

DECIMATION IN TIME FAST FOURIER TRANSFORM

(DIT - FFT)

In the following presentation, the number of samples are assumed as power of 2. i.e., $N = 2^v$, where $v \rightarrow$ fixed integer.

The decimation in time approach is one of breaking N -point transform into 2 $N/2$ transforms, then breaking each $N/2$ point transform into $N/4$ point transforms and continuing this decimation process until 2-point transforms are obtained. This technique is known as divide and conquer approach.

Let $x(n) = x(0), x(1), x(2), \dots, x(N-1)$.

Even indexed sequence : $x(0), x(2), \dots, x(N-2)$

odd indexed sequence : $x(1), x(3), \dots, x(N-1)$

The N -point DFT of $x(n)$ is

$$X(K) = \sum_{n=0}^{N-1} x(n) w_N^{kn}$$

$$\Rightarrow X(K) = \sum_{\substack{n=0 \\ n, \text{ even}}}^{N-2} x(n) \cdot w_N^{kn} + \sum_{\substack{n=1 \\ n, \text{ odd}}}^{N-1} x(n) \cdot w_N^{kn}$$

for the first decimation, put $n = 2r$ in the first summation and $n = 2r + 1$ in the second summation. This gives -

$$X(K) = \sum_{r=0}^{N/2-1} x(2r) \omega_N^{K \cdot 2r} + \sum_{r=0}^{N/2-1} x(2r+1) \omega_N^{K(2r+1)}$$

$$= \sum_{r=0}^{N/2-1} x(2r) \omega_N^{K \cdot 2r} + \omega_N^K \sum_{r=0}^{N/2-1} x(2r+1) \omega_N^{K \cdot 2r}$$

Since $\omega_N^{K \cdot 2r} = e^{j \frac{2\pi}{N} \cdot K \cdot 2r} = e^{j \frac{2\pi}{N/2} \cdot K \cdot r} = \omega_{N/2}^{K \cdot r}$

the above equation can be written as -

$$X(K) = \sum_{r=0}^{N/2-1} x(2r) \omega_{N/2}^{K \cdot r} + \omega_N^K \sum_{r=0}^{N/2-1} x(2r+1) \omega_{N/2}^{K \cdot r}$$

$$= G(K) + \omega_N^K H(K)$$

$$K = 0, 1, 2, \dots, \frac{N}{2} - 1.$$

where $G(K)$ and $H(K)$ are $N/2$ -point DFTs, are of even indexed and odd indexed sequences respectively for computing $x(K)$ for $K = \frac{N}{2}, \frac{N}{2} + 1, \dots, N-1$, the periodicity of $G(K)$ and $H(K)$ are exploited. It may be noted that $G(K)$ and $H(K)$ are periodic with a period equal to $N/2$. Thus we can write

$$X(K) = \begin{cases} G(K) + \omega_N^K H(K), & K = 0, 1, 2, \dots, \frac{N}{2} - 1. \\ G(K + \frac{N}{2}) + \omega_N^K H(K + \frac{N}{2}), & K = \frac{N}{2}, \frac{N}{2} + 1, \dots, N-1. \end{cases}$$

for $N=8$, the above equation becomes -

$$X(K) = \begin{cases} G(K) + W_8^K H(K) & , K = 0, 1, 2, 3. \\ G(K+4) + W_8^K H(K+4) & , K = 4, 5, 6, 7. \end{cases}$$

