

4. Job Scheduling + MST

Tuesday, February 23, 2021 8:14 AM

Job	Profit	Slot
A1	100	5
A2	90	4
A3	80	3
A4	70	2
A5	60	1
A6	50	5
A7	40	4

- Jobs = 0
- Allocate j to slot 1 for all pending elements until the slot is filled

A1 --> slot 0/1

A2:

X points to the largest filled slot --> slot 0/1

If $S[X] > 4(A2)$, then decrement x

Keep a count of the amount of time x is decremented

Job	Profit	Slot
A1	100	3
A2	90	5
A3	80	4
A4	70	2
A5	60	2
A6	50	1
A7	40	2

Complexity: $O(n)$

With dynamic programming: $O(n^2)$

Minimum Spanning Tree with Prim's algorithm (Greedy algorithm)

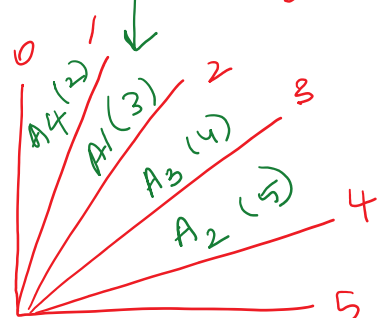
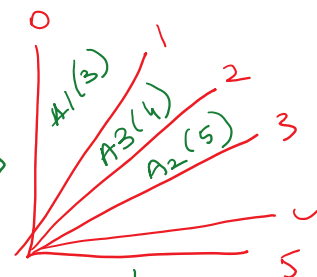
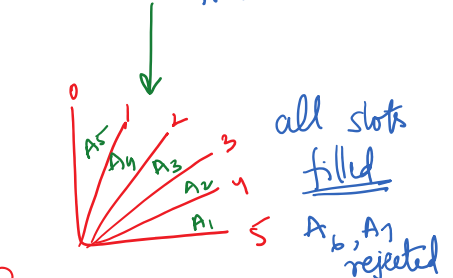
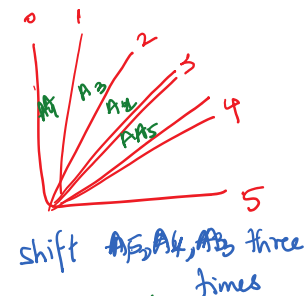
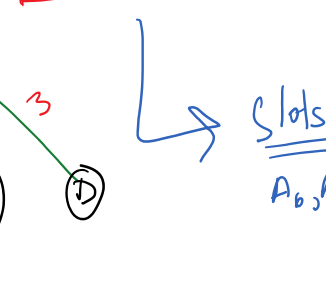
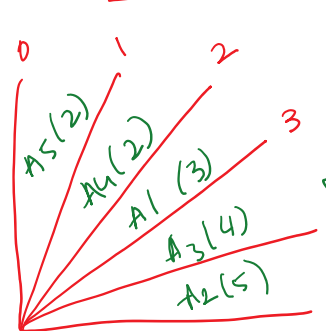
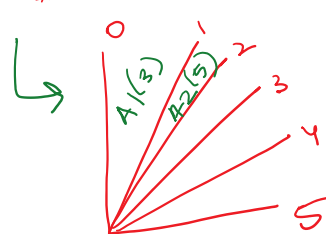
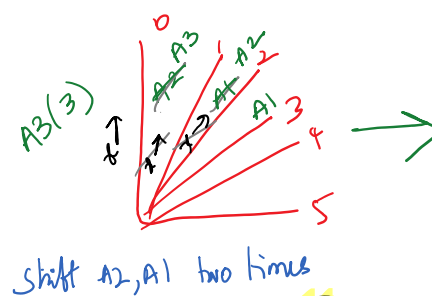
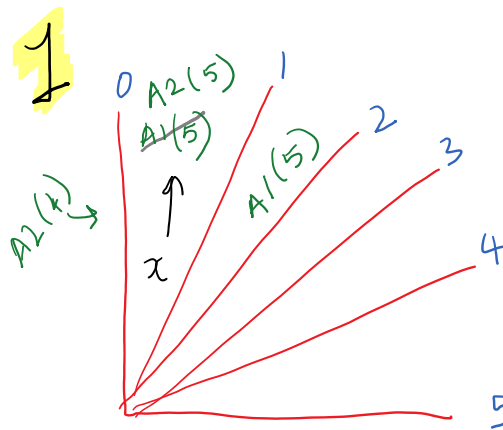
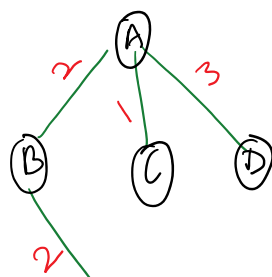
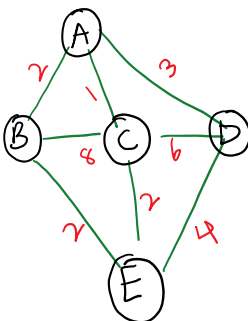
Given a graph $G = \{v, e\}$ (undirected)

