Chapter 2 Application Layer

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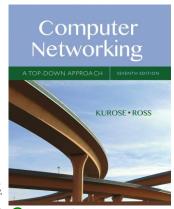
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Computer Networking: A Top Down Approach

7th edition Jim Kurose, Keith Ross Pearson/Addison Wesley April 2016

Application Layer 2-1

Chapter 2: outline

- 2.1 principles of network applications
- 2.2 Web and HTTP
- 2.3 electronic mail
 - SMTP, POP3, IMAP
- **2.4 DNS**

Chapter 2: application layer

our goals:

- conceptual, implementation aspects of network application protocols
 - transport-layer service models
 - client-server paradigm
 - peer-to-peer paradigm
 - content distribution networks

- learn about protocols by examining popular application-level protocols
 - HTTP
 - SMTP / POP3 / IMAP
 - DNS
- creating network applications
 - socket API

Application Layer 2-3

Some network apps

- e-mail
- web
- text messaging
- remote login
- P2P file sharing
- multi-user network games
- streaming stored video (YouTube, Hulu, Netflix)

- voice over IP (e.g., Skype)
- real-time video conferencing
- social networking
- search
- ..
- ...

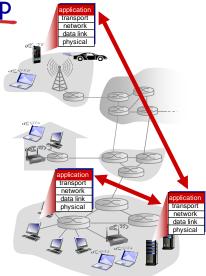
Creating a network app

write programs that:

- run on (different) end systems
- communicate over network
- e.g., web server software communicates with browser software

no need to write software for network-core devices

- network-core devices do not run user applications
- applications on end systems allows for rapid app development, propagation



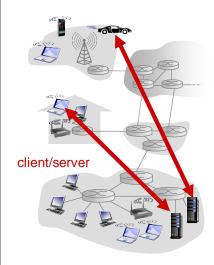
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Application architectures

possible structure of applications:

- client-server
- peer-to-peer (P2P)

Client-server architecture



server:

- always-on host
- permanent IP address
- data centers for scaling

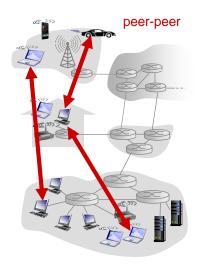
clients:

- communicate with server
- may be intermittently connected
- may have dynamic IP addresses
- do not communicate directly with each other

Application Layer 2-7

P2P architecture

- no always-on server
- arbitrary end systems directly communicate
- peers request service from other peers, provide service in return to other peers
 - self scalability new peers bring new service capacity, as well as new service demands
- peers are intermittently connected and change IP addresses
 - complex management



Processes communicating

process: program running within a host

- within same host, two processes communicate using inter-process communication (defined by OS)
- processes in different hosts communicate by exchanging messages

clients, servers

client process: process that initiates communication

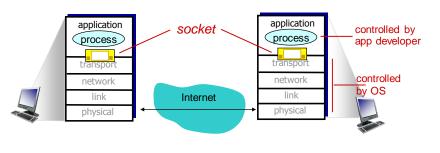
server process: process that waits to be contacted

 aside: applications with P2P architectures have client processes & server processes

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Sockets

- process sends/receives messages to/from its socket
- socket analogous to door
 - sending process shoves message out door
 - sending process relies on transport infrastructure on other side of door to deliver message to socket at receiving process



Addressing processes

- to receive messages, process must have identifier
- host device has unique 32bit IP address
- Q: does IP address of host on which process runs suffice for identifying the process?
 - A: no, many processes can be running on same host
- identifier includes both IP address and port numbers associated with process on host.
- example port numbers:
 - HTTP server: 80
 - mail server: 25
- to send HTTP message to gaia.cs.umass.edu web server:
 - IP address: 128.119.245.12
 - port number: 80
- more shortly...

