

Course Name:	Computer Organization and Architecture Laboratory	Semester:	III
Date of Performance:	14/08/25	Batch No:	B1
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Faculty Sign & Date:		Grade/Marks:	25

Experiment No: 02
Title: Restoring method of division

Aim and Objective of the Experiment:
<p>Aim: The basis of algorithm is based on paper and pencil approach and the operation involve repetitive shifting with addition and subtraction. So the main aim is to depict the usual process in the form of an algorithm.</p> <p>Objective: To implement the restoring division algorithm using shift, subtract/add operations to simulate binary division.</p>

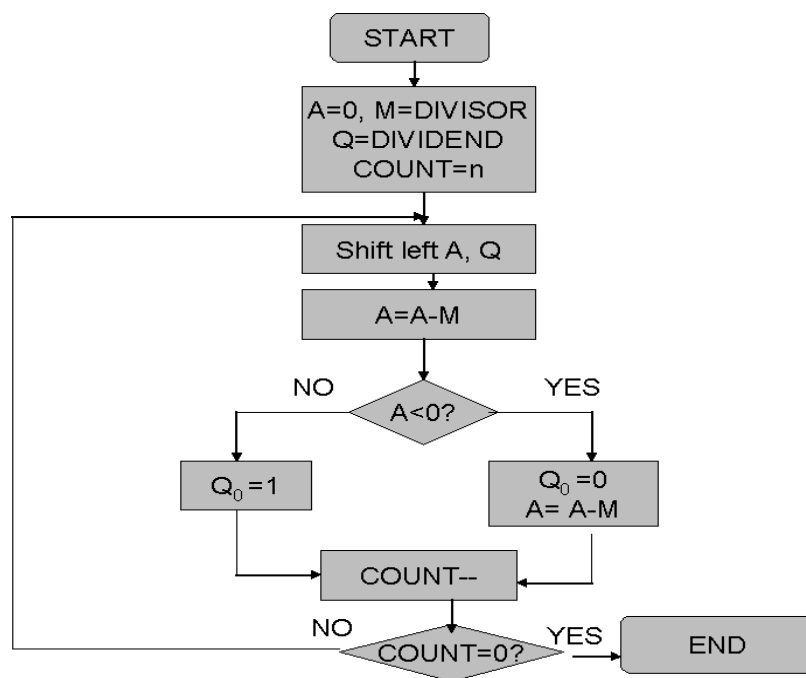
COs to be achieved:
<p>CO1: Describe and define the structure of a computer with buses structure and detail working of the ALU</p>

Books/Journals/Websites referred:
<ol style="list-style-type: none"> 1. Carl Hamacher, Zvonko Vranesic and Safwat Zaky, "Computer Organization", Fifth Edition, TataMcGraw-Hill. 2. William Stallings, "Computer Organization and Architecture: Designing for Performance", Eighth Edition, Pearson. 3. Dr. M. Usha, T. S. Srikanth, "Computer System Architecture and Organization", First Edition, Wiley-India.

Theory: The **restoring division algorithm** is a binary division technique that mimics long division used in decimal arithmetic. It operates on the principle of repeated subtraction and left-shifting. The dividend and divisor are stored in binary registers, and at each iteration, the contents of the accumulator and the quotient registers are left-shifted. The divisor is subtracted from the accumulator, and if the result is negative, the previous value of the accumulator is restored (hence the name "restoring"). This process is repeated for a number of steps equal to the bit length of the operands. The final values in the **quotient** and **accumulator** represent the result of the division and remainder, respectively.

It is primarily used in hardware-based division circuits due to its simple control logic but is slower than the **non-restoring** method due to the additional restoring step.

Flowchart for Restoring of Division:



Design Steps:

1. Start
2. Initialize A=0, M=Divisor, Q=Dividend and count=n (no of bits)
3. Left shift A, Q
4. If MSB of A and M are same

5. Then $A=A-M$
6. Else $A=A+M$
7. If MSB of previous A and present A are same
8. $Q_0=0$ & store present A
9. Else $Q_0=0$ & restore previous A
10. Decrement count.
11. If count=0 go to 11
12. Else go to 3
13. STOP

Example: - (Handwritten solved problem needs to be uploaded):-

$Q \div M$			0000
$7 \div 3$			1101
$M : 0011$		$Q : 0111$	1101
$-M : 1101$			0011
			0000
A	Q	count	0001
0000	0111	4	1101
0000	111 <input type="checkbox"/>		1110
1101	1110		+0011
0000	1110	3	0001
0001	110 <input type="checkbox"/>		0011
1110	1100		+1101
0001	1100	2	0000
0011	100 <input type="checkbox"/>		
0000	1001	1	0001
0001	001 <input type="checkbox"/>		+1101
1110	0010	0	0000 1110
$R = 0001$	$Q = 0010$		1100

Q M

7 ÷ 3

M : 0011 Q : 0111

-M : 1101

A	Q	count
0000	0111	4
0000	111 <input type="checkbox"/>	
1101	111 <u>0</u>	
0000	111 0	3
0001	110 <input type="checkbox"/>	
1110	110 <u>0</u>	
0001	110 0	2
0011	100 <input type="checkbox"/>	
0000	100 <u>1</u>	1
0001	001 <input type="checkbox"/>	
1110	001 <u>0</u>	0

R = 0001 Q = 0010

0000

1101

1101

0011

0000

0001

1101

1110

+ 0011

0001

0011

+ 1101

0000

0001

+ 1101

~~0000~~ 1110

1100

CODE:-

```

def compute_bits(num1,num2):

    a = max(num1,num2)

    # Calculate minimum bits to represent max number plus one bit for sign
    if needed

    # For simplicity just add some bits depending on range:

    if a <= 7:

        length = 4

    elif a <= 15:

        length = 5
  
```

```
elif a <= 21:

    length = 6

elif a <= 28:

    length = 7

elif a <= 35:

    length = 8

elif a <= 41:

    length = 9

elif a <= 49:

    length = 10

else:

    print("NOT ENOUGH TO COMPUTE YOUR RESULT")

    length = None

return length


def to_binary(num: int, bit_length: int) -> list:

    binary_num = [0] * bit_length

    i = bit_length - 1

    while num > 0 and i >= 0:

        binary_num[i] = num % 2

        num = num // 2

        i -= 1
```

```
    return binary_num

def binary_addition(b1, b2):

    carry = 0

    final = [0] * len(b1)

    for i in range(len(b1) - 1, -1, -1):

        total = b1[i] + b2[i] + carry

        final[i] = total % 2

        carry = total // 2

    return final

def to_negative(binary):

    inverted = [1 - bit for bit in binary]

    negative = binary_addition(inverted, ([0]*(len(binary)-1)) + [1])

    if len(negative) > len(binary):

        negative = negative[1:]

    return negative
```

```
def to_set_initials():

    question = input("Enter division (format m/q): ")

    nums = question.split('/')

    Dividend, Divisor = int(nums[0]), int(nums[1])

    bits_required = compute_bits(Dividend, Divisor)

    if bits_required is None:

        exit()

    M = to_binary(Divisor, bits_required)

    M_neg = to_negative(M)

    Q = to_binary(Dividend, bits_required)

    A = [0] * bits_required

    n = bits_required

    # print(f"M: {M}    M_neg: {M_neg}    Q: {Q}    A: {A}")

    return (M, M_neg, Q, A, n)

def to_restore_divide(M, M_neg, Q, A, n):

    while n != 0:

        # Left shift A and Q combined by 1 bit

        A = A[1:] + [Q[0]]
```

```
Q = Q[1:]

A = binary_addition(A, M_neg)

if A[0] == 1:

    Q = Q + [0]

    A = binary_addition(A, M)

else:

    Q = Q + [1]

n -= 1

return (A, Q)

def to_decimal(binary_list):

    decimal = 0

    length = len(binary_list)

    for i, bit in enumerate(binary_list):

        decimal += bit * (2 ** (length - i - 1))

    return decimal

my_values = to_set_initials()
```



```
remainder, quotient = to_restore_divide(*my_values)

print(f"The answer is : Remainder {to_decimal(remainder)}\nQuotient {to_decimal(quotient)}")
```

Post Lab Subjective/Objective type Questions:

1. Explain the role of the accumulator (A) in the restoring division algorithm and how it affects the quotient generation.

- A stores the partial remainder during division.
- Each step: shift A and quotient left, subtract divisor from A.
- If $A \geq 0$ after subtraction, set quotient bit to 1.
- If $A < 0$, restore A (add divisor back) and set quotient bit to 0.
- So, A helps decide each quotient bit by showing if divisor fits or not.

2. What are the advantages of restoring division over non restoring division?

- Simpler and easier to understand and implement.
- Always restores remainder when subtraction fails, ensuring correctness.
- Quotient bits are directly set by checking remainder sign.
- Less complex hardware compared to non-restoring division.

Conclusion:

PROBLEMS OUTPUT DEBUG CONSOLE TERMINAL PORTS

```
PS C:\Users\KJSCE\Desktop\cc> & C:/Users/KJSCE/AppData/Local/Programs/Python/Python312/python.exe c:/Users/KJSCE/Desktop/cc/restoring_div.py
Enter division (format m/q): 10/6
The answer is : Remainder 4
Quotient 1
PS C:\Users\KJSCE\Desktop\cc> & C:/Users/KJSCE/AppData/Local/Programs/Python/Python312/python.exe c:/Users/KJSCE/Desktop/cc/restoring_div.py
Enter division (format m/q): 15/4
The answer is : Remainder 3
Quotient 3
PS C:\Users\KJSCE\Desktop\cc> & C:/Users/KJSCE/AppData/Local/Programs/Python/Python312/python.exe c:/Users/KJSCE/Desktop/cc/restoring_div.py
Enter division (format m/q): 8/4
The answer is : Remainder 0
Quotient 2
PS C:\Users\KJSCE\Desktop\cc> |
```

Signature of faculty in-charge with Date: