).

**Table 2.1** Generic addressing modes

Name	Assembler syntax	Addressing function
Immediate	#Value	Operand = Value
Register	Ri	EA = Ri
Absolute (Direct)	LOC	EA = LOC
Indirect	(Ri) (LOC)	EA = [Ri] EA = [LOC]
Index	X(Ri)	EA = [Ri] + X
Base with index	Ri(Rj)	EA = [Ri] + [Rj]
Base with index and offset	X(Ri,Rj)	EA = [Ri] + [Rj] + X
Relative	X(PC)	EA = [PC] + X
Autoincrement	(Ri)+	EA = [Ri]; Increment Ri
Autodecrement	-(Ri)	Decrement Ri; EA = [Ri]

EA = effective address  
Value = a signed number

## IMPLEMENTATION OF VARIABLE AND CONSTANTS

- Variable** is represented by allocating a memory-location to hold its value.
- Thus, the value can be changed as needed using appropriate instructions.
- There are 2 accessing modes to access the variables:
  - 1) Register Mode
  - 2) Absolute Mode

### Register Mode

- The operand is the contents of a register.
- The name (or address)

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## ➤ ADDITIONAL ADDRESSING MODES

### 1) Auto Increment Mode

- Effective-address of operand is contents of a register specified in the instruction.
- After accessing the operand, the contents of this register are automatically incremented to point to the next item in a list.
- Implicitly, the increment amount is 1.
- This mode is denoted as
  - $[R_i]++$ ; where  $R_i$ =pointer-register.

### 1) Auto Decrement Mode

- The contents of a register specified in the instruction are first automatically decremented and are then used as the effective-address of the operand.
- This mode is denoted as
  - $-(R_i)$ ; where  $R_i$ =pointer-register.
- These 2 modes can be used together to implement an important data structure called a stack.

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Ex (1)

Add  $(R_2) + , R_0$

Initially

$R_2$	1040

$R_0$	10

103C	
1040	10
1044	20
1048	

$$\begin{aligned} (1040) &= 10 \\ \text{Add } (R_0) &= 10 \\ \hline R_0 &= 20H \end{aligned}$$

After execution

$R_2$	1044

$R_0$	20H

Sign in

Ex 27 Add  $-(R_2), R_0$

Initially

$R_2$	1040
-------	------

$R_0$	10
-------	----

103C	08
1040	10
1044	:
1048	:

$$\begin{aligned}(R_0) &= 10 \\ \text{Add } (103C) &= 08 \\ \hline R_0 &= 18\end{aligned}$$

(29)

After execution

$R_2$	103C	$R_0$	18
:	:	:	:

Registers R1 and R2 of a computer contains the decimal values 1200 and 4600. What is the effective- address of the memory operand in each of the following instructions?

- (a) Load 20(R1), R5
- (b) Move #3000, R5
- (c) Store R5, 30(R1, R2)
- (d) Add -(R2), R5
- (e) Subtract (R1)+, R5

# Basic Input/Output Operations



**ASSEMBLY LANGUAGE**

- We generally use symbolic-names to write a program.
- A complete set of symbolic-names and rules for their use constitute an **Assembly Language**.
- The set of rules for using the mnemonics in the specification of complete instructions and programs is called the **Syntax** of the language.
- Programs written in an assembly language can be automatically translated into a sequence of machine instructions by a program called an **Assembler**.
- The user program in its original alphanumeric text formal is called a **Source Program**, and the assembled machine language program is called an **Object Program**.

For example:

```
MOVE R0,SUM ;The term MOVE represents OP code for operation performed by instruction.  
ADD #5,R3 ;Adds number 5 to contents of register R3 & puts the result back into registerR3.
```

**ASSEMBLER DIRECTIVES**

- **Directives** are the assembler commands to the assembler concerning the program being assembled.
- These commands are not translated into machine opcode in the object-program.

	Memory address label	Operation	Addressing or data information
Assembler directives	SUM	EQU	200
		ORIGIN	204
	N	DATAWORD	100
	NUM1	RESERVE	400
		ORIGIN	100
Statements that	START	MOVE	N,R1



# I/O

- The data on which the instructions operate are not necessarily already stored in memory.
- Data need to be transferred between processor and outside world (disk, keyboard, etc.)
- I/O operations are essential, the way they are performed can have a significant effect on the performance of the computer.



# Program-Controlled I/O Example

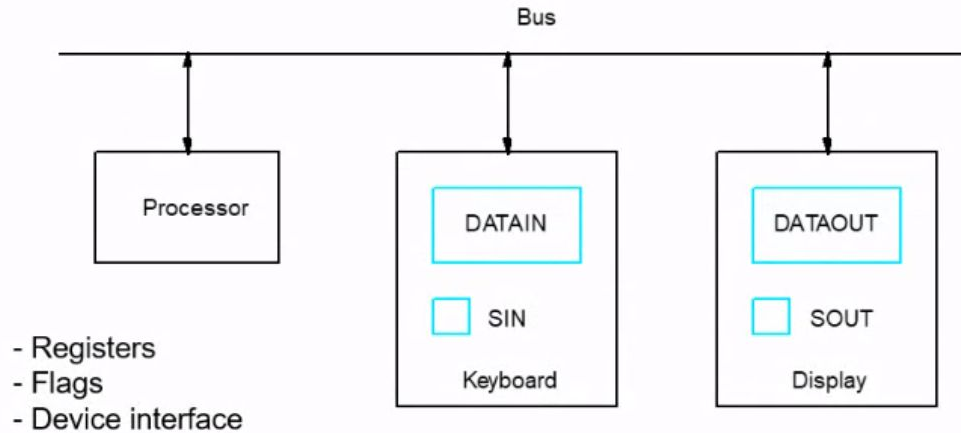


Figure 2.19 Bus connection for processor, keyboard, and display

# Program-Controlled I/O Example

- Machine instructions that can check the state of the status flags and transfer data:

READWAIT Branch to READWAIT if  $SIN = 0$   
Input from DATAIN to R1

WRITEWAIT Branch to WRITEWAIT if  $SOUT = 0$   
Output from R1 to DATAOUT

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