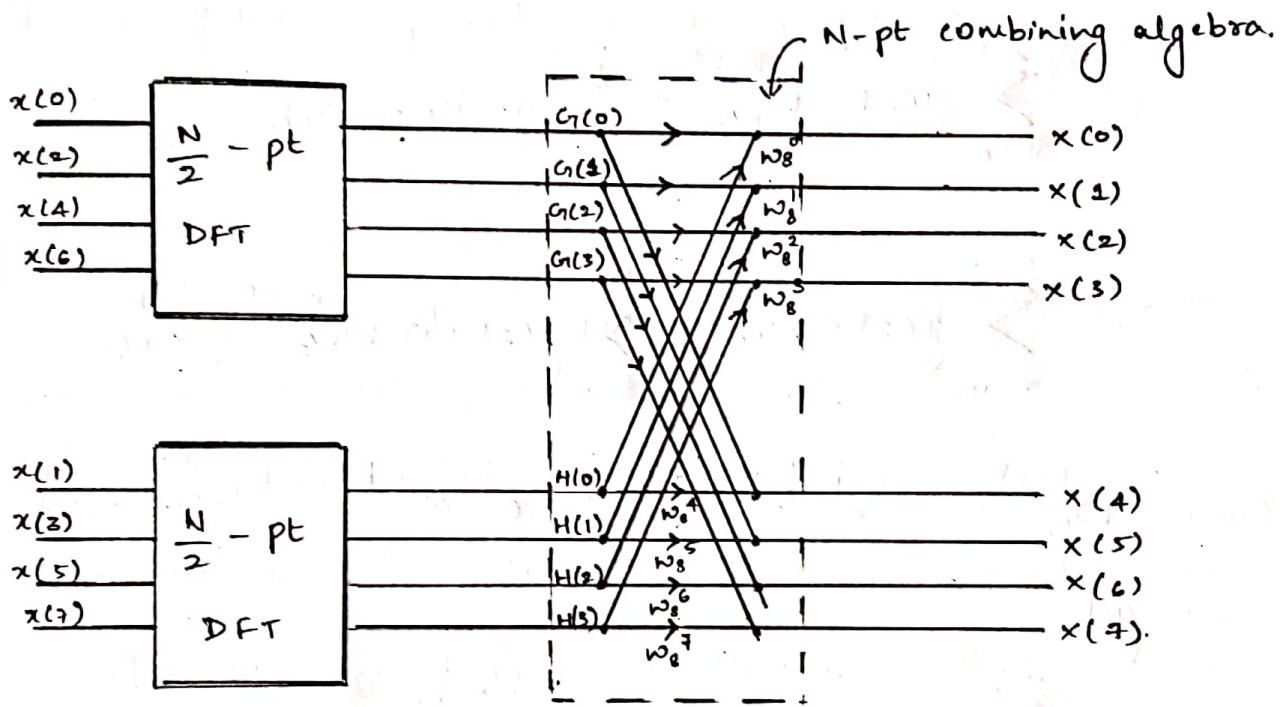


fig 1. Signal flow diagram after first decimation.

Total number of complex multiplications after first decimation is given by -

$$M_1 = \left(\frac{N}{2}\right)^2 + \left(\frac{N}{2}\right)^2 + N$$

The first two terms account for the computation of $N/2$ point DFT's while the last term accounts for N -point combining algebra.

Each of the $N/2$ - point sequences are further decimated into sequences of length equal to $N/4$.

$$G(k) = \sum_{l=0}^{N/2-1} g(l) \omega_{N/2}^{kl}$$

$$= \sum_{\gamma=0}^{\frac{N}{4}-1} g(2\gamma) \omega_{N/2}^{k \cdot 2\gamma} + \sum_{\gamma=0}^{\frac{N}{4}-1} g(2\gamma+1) \omega_{N/2}^{k(2\gamma+1)}$$

$$= \sum_{\gamma=0}^{\frac{N}{4}-1} g(2\gamma) \omega_{N/2}^{k \cdot 2\gamma} + \sum_{\gamma=0}^{\frac{N}{4}-1} g(2\gamma+1) \omega_{N/2}^{k \cdot 2\gamma} \cdot \omega_{N/2}^k$$

Since $\omega_{N/2}^{k \cdot 2\gamma} = e^{j \frac{2\pi}{N/2} \cdot k \cdot 2\gamma} = e^{j \frac{2\pi}{N/4} \cdot k \gamma} = \omega_{N/4}^{k\gamma}$,

we get

$$G(k) = \sum_{\gamma=0}^{\frac{N}{4}-1} g(2\gamma) \cdot \omega_{N/4}^{k\gamma} + \omega_{N/2}^k \sum_{\gamma=0}^{\frac{N}{4}-1} g(2\gamma+1) \cdot \omega_{N/4}^{k\gamma}$$

$$G(k) = A(k) + \omega_{N/2}^k B(k).$$

Since $A(k)$ and $B(k)$ are periodic with a period equal to $N/4$, we can write -

$$G(k) = \begin{cases} A(k) + \omega_{N/2}^k B(k) & k = 0, 1, \dots, \frac{N}{4}-1 \\ A(k + \frac{N}{4}) + \omega_{N/2}^k B(k + \frac{N}{4}) & k = \frac{N}{4}, \frac{N}{4}+1, \dots, \frac{N}{2}-1 \end{cases}$$

