Algorithms for Problem Solving

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February 4, 2025

1 Introduction

This report presents solutions to various algorithmic problems, including searching, graph traversal, dynamic programming, interval merging, and subarray sum optimization. Below is a brief problem description with representations, solution, results and conclusion. To make these results and conclusions reproducible, here is the github repository on which our code can be found https://github.com/sudo-001/TD-00-Search-Base-Software-Engineering.

2 Binary Search

2.1 Problem Representation

Binary search is an efficient algorithm for finding an element in a sorted list. It works by repeatedly dividing the search space in half until the target value is found or the search space is empty.

Example: Suppose we have the following sorted array:

If we are searching for the target 7, the Binary Search follows these steps:

- 1. Start with the full array: [1, 3, 5, 7, 9, 11, 13, 15]
- 2. Find the middle element (index 3, value = 7)
- 3. Since 7 is equal to the target, we return index 3.

Visual Representation:

Target found at index 3.

2.2 Solution

The algorithm follows these steps:

- 1. Initialize two pointers, low (start of the array) and high (end of the array).
- 2. Compute the middle index mid.
- 3. Compare the middle element with the target value:
 - If the middle element is equal to the target, return its index.
 - If the middle element is greater than the target, search in the left half.
 - Otherwise, search in the right half.
- 4. Repeat the process until the element is found or the search space is exhausted.

```
def binary_search(arr, target):
    left, right = 0, len(arr) - 1
                  if arr[mid] == target:
    return mid # Found the target
elif arr[mid] < target:
    left = mid + 1 # Target is in the right half</pre>
         arr = list(map(int, input("Enter the sorted array(separed by space): ").split()))
target = int(input("Enter the target : "))
result = binary_search(arr, target)
print("Result : Index {result}" if result != -1 else "Target not found")
except ValueError:
print("Please enter a valid integer")
                  vith open(filename, "r") as file:
   for line in file:
      parts = line.strip().split("->")
      sorted_list = list(map(int, parts[0].strip().split()))
      target = int(parts[1].strip())
         result = binary_search(sorted_list, target)
print(f"Sorted Table: {sorted_list}, Target: {target} -> Result: Index {result}" if result != -1 else "Target not found")
except FileNotFoundError:
    print(f"File {filename} not found.")
def main():
         while True:
    print("\n\n* MAIN MENU OF EXERCISE 1 *")
    print("1. Manual Test")
    print("2. Run a batch of tests")
    print("3. Exit")
                  if choice == "1":
    manual_test()
elif choice == "2":
    filename = input("Enter the file name of the batch of tests (eg: test_cases.txt) : ")
    batch_test(filename)
elif choice == "3":
    print("Exiting...\n\n")
    break
 if __name__ == "__main__":
    main()
```

2.3 Results

Sample test cases:

- Input: arr = [1, 3, 5, 7, 9], target = 5
- Output: Index = 2

2.4 Conclusion

Binary search significantly reduces search time compared to linear search, operating in $O(\log n)$ time complexity instead of O(n).

3 Graph Traversal (BFS and DFS)

3.1 Problem Representation

Graph traversal is a fundamental algorithmic technique used to explore nodes and edges of a graph. The two main traversal methods are:

- Breadth-First Search (BFS): Explores all neighbors of a node before moving to the next level.
- **Depth-First Search (DFS)**: Explores as deep as possible along one branch before backtracking.

A sample undirected graph:

Example of traversal starting from A:

- BFS: A -> B -> D -> C -> E -> F
- DFS: A -> B -> C -> E -> F -> D

3.2 Solution

The algorithm follows these steps as you can see on the pictures ??:

- 1. BFS uses a queue, exploring all neighbors before going deeper.
- 2. DFS uses a stack (or recursion), exploring one branch fully before backtracking.
- 3. Both methods ensure all nodes are visited.

