CS761 Artificial Intelligence

Tutorial 1: Agent Driven by Propositional Logic & A Basic Tutorial for Prolog Agent Driven by Propositional Logic

Recall: Knowledge-based Agent

The wumpus world scenario may be realised by a knowledge-based agent.

```
Knowledge-based Agent Program Framework
```

```
INPUT: percept, background knowledge KB, clock t (initially 0) OUTPUT: An action

Tell(KB, MakePerceptSentence(percept, t))

action \leftarrow Ask(KB, MakeActionQuery(percept, t))

return Tell(KB, MakeActionSentence(action, t)

t \leftarrow t + 1

return action
```

Two important operations:

- *Tell*: Add a sentence to KB.
 - We introduced proposition logic as a knowledge representation language.
 - A clause is a proposition of the form

$$(h_1 \lor h_2 \lor \cdots \lor h_m) \leftarrow (\ell_1 \land \ell_2 \land \cdots \land \ell_k)$$

- A propositional knowledge base is a set of clauses.
- Ask: Reason if a sentence is entailed by the KB
 - A model of a propositional knowledge base KB is an interpretation that satisfies KB.
 - A proposition g is called a <u>logical consequence</u> of a knowledge base KB, written as

$$KB \models g$$

if *g* is true in every model of KB.

An inference engine decides for any KB, a set of percept atoms
 Percept, and proposition g, whether

$$KB \cup Percept \models g$$

A Propositional Wumpus World Agent

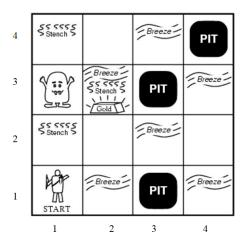
We now bring together all points above in order to construct a wumpus world agent using propositional logic.

We need to handle the following questions:

- World model:
 - How does the world evolves independently of the agent?
 - How do the agent's actions affect the world?
- Internal state of the agent:
 - How does the agent perceive the environment?
 - How does the agent update its knowledge base?
- Inference:

How does the agent use logical inference to make decisions?

The wumpus world:



States: Static knowledge

Define the atomic propositions: For any $i, j \in \{1, 2, 3, 4\}^2$

- $B_{i,j}$: Breeze can be sensed at location (i, j)
- $P_{i,j}$: Location (i, j) has a pit in it
- $W_{i,j}$: Location (i, j) has a wumpus in it

In KB, add the static knowledge axioms:

- Starting location: $\neg P_{1,1}$, $\neg W_{1,1}$
- $\bullet \ B_{1,1} \leftrightarrow (P_{1,2} \vee P_{2,1})$
- $\bullet \ S_{1,1} \leftrightarrow (W_{1,2} \lor W_{2,1})$
- There is at least one wumpus:
 - $\bullet \ \, W_{1,1} \vee W_{1,2} \vee \cdots \vee W_{4,3} \vee W_{4,4} \\$
- and others ...

States: Dynamic Knowledge

Define the atomic propositions: For $t \in \mathbb{N}$, $i, j \in \{1, 2, 3, 4\}$,

- $WumpusAlive^t$: Wumpus is alive at time t
- $OK_{i,j}^t$: It is OK to move to position (i, j) at time t
- $L_{i,i}^t$: Agent is located at (i, j) at time t
- $FacingEast^t$: The agent is facing east at time t

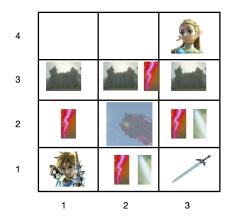
In KB, add the following dynamic knowledge axioms:

- $L_{1,1}^0$
- FacingEast⁰
- $\bullet \ OK_{i,j}^t \leftrightarrow (\neg P_{i,j} \land \neg (W_{i,j} \land WumpusAlive^t))$
- and maybe others

Q2 in Assignment 2

For any $i \in \{1, 2, 3\}$ and $j \in \{1, 2, 3, 4\}$

- $C_{i,j}$: Calamity can be sensed at location (i, j)
- $G_{i,j}$: Ganon is at location (i, j)
- $S_{i,j}$: Master sword is at location (i, j)
- HoldingSword: Link has picked up the master sword



Q2 in Assignment 2

Questions:

- (a) Use atomic propositions given above to represent the following static knowledge axioms:
 - (1) Ganon must be at one location. (You do NOT need to assume whether Ganon is unique or not.)
 - (2) Ganon and master sword are not at the same location.
 - (3) Calamity cannot be sensed at two adjacent locations at the same time.
- (b) Use atomic propositions given above to define dynamic knowledge axiom $OK_{i,j}$ that denotes that it is safe to move to location (i, j).

Note: The explanation for how a formulas is derived is **NOT** required.

A Basic Tutorial for Prolog

Imperative v.s. Declarative Programming

Imperative programs: describe the steps of computation, i.e., how to produce an output.

Example. The procedure to sum a list of integers.

```
int sum(int[] list){
   int result = 0;
   for (int i=0; i<list.length;++i){
      result += list[i];
   }
   return result;
}</pre>
```

Declarative programs: describe what output you want, rather than how to produce the output.

Example. The definition of the sum of a list of integers.

Logic Programming

Logic programming is a declarative programming language paradigm that aims to separate a program into its logic component and its control component:

- the logic component applies logic as a knowledge description language to describe what to compute;
- the control component applies inference to solve the problem.

