



# An Analysis of TAGE and L-TAGE Branch Predictors

Team Amdahl's Vandals



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# Outline

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# Introduction

Branch predictors:

- Critical role in achieving effective performance
- 1% improvement in prediction accuracy can save millions of cycles in large programs

A close-up photograph of a microchip on a circuit board, overlaid with a teal gradient. The chip is black with gold pins and has white text: "CHRONTEL", "7007A-T", "RUC 9949", and "M9J19".

## More efficient Branch Predictors?

- History repeats itself
- Multiple history lengths can capture correlations from both remote branch outcomes and recent history
- Compute final prediction using multiple isolated predictor components



## More efficient Branch Predictors?

- **O-GEHL Predictor:** uses adder tree to combine predictions from multiple components
- Geometric series gives history lengths to be utilised
- Allows long history lengths to be exploited and recent branch outcome correlations to be captured



# The TAGE Predictor

Tagged GEometric  
history length predictor

- Tagless bimodal base predictor + partially tagged components indexed using history lengths in GP
- Prediction by either base predictor or tag match on a tagged component
- Multiple hits?
  - If tag-matching table with longest history has strong prediction, use it
  - Else use tag-matching table with second-longest history (alternate prediction)





# The TAGE Predictor

TAged GEometric  
history length predictor

- New and efficient predictor update algorithm used
- Makes partial tagging more efficient than adder for final prediction computation
- Outperforms O-GEHL at equivalent storage budgets and predictor complexity

# Principles and Explanation - 1k ft. view

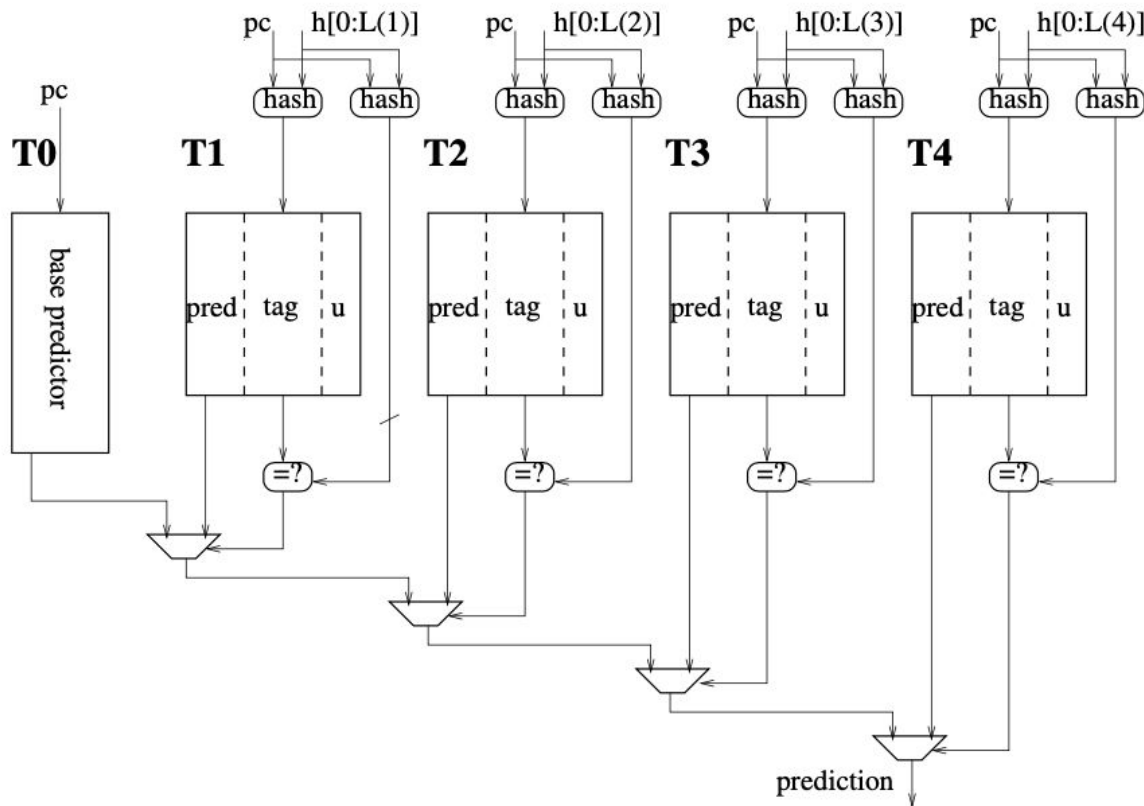
- Geometric history lengths
- Update rules
- Rationale behind them





