

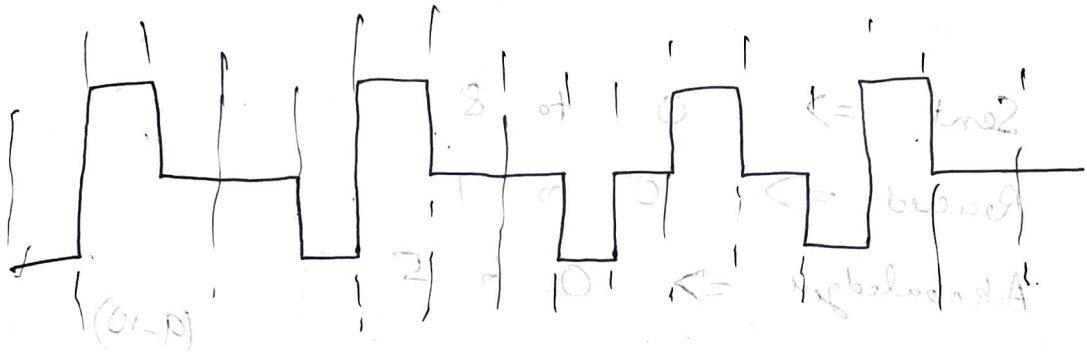
(9.)

(a.)

$$0100 = \text{mod}$$

$$1011 = \text{transitions}$$

$$1011 \oplus 0101 = 1110 = 14 \text{ mod } 16$$



$$\text{Total transitions} = 12$$



(b)

$$0111$$

$$+ 1111$$

$$\hline 10110$$

$$0110$$

$$+ 0001$$

$$\hline 0111$$

$$010111$$

$$+ 1010$$

$$\hline 10001$$

$$0001$$

$$+ 0001$$

$$\hline 0010$$

$$\text{Sum} = 0010$$

$$1's \text{ complement} = 1101$$

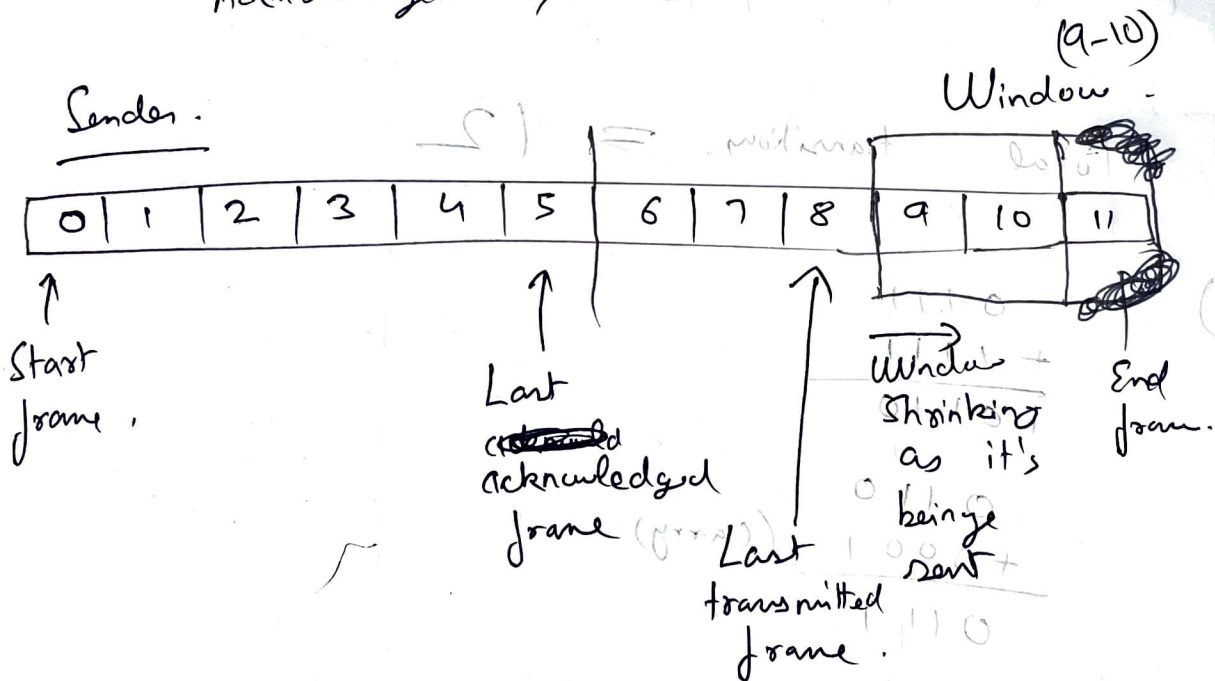
$$\text{Final bit stream} = 011111110101101$$