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App-layer protocol defines

- types of messages exchanged,
 - e.g., request, response
- message syntax:
 - what fields in messages & how fields are delineated
- message semantics
 - meaning of information in fields
- rules for when and how processes send & respond to messages

open protocols:

- defined in RFCs
- allows for interoperability
- e.g., HTTP, SMTP

proprietary protocols:

e.g., Skype

What transport service does an app need?

data integrity

- some apps (e.g., file transfer, web transactions) require
 100% reliable data transfer
- other apps (e.g., audio) can tolerate some loss

timing

 some apps (e.g., Internet telephony, interactive games) require low delay to be "effective"

throughput

- some apps (e.g., multimedia) require minimum amount of throughput to be "effective"
- other apps ("elastic apps") make use of whatever throughput they get

security

encryption, data integrity,

Application Layer 2-13

Transport service requirements: common apps

	application	data loss	throughput	time sensitive
'				
	file transfer	no loss	elastic	no
V	e-mail	no loss	elastic	no
	Veb documents	no loss	elastic	no
real-ti	me audio/video	loss-tolerant	audio: 5kbps-1Mbps	yes, 100's
			video:10kbps-5Mbps	s msec
sto	red audio/video	loss-tolerant	same as above	
in	eractive games	loss-tolerant	few kbps up	yes, few secs
	text messaging	no loss	elastic	yes, 100's
_				msec
				yes and no

Internet transport protocols services

TCP service:

- reliable transport between sending and receiving process
- flow control: sender won't overwhelm receiver
- congestion control: throttle sender when network overloaded
- does not provide: timing, minimum throughput guarantee, security
- connection-oriented: setup required between client and server processes

UDP service:

- unreliable data transfer between sending and receiving process
- does not provide: reliability, flow control, congestion control, timing, throughput guarantee, security, or connection setup,
- Q: why bother? Why is there a UDP?

Application Layer 2-15

Internet apps: application, transport protocols

application	application layer protocol	underlying transport protocol
e-mail	SMTP [RFC 2821]	TCP
remote terminal access	Telnet [RFC 854]	TCP
Web	HTTP [RFC 2616]	TCP
file transfer	FTP [RFC 959]	TCP
streaming multimedia	HTTP (e.g., YouTube),	TCP or UDP
_	RTP [RFC 1889]	
Internet telephony	SIP, RTP, proprietary	
, ,	(e.g., Skype)	TCP or UDP

Securing TCP

TCP & UDP

- no encryption
- cleartext passwds sent into socket traverse Internet in cleartext

SSL

- provides encrypted TCP connection
- data integrity
- end-point authentication

SSL is at app layer

apps use SSL libraries, that "talk" to TCP

SSL socket API

- cleartext passwords sent into socket traverse Internet encrypted
- see Chapter 8

Application Layer 2-17

Chapter 2: outline

- 2.1 principles of network applications
- 2.2 Web and HTTP
- 2.3 electronic mail
 - SMTP, POP3, IMAP
- **2.4 DNS**

Web and HTTP

First, a review...

- web page consists of objects
- object can be HTML file, JPEG image, Java applet, audio file,...
- web page consists of base HTML-file which includes several referenced objects
- each object is addressable by a URL, e.g.,

www.someschool.edu/someDept/pic.gif

host name

path name

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HTTP overview

HTTP: hypertext transfer protocol

- Web's application layer protocol
- client/server model
 - client: browser that requests, receives, (using HTTP protocol) and "displays" Web objects
 - server: Web server sends (using HTTP protocol) objects in response to requests



HTTP overview (continued)

uses TCP:

- client initiates TCP connection (creates socket) to server, port 80
- server accepts TCP connection from client
- HTTP messages

 (application-layer protocol messages) exchanged
 between browser (HTTP client) and Web server
 (HTTP server)
- TCP connection closed

HTTP is "stateless"

 server maintains no information about past client requests

- aside

protocols that maintain "state" are complex!