3.3 Results

Sample test cases:

- Input Graph: A -- B, A -- D, B -- C, B -- E, E -- F
- BFS from A: A -> B -> D -> C -> E -> F
- DFS from A: A -> B -> C -> E -> F -> D

Github:

Figure 1: BFS/DFS code snippet (Graph)

Figure 2: BFS/DFS code snippet (Main)

3.4 Conclusion

Graph traversal techniques like BFS and DFS are essential for solving many real-world problems, including pathfinding and network analysis.

4 0/1 Knapsack Problem

4.1 Problem Representation

The 0/1 Knapsack problem is a classic optimization problem where the goal is to select a subset of items to maximize the total value while staying within a weight limit. Each item has a value and a weight, and we must decide which items to include in the knapsack such that the total weight does not exceed the maximum weight.

The problem can be defined as follows:

- We are given a set of items, where each item has a value and a weight.
- We are also given a maximum weight capacity of the knapsack.
- The objective is to maximize the total value of the selected items, subject to the constraint that the total weight does not exceed the knapsack's capacity.

Example: Suppose we have the following items:

Items =
$$[(60, 10), (100, 20), (120, 30)]$$

Max weight = 50

Where each tuple represents (value, weight). In this case, the goal is to select a subset of items such that the total weight is less than or equal to 50, while maximizing the total value.

4.2 Solution

The 0/1 Knapsack problem can be solved using dynamic programming. The algorithm proceeds by constructing a table to store the maximum value that can be achieved for each weight capacity and item combination. The steps are as follows:

- 1. Create a 2D array dp where dp[i][w] represents the maximum value that can be obtained using the first i items and with a weight capacity of w.
- 2. For each item i, and each weight w, calculate the maximum value by either including or excluding the item:
 - If the current item cannot fit (i.e., weight ; w), then dp[i][w] = dp[i-1][w].
- If the current item can fit, then we take the maximum of either excluding or including the item: dp[i][w] = max(dp[i-1][w], dp[i-1][w-weight] + value).
- 3. Once the table is filled, the maximum value will be located in dp[n] [max_weight], where n is the number of items and max_weight is the maximum weight capacity.
- 4. To find the selected items, trace back from dp[n][max_weight] and check which items were included.

```
n = len(items)
                    for w in range(max_weight + 1):
                        if weight > w:
                             dp[i][w] = max(dp[i - 1][w], dp[i-1][w-weight] + value)
                        selected_items.append(items[i-1])
                return dp[n][max_weight], selected_items
                items=[]
                n=int(input("Enter the number of inputs : "))
                for in range(n):
                    items.append((value, weight))
                max weight = int(input("Enter the maximum weight capacity: "))
                max value, selected_items = knapsack(items, max_weight)
print(f"\n Maximum value: {max_value}")
                print(f"\n Selected items: {selected_items}")
                    with open(filename, "r") as file:
                             items = [tuple(map(int, item.split(','))) for item in parts[0].split(";")]
                             max weight=int(parts[1])
                             print(f"Testing case: {items} with max weight {max weight}")
                            max_value, selected_items = knapsack(items, max_weight)
print(f"* Maximum value: {max_value}")
                            print(f"* Selected items: {selected items}")
                   print(f"File {filename} not found.")
                while True:
                    choice = input("Choose an option (1-3): ")
                         filename = input("Enter the file name of the batch of tests (eg: knapsack items.txt) : ")
                         load from file(filename)
                        print("Exiting...\n\n")
           if <u>__name__</u> == "_<u>_main__</u>":
__main()
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http
```

4.3 Results

Sample test cases:

```
• Input: Items = [(60, 10), (100, 20), (120, 30)], Max weight = 50
```

• Output: Maximum value = 220

• Selected items: [(100, 20), (120, 30)]

The dynamic programming approach efficiently computes the maximum value by systematically filling in the DP table. Here is a step-by-step breakdown of the DP table computation for the example input:

```
Iteration 1: dp[1][w] = max(dp[0][w], dp[0][w - weight[1]] + value[1])
Iteration 2: dp[2][w] = max(dp[1][w], dp[1][w - weight[2]] + value[2])
Iteration 3: dp[3][w] = max(dp[2][w], dp[2][w - weight[3]] + value[3])
```

4.4 Conclusion

The 0/1 Knapsack problem is a classical optimization problem that can be efficiently solved using dynamic programming with a time complexity of $O(n \times W)$, where n is the number of items and W is the maximum weight capacity. This approach guarantees finding the optimal solution while minimizing computational overhead compared to brute force methods.

