# Block Ciphers: Modes of Operation

2019. 3. 19

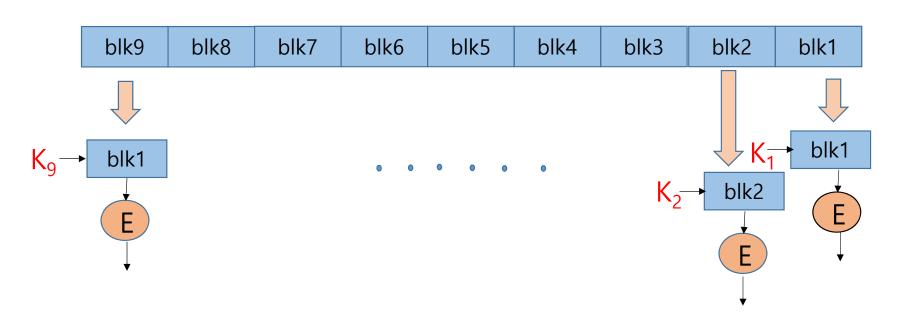
#### **Contents**

- Introduction to crypto
- Symmetric-key cryptography
  - Stream ciphers
  - Block ciphers
  - Block cypher operation modes
- Public-key cryptography
  - RSA
  - ECC
  - Digital signature
  - Public key Infrastructure

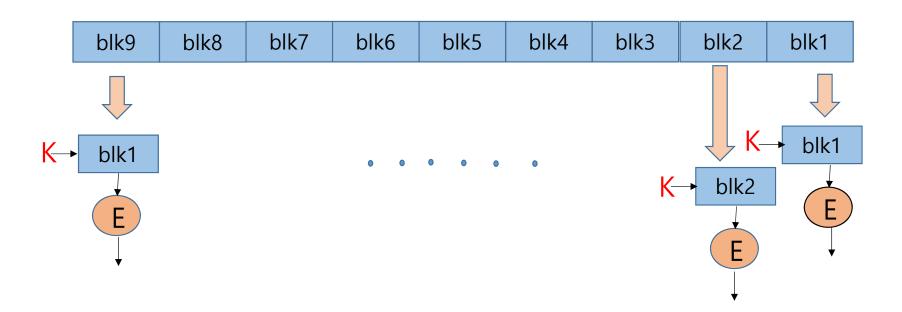
- Cryptographic hash function
  - Attack complexity
  - Hash Function algorithm
- Integrity and Authentication
  - Message authentication code
  - GCM
  - Digital signature
- Key establishment
  - server-based
  - Public-key based
  - Key agreement (Diffie-Hellman)

### Encryption of multiple blocks

- What if a file have multiple block?
  - If we use different keys for each block, encryption is like one time pad(OTP).



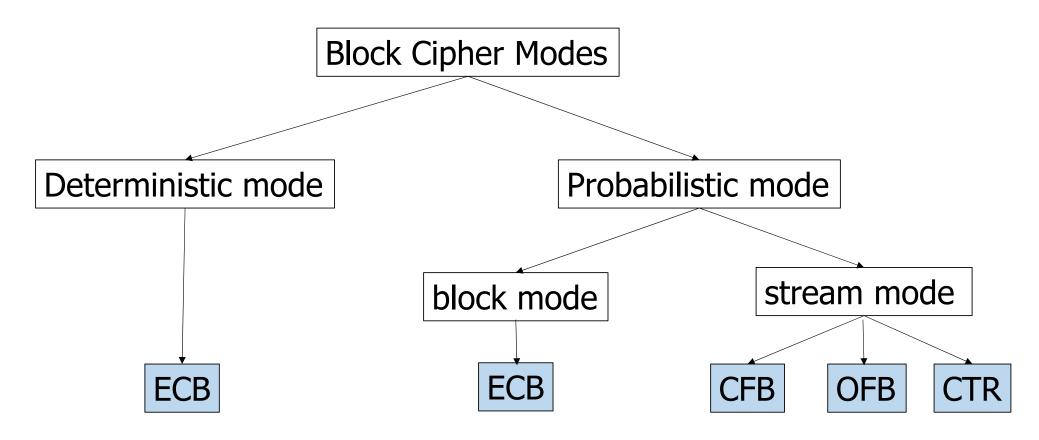
• What if we use the same key for all the blocks of the file? Are there any problems?



