5. Shift Reduce Parser 2

Monday, March 1, 2021 8:28 AM Arguented Grammar I, loswe (E'→E) $T \rightarrow F$ $F \rightarrow id \mid (E)$ E'→E £ → · E +T E - E-T E カ·ナ E +T T -> TKF I, closure (Io, E) E'→E. I3 closure (I, F) E→E.+T T→F. I2 closure (Io,T) I4 closure (I., id) E - T. F -> id. T → T. * F Is closure (I., (') F -> (· E) E → · E+T E → .T T → ·T*F F -> .(E) goo $(I_0,F) \rightarrow I_3$ T→ F. goto (I., (') > T4 of till you cannot move E → E.+ T to a new item set F → (. E) $goto(I_o, T) \rightarrow I_2$ E → · E + T $E \rightarrow T$. $T \rightarrow T \cdot *F$ [- → · T T → T*F F -> -(E) F-id $F \rightarrow id$. F > F+.T F -> (E)

$$F \rightarrow id$$

$$F \rightarrow id$$

$$goto(T_1, +) \rightarrow T_6$$

$$E \rightarrow E + . T$$

$$T \rightarrow . T + F$$

$$T \rightarrow . F$$

$$F \rightarrow . (E)$$

$$F \rightarrow . id$$

$$goto(T_4, T)$$

$$F \rightarrow T.$$

$$F \rightarrow$$

$$F \to (E.)$$

$$E \to E. + T$$

$$Joho (I_{4}, id)$$

$$F \Rightarrow id. \quad \Im I_{5}$$

$$Goho (I_{5}, T)$$

$$E \to E + T.$$

$$T \to T. * F$$

$$Goho (I_{6}, F)$$

$$T \to F. \quad \Im I_{5}$$

$$Goho (I_{6}, id)$$

$$F \Rightarrow id. \quad \Im I_{5}$$

$$Goho (I_{6}, id)$$

$$T = T * F.$$

$$Goho (I_{8}, T)$$

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$$Goho (I_{8}, T)$$

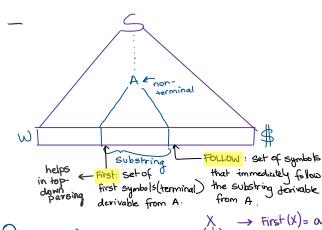
$$Goho (I_{9}, *)$$

$$F \to (E).$$

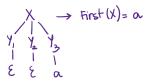
*DFA of goto function -> DFA

Parse Table

(1) FIRST & FOLLOWS



First (x) = $\{x\}$ $x \rightarrow y_1 y_2 \dots y_n$



Example:

Q 5→Az First (A) 4€3 A→ € First (S) = { 2} U) Hirst (x) = 7×5

X → Y, Y2 Yn

E E a

D First (X) → First (Y1) except E (empty string) U

First (Y2) except E U

First (Yn) except E

3 X → E; First (x) = {E}