Similarly, we can write -

$$H(K) = \begin{cases} C(K) + W_{N/2}^K D(K) & K = 0, 1, \dots, \frac{N}{4} - 1. \\ C(K + \frac{N}{4}) + W_{N/2}^K D(K + \frac{N}{4}) & K = \frac{N}{4}, \frac{N}{4} + 1, \dots, \frac{N}{2} - 1. \end{cases}$$

The above equations for $N=8$ become the following -

$$G(K) = \begin{cases} A(K) + W_4^K B(K) & K = 0, 1. \\ A(K+2) + W_4^K B(K+2) & K = 2, 3. \end{cases}$$

$$H(K) = \begin{cases} C(K) + W_4^K D(K) & K = 0, 1. \\ C(K+2) + W_4^K D(K+2) & K = 2, 3. \end{cases}$$

Continuing this process, each $\frac{N}{4}$ points transformation is broken into 2 $\frac{N}{8}$ point transforms.

Since $N = 2^p$, this process can be continued until there are $\log_2 N$ stages. It may be noted that in each stage there are $N/2$ butterflies and each butterfly has 2 complex multiplications. Therefore, after final decimation, we have -

$$2 \times \frac{N}{2} \times \log_2 N = N \log_2 N \text{ complex multiplications.}$$

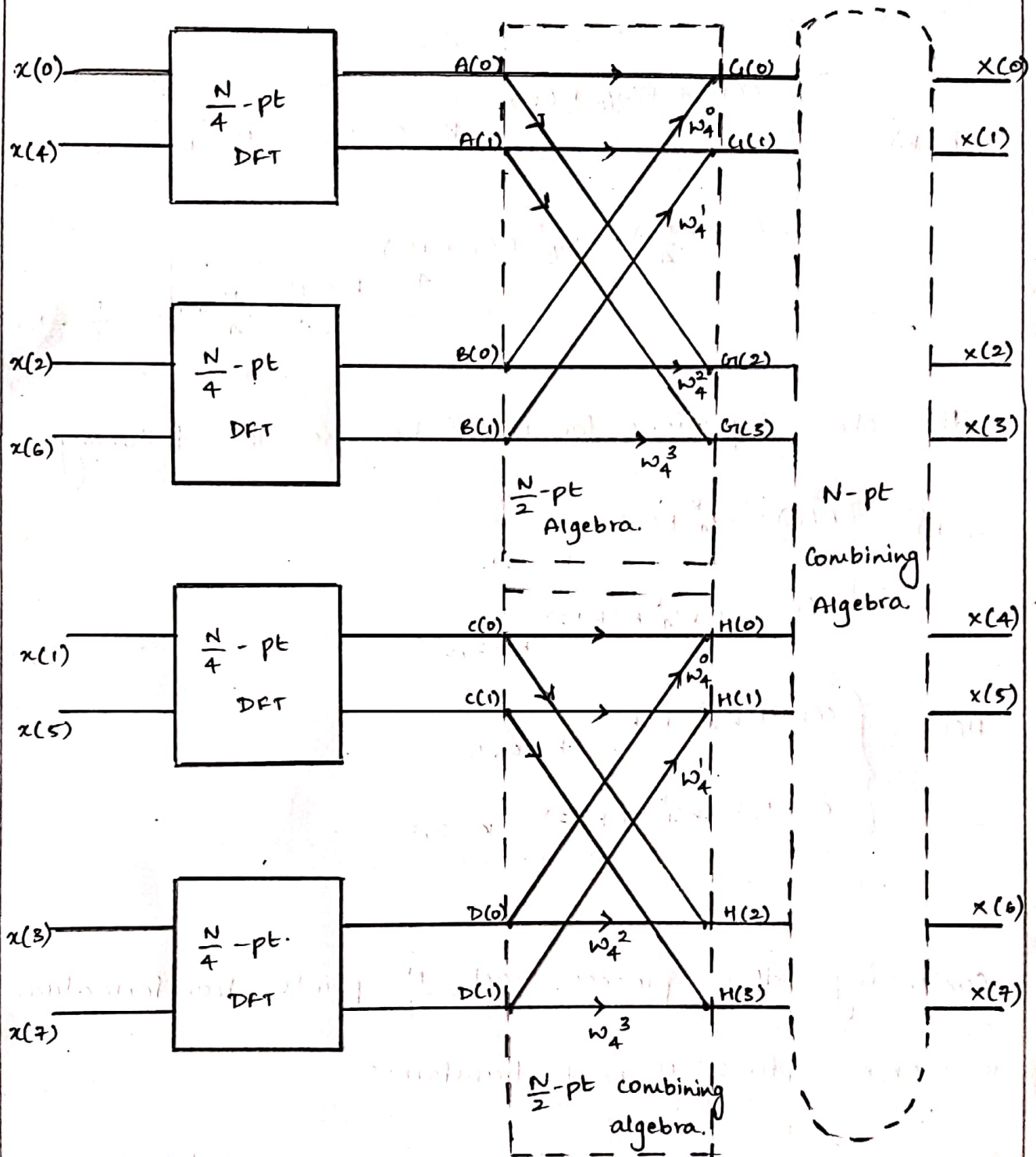


fig 2 : Signal flow diagram after second decimation.

