

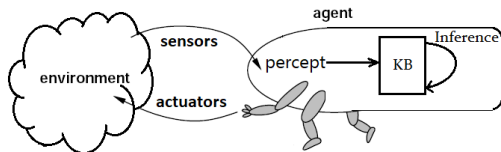
CS761 Artificial Intelligence

First-order Logic Inference: Logic Programming

Recall: Logical agents

A **knowledge-based agent** implements the following:

- **Knowledge base**: Expressing knowledge (i.e., constraints) in a logical language
- **Inference engine**: Performing reasoning to solve problems.



The “**declarative ideal**” of Robert Kowalski (1979):

$$\text{Algorithm} = \text{Logic} + \text{Control}$$



Imperative v.s. Declarative Programming

Imperative programs: describe the steps of computation, i.e., how to produce an output.

Example. The procedure to sum a list of integers.

```
int sum(int[] list){
    int result = 0;
    for (int i=0; i<list.length;++i){
        result += list[i];
    }
    return result;
}
```

Declarative programs: describe what output you want, rather than how to produce the output.

Example. The definition of the sum of a list of integers.

```
sum([],0).           % sum of an empty list is 0
sum([FirstItem | Rest], Sum) :-
    sum(Rest,SumRest), Sum is FirstItem + SumRest.
                    % sum of a non-empty list is the first item +
                    % sum of the rest of the list.
```

Logic Programming

Logic programming is a declarative programming language paradigm that aims to separate a program into its **logic component** and its **control component**:

- the logic component applies logic as a knowledge description language to describe **what to compute**;
- the control component applies inference to **solve the problem**.

Prolog is one of the most important logic programming languages:

- Developed in Marseille, France in 1972.
- First-order logic as its knowledge representation language.
- Widely used in theorem proving, expert systems, automated planning, and natural language processing.
- Download and install the client (Windows, Linux, **MacOS (not recommended)**):

<https://www.swi-prolog.org/download/stable>

Online version (recommended):

<https://swish.swi-prolog.org/>



SWI Prolog

Recall example. *“The law says that it is a crime for a New Zealander to sell alcohol to minors. The girl Lucy is 15 years old and has some beers. All of the beers were sold to her by David, who is a New Zealander.”* Prove that David is guilty.

- *“it is a crime for a New Zealander to sell alcohol to minors”*

$$(1) \text{ Crime}(x) \leftarrow \text{NZ}(x) \wedge \text{Alcohol}(y) \wedge \text{Sells}(x, y, z) \wedge \text{Minor}(z)$$

- *“Lucy is 15 years old and has some beers”*

$$(2) \text{ Owns}(\text{Lucy}, B), (3) \text{ Beers}(B), (4) \text{ Under17}(\text{Lucy})$$

- *“All of the beers were sold to her by David”*

$$(5) \text{ Sells}(\text{David}, x, \text{Lucy}) \leftarrow \text{Beers}(x) \wedge \text{Owns}(\text{Lucy}, x)$$

- Beers are alcohol

$$(6) \text{ Alcohol}(x) \leftarrow \text{Beers}(x)$$

- Anyone younger than 17 years old is a minor.

$$(7) \text{ Minor}(x) \leftarrow \text{Under17}(x)$$

- *“David, who is a New Zealander”*

$$(8) \text{ NZ}(\text{David})$$

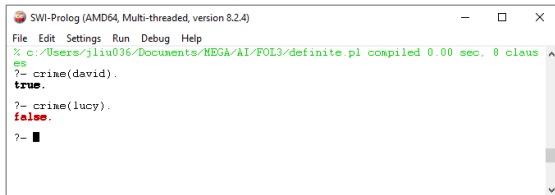
- Query: $\text{Ask}(\text{KB}, \text{Crime}(\text{David}))$

Prolog program:

```
crime(X) :- nz(X), alcohol(Y), sells(X,Y,Z), minor(Z).
owns(lucy,b).
beers(b).
under17(lucy).
sells(david,X,lucy) :- beers(X), owns(lucy,X).
alcohol(X) :- beers(X).
minor(X) :- under17(X).
nz(david).
```

Query:

```
?- crime(david).
true.
?- crime(lucy).
false.
```

A screenshot of the SWI-Prolog IDE window. The title bar reads "SWI-Prolog (AMD64, Multi-threaded, version 8.2.4)". The menu bar includes "File", "Edit", "Settings", "Run", "Debug", and "Help". The main text area shows the compilation status: "% c:/Users/jliu036/Documents/MEGA/AI/FOL3/definite.pl compiled 0.00 sec, 8 claus". Below this, the query results are displayed: "?- crime(david)." followed by "true." in green, and "?- crime(lucy)." followed by "false." in red. The prompt "?- " is followed by a black cursor block.

```
SWI-Prolog (AMD64, Multi-threaded, version 8.2.4)
File Edit Settings Run Debug Help
% c:/Users/jliu036/Documents/MEGA/AI/FOL3/definite.pl compiled 0.00 sec, 8 claus
?- crime(david).
true.
?- crime(lucy).
false.
?- 
```

Prolog Program

Basic Syntax. A Prolog program implements a first-order knowledge base:

- **Facts:** Atoms
E.g. `rainy(dunedin). cold(dunedin). play.`
- **Rules:** `Head :- Body1, Body2, ..., Bodyk.`
E.g. `snowy(X) :- rainy(X), cold(X).`

Note.

- **Constant, function, relation names** all start with small case letters. E.g. `dunedin, david, sum, sells, 'ET'`.
- **Variable names** start with capital letters or `_`. E.g. `X, Lucy, _x`
- Left implication: `":-`
- Conjunction: `","`

Rules. A **rule** in Prolog

$\text{Head} \text{ :- } \text{Body}_1, \text{Body}_2, \dots, \text{Body}_k.$

represents a logical implication.

- **Universal quantification:** Variables in the head are universally quantified.

```
cold(wellington).  
rainy(auckland).  
snowy(X) :- rainy(X), cold(X).
```

The last statement can be viewed as $\forall x: (\text{rainy}(x) \wedge \text{cold}(x)) \rightarrow \text{snowy}(x).$

- **Existential quantification:** Variables that only appear in the body are existentially quantified.

```
in(auckland, northIs).  
in(wellington, northIs).  
sameIsland(X, Z) :- in(X, Y), in(Z, Y).
```

The last statement can be viewed as

$\forall x, z: [\exists y: \text{in}(x, y) \wedge \text{in}(z, y)] \rightarrow \text{sameIsland}(x, z).$

Queries. Queries start with “?-”

```
rainy(dunedin).  
cold(dunedin).  
rainy(wellington).  
cold(wellington).  
rainy(auckland).  
snowy(X) :- rainy(X), cold(X).
```

- True/false queries:

```
E.g. ?- snowy(dunedin).  
      true.
```

- Unification in queries:

```
E.g.  
    ?- snowy(C).  
       C=dunedin ;  
       C=wellington.
```

Recursive rules. Prolog allows recursive definitions of predicates:

Example.

```
in(hamilton,waikato).  
in(waikato,northIs).  
belong(X,Y) :- in(X,Y).  
belong(X,Y) :- in(X,Z), belong(Z,Y).  
  
?- belong(hamilton,northIs).  
true.
```

Negation as failure. $\backslash +$ denotes the negation symbol in Prolog.
E.g.

```
man(jim).  
man(fred).  
woman(X) :- \+(man(X)).  
?- woman(jim).  
false.  
?- woman(X).  
false.
```

Note.

- $\backslash+(literal(x))$ is true whenever it is not possible to prove $literal(x)$ to be true, i.e.,

$\backslash+(literal(x))$ returns true $\Leftrightarrow literal(x)$ does not exist in KB.

- $\backslash+(literal(X))$ is true whenever it is not possible to unify X with a constant that makes $literal(X)$ true, i.e.,

$\backslash+(literal(X))$ returns true $\Leftrightarrow \neg \exists x: literal(x)$.

Example. [a simple KB in Prolog]

```
parent(ann,bob).
parent(abe,bob).
parent(bob,dan).
parent(cat,dan).
parent(ann,ema).
parent(dan,fay).
male(abe).
male(bob).
male(dan).
female(X) :- \+(male(X)).

father(X,Y) :- parent(X,Y), male(X).
mother(X,Y) :- parent(X,Y), female(X).
son(X,Y) :- parent(Y,X), male(X).
sister(X,Z) :- parent(Y,X), parent(Y,Z), female(X).
aunt(X,Z) :- sister(X,Y), parent(Y,Z).
grandfather(X,Z) :- father(X,Y), parent(Y,Z).

ancestor(X,Y) :- parent(X,Y).
ancestor(X,Z) :- parent(X,Y), ancestor(Y,Z).
```

Prolog Data Structure: Lists

Prolog defines a list using the `[]` notation:

- `[a, b, c]`
- `[a | [b, c]], [a, b | [c]], [a, b, c | []]`

Checking list membership:

```
member(X, [X | _]).  
member(X, [_ | T]) :- member(X, T).
```

Checking if a list is sorted:

```
sorted([]).                % empty list is sorted  
sorted([_]).               % singleton is sorted  
sorted([A, B | T]) :- A <= B, sorted([B | T]). % <= is a built-in predicate
```

Appending elements to a list:

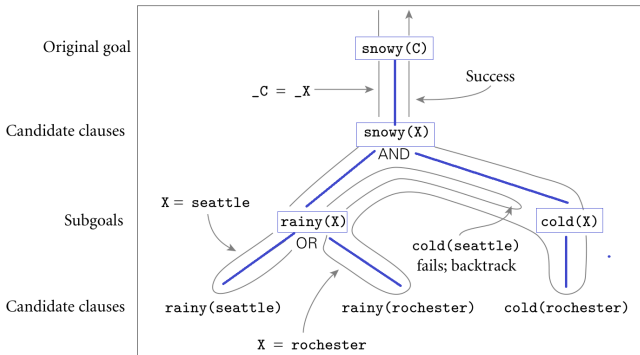
```
append([], A, A).  
append([H | T], A, [H | L]) :- append(T, A, L).  
?- append([a, b, c], [d, e], L).  
L = [a, b, c, d, e].  
?- append(X, [d, e], [a, b, c, d, e]).  
X = [a, b, c].
```

Inference in Prolog

Backward chaining. Start with query and work backward, attempting to “unresolve” it into a set of pre-existing clauses.

- The backward chaining search can be described by a **tree of subgoals**.
- Prolog explores this tree using **backtracking**: **depth-first, from top line to bottom line**.

```
rainy(seattle).  
rainy(rochester).  
cold(rochester).  
snowy(X) :- rainy(X), cold(X).
```



However, Prolog is not a fully logical programming language. We are going to discuss three important “imperative features” of Prolog:

- Ordering of literals
- Cut
- Arithmetic

Imperative feature 1: Ordering

- During search, Prolog considers clauses in the **fixed order** from first to last.
- Prolog programs **do not** have the same semantics as first-order logic.
- Programmers must **be careful with the order of clauses** to ensure recursive programs will terminate.

Example. The following code may find all pairs (X,Y) with path(X,Y).

```
edge(a, b). edge(b, c). edge(c, d).  
edge(d, e). edge(b, e). edge(d, f).  
path(X, X).  
path(X, Y) :- edge(Z, Y), path(X, Z).  
?- path(a, a).  
true.
```

The following code will produce **an error** with path(X,Y).

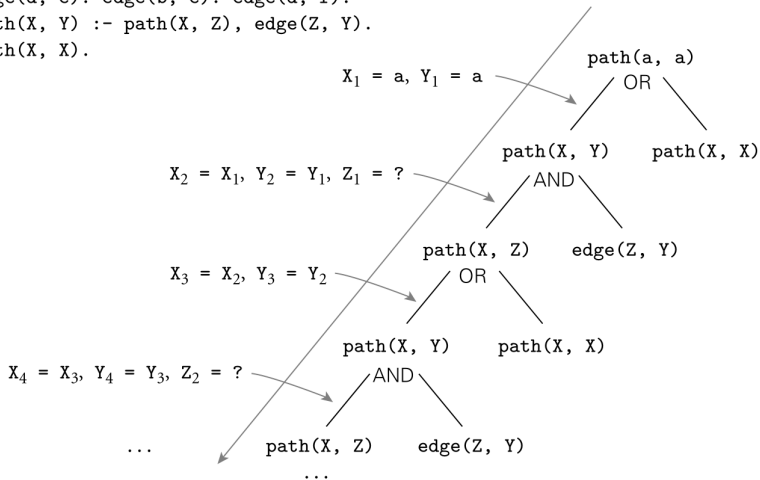
```
edge(a, b). edge(b, c). edge(c, d).  
edge(d, e). edge(b, e). edge(d, f).  
path(X, Y) :- path(X, Z), edge(Z, Y).  
path(X, X).  
?- path(a, a).  
ERROR (Memory Limit)
```



```

edge(a, b). edge(b, c). edge(c, d).
edge(d, e). edge(b, e). edge(d, f).
path(X, Y) :- path(X, Z), edge(Z, Y).
path(X, X).

```



Imperative feature 2: Cut

The **cut !** is a zero-argument predicate:

- As a subgoal, it is always satisfied.
- **Side effect:** It commits the interpreter to whatever choices have been made since unifying the parent goal with the left-hand side of the current rule, including the choice of that unification itself.

Example.

```
member(X, [X | _]).  
member(X, [_ | T]) :- member(X, T).  
prime_candidate(X) :- member(X, candidates), prime(X).  
?- prime_candidate(6).
```

Assumption:

- `prime(X)` is expensive. and `prime(6)` is false.
- `candidates` is the list `[6,5,3,6,2,6,6]`.

The search may execute `prime(6)` four times, every time with the same result.

To cut down search time:

```
member(X, [X | _]) :- !.                % Rule 1  
member(X, [_ | T]) :- member(X, T).     % Rule 2
```

Cut: When `member(X,L)` unifies with the head of Rule 1, it will no longer be unified with Rule 2 in further search.

Imperative feature 3: Arithmetic

Arithmetic symbols are predicate symbols, not arithmetic operators:

- $2+3$ is a two-argument predicate $+(2, 3)$, not 5

```
?- (2 + 3) = 5.  
false.
```

- `is` is a built-in predicate that unifies the first argument with the arithmetic value of the second argument, which must be instantiated.

```
?- is(5, 2+3).  
true.  
?- X is 2+3.  
X = 5.  
?- is(2+3, 5).  
false.  
?- 5 is Y+3.  
ERROR.
```

- Enumerate all natural numbers:

```
natural(1).  
natural(N) :- natural(M), N is M+1.  
my_loop(N) :- natural(I), write(I), nl, I=N, !.  
                % nl is 'new line' predicate
```

Extended Example: Tic-Tac-Toe

Game rule.

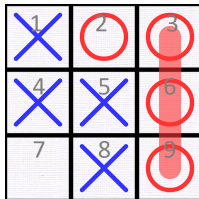
- Play on a 3×3 grid.
- Two players, X (agent) and O, take turns placing markers in empty squares.
- Any player wins if they place three markers in a (horizontal, vertical, or diagonal) line.

We will develop a Prolog program for Tic-Tac-Toe.

- **Squares:** 1, 2, 3, 4, 5, 6, 7, 8, 9
- **Agent perception:** $x(i)$, $o(i)$ indicating the markers in square i .

```
full(A) :- x(A).    full(A) :- o(A).  
empty(A) :- \+(full(A))    % the built-in \+ succeeds if  
                           % its argument cannot be proven.
```

- **Agent action:** $move(A)$. Define the next move of the Prolog agent.



Knowledge base:

- Winning conditions:

```
ordered_line(1, 2, 3). ordered_line(4, 5, 6).  
ordered_line(7, 8, 9). ordered_line(1, 4, 7).  
ordered_line(2, 5, 8). ordered_line(3, 6, 9).  
ordered_line(1, 5, 9). ordered_line(3, 5, 7).  
line(A, B, C) :- ordered_line(A, B, C).  
line(A, B, C) :- ordered_line(A, C, B).  
line(A, B, C) :- ordered_line(B, A, C).  
line(A, B, C) :- ordered_line(B, C, A).  
line(A, B, C) :- ordered_line(C, A, B).  
line(A, B, C) :- ordered_line(C, B, A).
```

- `move(A) :- good(A), empty(A).`

We will need to define “good” moves, and use **ordering** to indicate preference among good moves.

Good move 1. Winning.

```
win(A) :- x(B), x(C), line(A,B,C).
```

Good move 2. Prevent opponent from winning.

```
block_win(A) :- o(B), o(C), line(A,B,C).
```

Good move 3. Create a “split” situation

```
split(A) :- x(B), x(C), different(B,C),  
            line(A,B,D), line(A,C,E), empty(D), empty(E).  
same(A,A).  
different(A,B) :- \+(same(A,B)).    % negation as failure
```

O	1	2	3
4	X	6	O
7	8	9	X

Good move 4. Get two in a row in such a way that the opponent's blocking move won't build towards three in a row.

```
strong_build(A) :- x(B), line(A,B,C), empty(C), \+(risky(C)).  
risky(C) :- o(D), line(C,D,E), empty(E).
```

Good move 5. Get two in a row in such a way that the opponent's blocking move won't become a split.

```
weak_build(A) :- x(B), line(A,B,C), empty(C), \+(double_risky(C)).  
double_risky(C) :- o(D), o(E), different(D,E), line(C,D,F),  
    line(C,E,G), empty(F), empty(G).
```

Strategy for the Prolog agent.

```
good(A) :- win(A).      good(A) :- block_win(A).  
good(A) :- split(A).    good(A) :- strong_build(A).  
good(A) :- weak_build(A).
```

Last resort. If none of the above, pick an unoccupied square, giving priority to the centre, the corners, and the sides in order.

```
good(5).      % centre  
good(1).      good(3).      good(7).      good(9).      % corners  
good(2).      good(4).      good(6).      good(8).      % sides
```

Summary of The Topic

The following are the main knowledge points covered:

- **Declarative programming** is different from **imperative programming** in that a declarative program describes what output you want, rather than how to produce the output.
- **Logic programming**: A declarative programming paradigm that aim to decompose a programming language into a logic component (to represent knowledge), and a control component (to perform inference).
- **Prolog**: is an important logic programming language that aims to implement a first-order knowledge base.
- **Prolog inference**: Backward chaining starts from the goal and works backwards while traversing the subgoal tree using backtracking.
- Prolog does not fully correspond to FO logic:
 - Ordering
 - Arithmetics
 - Cut