5 Merge Intervals Problem

5.1 Problem Representation

The Merge Intervals problem involves merging overlapping intervals. Given a list of intervals, each interval is represented as a tuple (start, end). The goal is to merge the intervals where they overlap, and return a list of the non-overlapping intervals that cover all the intervals in the input.

The problem can be described as:

- We are given a collection of intervals.
- Each interval is represented by two integers, the start time and the end time.
- Our task is to merge any overlapping intervals and return the resulting list of intervals.

Example: Given the following intervals:

Intervals =
$$[(1, 3), (2, 4), (5, 7), (6, 8)]$$

The merged intervals should be:

$$Merged = [(1, 4), (5, 8)]$$

5.2 Solution

The algorithm to solve the Merge Intervals problem works as follows:

- 1. First, we sort the intervals by their starting times.
- 2. We initialize an empty list merged and add the first interval to it.
- 3. Then, for each subsequent interval, we check if it overlaps with the last interval in merged. If it does, we merge the two intervals by updating the end time of the last interval in merged.
- 4. If the current interval does not overlap with the last interval in merged, we add it as a new interval.
- 5. Finally, the merged list will contain all the non-overlapping intervals.

5.3 Results

For the input:

- Input: Intervals = [(1, 3), (2, 4), (5, 7), (6, 8)]
- Output: Merged intervals = [(1, 4), (5, 8)]

The algorithm successfully merges the overlapping intervals into two intervals: (1, 4) and (5, 8).

Here is a step-by-step breakdown of the merging process:

- Sort intervals: [(1, 3), (2, 4), (5, 7), (6, 8)].
- Initialize merged = [(1, 3)].
- Check interval (2, 4): Overlaps with (1, 3), so merge into (1, 4).
- Check interval (5, 7): Does not overlap with (1, 4), so add to merged.
- Check interval (6, 8): Overlaps with (5, 7), so merge into (5, 8).
- Final result: [(1, 4), (5, 8)].

5.4 Conclusion

The Merge Intervals problem can be efficiently solved by sorting the intervals and then merging the overlapping ones. This approach has a time complexity of $O(n \log n)$, where n is the number of intervals, due to the sorting step. The merging process itself takes linear time, making the solution both efficient and scalable.

```
return []
        merged = [intervals[0]]
             last = merged[-1]
             if current[0] <= last[1]: # Overlapping case
  merged[-1] = (last[0], max(last[1], current[1]))</pre>
         return merged
        n = int(input("Enter the number of intervals: "))
         for _ in range(n):
    start, end = map(int, input("Enter start and end time (separated by space): ").split())
         print(f"\n* Merged Intervals: {merged}")
                       print(f"\nTesting case: {intervals}")
                      merged = merge_intervals(intervals)
                       print(f"* Merged Intervals: {merged}")
        except FileNotFoundError:
    print(f"File {filename} not found.")
             print("\n\n* MAIN MENU EXERCISE 4*")
             print("1. Enter intervals manually")
print("2. Load test cases from a file")
             elif choice == '2':
                  load_from_file(filename)
                  print("Exciting...")
                  break
                 print("Invalid choice, please enter a number between 1 and 3.")
    if name == " main ":
```

6 Maximum Subarray Sum (Kadane's Algorithm)

6.1 Problem Representation

The Maximum Subarray Sum problem asks for finding the contiguous subarray (containing at least one number) within a one-dimensional array of numbers that has the largest sum. Kadane's Algorithm provides an efficient solution to this problem.

The problem can be described as follows:

- We are given an array of integers, which can include both positive and negative numbers.
- We need to find the contiguous subarray that has the largest sum.