- past history (state) must be maintained
- if server/client crashes, their views of "state" may be inconsistent, must be reconciled

Application Layer 2-21

HTTP connections

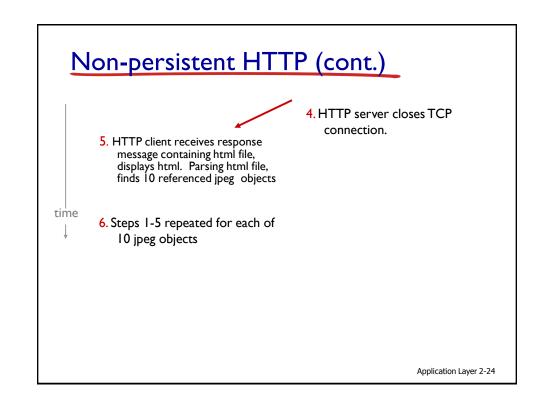
non-persistent HTTP

- at most one object sent over TCP connection
 - connection then closed
- downloading multiple objects required multiple connections

persistent HTTP

 multiple objects can be sent over single TCP connection between client, server

Non-persistent HTTP suppose user enters URL: (contains text, www.someSchool.edu/someDepartment/home.index references to 10 jpeg images) la. HTTP client initiates TCP connection to HTTP server Ib. HTTP server at host (process) at www.someSchool.edu waiting www.someSchool.edu on port 80 for TCP connection at port 80. "accepts" connection, notifying 2. HTTP client sends HTTP request message (containing URL) into 3. HTTP server receives request TCP connection socket. message, forms response Message indicates that client message containing requested wants object object, and sends message into someDepartment/home.index its socket time Application Layer 2-23



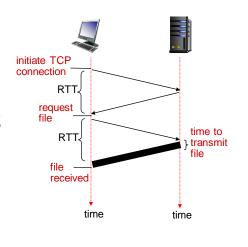
Non-persistent HTTP: response time

RTT (definition): time for a small packet to travel from client to server and back

HTTP response time:

- one RTT to initiate TCP connection
- one RTT for HTTP request and first few bytes of HTTP response to return
- file transmission time
- non-persistent HTTP response time =

2RTT+ file transmission time



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Persistent HTTP

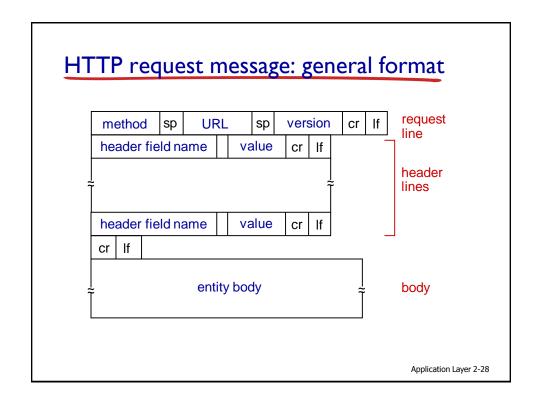
non-persistent HTTP issues:

- requires 2 RTTs per object
- OS overhead for each TCP connection
- browsers often open parallel TCP connections to fetch referenced objects

persistent HTTP:

- server leaves connection open after sending response
- subsequent HTTP messages between same client/server sent over open connection
- client sends requests as soon as it encounters a referenced object
- as little as one RTT for all the referenced objects

HTTP request message two types of HTTP messages: request, response HTTP request message: ASCII (human-readable format) carriage return character line-feed character request line (GET, POST, GET /index.html HTTP/1.1\r\n Host: www-net.cs.umass.edu\r\n **HEAD** commands) User-Agent: Firefox/3.6.10\r\n Accept: text/html,application/xhtml+xml\r\n header Accept-Language: en-us,en;q=0.5\r\n lines Accept-Encoding: gzip,deflate\r\n Accept-Charset: ISO-8859-1,utf-8;q=0.7\r\n carriage return, Keep-Alive: 115\r\n line feed at start Connection: keep-alive\r\n of line indicates \r\n end of header lines * Check out the online interactive exercises for more examples: http://gaia.cs.umass.edu/kurose_ross/interactive/ Application Layer 2-27



Uploading form input

POST method:

- web page often includes form input
- input is uploaded to server in entity body

URL method:

- uses GET method
- input is uploaded in URL field of request line:

www.somesite.com/animalsearch?monkeys&banana

Application Layer 2-29

Method types

HTTP/1.0:

- GET
- POST
- HEAD
 - asks server to leave requested object out of response

HTTP/I.I:

- GET, POST, HEAD
- PUT
 - uploads file in entity body to path specified in URL field
- DELETE
 - deletes file specified in the URL field

HTTP response message

```
status line
(protocol -
               HTTP/1.1 200 OK\r\n
status code
               Date: Sun, 26 Sep 2010 20:09:20 GMT\r\n
status phrase)
               Server: Apache/2.0.52 (CentOS) \r\n
               Last-Modified: Tue, 30 Oct 2007 17:00:02
                  GMT\r\n
               ETag: "17dc6-a5c-bf716880"\r\n
     header
               Accept-Ranges: bytes\r\n
       lines
               Content-Length: 2652\r\n
               Keep-Alive: timeout=10, max=100\rn
                Connection: Keep-Alive\r\n
                Content-Type: text/html; charset=ISO-8859-
                  1\r\n
data, e.g.,
                \r\n
requested
               data data data data ...
HTML file
```

HTTP response status codes

- status code appears in 1st line in server-toclient response message.
- some sample codes:

* Check out the online interactive exercises for more examples: http://gaia.cs.umass.edu/kurose_ross/interactive/

- 200 OK
 - request succeeded, requested object later in this msg
- 301 Moved Permanently
 - requested object moved, new location specified later in this msg (Location:)
- 400 Bad Request
 - request msg not understood by server
- 404 Not Found
 - requested document not found on this server
- 505 HTTP Version Not Supported

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Trying out HTTP (client side) for yourself

I. Telnet to your favorite Web server:

telnet gaia.cs.umass.edu 80 opens TCP connection to port 80 (default HTTP server port) at gaia.cs.umass.edu. anything typed in will be sent to port 80 at gaia.cs.umass.edu

2. type in a GET HTTP request:

GET /kurose_ross/interactive/index.php HTTP/1.1

Host: gaia.cs.umass.edu

by typing this in (hit carriage return twice), you send this minimal (but complete)

GET request to HTTP server

3. look at response message sent by HTTP server! (or use Wireshark to look at captured HTTP request/response)

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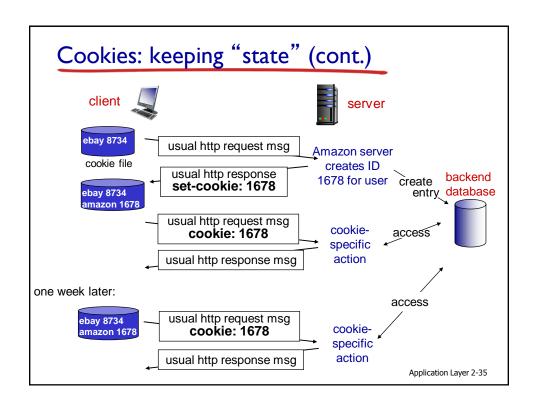
User-server state: cookies

many Web sites use cookies four components:

- I) cookie header line of HTTP response message
- 2) cookie header line in next HTTP request message
- cookie file kept on user's host, managed by user's browser
- 4) back-end database at Web site

example:

- Susan always access Internet from PC
- visits specific e-commerce site for first time
- when initial HTTP requests arrives at site, site creates:
 - unique ID
 - entry in backend database for ID



Cookies (continued)

what cookies can be used for:

- authorization
- shopping carts
- recommendations
- user session state (Web e-mail)

cookies and privacy:

- cookies permit sites to learn a lot about you
- you may supply name and e-mail to sites

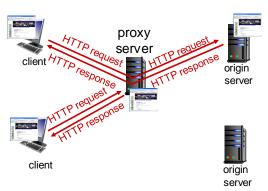
how to keep "state":

- protocol endpoints: maintain state at sender/receiver over multiple transactions
- cookies: http messages carry state

Web caches (proxy server)

goal: satisfy client request without involving origin server

- user sets browser: Web accesses via cache
- browser sends all HTTP requests to cache
 - object in cache: cache returns object
 - else cache requests object from origin server, then returns object to client



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More about Web caching

- cache acts as both client and server
 - server for original requesting client
 - client to origin server
- typically cache is installed by ISP (university, company, residential ISP)

why Web caching?