(C)

Sent  $\Rightarrow$  0 to 8

Received  $\Rightarrow$  0 to 7

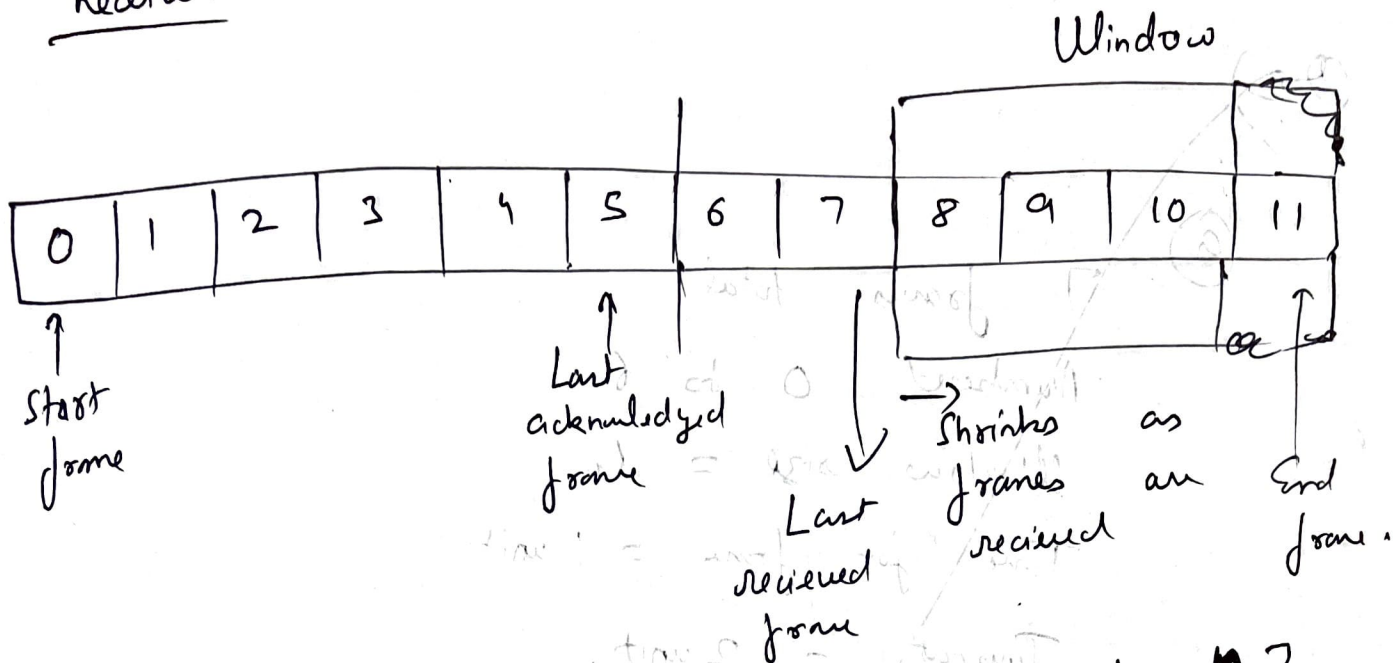
Acknowledged  $\Rightarrow$  0 to 5



$\Rightarrow$  Window size is at max 5 so all the remaining frames (9-10) are in the window. 3 are still unacknowledged.

