

# *APPLIED GRAPH THEORY AND ALGORITHMS (CSC4066)*

## *Johnson Algorithm*



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# Johnson Algorithm

- The problem is to find the shortest path between every pair of vertices in a given weighted directed graph and weight may be negative
- It is suitable for sparse graph
- Johnson's Algorithm uses both Dijkstra's Algorithm and Bellman-Ford Algorithm
- Johnson's Algorithm uses the technique of "reweighting."



# Johnson Algorithm

- If all edge weights  $w$  in a graph  $G = (V, E)$  are nonnegative, we can find the shortest paths between all pairs of vertices by running Dijkstra's Algorithm once from each vertex.
- If  $G$  has negative - weight edges, we compute a new - set of non - negative edge weights that allows us to use the same method.

# Johnson Algorithm

- Reweighing is done through the formula

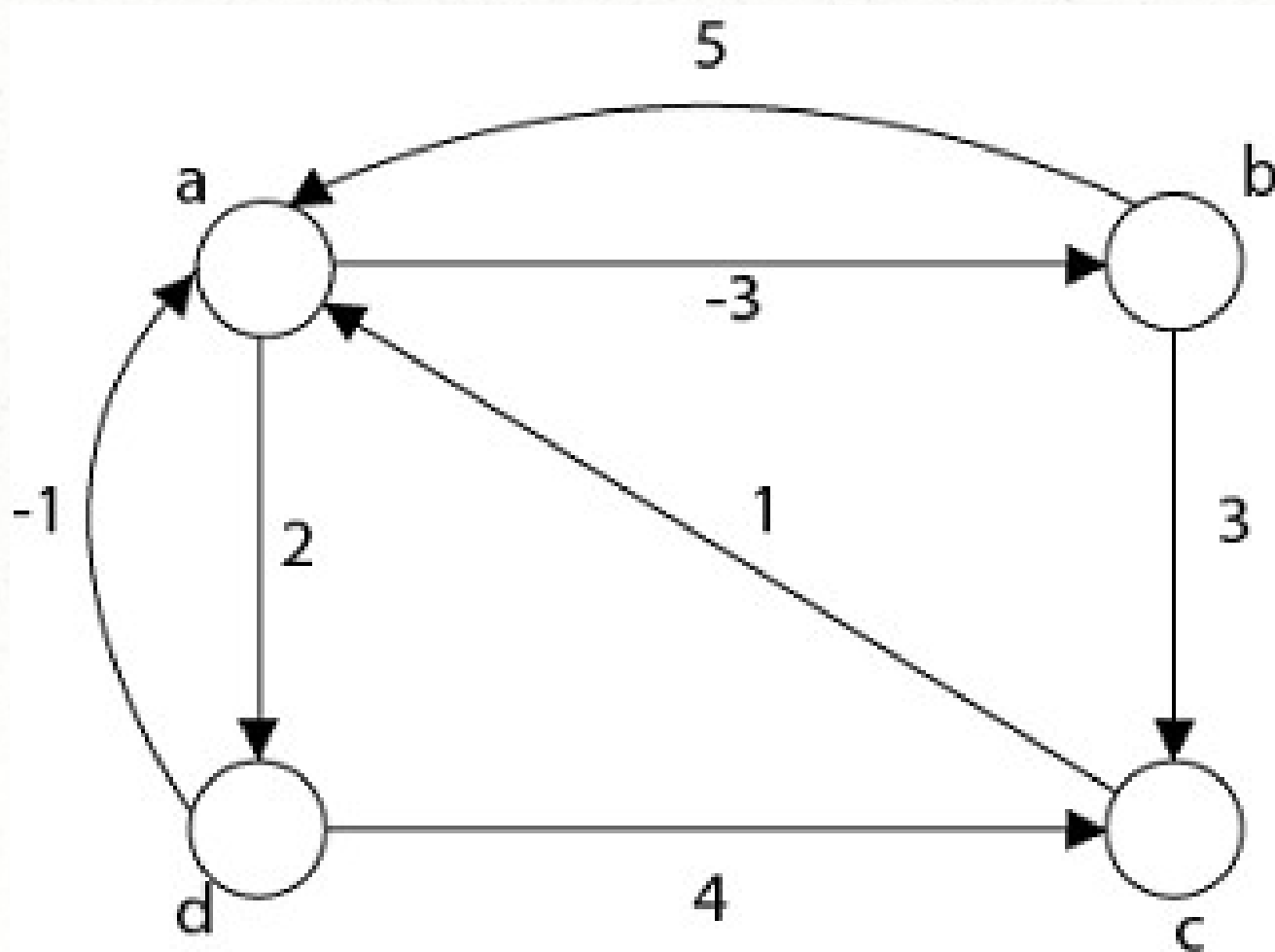
$$w(u, v) = w(u, v) + h(u) - h(v)$$

### JOHNSON (G)

```
1. Compute  $G'$  where  $V[G'] = V[G] \cup \{s\}$  and  
    $E[G'] = E[G] \cup \{(s, v) : v \in V[G]\}$   
  
2. If  $BELLMAN-FORD(G', w, s) = FALSE$   
   then "input graph contains a negative weight cycle"  
   else  
     for each vertex  $v \in V[G']$   
       do  $h(v) \leftarrow \delta(s, v)$   
       Computed by Bellman-Ford algorithm  
     for each edge  $(u, v) \in E[G']$   
       do  $w(u, v) \leftarrow w(u, v) + h(u) - h(v)$   
     for each vertex  $u \in V[G]$   
       do run DIJKSTRA( $G, w, u$ ) to compute  
          $\delta(u, v)$  for all  $v \in V[G]$   
       for each vertex  $v \in V[G]$   
         do  $d_{uv} \leftarrow \delta(u, v) + h(v) - h(u)$   
Return D.
```



