

Course Name:	Computer Organization and Architecture Laboratory	Semester:	III
Date of Performance:	14/08/25	Batch No:	B1
Faculty Name:	Prof. Sheetal Pereira	Roll No:	16010124080
Faculty Sign & Date:		Grade/Marks:	25

Experiment No: 04
Title: Non-Restoring Method of Division

Aim and Objective of the Experiment:
<p>Aim: The basis of algorithm is based on paper and pencil approach and the operation involve repetitive shifting with addition and subtraction. So the main aim is to depict the usual process in the form of an algorithm.</p> <p>Objective: To study and implement Non Restoring method of division</p>

COs to be achieved:
CO1: Describe and define the structure of a computer with buses structure and detail working of the ALU

Books/Journals/Websites referred:
<ol style="list-style-type: none"> 1. Carl Hamacher, Zvonko Vranesic and Safwat Zaky, "Computer Organization", Fifth Edition, TataMcGraw-Hill. 2. William Stallings, "Computer Organization and Architecture: Designing for Performance", Eighth Edition, Pearson. 3. Dr. M. Usha, T. S. Srikanth, "Computer System Architecture and Organization", First Edition, Wiley-India.

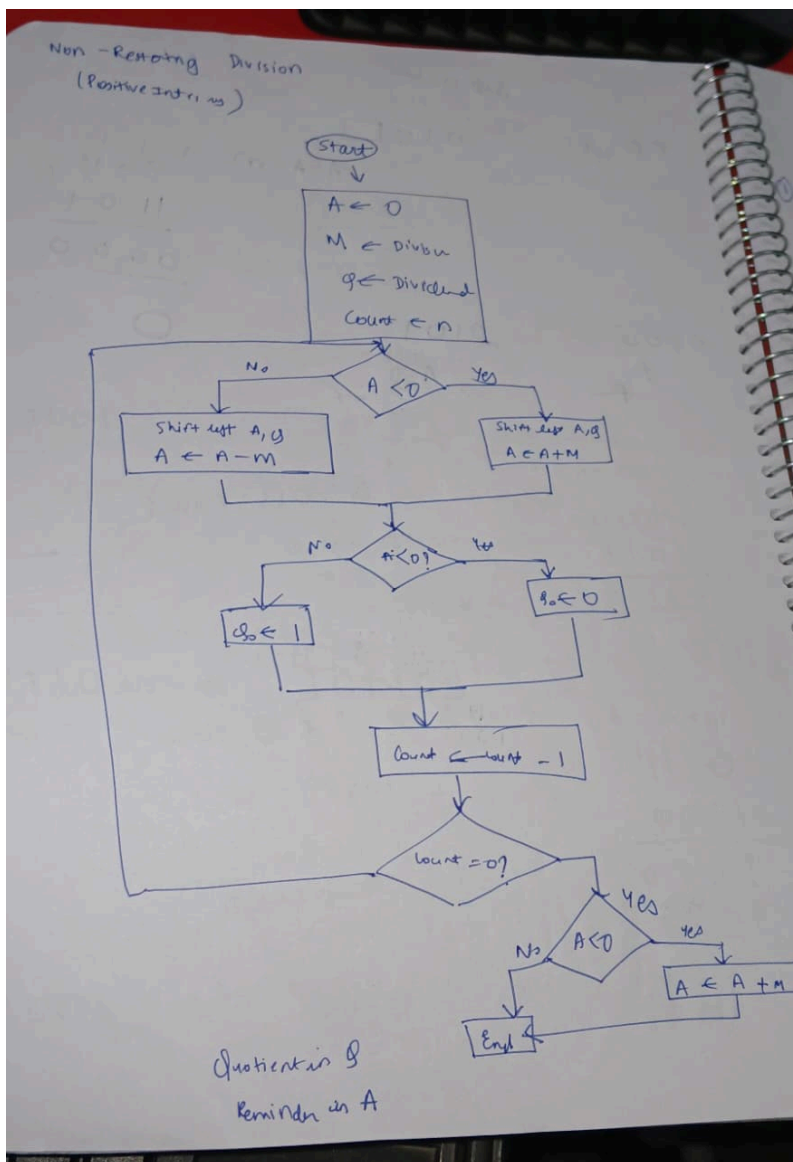
Theory: The non-restoring division algorithm is a faster alternative to the restoring method for binary division. Instead of "restoring" the accumulator when a subtraction results in a negative

value, it continues to the next step by adding the divisor back only when needed. This reduces the number of operations per iteration, improving performance.