- reduce response time for client request
- reduce traffic on an institution's access link
- Internet dense with caches: enables "poor" content providers to effectively deliver content (so too does P2P file sharing)

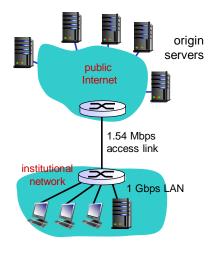
Caching example:

assumptions:

- avg object size: I00K bits
- avg request rate from browsers to origin servers: I 5/sec
- avg data rate to browsers: I.50 Mbps
- RTT from institutional router to any origin server: 2 sec
- access link rate: 1.54 Mbps

consequences:

- LAN utilization: 15% _problem!
- access link utilization 99%
- total delay = Internet delay + access delay + LAN delay
 - = 2 sec + minutes + usecs



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Caching example: fatter access link

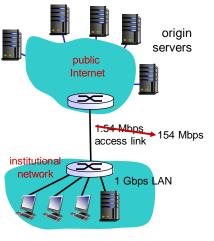
assumptions:

- avg object size: I 00K bits
- avg request rate from browsers to origin servers:15/sec
- avg data rate to browsers: 1.50 Mbps
- RTT from institutional router to any origin server: 2 sec
- access link rate: 1.54 Mbps
 154 Mbps

consequences:

- LAN utilization: 15%
- access link utilization = 99% 9.9%
- total delay = Internet delay + access delay + LAN delay
 - = 2 sec + minutes + usecs msecs

Cost: increased access link speed (not cheap!)



Caching example: install local cache

assumptions:

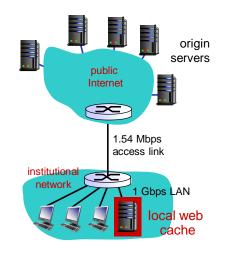
- avg object size: I00K bits
- avg request rate from browsers to origin servers:15/sec
- avg data rate to browsers: I.50 Mbps
- RTT from institutional router to any origin server: 2 sec
- access link rate: 1.54 Mbps

consequences:

- LAN utilization: 15%
- access link utilization = ?
- total delay = ?

How to compute link utilization, delay?

Cost: web cache (cheap!)

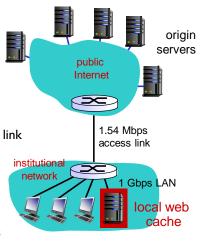


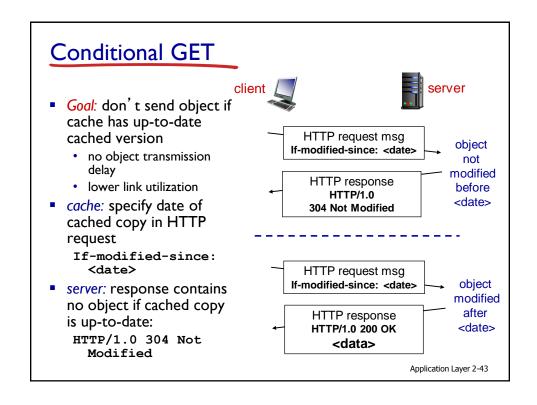
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Caching example: install local cache

Calculating access link utilization, delay with cache:

- suppose cache hit rate is 0.4
 - 40% requests satisfied at cache, 60% requests satisfied at origin
- access link utilization:
 - 60% of requests use access link
- data rate to browsers over access link
 - = 0.6*1.50 Mbps = .9 Mbps
 - utilization = 0.9/1.54 = .58
- total delay
 - = 0.6 * (delay from origin servers) +0.4
 * (delay when satisfied at cache)
 - $= 0.6 (2.01) + 0.4 (\sim msecs) = \sim 1.2 secs$
 - less than with 154 Mbps link (and cheaper too!)





