# COMP0174 Practical Program Analysis Available Expressions Analysis

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# Four Classic Analyses

	Forward	Backward
Must	Available Expressions	Very Busy Expressions
May	Reaching Definitions	Live Variables

#### Available Expressions

**Definition**. For each program point, which expressions must have already been computed, and not later modified, on all paths to the program point.

It is a forward must analysis.

**Application**: optimization (don't recompute expressions that are still available).

```
[x \coloneqq a + b]^{1};
[y \coloneqq a * b]^{2};
while [y > a + b]^{3}do
[a \coloneqq a + 1]^{4};
[x \coloneqq a + b]^{5};
```

The expression a + b is available every time execution reaches the condition 3, therefore the expression need not be recomputed.

### Killed Expression

An expression is killed in a block if any of the variables used in the expression are modified in the block:

$$kill_{AE}([x \coloneqq a]^l) = \{a' \in AExp_* | x \in Vars(a')\}$$
  
 $kill_{AE}([skip]^l) = \emptyset$   
 $kill_{AE}([b]^l) = \emptyset$ 

where  $AExp_*$  are all expressions in the program

#### Killed Expression

An expression is killed in a block if any of the variables used in the expression are modified in the block:

$$kill_{AE}\big([x\coloneqq a]^l\big) = \{a'\in AExp_*|\ x\in Vars(a')\}$$
 
$$kill_{AE}\big([skip]^l\big) = \emptyset$$
 Assignment statement: 
$$kill_{AE}\big([b]^l\big) = \emptyset$$
 kills all expressions that

kills all expressions that use variable x assigned in the block because they have to be recomputed again

where  $AExp_*$  are all expressions in the program

#### Generated Expression

A generated expression is an expression that is evaluated in the block and where none of the variables used in the expression are later modified in the block:

$$gen_{AE}([x \coloneqq a]^l) = \{a' \in AExp(a) \mid x \notin Vars(a')\}$$

$$get_{AE}([skip]^l) = \emptyset$$

$$gen_{AE}([b]^l) = AExp(b)$$

## Analysis

The goal of the analysis is to compute the largest set satisfying the equation for  $AE_{entry}$ :

$$AE_{entry}(l) = \begin{cases} \emptyset & if \ l = init(program) \\ \cap \{AE_{exit}(l') \mid (l', l) \in flow(program)\} \ otherwise \end{cases}$$

$$AE_{exit}(l) = \left(AE_{entry}(l) \setminus kill(B^l)\right) \cup gen_{AE}(B^l)$$

where  $B^l \in blocks(program)$ 

$[x \coloneqq a + b]^1;$	
$[y \coloneqq a * b]^2;$	
while $[y > a + b]^3$	do
$[a \coloneqq a + 1]^4$ ;	
$[x \coloneqq a + b]^5$ ;	

l	$kill_{AE}(l)$	$gen_{AE}(l)$
1	Ø	${a+b}$
2	Ø	$\{a*b\}$
3	Ø	${a+b}$
4	$\{a+b,a*b,a+1\}$	Ø
5	Ø	${a+b}$

```
AE_{entry}(1) = \emptyset

AE_{entry}(2) = AE_{exit}(1)

AE_{entry}(3) = AE_{exit}(2) \cap AE_{exit}(5)

AE_{entry}(4) = AE_{exit}(3)

AE_{entry}(5) = AE_{exit}(4)

AE_{entry}(5) = AE_{exit}(4)

AE_{entry}(5) = AE_{exit}(4)

AE_{exit}(1) = AE_{entry}(1) \cup \{a + b\}

AE_{exit}(2) = AE_{entry}(3) \cup \{a + b\}

AE_{exit}(4)

AE_{exit}(4)

AE_{exit}(5) = AE_{entry}(5) \cup \{a + b\}
```

Equations for entry and exit functions:

$$AE_{entry}(1) = \emptyset \qquad AE_{exit}(1) = AE_{entry}(1) \cup \{a + b\} \\ AE_{entry}(2) = AE_{exit}(1) \qquad AE_{exit}(2) = AE_{entry}(2) \cup \{a * b\} \\ AE_{entry}(3) = AE_{exit}(2) \cap AE_{exit}(5) \qquad AE_{exit}(3) = AE_{entry}(3) \cup \{a + b\} \\ AE_{entry}(4) = AE_{exit}(3) \qquad AE_{exit}(4) = AE_{entry}(4) \setminus \{a + b, a * b, a + 1\} \\ AE_{entry}(5) = AE_{exit}(4) \qquad AE_{exit}(5) = AE_{entry}(5) \cup \{a + b\}$$

AE at the entry of Block 3 = AE available at the exit of Block 2 (when entering the loop for the first time) and of Block 5 (when coming back from the exit of the loop)

$[x \coloneqq a + b]^1;$	
$[y \coloneqq a * b]^2;$	
while $[y > a + b]^3$	do
$[a \coloneqq a + 1]^4$ ;	
$[x \coloneqq a + b]^5$ ;	

l	$AE_{entry}(l)$	$AE_{exit}(l)$
1	Ø	${a+b}$
2	${a+b}$	$\{a+b, a*b\}$
3	${a+b}$	${a+b}$
4	${a+b}$	Ø
5	Ø	${a+b}$