#### Modes of operation

- Block cipher modes of operation
  - ECB: Electronic code book
  - CBC: Cipher block chaining
  - CFB: Cipher feedback
  - OFB: Output feedback
  - CTR: Counter mode
  - and more

#### Classification of operation modes



#### **ECB**

- Mapping between blocks of plaintext and ciphertext is fixed as long as the key is same. (deterministic)
- It is like a traditional code book.

Key = Ki	
$P_0$	$C_0$
P <sub>1</sub>	C <sub>1</sub>
P <sub>2</sub>	C <sub>2</sub>
P <sub>3</sub>	$C_3$
P <sub>4</sub>	C <sub>4</sub>
•••	•••

### **Advantages of ECB**

- Block synchronization is unnecessary.
  - Receiver can decrypt the received blocks regardless of receiving other blocks.
- Bit errors affect only corresponding block, not succeeding blocks.
- Encryption/decryption processes can be parallelized.

#### ECB weakness

- Suppose  $P_i = P_i$
- Then  $C_i = C_j$  and an attacker knows  $P_i = P_j$
- This gives the attacker some information, even if he does not know P<sub>i</sub> or P<sub>j</sub>
- He might know P<sub>i</sub>
- Is this a serious issue?

#### Substitution attack

Consider the following plaintext.

"Abel loves Bob. Cain hates Tom."

Suppose the block size is 64-bits:

- Then, the cipher texts are  $C_0$ ,  $C_1$ ,  $C_2$ ,  $C_3$ .
- Attacker reordered the cipher text blocks: C<sub>0</sub>,C<sub>3</sub>,C<sub>2</sub>,C<sub>1</sub>
- Then the decrypted plaintext is:

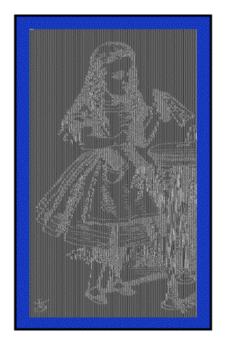
"Abel loves Tom. Cain Hates Bob. "

Still, attacker does not know contents about the ciphertext.

### An Example of ECB encryption

Alice's uncompressed image, and ECB encrypted (TEA)

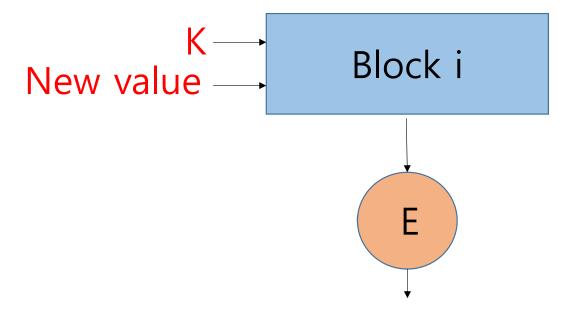




• Why does this happen?

(source: Information Security of M. Stamp)

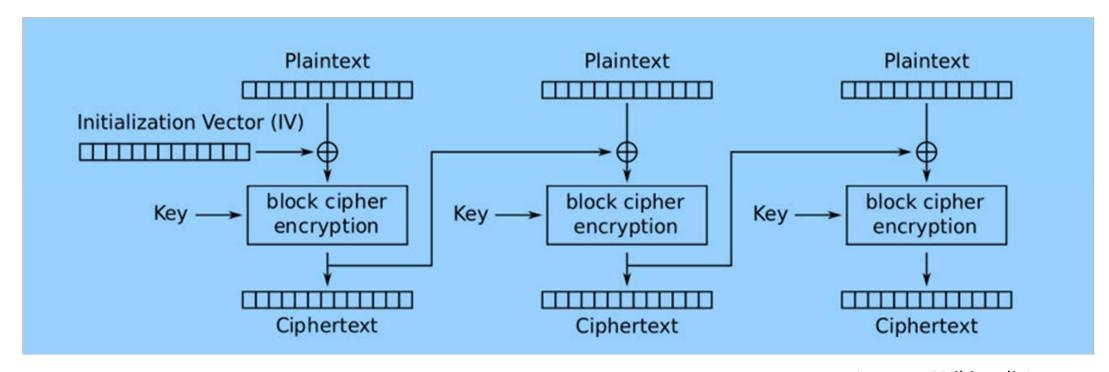
# Fix the problem



### CBC(Cipher block chaining) Encryption

$$C_0 = E_K(IV \oplus P_0), C_1 = E_K(C_0 \oplus P_1), C_2 = E_K(C_1 \oplus P_2),...$$

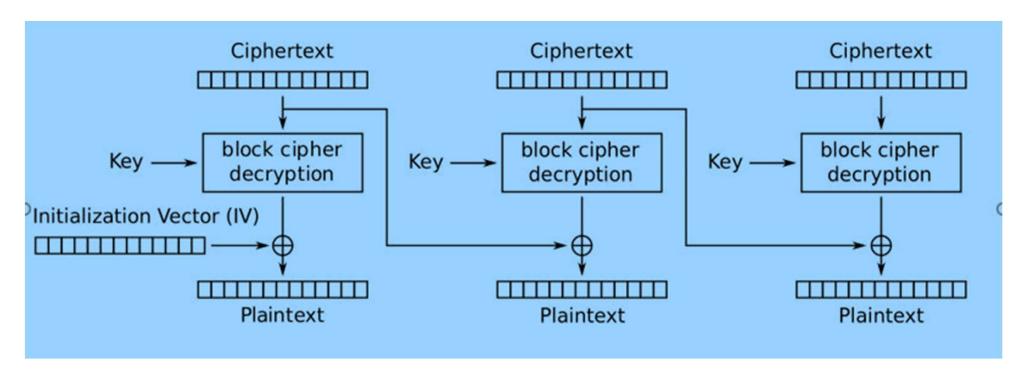
$$C_2 = E_K(C_1 \oplus P_2), \dots$$



(source: Wikipedia)

### **CBC** Decryption

$$P_0 = D_K(C_0) \oplus IV$$
,  $P_1 = D_K(C_1) \oplus C_0$ ,  $P_2 = D_K(C_2) \oplus C_1$ ,...



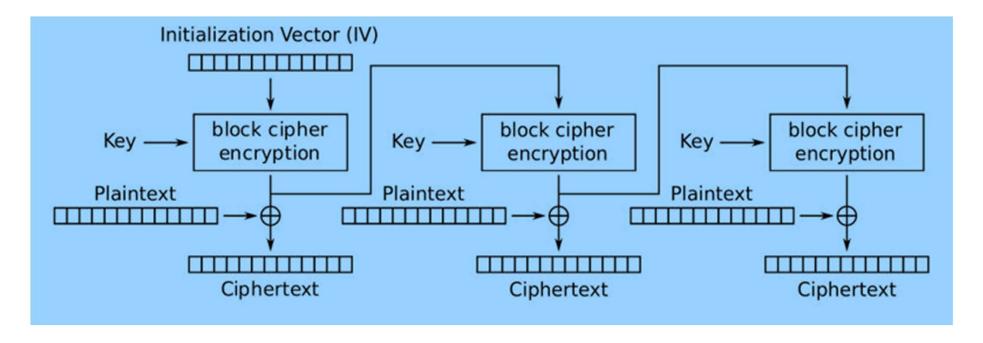
(source: Wikipedia)

#### **CBC**

- CBC encryption is probabilistic.
  - If we use new IV every time we encrypt, two ciphertexts of the same plaintext blocks are completely different.
- IV should be nonce. (should be used only once)
- But it should not be secret.(doesn't need to be)

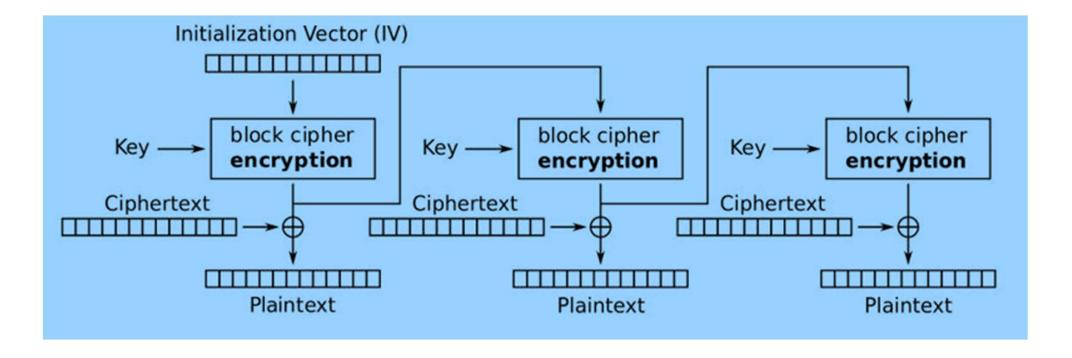
### **OFB Encryption**

 $I_0 = IV$ ,  $O_i = E_K(I_i)$ ,  $I_i = O_{i-1}$ ,  $C_i = P_i \oplus O_i$ ,  $P_i = C_i \oplus O_i$ ... It works like the stream cipher. The key stream is generated block by block, not bitwise.



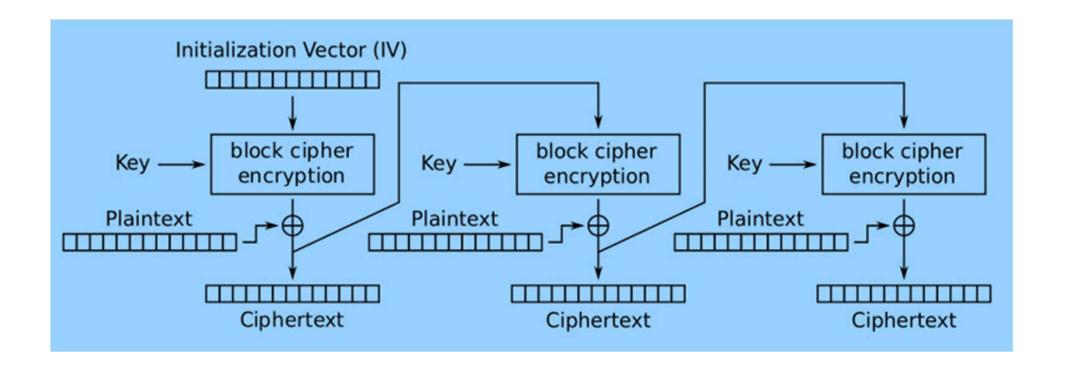
### **OFB Decryption**

Note that when decrypting, block cipher uses the encryption.

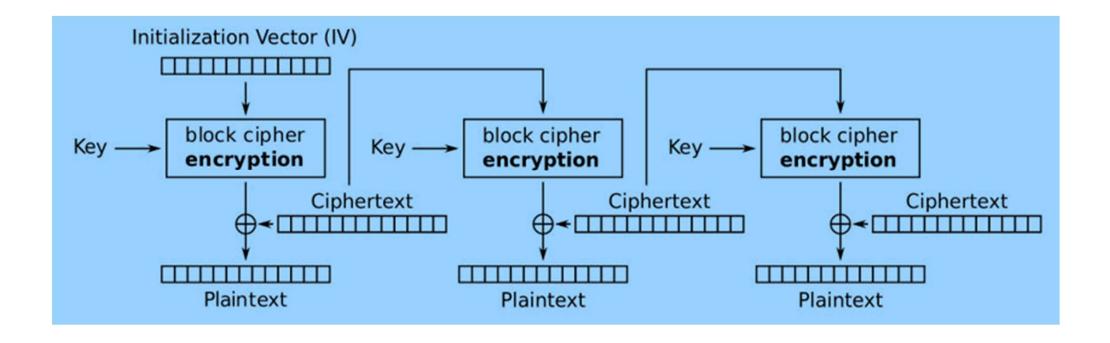


#### **CFB Encryption**

$$C_0 = E_K(IV) \oplus P_0, C_i = E_K(C_{i-1}) \oplus P_i, P_i = E_K(C_{i-1}) \oplus C_i,$$

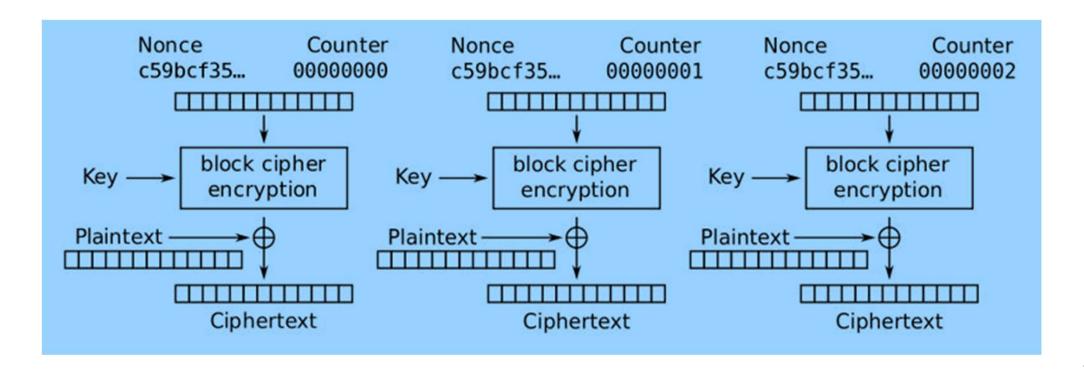


# **CFB Decryption**

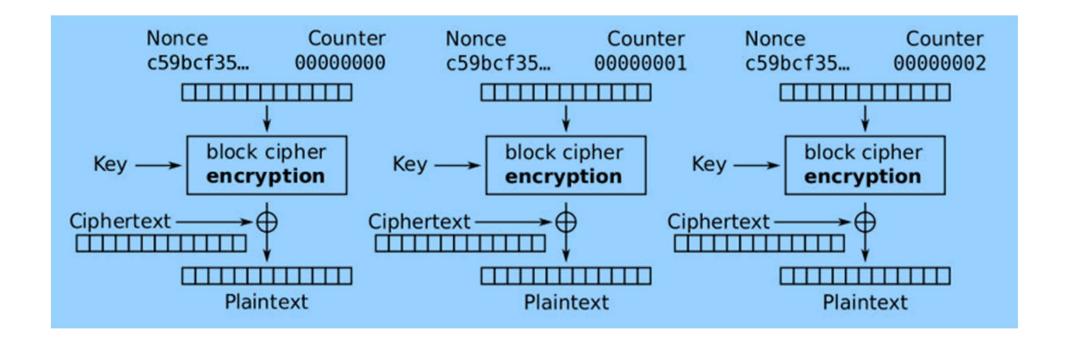


### **CTR Encryption**

 $C_i = E_K(IV \parallel CTR_i) \oplus P_i$ ,  $P_i = E_K(IV \parallel CTR_i) \oplus C_i$ ,...



# **CTR** Decryption



### Advantage of CTR

• The encryption/decryption of all blocks can be processed in parallel.

#### Question

- Why are there so many modes operations? Which one can be recommended for your use?
- IV should be nonce. How can we generate IVs every time new message blocks are sent?
- In doing this block operation, at the same time can we do the integrity and authentication check of the message?
  - i.e., can we verify that (1) the message is really created by Alice, and (2) the ciphertext was not tampered during transmission?