Chapter 2: outline

- 2.1 principles of network applications
- 2.2 Web and HTTP
- 2.3 electronic mail
 - SMTP, POP3, IMAP
- **2.4 DNS**

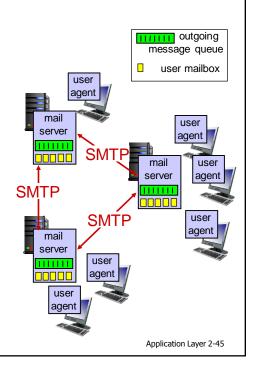
Electronic mail

Three major components:

- user agents
- mail servers
- simple mail transfer protocol: SMTP

User Agent

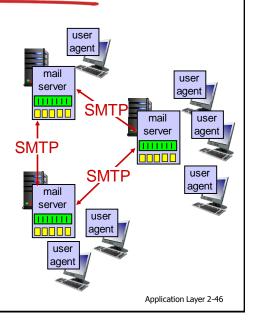
- a.k.a. "mail reader"
- composing, editing, reading mail messages
- e.g., Outlook, Thunderbird, iPhone mail client
- outgoing, incoming messages stored on server



Electronic mail: mail servers

mail servers:

- mailbox contains incoming messages for user
- message queue of outgoing (to be sent) mail messages
- SMTP protocol between mail servers to send email messages
 - client: sending mail server
 - "server": receiving mail server



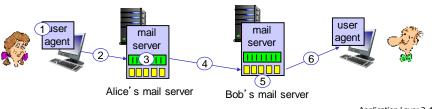
Electronic Mail: SMTP [RFC 2821]

- uses TCP to reliably transfer email message from client to server, port 25
- direct transfer: sending server to receiving server
- three phases of transfer
 - handshaking (greeting)
 - · transfer of messages
 - closure
- command/response interaction (like HTTP)
 - commands: ASCII text
 - · response: status code and phrase
- messages must be in 7-bit ASCI

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Scenario: Alice sends message to Bob

- I) Alice uses UA to compose message "to" bob@someschool.edu
- Alice's UA sends message to her mail server; message placed in message queue
- client side of SMTP opens TCP connection with Bob's mail server
- SMTP client sends Alice's message over the TCP connection
- 5) Bob's mail server places the message in Bob's mailbox
- 6) Bob invokes his user agent to read message



Sample SMTP interaction

```
S: 220 hamburger.edu
C: HELO crepes.fr
S: 250 Hello crepes.fr, pleased to meet you
C: MAIL FROM: <alice@crepes.fr>
S: 250 alice@crepes.fr... Sender ok
C: RCPT TO: <bob@hamburger.edu>
S: 250 bob@hamburger.edu ... Recipient ok
C: DATA
S: 354 Enter mail, end with "." on a line by itself
C: Do you like ketchup?
C: How about pickles?
C: .
S: 250 Message accepted for delivery
C: QUIT
S: 221 hamburger.edu closing connection
```

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Try SMTP interaction for yourself:

- telnet servername 25
- see 220 reply from server
- enter HELO, MAIL FROM, RCPT TO, DATA, QUIT commands

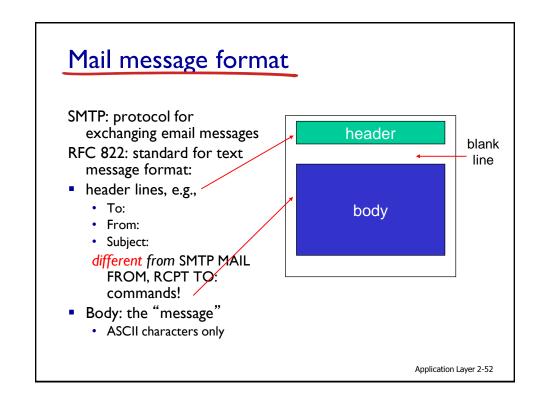
above lets you send email without using email client (reader)

SMTP: final words

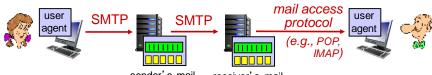
- SMTP uses persistent connections
- SMTP requires message (header & body) to be in 7-bit ASCII
- SMTP server uses
 CRLF.CRLF to
 determine end of message

comparison with HTTP:

- HTTP: pull
- SMTP: push
- both have ASCII command/response interaction, status codes
- HTTP: each object encapsulated in its own response message
- SMTP: multiple objects sent in multipart message



Mail access protocols



sender's mail receiver's mail server server

- SMTP: delivery/storage to receiver's server
- mail access protocol: retrieval from server
 - POP: Post Office Protocol [RFC 1939]: authorization, download
 - IMAP: Internet Mail Access Protocol [RFC 1730]: more features, including manipulation of stored messages on server
 - HTTP: gmail, Hotmail, Yahoo! Mail, etc.

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POP3 protocol

authorization phase

- client commands:
 - user: declare username
 - pass: password
- server responses
 - +OK
 - -ERR

transaction phase, client:

- list: list message numbers
- retr: retrieve message by number
- dele: delete
- quit

S: +OK POP3 server ready

C: user bob

S: +OK

C: pass hungry

 $S: \ + OK \ \,$ user successfully logged on

C: list

s: 1 498

s: 2 912

s: .

C: retr 1

S: <message 1 contents>

s: .

C: dele 1

C: retr 2

S: <message 1 contents>

s: .

C: dele 2

C: quit

S: +OK POP3 server signing off

POP3 (more) and IMAP

more about POP3

- previous example uses POP3 "download and delete" mode
 - Bob cannot re-read email if he changes client
- POP3 "download-andkeep": copies of messages on different clients
- POP3 is stateless across sessions

IMAP

- keeps all messages in one place: at server
- allows user to organize messages in folders
- keeps user state across sessions:
 - names of folders and mappings between message IDs and folder name

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Chapter 2: outline

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- **2.4 DNS**

DNS: domain name system

people: many identifiers:

SSN, name, passport #

Internet hosts, routers:

- IP address (32 bit) used for addressing datagrams
- "name", e.g., www.yahoo.com used by humans
- Q: how to map between IP address and name, and vice versa?

Domain Name System:

- distributed database implemented in hierarchy of many name servers
- application-layer protocol: hosts, name servers communicate to resolve names (address/name translation)
 - note: core Internet function, implemented as applicationlayer protocol
 - complexity at network's "edge"

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DNS: services, structure

DNS services

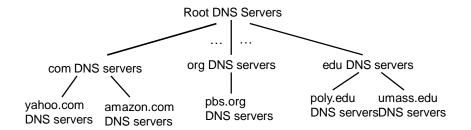
- hostname to IP address translation
- host aliasing
 - · canonical, alias names
- mail server aliasing
- load distribution
 - replicated Web servers: many IP addresses correspond to one name

why not centralize DNS?

- single point of failure
- traffic volume
- distant centralized database
- maintenance

A: doesn't scale!

DNS: a distributed, hierarchical database



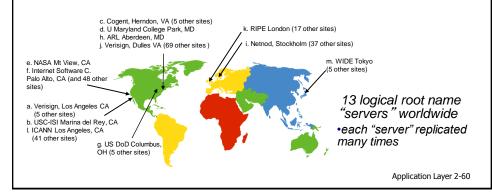
client wants IP for www.amazon.com; 1st approximation:

- client queries root server to find com DNS server
- client queries .com DNS server to get amazon.com DNS server
- client queries amazon.com DNS server to get IP address for www.amazon.com

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DNS: root name servers

- contacted by local name server that can not resolve name
- root name server:
 - contacts authoritative name server if name mapping not known
 - gets mapping
 - returns mapping to local name server



TLD, authoritative servers

top-level domain (TLD) servers:

- responsible for com, org, net, edu, aero, jobs, museums, and all top-level country domains, e.g.: uk, fr, ca, jp
- Network Solutions maintains servers for .com TLD
- · Educause for .edu TLD

authoritative DNS servers:

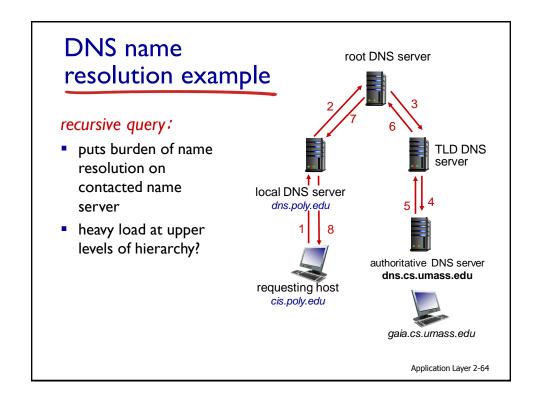
- organization's own DNS server(s), providing authoritative hostname to IP mappings for organization's named hosts
- · can be maintained by organization or service provider

Application Layer 2-61

Local DNS name server

- does not strictly belong to hierarchy
- each ISP (residential ISP, company, university) has one
 - also called "default name server"
- when host makes DNS query, query is sent to its local DNS server
 - has local cache of recent name-to-address translation pairs (but may be out of date!)
 - acts as proxy, forwards query into hierarchy

DNS name root DNS server resolution example host at cis.poly.edu TLD DNS server wants IP address for gaia.cs.umass.edu local DNS server iterated query: dns.poly.edu contacted server replies with name of server to contact authoritative DNS server "I don't know this dns.cs.umass.edu requesting host name, but ask this cis.poly.edu server" gaia.cs.umass.edu Application Layer 2-63



DNS: caching, updating records

- once (any) name server learns mapping, it caches mapping
 - cache entries timeout (disappear) after some time (TTL)
 - TLD servers typically cached in local name servers
 - · thus root name servers not often visited
- cached entries may be <u>out-of-date</u> (best effort name-to-address translation!)
 - if name host changes IP address, may not be known Internet-wide until all TTLs expire
- update/notify mechanisms proposed IETF standard
 - RFC 2136

Application Layer 2-65

DNS records

DNS: distributed database storing resource records (RR)

RR format: (name, value, type, ttl)

type=A

- name is hostname
- value is IP address

type=NS

- name is domain (e.g., foo.com)
- value is hostname of authoritative name server for this domain

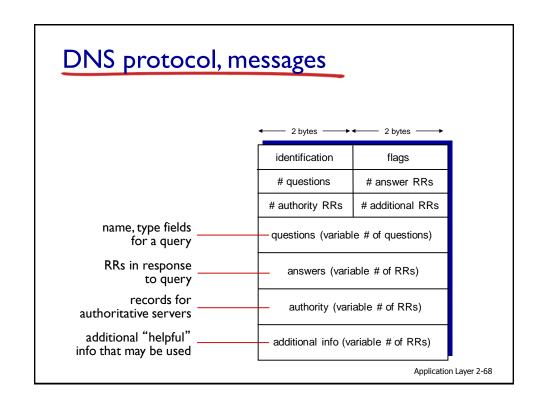
type=CNAME

- name is alias name for some "canonical" (the real) name
- www.ibm.com is really servereast.backup2.ibm.com
- value is canonical name

type=MX

 value is name of mailserver associated with name

DNS protocol, messages query and reply messages, both with same message format 2 bytes -2 bytes identification flags message header identification: 16 bit # for # questions # answer RRs query, reply to query uses # authority RRs # additional RRs same# flags: questions (variable # of questions) query or reply recursion desired answers (variable # of RRs) recursion available authority (variable # of RRs) reply is authoritative additional info (variable # of RRs) Application Layer 2-67



Inserting records into DNS

- example: new startup "Network Utopia"
- register name networkuptopia.com at DNS registrar (e.g., Network Solutions)
 - provide names, IP addresses of authoritative name server (primary and secondary)
 - registrar inserts two RRs into .com TLD server: (networkutopia.com, dns1.networkutopia.com, NS) (dns1.networkutopia.com, 212.212.212.1, A)
- create authoritative server type A record for www.networkuptopia.com; type MX record for networkutopia.com

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Chapter 2: outline

- 2.1 principles of network applications
- 2.2 Web and HTTP
- 2.3 electronic mail
 - SMTP, POP3, IMAP
- **2.4 DNS**