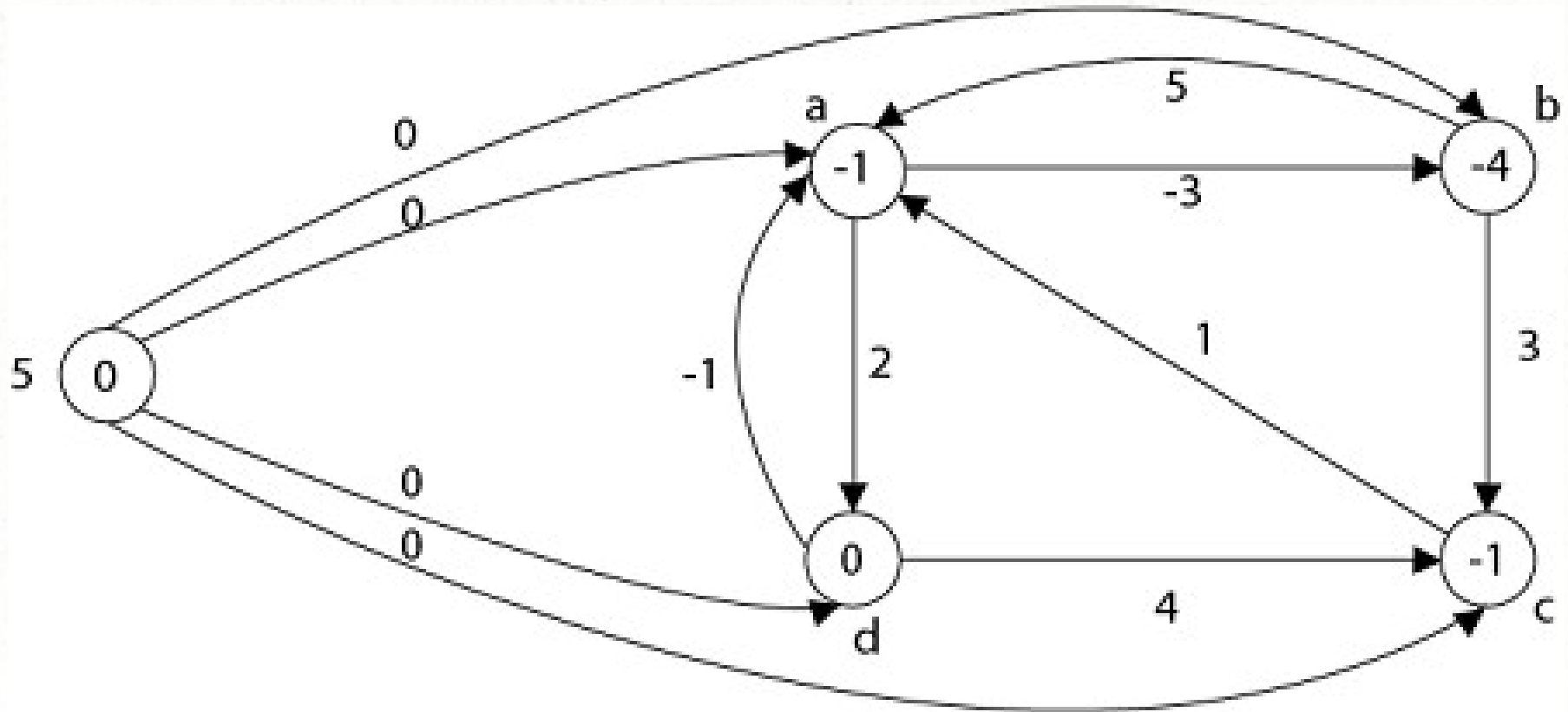
# Johnson Algorithm



# Johnson Algorithm

- Step1: Take any source vertex's' outside the graph and make distance from's' to every vertex '0'.
- Step2: Apply Bellman-Ford Algorithm and calculate minimum weight on each vertex.

# Johnson Algorithm





# Johnson Algorithm

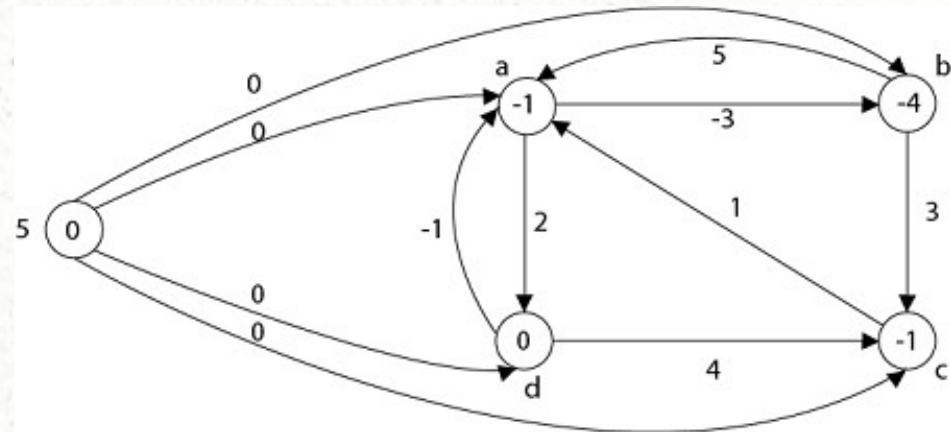
- Step3:

$$\begin{aligned}w(a, b) &= w(a, b) + h(a) - h(b) \\&= -3 + (-1) - (-4) \\&= 0\end{aligned}$$

$$\begin{aligned}w(b, a) &= w(b, a) + h(b) - h(a) \\&= 5 + (-4) - (-1) \\&= 2\end{aligned}$$