$$5-2 = 3 \rightarrow 9 \text{ and } 10$$

# Receiver.

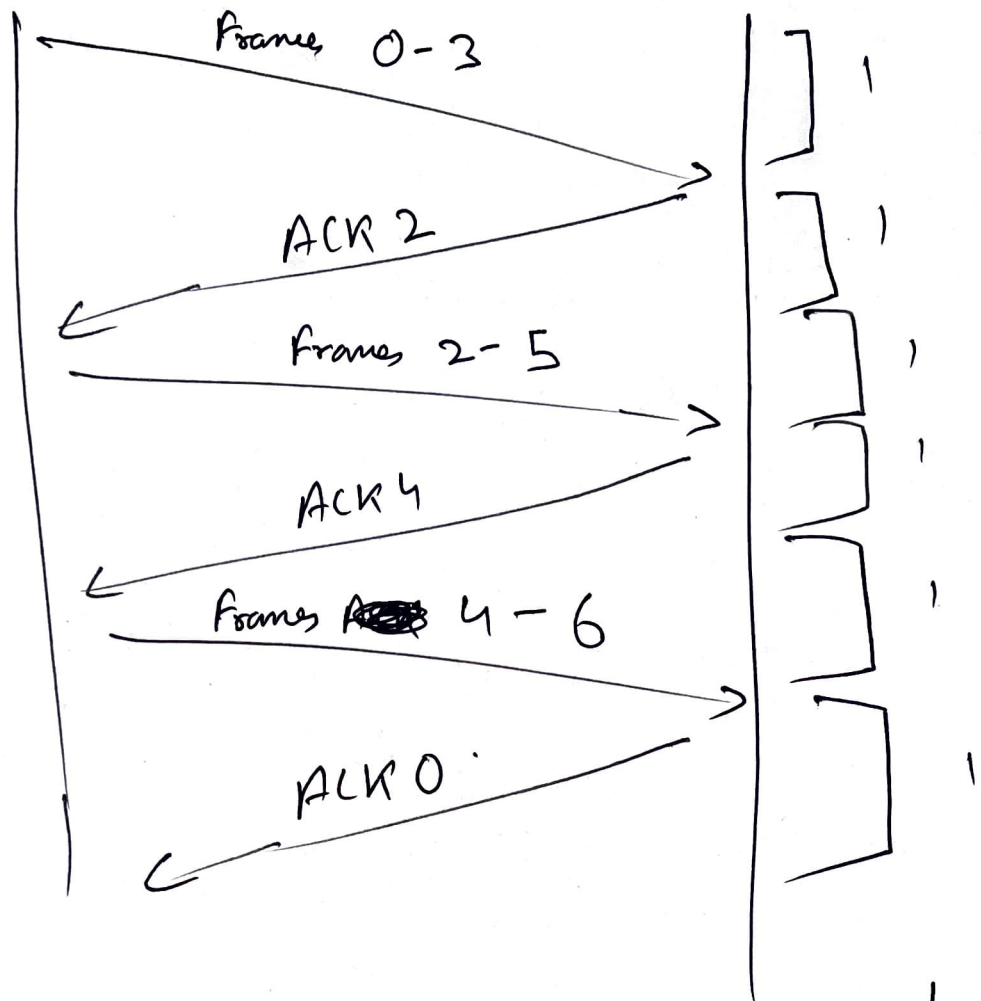


$\Rightarrow$  Max window size is 5 and 3 frames are remaining in the window. 2 are not ack. total number of frames in window (8-10)

$5 - 2 = 3$  frames in window (8-10)

Q2)

⑨



⑩

Selection repeat ARO is used.

if

frame loss probability is low

Go back N will have to send everything again which is inefficient

hence selection repeat ARO is better.

93 (a)

$T \rightarrow R$

100 mbps.

$V = 2 \times 10^8$  m/s

dist = 500 m.

$$\text{Propagation delay} = \frac{500}{2 \times 10^8} \times 10^6 \text{ lls} \\ = 2.5 \text{ lls.}$$

$$L = \frac{2DR}{V} \\ = \frac{2 \times 500 \times 100 \times 10^6}{2 \times 10^8} \\ = 500 \text{ bits.}$$

$$\text{Transmission Time} = \frac{500}{100 \times 10^6} \times 10^6 = 5 \text{ lls.}$$

7 lls have passed and it is

transmitting, hence

actual packet length > min packet length.

Else T would have transmitted in ~~5 lls~~ 5 lls itself.

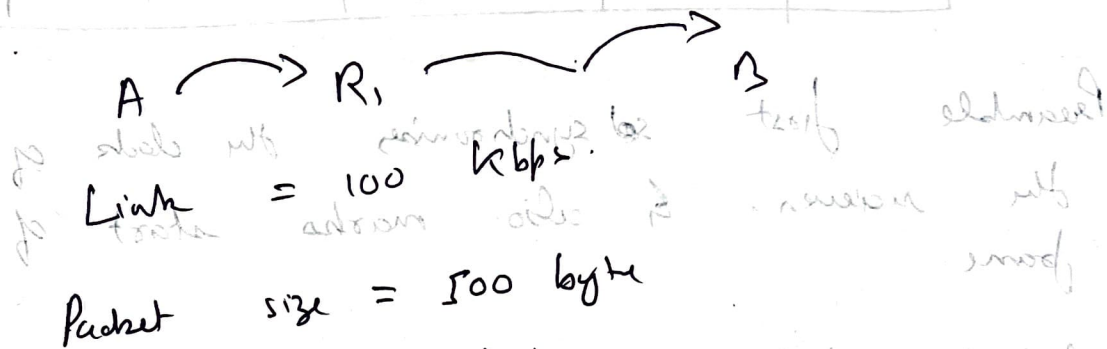
Yes, T may be able to transmit.  
and even if a collision happens,  
it will be detected. and then  
T will be able to resend after  
the collision.



Q4)

(a)

Total data	=	5000 bytes	
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Header = 20 byte

Time for packet to hop = 100 ms

No. of packets =  $\frac{5000}{500} = 10$

100 ms = Transmission delay + propagation delay

For all packets

A to R<sub>1</sub>

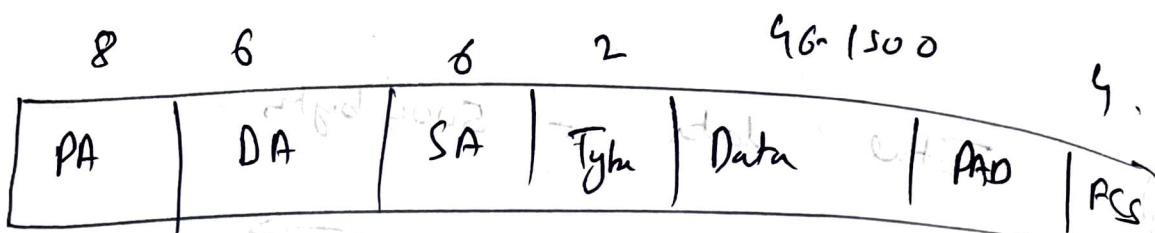
time =  $10 \times 100 \text{ ms} = 1 \text{ sec}$

R<sub>1</sub> to B also

time = 1 sec

Total = 2 sec

(b)



Preamble first synchronises the clock of the receiver, & also marks start of frame.

Destination address is matched, if not then discarded. Else proceed.

FCS checked for data integrity.

Ethernet specific header is stripped off and the packet is passed to network layer.

Q3

6



A sends RTS to B

B sends CTS

However due to collision with  
some packet. possible CTS that  
A might reach A but does not  
reach C. So A sends the data to  
B. C will also send the data.  
as it did not get CTS.  
Hence collision will happen at B.