Prolog is one of the most important logic programming languages:

- Developed in Marseille, France in 1972.
- First-order logic as its knowledge representation language.
- Widely used in theorem proving, expert systems, automated planning, and natural language processing.
- Download and install the client (Windows, Linux, MacOS (not recommended)):



https://www.swi-prolog.org/download/stable Online version (recommended): https://swish.swi-prolog.org/ **Example.** "The law says that it is a crime for a New Zealander to sell alcohol to minors. The girl Lucy is 15 years old and has some beers. All of the beers were sold to her by David, who is a New Zealander." Prove that David is guilty.

- "it is a crime for a New Zealander to sell alcohol to minors"
 - (1) $Crime(x) \leftarrow NZ(x) \land Alcohol(y) \land Sells(x, y, z) \land Minor(z)$
- "Lucy is 15 years old and has some beers"

"All of the beers were sold to her by David"

(5)
$$Sells(David, x, Lucy) \leftarrow Beers(x) \land Owns(Lucy, x)$$

Beers are alcohol

$$(6) Alcohol(x) \leftarrow Beers(x)$$

• Anyone younger than 17 years old is a minor.

$$(7) Minor(x) \leftarrow Under 17(x)$$

"David, who is a New Zealander"

Query: Ask(KB, Crime(David))

Prolog program:

```
crime(X) :- nz(X), alcohol(Y), sells(X,Y,Z), minor(Z).
owns(lucy,b).
beers(b).
under17(lucy).
sells(david,X,lucy) :- beers(X), owns(lucy,X).
alcohol(X) :- beers(X).
minor(X) :- under17(X).
nz(david).
```

Query:

```
?- crime(david).
true.
?- crime(lucy).
false.
```

Prolog Program

Basic Syntax. A Prolog program implements a first-order knowledge base:

```
Facts: AtomsE.g. rainy(dunedin). cold(dunedin). play.
```

```
    Rules: Head: Body<sub>1</sub>, Body<sub>2</sub>,..., Body<sub>k</sub>.
    E.g. snowy(X): rainy(X), cold(X).
```

Note.

- Constant, function, relation names all start with small case letters. E.g. dunedin, david, sum, sells, "ET".
- Variable names start with capital letters or "_". E.g. X, Lucy, _x
- Left implication: ":-"
- Conjunction: ","

Rules. A rule in Prolog

```
Head :- Body_1, Body_2, ..., Body_k.
```

represents a logical implication.

 Universal quantification: Variables in the head are universally quantified.

```
cold(wellington).
rainy(auckland).
snowy(X) :- rainy(X), cold(X).
```

The last statement can be viewed as $\forall x : (rainy(x) \land cold(x)) \rightarrow snowy(x)$.

 Existential quantification: Variables that only appear in the body are existentially quantified.

```
in(auckland, northIs).
in(wellington, northIs).
sameIsland(X, Z) :- in(X, Y), in(Z, Y).
```

The last statement can be viewed as

```
\forall x, z \colon [\exists y \colon in(x, y) \land in(z, y)] \rightarrow sameIsland(x, z).
```

```
Queries. Queries start with "?-"
    rainy(dunedin).
    cold(dunedin).
    rainy(wellington).
    cold(wellington).
    rainy(auckland).
    snowy(X) := rainy(X), cold(X).
  True/false queries:
     E.g. ?- snowy(dunedin).
         true.
  • Unification in queries:
     E.g.
         ?- snowy(C).
         C=dunedin:
         C=wellington.
```

Recursive rules. Prolog allows recursive definitions of predicates:

Example.

```
in(hamilton,waikato).
in(waikato,northIs).
belong(X,Y) :- in(X,Y).
belong(X,Y) :- in(X,Z), belong(Z,Y).
?- belong(hamilton,northIs).
true.
```

```
Negation as failure. \+ denotes the negation symbol in Prolog. E.g.
```

```
man(jim).
man(fred).
woman(X) :- \+(man(X)).
?- woman(jim).
false.
?- woman(X).
```

Note.

 \+(literal(x)) is true whenever it is not possible to prove literal(x) to be true, i.e.,

```
+(literal(x)) returns true \Leftrightarrow literal(x) does not exist in KB.
```

• \+(literal(X)) is true whenever it is not possible to unify X with a constant that makes literal(X) true, i.e.,

```
\+(literal(X)) returns true \Leftrightarrow \neg \exists x : literal(x).
```

Example. [a simple KB in Prolog]

```
parent(ann,bob).
parent(abe,bob).
parent(bob,dan).
parent(cat,dan).
parent(ann.ema).
parent(dan, fay).
male(abe).
male(bob).
male(dan).
female(X) :- +(male(X)).
father(X,Y) := parent(X,Y), male(X).
mother(X,Y) := parent(X,Y), female(X).
son(X.Y) := parent(Y.X). male(X).
sister(X,Z) := parent(Y,X), parent(Y,Z), female(X).
aunt(X.Z) :- sister(X.Y). parent(Y.Z).
grandfather(X,Z) := father(X,Y), parent(Y,Z).
ancestor(X,Y) :- parent(X,Y).
ancestor(X,Z) := parent(X,Y), ancestor(Y,Z).
```

Q1 in Assignment 2

Implement a Prolog predicate mapcolouring(N,A,WA,B,G,T,M,H,WE) where the variable X should be the corresponding word as indicated in the following map. Submit your answer as mapcolouring.pl file.

```
% Knowledge bases
 3 different (red, green).
 4 different(red, blue).
 5 different (green, blue).
 6 different (green, red).
 7 different(blue, red).
 8 different(blue, green).
                                                            WA
    Your implementation
   mapcolouring(N,A,WA,B,G,T,M,H,WE):-
13
       % Your code goes here
14
15
```