The total signal flow diagram after final decimation is as shown below. -

Position Index.

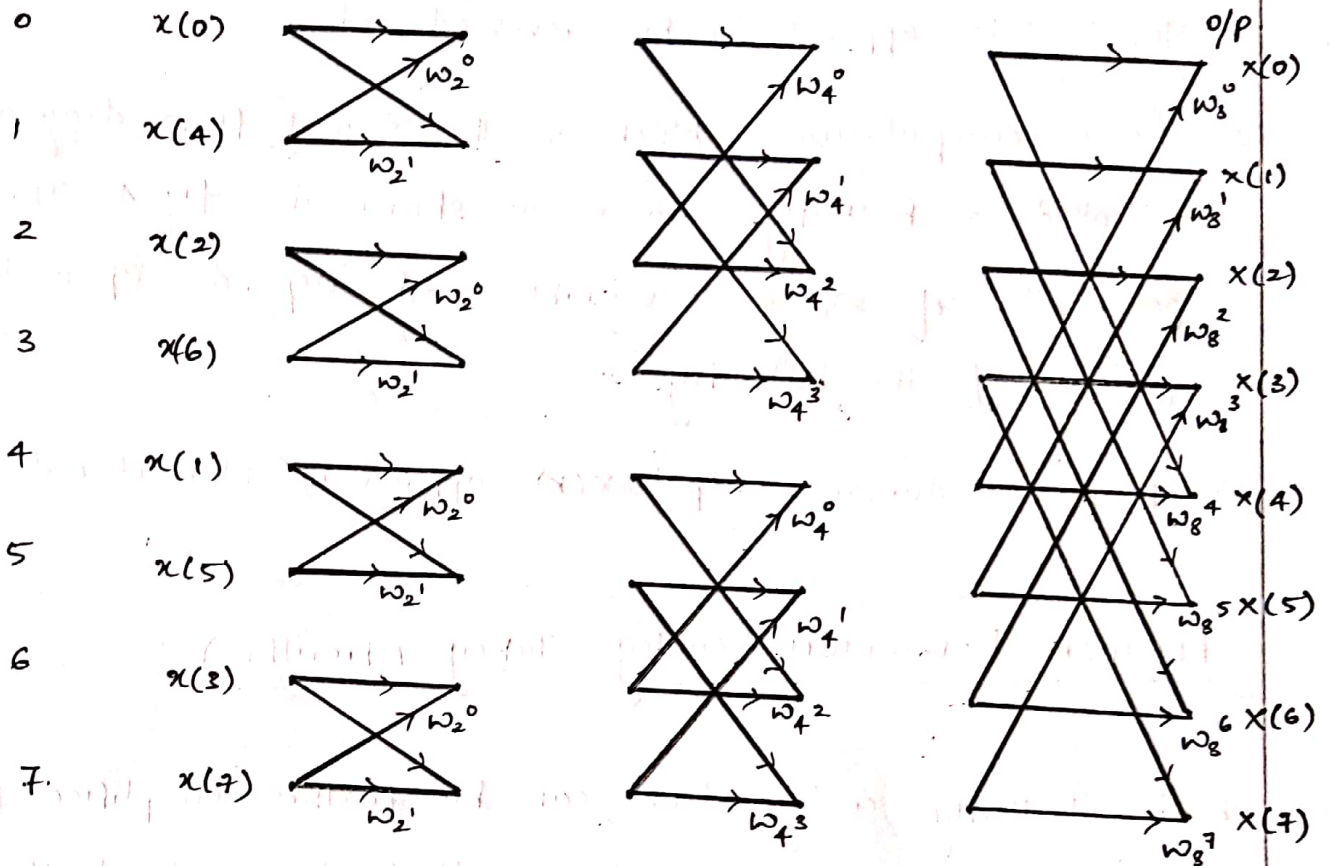


fig 3. : Signal flow diagram after final decimation.