Example: Given the following array:

Array =
$$[-2, 1, -3, 4, -1, 2, 1, -5, 4]$$

The maximum sum is 6, and the subarray corresponding to that sum is:

Subarray =
$$[4, -1, 2, 1]$$

6.2 Solution

Kadane's Algorithm is an efficient way to solve the Maximum Subarray Sum problem with a time complexity of O(n). The basic idea of the algorithm is to traverse the array while keeping track of the current sum and the maximum sum found so far. Here's how it works:

- 1. Initialize two variables: max_sum (to store the maximum sum found so far) and current_sum (to store the sum of the current subarray).
- 2. Iterate through the array:
 - If the current element is greater than the sum of the current element and the previous subarray (current_sum), start a new subarray from the current element.
 - Otherwise, add the current element to current_sum.
 - If current_sum exceeds max_sum, update max_sum and store the indices of the subarray that gives the maximum sum.
- 3. Finally, return the max_sum and the corresponding subarray.

6.3 Results

For the input array:

- Input: Array = [-2, 1, -3, 4, -1, 2, 1, -5, 4]
- Output: Maximum Sum = 6, Subarray = [4, -1, 2, 1]

The algorithm successfully finds the subarray with the maximum sum and returns both the sum and the subarray.

```
if not arr:
                   return 0, []
               max sum = current sum = arr[0]
               start = end = s = 0
               for i in range(1, len(arr)):
                       current sum = arr[i]
                       current sum += arr[i]
                   if current sum > max sum:
                        start, end = s, i
               return max_sum, arr[start:end+1]
               arr = list(map(int, input("Enter the array (separated by space): ").split()))
               print(f"\n* Maximum Sum: {max_sum}")
               print(f"* Subarray: {subarray}")
                   with open(filename, 'r') as file:
                        for line in file:
                            print(f"\nTesting case: {arr}")
                            print(f"* Maximum Sum: {max sum}")
                            print(f"* Subarray: {subarray}")
                   print(f"A File {filename} not found.")
               while True:
                   print("\n* MAIN MENU Exercise 5*")
                   print("1. Enter an array manually")
                   choice = input("Choose an option (1-3): ")
                   if choice == '1':
                   elif choice == '2':
    filename = input("Enter the filename (e.g., subarrays.txt): ")
                       load_from_file(filename)
                       print("\overling Exciting...\n\n")
                       break
                       print("A Invalid choice, please enter 1-3.")
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Githul 5 if __name__ == "__main__": 12
https://githubmainf/sudo-001/TD-00-Search-Base-Software-Engineering
```

6.4 Conclusion

Kadane's Algorithm is an efficient and optimal solution to the Maximum Subarray Sum problem, with a time complexity of O(n). By traversing the array only once and keeping track of the current and maximum sums, the algorithm can find the solution in linear time. This makes it highly suitable for large arrays and real-time applications.

7 Conclusion

The algorithms implemented in this report showcase their effectiveness in addressing specific computational problems with optimal performance. Binary Search, with its $O(\log n)$ time complexity, proved highly efficient for searching in sorted data, significantly reducing search time compared to linear search methods. The graph traversal algorithm, operating in O(V+E), successfully tackled graph-related problems, efficiently exploring all nodes and edges in a graph while maintaining optimal performance. Dynamic Programming optimized the solution for the 0/1 Knapsack problem, reducing the computational complexity from an exponential brute-force approach to a polynomial one, thereby enabling the solution of larger instances in a reasonable time frame. Kadane's Algorithm, with a time complexity of O(n), demonstrated its power in solving the Maximum Subarray Problem, quickly identifying the optimal subarray even for large arrays. Overall, the effectiveness of these algorithms is evident in their ability to solve complex problems efficiently, making them valuable tools in various domains, including search optimization, graph analysis, and dynamic programming problems. The results affirm the importance of choosing the right algorithm for problem-solving, as each one significantly improves computational efficiency and leads to faster, more scalable solutions.