# TAGE

A base predictor +  
several tagged predictor  
components indexed  
with increasing history  
lengths





# Geometric History Length Prediction

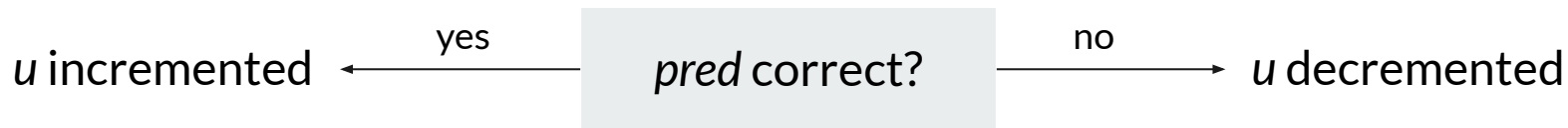
- $M$  distinct predictor tables
- Indexed with hash functions of branch address and global branch history
- Distinct history lengths used for computing index of distinct tables
- Base table  $T_0$  indexed using branch address only
- For table  $T_i$  ( $i = 1$  to  $M-1$ ), history length:

$$L(i) = (int)(\alpha^{i-1} \times L(1) + 0.5)$$



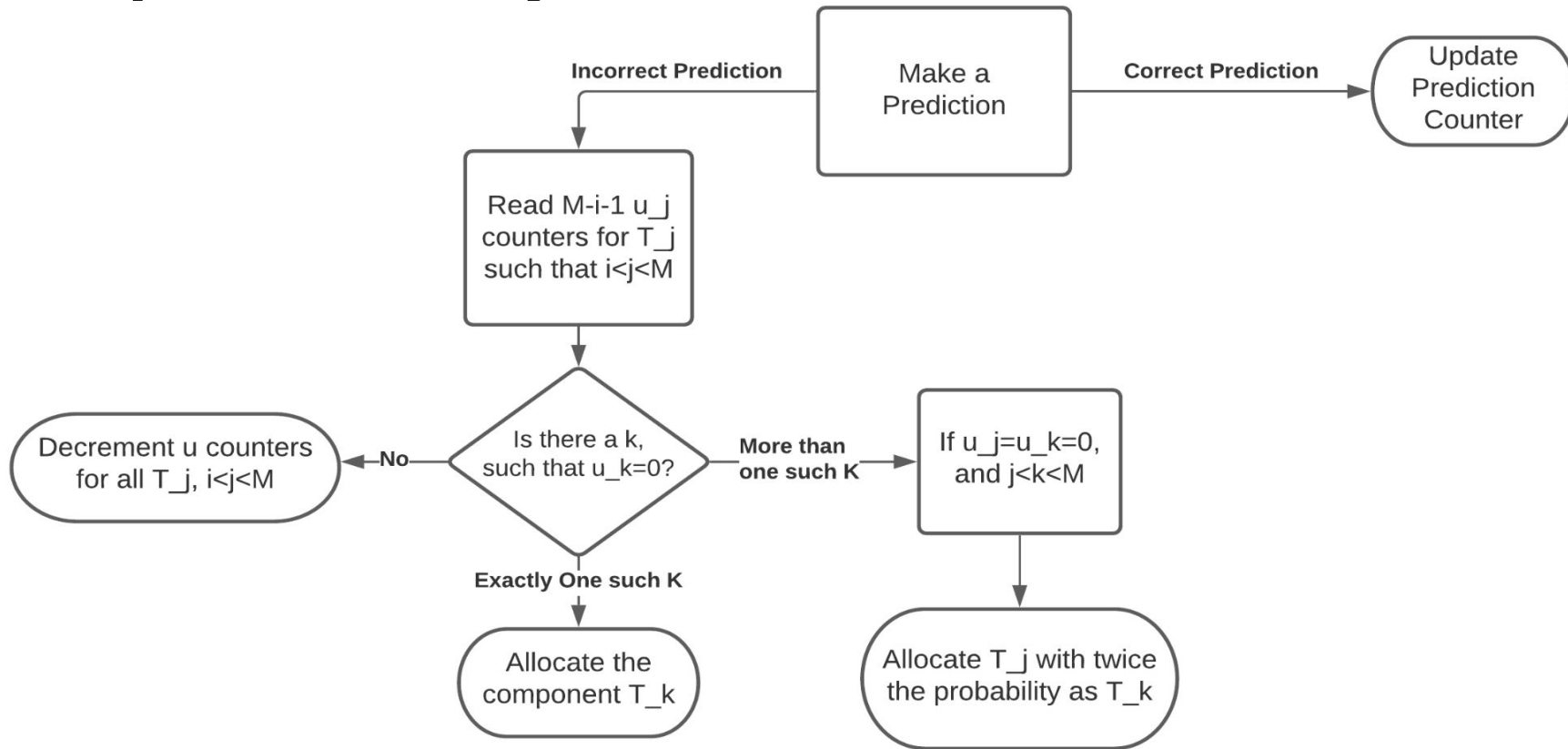
# Update Policy

- Each tagged component entry has a useful counter  $u$
- Updated when alternate prediction (second best)  $\neq$  final prediction  $pred$ .



- MSBs and LSBs of all  $u$  counters periodically reset to 0.

# Update Policy





## Why this update policy?

- Minimizes perturbation caused by single occurrence of a branch
- At most one tagged entry is allocated on a misprediction
- The useful counter  $u$  helps mimic a pseudo LRU policy



# L-TAGE Predictor

- An X component TAGE predictor combined with a Y-entry loop predictor
- Loops are detected based on repeated access to the tables for branch prediction





# Loop Predictor Component

- Identifies regular loops with constant number of iterations
- Provides global prediction when loop has been executed 3 times successively with same number of iterations
- Replacement policy is based on the age
- Used when confidence of loop prediction is high enough



# Why loop prediction?

- Can capture regular loop behaviors, which TAGE cannot, using very limited storage
- Can predict the number of times a loop will iterate
- Allows several basic blocks to be prefetched accurately
- Reduces number of instruction cache and memory requests

# So what did we do?

- We implemented TAGE and L-TAGE in ChampSim

Find the code here: <https://github.com/sudoRicheek/amdahls-vandals>

To see some sample results clone the repo and:

- Load the traces [server\_001, server\_002, server\_003, server\_004, server\_009] into traces/
- cd ChampSim/
- ./run\_cases.sh

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The generated results can be found in results\_10M/

# What did we observe?

- Comparing performance across predictors
- Varying number of components for TAGE and L-TAGE
- Varying history lengths
- Varying tag widths

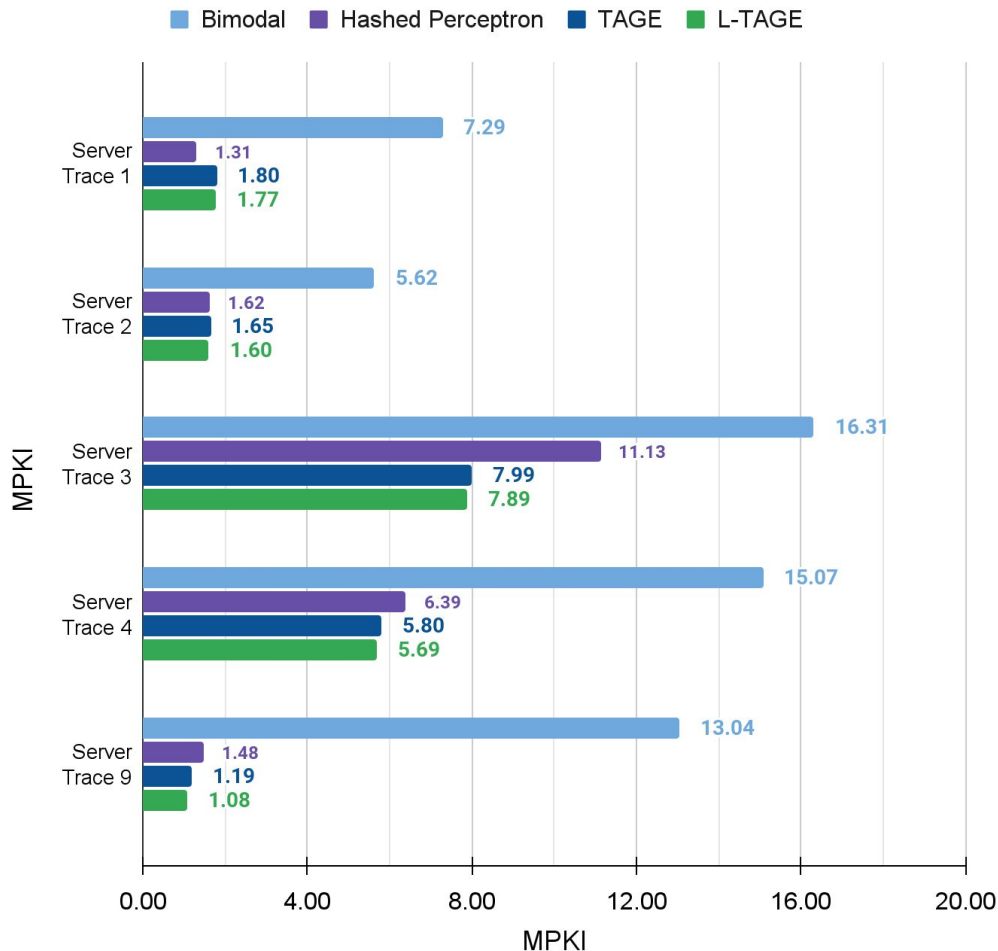
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# Comparing Predictors



In all cases, our TAGE and L-TAGE perform better than Bimodal

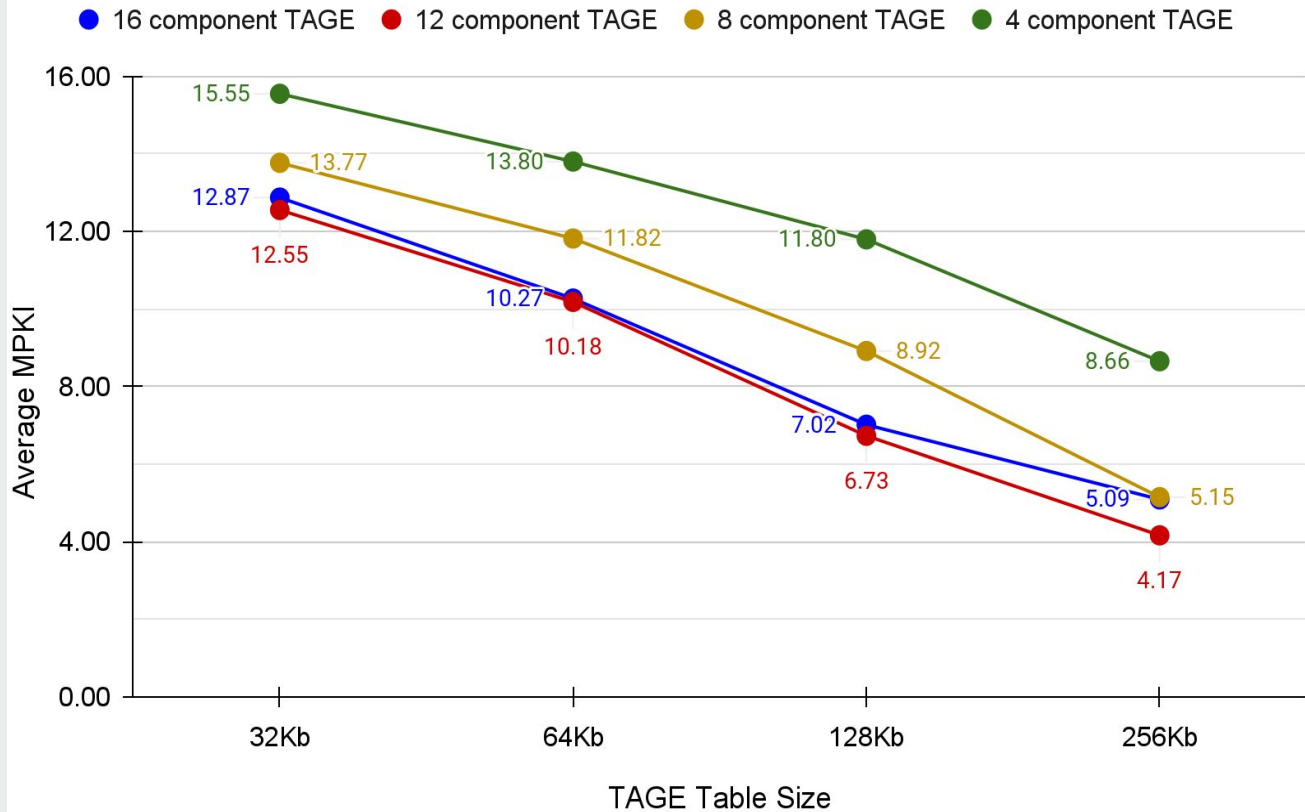
In most cases, it also performs better than Hashed Perceptron



# TAGE: Varying number of components



- On increasing Table Size, accuracy of all predictors increase
- On increasing components upto 12, there is significant decrease in MPKI
- On increasing components from 12 to 16, average MPKI begins to stagnate

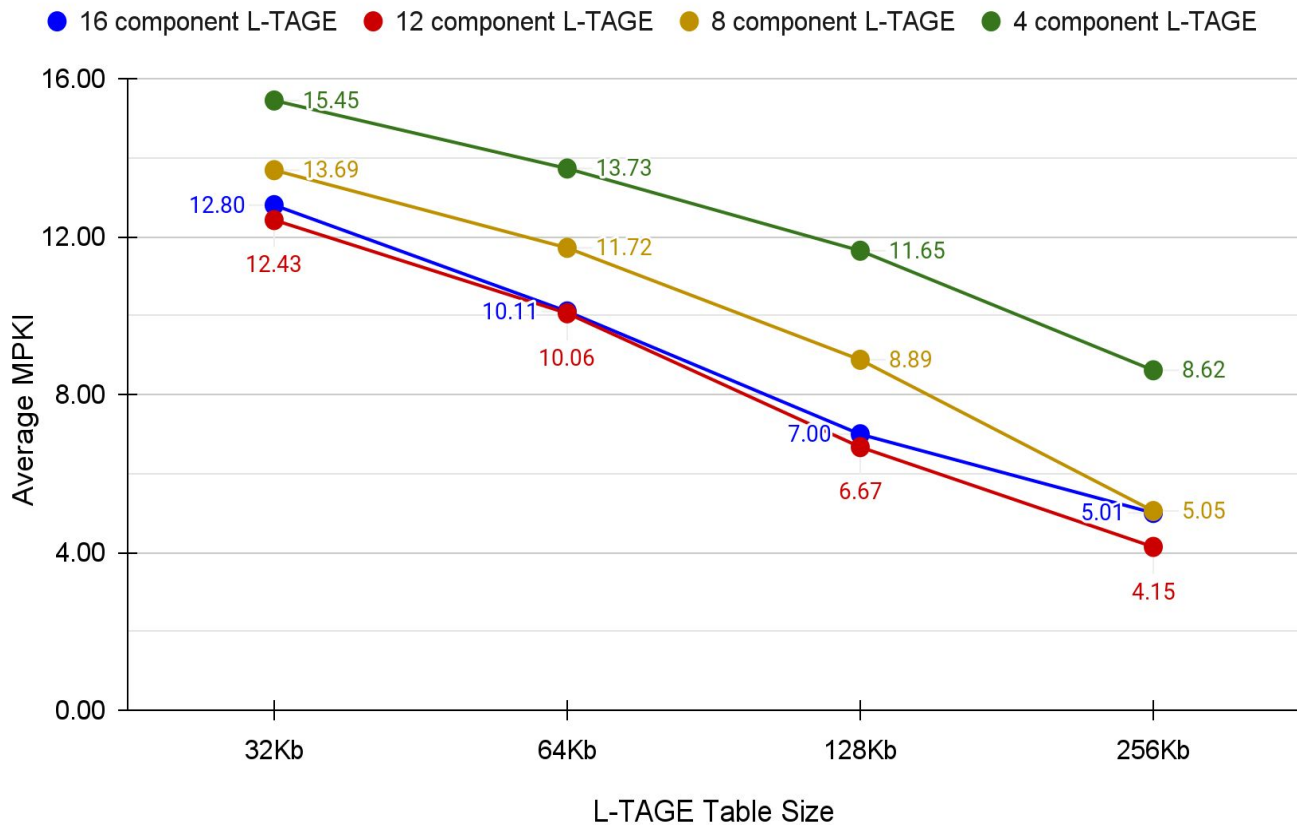


MPKI average taken over 5 server caches



# L-TAGE: Varying number of components

- Again, on increasing Table Size, accuracy of all predictors increase
- On increasing components upto 12, there is significant decrease in MPKI
- On increasing components from 12 to 16, MPKI begins to stagnate



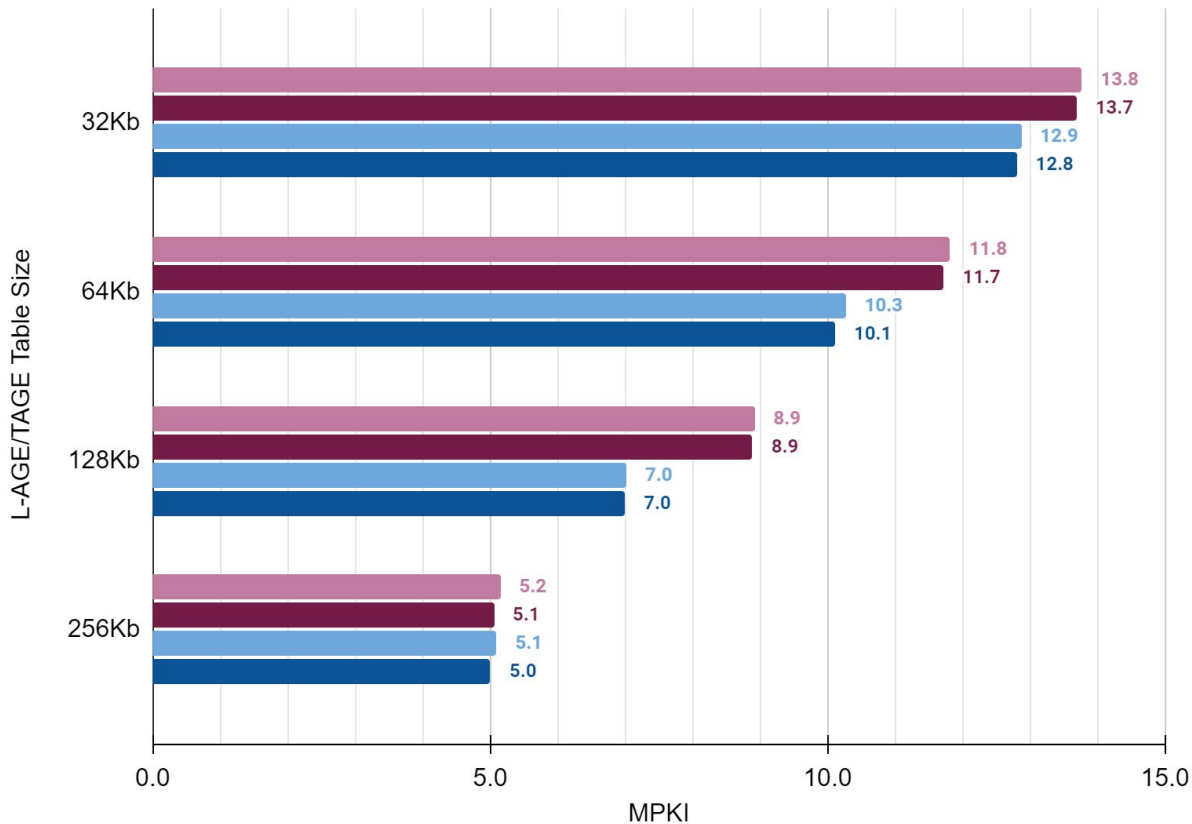
MPKI average taken over 5 server caches

# TAGE vs. L-TAGE



- We again see a clear trend of decreasing MPKI with increasing Table Size.
- A 16 table TAGE and L-TAGE performs much better than its 8 table counterpart.
- n-component L-TAGE performs better than n-component TAGE over all the tabulated Table sizes.

8 component TAGE   8 component L-TAGE   16 component TAGE   16 component L-TAGE



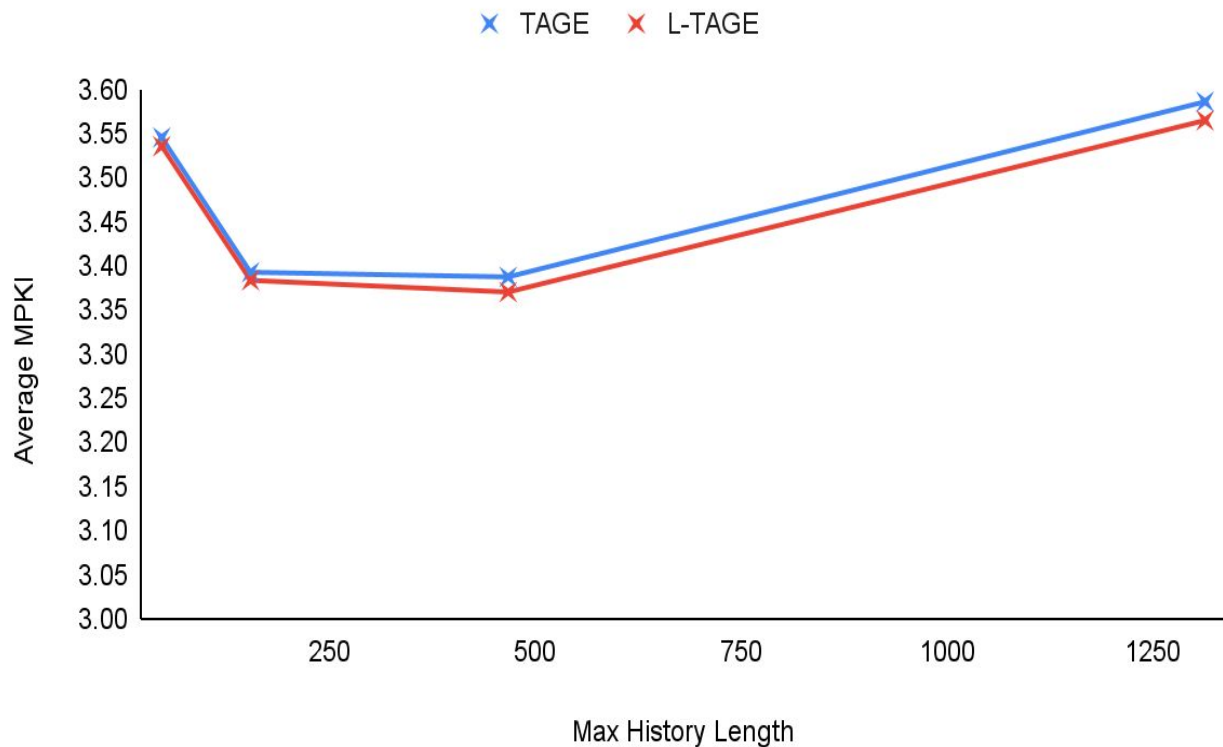
MPKI average taken over 5 server caches

# Varying History Length



- With increase in history length, initially MPKI decreases (hence a larger history helps)
- As we increase further the maximum history length, MPKI stagnates and then increases marginally.
- TAGE and L-TAGE are robust to perturbations in history length.
- Marginal increase could be because server traces don't require very long history