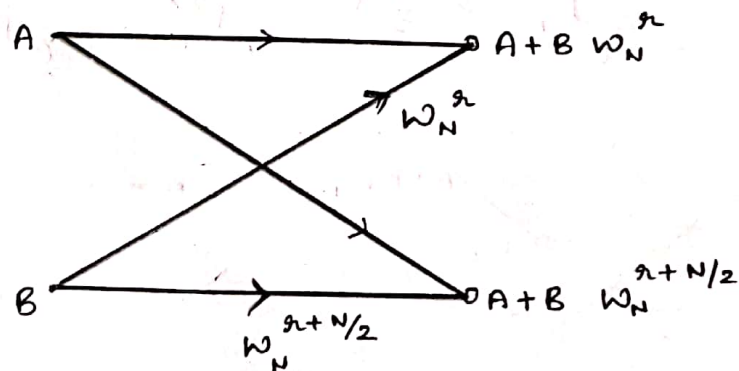


fig 4 : A sample Butterfly.

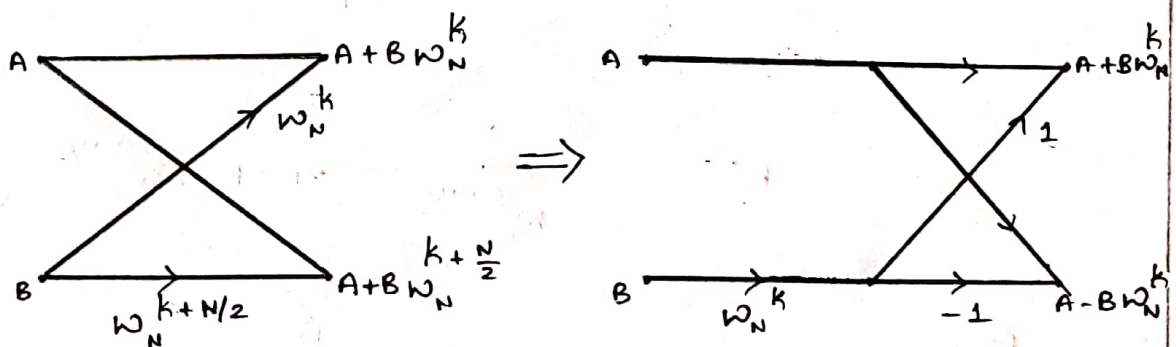
The following observations are made from the signal flow graph shown in fig 3 :

- 1) Input data appears in bit-reversed order.
- 2) Basic computational block in the signal flow diagram is called a Butterfly and is as shown in fig 4. The power 'a' of ω_N is a variable and depends upon the position of the butterfly.
- 3) Frequency domain output $x(k)$ appears in Normal order.