Objective: Find a tree that connects all the vertices in a graph without forming a cycle with a minimal cost of edges

n nodes --> n-1 edges

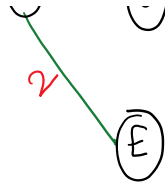
Adjacency matrix: $O(n^2)$

Adjacency list: $O(n)$



Slots filled
A6, A7 and subsequent jobs are rejected

(E)



Every time you insert an element into the graph, insert the element into a minheap to decrease complexity to $O(n \log n)$

i	Near(i)	Distance
1	1 \rightarrow 0	
2	1	28
3	6	inf
4	6	inf
5	6	25
6	6 \rightarrow 0	
7	6	inf

If $i == \text{near}(i)$, $\text{cost}(i) = 0$

Check 2,3,4,5,7 nearer to 1 and 6?

If $(\text{dist}(i,1) < \text{dist}(i,6))$

Add 1

Else

Add 6

If both are infinity, 6 gets added

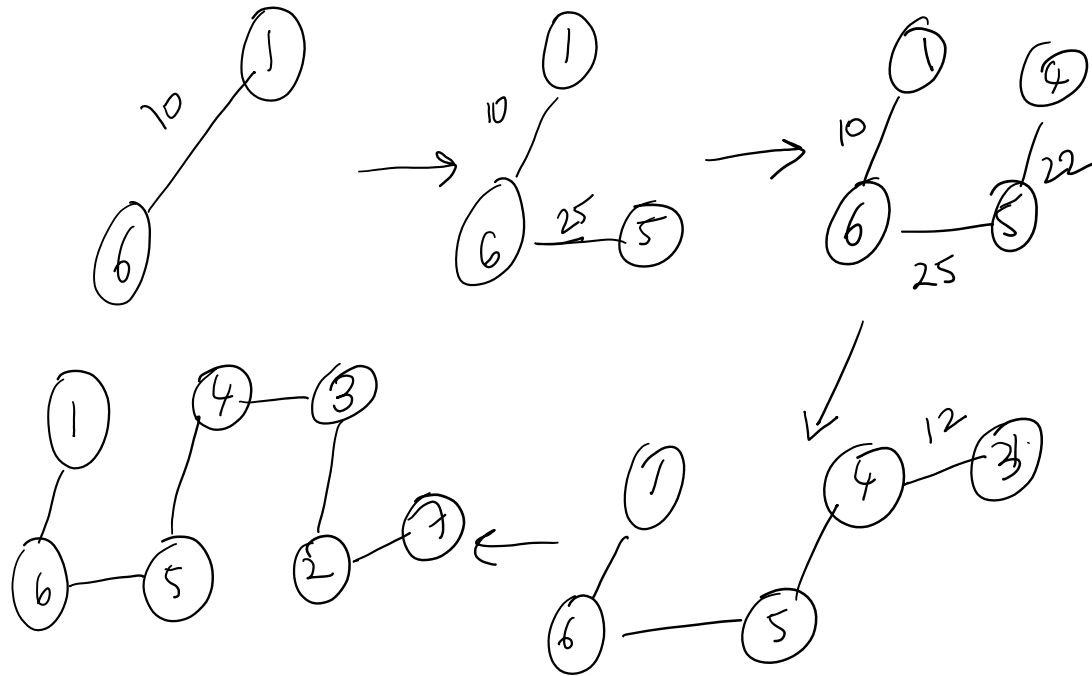
1	0	
2	1	28
3	5	inf
4	5	22
5	0	
6	0	
7	5	24

Check 2,3,4,7 nearer to 1,5,6?

Since 2 is nearer to 1 in the last step, only compare with 1 and 5

1	0	
2	1	28
3	4	12
4	0	
5	0	
6	0	
7	4	18

1	0	
2	3	16
3	0	
4	0	
5	0	
6	0	
7	5	18



1	0	
2	3	16
3	0	
4	0	
5	0	
6	0	
7	2	14