$$w(b, c) = w(b, c) + h(b) - h(c)$$

$$\begin{aligned}w(b, c) &= 3 + (-4) - (-1) \\&= 0\end{aligned}$$



# Johnson Algorithm

$$w(c, a) = w(c, a) + h(c) - h(a)$$

$$w(c, a) = 1 + (-1) - (-1) \\ = 1$$

$$w(d, c) = w(d, c) + h(d) - h(c)$$

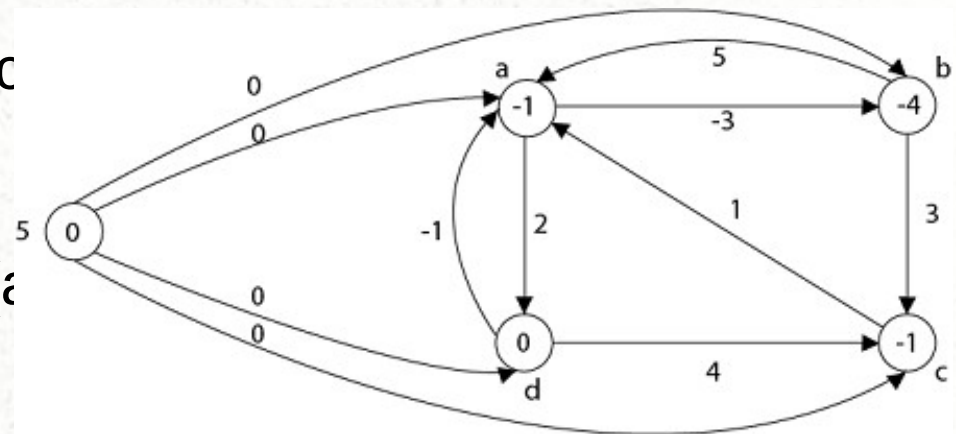
$$w(d, c) = 4 + 0 - (-1) \\ = 5$$

$$w(d, a) = w(d, a) + h(d) - h(a)$$

$$w(d, a) = -1 + 0 - (-1) \\ = 0$$

$$w(a, d) = w(a, d) + h(a) - h(d)$$

$$w(a, d) = 2 + (-1) - 0 = 1$$



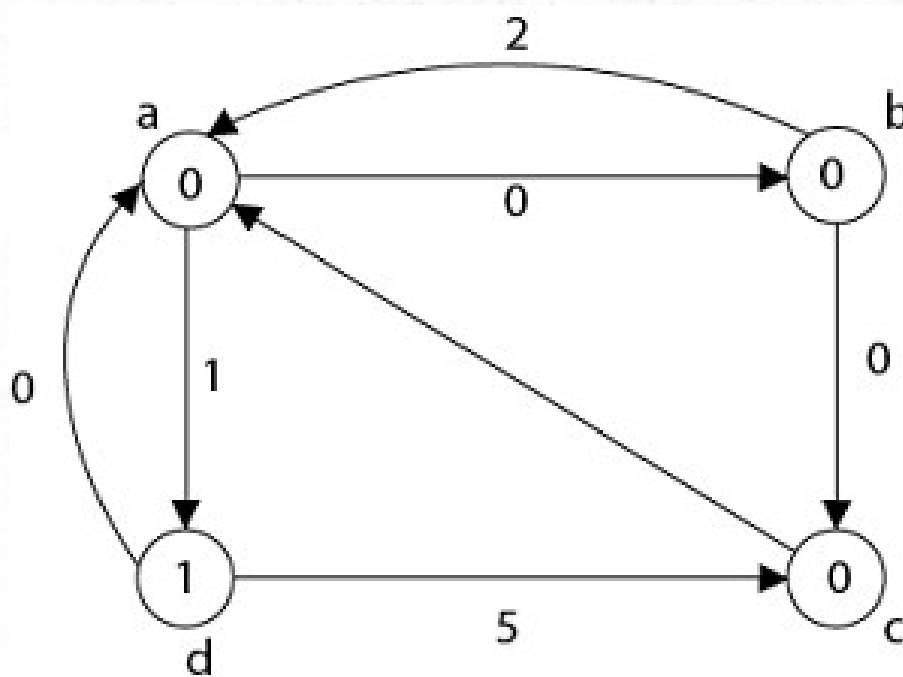
# Johnson Algorithm

- Step 4: Now all edge weights are positive and now we can apply Dijkstra's Algorithm on each vertex and make a matrix corresponds to each vertex in a graph