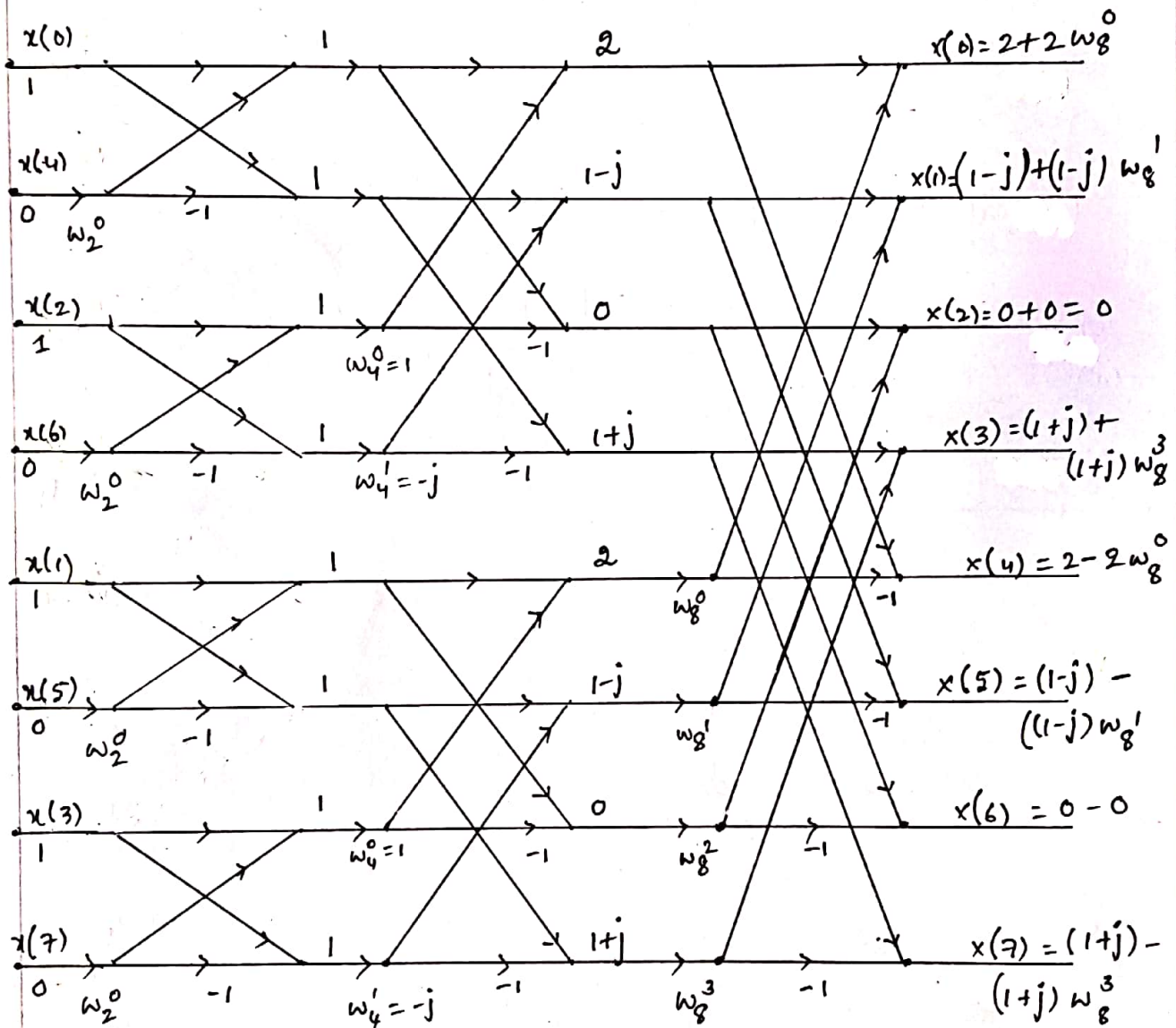
FURTHER REDUCTION. (Cooley - Tukey Algorithm).

Basic Butterfly configuration can be further simplified to reduce the number of complex multiplications per butterfly by one.

$$\omega_N^{k + \frac{N}{2}} = e^{j \frac{2\pi}{N} (k + \frac{N}{2})} = e^{j \frac{2\pi k}{N}} \cdot e^{j\pi} = -\omega_N^k$$



Compute 8-point DFT of the sequence $x(n) = (1, 1, 1, 1, 0, 0, 0, 0)$ using decimation in time radix-2 FFT algorithm.



$$x(0) = 2 + 2w_8^0 = 2 + 2 = \underline{\underline{4}}$$

$$x(1) = (1-j) + (1-j)w_8^1 = 1-j + (1-j)\left(\frac{1}{\sqrt{2}} - j\frac{1}{\sqrt{2}}\right)$$

$$x(1) = \underline{\underline{1-j2.414}}$$

$$x(2) = \underline{\underline{0}}$$

$$x(3) = (1+j) + (1+j)w_8^3 = 1+j + (1+j)\left(-\frac{1}{\sqrt{2}} - j\frac{1}{\sqrt{2}}\right)$$

$$= 1+j - \frac{1}{\sqrt{2}} - j\frac{1}{\sqrt{2}} - j\frac{1}{\sqrt{2}} + \frac{1}{\sqrt{2}}$$

$$x(3) = \underline{\underline{1+j0.414}}$$

$$x(4) = 2 - 2\omega_8^0 = 2 - 2 = 0$$

$$\begin{aligned} x(5) &= (1-j) - (1-j)\omega_8^1 = 1-j - (1-j)\left(\frac{1}{\sqrt{2}} - j\frac{1}{\sqrt{2}}\right) \\ &= 1-j - \frac{1}{\sqrt{2}} + j\frac{1}{\sqrt{2}} + j\frac{1}{\sqrt{2}} + \frac{1}{\sqrt{2}} \\ &= \underline{\underline{1 - j0.414}} \end{aligned}$$

$$x(6) = 0$$

$$x(7) = (1+j) - (1+j)\omega_8^3 = 1 + j2.414$$

2) Compute 8-point DFT of the sequence.
 $x(n) = (1, 0, 1, 0, 1, 0, 1, 0)$ using decimation in time
 radix-2 FFT algorithm.