MPKI v/s Maximum History Length

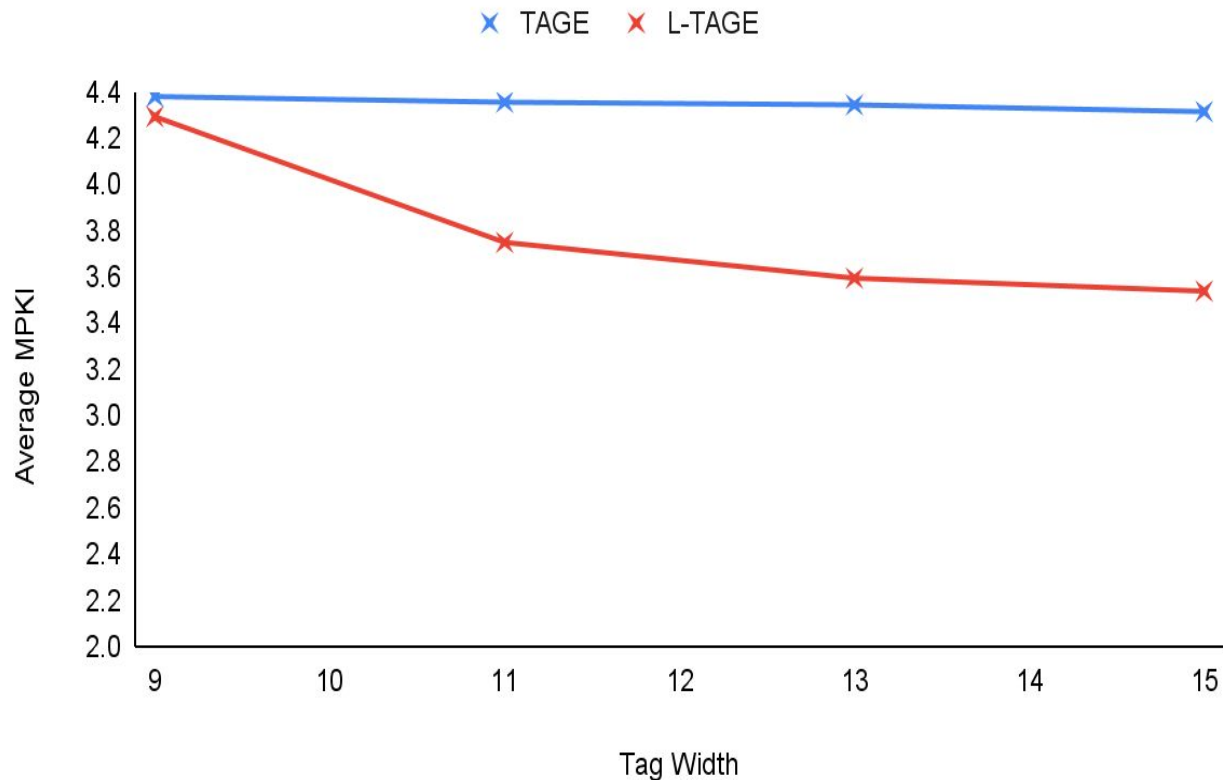


MPKI average taken over 5 server caches

# Varying Tag Width

- Using a larger tag width will need more storage for tables
- Too small tag widths will result in false tag matches, leading to mispredictions.
- Results show that an increase in tag width stagnates MPKI after a certain point

MPKI v/s Tag Width



MPKI average taken over 5 server caches

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## What we learnt



# Conclusion

- We first saw how TAGE greatly improved the efficacy of geometric length based branch predictors like O-GEHL
- We built an intuition to justify history lengths in GP, Update Rule, and the Loop Predictor in L-TAGE
- We finally present our code and show extensive results by studying each major component of L-TAGE and TAGE in detail.





# Contributions

Roll No.	Name	Contribution
190260036	Richeek Das	Code, Inferences, Optimal Parameters
190050002	Aakriti	Graphs generation, Inferences, Results video
190050020	Ankit Kumar Misra	Code, TAGE theory in PPT, Results video
190050111	Shabnam Sahay	Theory, PPT formatting and editing, Theory video
190050119	Sumit Jain	TAGE Theory, Results, Graphs, Inferences



# References

[A case for \(partially\) tagged Geometric History Length Branch Prediction](#)

[The L-TAGE branch predictor](#)

<https://github.com/boubinlg/BranchPrediction>

<https://github.com/2Bor2C/tage>

[Championship Branch Prediction \(CBP\) 2016](#)