The algorithm works by checking the sign of the accumulator after each operation. If it's positive, it subtracts the divisor in the next step; if negative, it adds the divisor. The quotient bit is determined accordingly. After all iterations, if the accumulator is negative, a final correction step is performed by adding the divisor once.

This method is more efficient in hardware implementations as it avoids unnecessary restoring operations.

Flowchart for Restoring of Division (Students need to draw):-



Design Steps:

1. Start
2. Initialize:
 - $A = 0$ (Accumulator)
 - $M = \text{Divisor}$
 - $Q = \text{Dividend}$
 - $\text{count} = n$ (number of bits)
3. Left shift (A, Q) combined by 1 bit
4. If MSB of previous A = 0 (A was positive or zero)
Then $A = A - M$
Else (A was negative)
 $A = A + M$
5. If MSB of A = 0
Set $Q_0 = 1$
Else
Set $Q_0 = 0$
6. Decrement count by 1
7. If $\text{count} \neq 0$, go to step 3
8. If $A < 0$ (negative remainder), then
 $A = A + M$ (final correction step)
9. Stop

Example: - (Handwritten solved problem needs to be uploaded):-

$Q \div M$ $M \Rightarrow 0011$ $Q : 0111$
 $7 \div 3$ $-M \Rightarrow 1101$

A	Q	n
0000	0111	4
1101	0111 [0]	3
1011	110 [0]	2
1110	110 [0]	1
1101	100 [0]	0
0000	100 [1]	
0001	001 [0]	
1110	001 [0]	

0001 0010
R=1 Q=2

$Q \div M$ $M \Rightarrow 000111$ $Q : 010100$
 $20 \div 7$ $-M \Rightarrow 111001$

A	Q	n
000000	010100	6
111001	10100 [0]	5
110011	01000 [0]	4
111010	01000 [0]	3
110000	10000 [0]	2
110111	00000 [0]	1
111110	00000 [1]	
000000	00000 [1]	
000110	00001 [0]	
111111	00001 [0]	

