# Contents

| 1  | 01.19.21 |   |  |  |  |
|----|----------|---|--|--|--|
|    | 1.1      | Tuples & DFAs   |  |  |  |
|    | 1.2      | Domains & Codomains   |  |  |  |
|    | 1.3      | Strings   |  |  |  |
|    | 1.4      | <b>TODO</b> Review Recursive Definitions  |  |  |  |
|    | 1.5      | Languages   |  |  |  |
| 2  | 01.14.21 |   |  |  |  |
|    | 2.1      | Automaton (automata)  |  |  |  |
|    | 2.2      | The Mathematics of Automata   |  |  |  |
|    |          | 2.2.1 Mathematicians & History  |  |  |  |
|    |          | 2.2.2 Sequential Logic  |  |  |  |
|    | 2.3      | Necessary Review  |  |  |  |
|    | 2.4      | Functions   |  |  |  |
|    | 2.5      | <b>TODO</b> Types of Functions - Definition & Logical Statement .   |  |  |  |
|    | 2.6      | Finite Automaton (Finite State Machine)   |  |  |  |
|    |          |   |  |  |  |
| 1  | 0.       | 1.19.21   |  |  |  |
| 1. | 1 7      | Tuples & DFAs   |  |  |  |
|    | • T      | uples are sequences which are always finite in length   |  |  |  |
|    | • T      | he deterministic finite automaton shown is a 5-tuple:   |  |  |  |
|    |          | 1. Q: finite nonempty set of states   |  |  |  |
|    |          | - state: configuration of logic of a machine  |  |  |  |
|    |          | 2. $\Sigma$ (Sigma) - input alphabet  |  |  |  |
|    |          | <ul> <li>alphabet: a finite, nonempty set of symbols where symbols<br/>are an object of length 1</li> </ul> |  |  |  |
|    |          | 3. $\delta$ (Delta) - transition function   |  |  |  |
|    |          | 4. $Q_0 \in Q$ - starting state   |  |  |  |
|    |          | 5. $F \subset Q$ - set of final states  |  |  |  |
|    | • F      | or this deterministic finite automaton,   |  |  |  |
|    |          | $-\delta: \mathbb{Q} \times \Sigma \to \mathbb{Q}_2$  |  |  |  |

Represented as a table,

| Step | State | Input | Transition            |
|------|-------|-------|-----------------------|
| 1    | $Q_1$ | 1     | $Q_1 \to Q_2$         |
| 2    | $Q_2$ | 0     | $Q_2 \to Q_1$         |
| 3    | $Q_1$ | 1     | $Q_1 \to Q_2$         |
| 4    | $Q_2$ | 1     | $Q_2 \rightarrow Q_2$ |

# 1.2 Domains & Codomains

- Domain: set of all possible function inputs
- Codomain: set of all possible outputs

# 1.3 Strings

- In computer science, strings are character arrays
- In mathematics, strings are sequences of symbols
- Specifically a string over an alphabet,  $\Sigma$ , is a sequence of symbols belonging to  $\Sigma$
- $\epsilon$  is the empty string
- Concatenation: If  $w_1, w_2 \in \Sigma, w_1 \cdot w_2 = w_1 w_2$
- If  $c \in \Sigma$ , then  $\epsilon \cdot c = c \cdot \epsilon = c$

# 1.4 TODO Review Recursive Definitions

- Base step: a step that can not be broken down any further, a fact that is always true regardless of the input
- Recursive step:
- Defining the length of a string over  $\Sigma$ 
  - Base:  $|\epsilon| = 0$
  - Recursive:
    - \* let w be a string over  $\Sigma$ , and  $c \in \Sigma$
    - \* then  $|\mathbf{w} \cdot \mathbf{c}| = |\mathbf{w}| + 1$
- Using this to define |1011|,

- 1.  $|1011| = |101 \cdot 1| = |101| + 1 =$
- 2.  $|10 \cdot 1| + 1 = |10| + 1 + 1 =$
- 3.  $|1 \cdot 0| + 1 + 1 = |1| + 1 + 1 + 1 = |1|$
- 4.  $|\epsilon \cdot 1| + 1 + 1 + 1 =$
- 5.  $|\epsilon| + 1 + 1 + 1 + 1 =$
- 6. 0+1+1+1+1=4

# 1.5 Languages

- Languages over  $\Sigma$  a set of finite strings over  $\Sigma$
- $\bullet$  Langauges recognized by an automaton, M, L(M) is the language accepted by M
- $\emptyset$  is the empty language
- $\epsilon \neq \emptyset$
- $\epsilon \neq \{\epsilon\}$
- $\epsilon$  is not a symbol in any alphabet

# $2 \quad 01.14.21$

# 2.1 Automaton (automata)

- Self running machine requiring a continuous power source
  - Historically used power sources include water, steam, and electricity
- Course revolves around defining the mathematics powering machines

#### 2.2 The Mathematics of Automata

# 2.2.1 Mathematicians & History

- Cantor defines sets as collections of objects
- Cantor also argues that infinites can be of different magnitudes there are infinitely more real numbers than natural numbers
- Goedel eventually derives his incompleteness theorem

- No logical system that contains the natural numbers can prove its own soundness
- Every sound logical system containing the natural numbers contains valid statements that cannot be proved or disproved
- In 1936, Turing proves The Halting Problem is not decidable, it is impossible
  - The Halting Problem is an algorithm that can analyze any other algorithm and determine whether or not it goes into an infinite loop
- Turing creates the turing machine as an object consisting of sets and processes wherein the object can use any finite process to complete an action.
- Turing machine sets the basis for a computer, which leads to a series of important questions:
  - What can & can't a machine do?
  - What does it mean for a problem of be harder than another?
  - What does it mean for a machine to be more powerfule than another?

#### 2.2.2 Sequential Logic

- Sentential Logic- based on boolean results
  - Predicated on AND, OR, NOT
  - XOR, XAND, etc. can be derived using the above

#### 2.3 Necessary Review

- Textbook Ch. 0
- Logic Statements
- Set Theory
- Functions

#### 2.4 Functions

- Functions something that maps objects from one set to another
- Given f:  $a \rightarrow b$ ;
  - Everything in a is mapped to something in b
    - \* For every x, such that x is an element of a, there exists a y, such that y is an element of b
  - No one point in the domain can be mapped to two different points in the codomain
    - \* Logically, you can't have a function that takes in one input and returns two different outputs
    - \* If f maps  $x \to y1$  and y2, y1 = y2
    - - $\forall$  x  $\in$  A  $y_1,y_2 \in$  B  $[f(x)=y_1 \land f(x)=y_2 \rightarrow y_1=y_2]$

# 2.5 TODO Types of Functions - Definition & Logical Statement

- Injective Functions
- Surjective Functions
- Proof by Induction  $(\forall)$
- Proof by Contradiction (¬∃)

# 2.6 Finite Automaton (Finite State Machine)

- States are logical confirgurations
- States are generally based upon input
- Purpose of a state machine is to make a yes/no decision