# Johnson Algorithm

- Case 1: 'a' as a source vertex



a, a      0

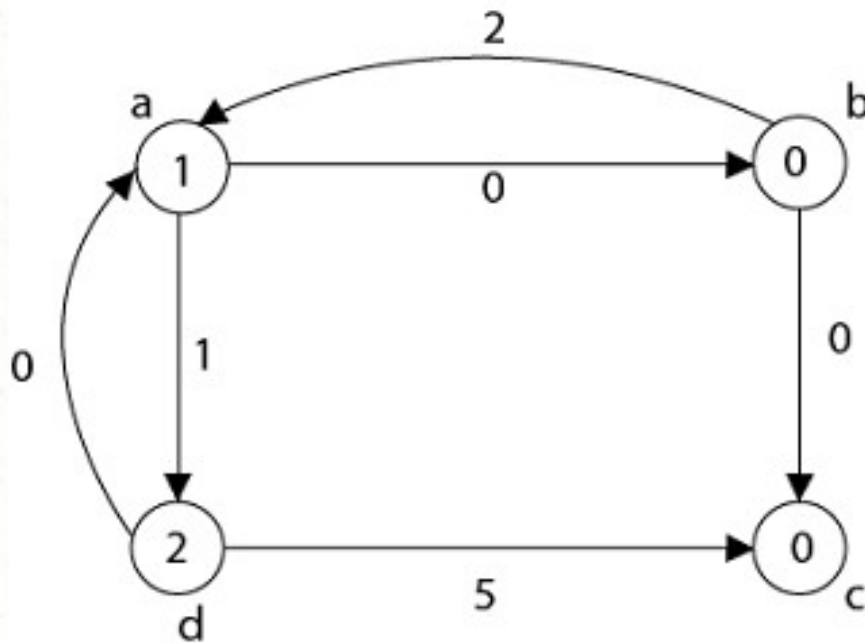
a, b      0

a, c      0

a, d      1

# Johnson Algorithm

- Case 2: 'b' as a source vertex



b, a

2

b, b

0

b, c

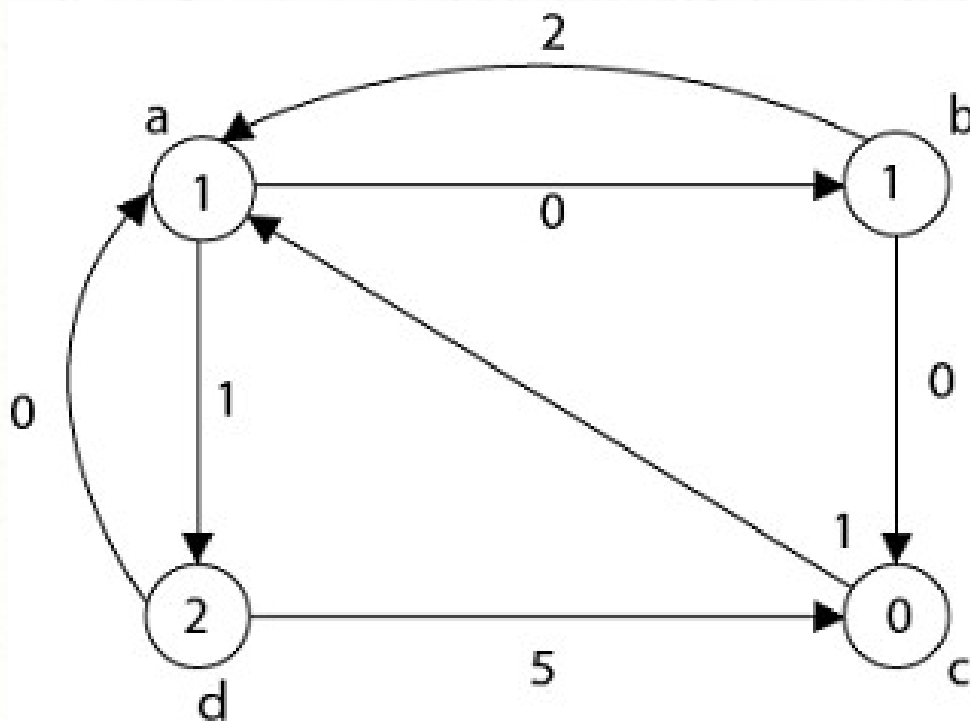
0

b, d

2

# Johnson Algorithm

- Case 3: 'c' as a source vertex



c, a

1

c, b

1

c, c

0

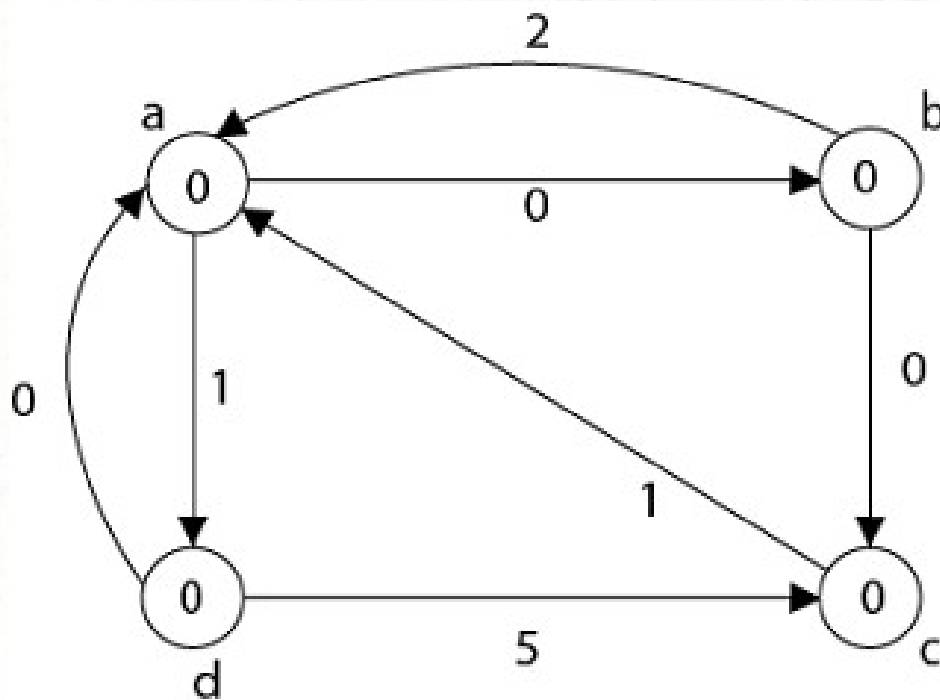
c, d

2



# Johnson Algorithm

- Case4: 'd' as source vertex



d, a

0

d, b

0

d, c

0

d, d

0

# Johnson Algorithm

	<b>a</b>	<b>b</b>	<b>c</b>	<b>d</b>
<b>a</b>	0	0	0	1
<b>b</b>	1	0	0	2
<b>c</b>	1	1	0	2
<b>d</b>	0	0	0	0

# Johnson Algorithm

- Step 5:

- $d_{uv} \leftarrow \delta(u, v) + h(v) - h(u)$

- $d(a, a) = 0 + (-1) - (-1) = 0$

- $d(a, b) = 0 + (-4) - (-1) = -3$

- $d(a, c) = 0 + (-1) - (-1) = 0$

- $d(a, d) = 1 + (0) - (-1) = 2$

- $d(b, a) = 1 + (-1) - (-4) = 4$

- $d(b, b) = 0 + (-4) - (-4) = 0$

- $d(c, a) = 1 + (-1) - (-1) = 1$

- $d(c, b) = 1 + (-4) - (-1) = -2$

- $d(c, c) = 0$

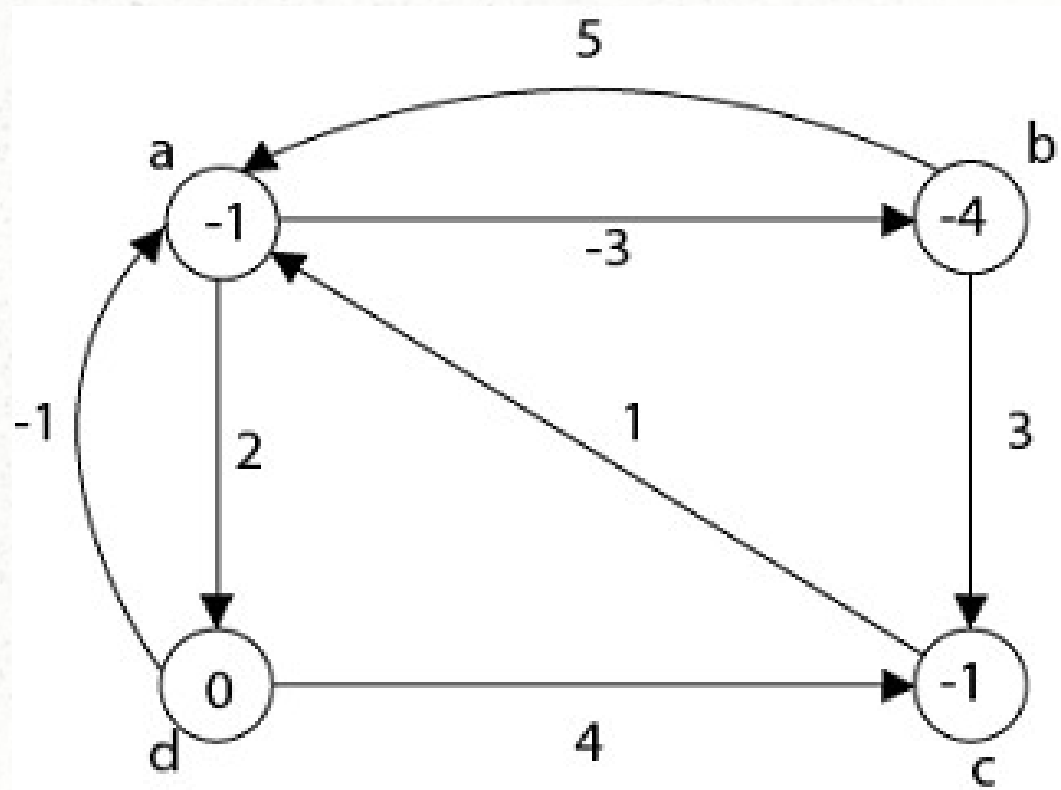
- $d(c, d) = 2 + (0) - (-1) = 3$

- $d(d, a) = 0 + (-1) - (0) = -1$

- $d(d, b) = 0 + (-4) - (0) = -4$

- $d(d, c) = 0 + (-1) - (0) = -1$

- $d(d, d) = 0$





# Johnson Algorithm

	<b>a</b>	<b>b</b>	<b>c</b>	<b>d</b>
<b>a</b>	0	-3	0	2
<b>b</b>	4	0	3	6
<b>c</b>	1	-2	0	3
<b>d</b>	-1	-4	-1	0

# Johnson Algorithm

- Time Complexity
- Time complexity of Floyd Warshal's Algorithm is  $O(V^3)$
- The time complexity of Bellman ford algorithm is  $O(VE)$
- Using Max Heap, Dijkstra's complexity is  $O(E \log V)$ , otherwise it is  $O(V^2)$
- The time complexity of Johnson algorithm is  $O(VE \log V)$

***Thank you***  
***Any Question***  
***??????????***