CODE:-

```
def compute_bits(num1,num2):  
  
    a = max(num1,num2)  
  
    # Calculate minimum bits to represent max number plus one bit for sign  
    if needed  
  
    # For simplicity just add some bits depending on range:  
  
    if a <= 7:  
  
        length = 4  
  
    elif a <= 15:  
  
        length = 5  
  
    elif a <= 21:  
  
        length = 6  
  
    elif a <= 28:  
  
        length = 7  
  
    elif a <= 35:  
  
        length = 8  
  
    elif a <= 41:  
  
        length = 9  
  
    elif a <= 49:  
  
        length = 10
```

```
else:

    print("NOT ENOUGH TO COMPUTE YOUR RESULT")

    length = None

return length


def to_binary(num: int, bit_length: int) -> list:

    binary_num = [0] * bit_length

    i = bit_length - 1

    while num > 0 and i >= 0:

        binary_num[i] = num % 2

        num = num // 2

        i -= 1

    return binary_num


def binary_addition(b1, b2):

    carry = 0

    final = [0] * len(b1)

    for i in range(len(b1) - 1, -1, -1):

        total = b1[i] + b2[i] + carry
```

```
        final[i] = total % 2

        carry = total // 2

    return final

def to_negative(binary):

    inverted = [1 - bit for bit in binary]

    negative = binary_addition(inverted, ([0]*(len(binary)-1)) + [1])

    if len(negative) > len(binary):

        negative = negative[1:]

    return negative

def to_set_initials():

    question = input("Enter division (format m/q): ")

    nums = question.split('/')

    Dividend, Divisor = int(nums[0]), int(nums[1])

    bits_required = compute_bits(Dividend, Divisor)

    if bits_required is None:

        exit()

    M = to_binary(Divisor, bits_required)

    M_neg = to_negative(M)
```

```
Q = to_binary(Dividend, bits_required)

A = [0] * bits_required

n = bits_required

# print(f"M: {M}    M_neg: {M_neg}    Q: {Q}    A: {A}")

return (M, M_neg, Q, A, n)

def non_restoring(M, M_neg, Q, A, n):
    while (n!=0):
        if A[0]==0:
            A = A[1:] + [Q[0]]
            Q = Q[1:]
            A= binary_addition(A, M_neg)

        else:
            A = A[1:] + [Q[0]]
            Q = Q[1:]
            A= binary_addition(A, M)

        if A[0]==0:
            Q=Q+[1]

        else:
```



```
        Q=Q+[0]

        n=n-1

    if n==0:

        if A[0]==1:

            A=binary_addition(A,M)

        else:

            pass

    return (Q,A)

def to_decimal(binary_list):

    decimal = 0

    length = len(binary_list)

    for i, bit in enumerate(binary_list):

        decimal += bit * (2 ** (length - i - 1))

    return decimal

my_values = to_set_initials()

quotient,remainder = non_restoring(*my_values)
```

```
print(f"The answer is : Remainder {to_decimal(remainder)}\nQuotient  
{to_decimal(quotient)}")
```

Post Lab Subjective/Objective type Questions:

- **Why is the non-restoring division algorithm considered faster than the restoring method in binary division?**
 - Non-restoring division does not restore the remainder after a negative subtraction.
 - Instead of doing an extra addition to fix the remainder every time subtraction fails, it just keeps going with the next step.
 - This removes the restore (add-back) step present in restoring division, saving time.
 - Fewer operations per cycle → faster overall division.
- **What is the purpose of the final correction step in the non-restoring division algorithm?**
 - Because non-restoring division doesn't restore after every negative remainder, the final remainder might be negative at the end.
 - The final correction step adds the divisor back if the remainder is negative.
 - This ensures the final remainder is positive and correct.
 - It fixes any leftover error caused by skipping restores during the process.

- **Solve 10/3 using Non-Restoring algorithm for division operation?**

PROBLEMS OUTPUT DEBUG CONSOLE TERMINAL PORTS

```
PS C:\Users\KJSCE\Desktop\cc> & C:/Users/KJSCE/AppData/Local/Programs/Python/Python312/python.exe c:/Users/KJSCE/Desktop/cc/non_restoring.py
Enter division (format m/q): 10/3
The answer is : Remainder 1
Quotient 3
PS C:\Users\KJSCE\Desktop\cc> |
```

Conclusion:

PROBLEMS OUTPUT DEBUG CONSOLE TERMINAL PORTS

```
PS C:\Users\KJSCE\Desktop\cc> & C:/Users/KJSCE/AppData/Local/Programs/Python/Python312/python.exe c:/Users/KJSCE/Desktop/cc/non_restoring.py
Enter division (format m/q): 17/5
The answer is : Remainder 2
Quotient 3
PS C:\Users\KJSCE\Desktop\cc> & C:/Users/KJSCE/AppData/Local/Programs/Python/Python312/python.exe c:/Users/KJSCE/Desktop/cc/non_restoring.py
Enter division (format m/q): 14/6
The answer is : Remainder 2
Quotient 2
PS C:\Users\KJSCE\Desktop\cc> & C:/Users/KJSCE/AppData/Local/Programs/Python/Python312/python.exe c:/Users/KJSCE/Desktop/cc/non_restoring.py
Enter division (format m/q): 12/2
The answer is : Remainder 0
Quotient 6
PS C:\Users\KJSCE\Desktop\cc> |
```

Signature of faculty in-charge with Date: