

DS 288: Numerical Methods

AUG. 24- 2021

RECAP:

ERROR GROWTH: EXPONENTIAL
& LINEAR ↴

EXAMPLES OF RECURSIVE EQUATIONS
(OR DIFFERENCE EQUATIONS)

↳ [1.3]

Ex:-
① $P_n = \frac{1}{3} P_{n-1}$ ② $P_n = \frac{10}{3} P_{n-1} - P_{n-2}$

CAN WE EXPLAIN THE ERROR
GROWTH BASED ON THESE
RELATIONS

(ASSUME CONSTANTS ARE EXACT)

$$\textcircled{a} \quad P_n = \frac{1}{3} P_{n-1} \quad (\text{exact})$$

$$= \widehat{P}_n = \frac{1}{3} \widehat{P}_{n-1} \quad (\text{COMPUTER})$$

$$\underbrace{P_n - \widehat{P}_n}_{\epsilon_n} = \frac{1}{3} (P_{n-1} - \widehat{P}_{n-1})$$

$\epsilon_n \rightarrow \text{ERROR AT STEP } n$

$$\epsilon_n = \frac{1}{3} \epsilon_{n-1}$$

$$\epsilon_1 = \frac{1}{3} \epsilon_0$$

$$\epsilon_2 = \frac{1}{3} \epsilon_1 = \left(\frac{1}{3}\right) \left(\frac{1}{3} \epsilon_0\right) = \left(\frac{1}{3}\right)^2 \epsilon_0$$

$$\epsilon_3 = \frac{1}{3} \epsilon_2 = \frac{1}{3} \left(\frac{1}{3}\right)^2 \epsilon_0 = \left(\frac{1}{3}\right)^3 \epsilon_0$$

⋮

$$\epsilon_n = \left(\frac{1}{3}\right)^n \epsilon_0$$

$$\text{AS } n \rightarrow \infty \quad \cancel{\left(\frac{1}{3}\right)^n \rightarrow 0}$$

ERRORS INTRODUCED WILL DECAY

NOW RELAX CONDITION THAT
CONSTANTS ARE EXACT

$$\frac{1}{3} = \frac{1}{3} - \delta \quad \begin{matrix} \delta \rightarrow \text{ROUND OFF} \\ \text{ERROR} \end{matrix}$$

In REPRESENTING
0.3333...

$$P_n = \frac{1}{3} P_{n-1} \text{ (EXACT)} \\ \hat{P}_n = (\frac{1}{3} - \delta) \hat{P}_{n-1} \text{ (COMPUTER)}$$

$$\underline{P_n - \hat{P}_n = \frac{1}{3} (P_{n-1} - \hat{P}_{n-1}) + \delta \hat{P}_{n-1}}$$

$$\epsilon_n = \frac{1}{3} \epsilon_{n-1} + \delta \hat{P}_{n-1}$$

$$\epsilon_1 = \frac{1}{3} \epsilon_0 + \delta \hat{P}_0$$

$$\epsilon_2 = \left(\frac{1}{3}\right)^2 \epsilon_0 + \frac{1}{3} \delta \hat{P}_0 + \delta \hat{P}_1$$

$$\epsilon_3 = \left(\frac{1}{3}\right)^3 \epsilon_0 + \left(\frac{1}{3}\right)^2 \delta \hat{P}_0 + \frac{1}{3} \delta \hat{P}_1 + \delta \hat{P}_2 \\ \vdots \quad \vdots \quad \vdots \quad \ddots \quad \ddots$$

$$E_n = \underbrace{\left(\frac{1}{3}\right)^n \epsilon_0}_{\text{n TERMS EACH ONE}} + \underbrace{\left(\frac{1}{3}\right)^{n-1} \delta p_0 + \dots + \delta \hat{p}_{n-1}}$$

n TERMS EACH ONE
BEING SMALLER THAN $\delta \hat{p}$

$$E_n \leq \left(\frac{1}{3}\right)^n \epsilon_0 + n \delta \hat{p}_0$$

$$\lim_{n \rightarrow \infty} E_n \leq \cancel{\left(\frac{1}{3}\right)^n \epsilon_0} + \underbrace{n \delta \hat{p}_0}_{\Downarrow}$$

ERRORS INTRODUCED
BOUNDED BY LINEAR GROWTH

(b) ASSUME CONSTANTS ARE EXACT.

$$p_n = \frac{10}{3} p_{n-1} - p_{n-2} \rightarrow \text{EXACT}$$

$$\hat{p}_n = \frac{10}{3} \hat{p}_{n-1} - \hat{p}_{n-2} \rightarrow \text{COMPUTER}$$

$$\underline{E_n = \frac{10}{3} E_{n-1} - E_{n-2} \rightarrow \text{DEFECT}}$$

$E_n \propto \epsilon_0, \epsilon_1$ DEFECT-EQUATION

HAND OUT ON
DIFFERENCE EQUATIONS

Solutions to Difference Equations: Review*

First let us look at the second-order homogeneous ordinary differential equation:

$$\underbrace{\frac{d^2\theta}{dt^2} + 2\frac{d\theta}{dt} + 2\theta = 0}_{\rightarrow \theta(t)} \quad (1)$$

with I.C.s

$$\begin{aligned}\theta(0) &= 1 \\ \frac{d\theta}{dt}(0) &= 0\end{aligned}$$

We seek a solution of the form $\theta = Ce^{\lambda t}$. Plugging into equation (1)

$$\begin{aligned}\underbrace{C \left[\lambda^2 e^{\lambda t} + 2\lambda e^{\lambda t} + 2e^{\lambda t} \right]}_{\theta(t)} &= 0 \\ Ce^{\lambda t} \left[\lambda^2 + 2\lambda + 2 \right] &= 0 \\ \theta \left[\lambda^2 + 2\lambda + 2 \right] &= 0\end{aligned} \quad (2)$$

Equation (2) requires that either $\theta(t) = 0$ for all t , which is not very interesting as a solution, or that $\lambda^2 + 2\lambda + 2 = 0$. This quadratic equation is known as the *auxiliary* or *characteristic equation* and has two solutions or roots, $\lambda_1 = -1 + j$ and $\lambda_2 = -1 - j$. So the general solution to equation (1) is

$$\underbrace{\theta(t) = C_1 e^{\lambda_1 t} + C_2 e^{\lambda_2 t}}_{\text{C}} \quad (3)$$

Now we determine C_1 and C_2 from the I.C.s

$$\begin{aligned}1 &= C_1 + C_2 \\ 0 &= \lambda_1 C_1 + \lambda_2 C_2\end{aligned} \quad \begin{aligned}\text{solve} \\ \implies\end{aligned} \quad \begin{aligned}C_1 &= \frac{\lambda_2}{\lambda_2 - \lambda_1} \\ C_2 &= \frac{\lambda_1}{\lambda_1 - \lambda_2}\end{aligned}$$

The solution to equation (1) is, therefore,

$$\underbrace{\theta(t) = \frac{\lambda_2}{\lambda_2 - \lambda_1} e^{\lambda_1 t} + \frac{\lambda_1}{\lambda_1 - \lambda_2} e^{\lambda_2 t}}_{\text{C}} \quad (4)$$

Now let us look at a difference equation

$$\underbrace{y_{n+2} - \frac{5}{6} y_{n+1} + \frac{1}{6} y_n = 0}_{\rightarrow y_n} \quad (5)$$

with I.C.s

$$\begin{aligned}y_0 &= 0 \\ y_1 &= 1\end{aligned}$$

$$y_n = \frac{5}{6} y_{n-1} - \frac{1}{6} y_{n-2}$$

* Credit: Prof. Simon Shepard, Dartmouth College

Analogously, we seek a solution of the form $y_n = C\lambda^n$. Plugging into equation (5)

$$\begin{aligned} C \lambda^{n+2} - \frac{5}{6}C \lambda^{n+1} + \frac{1}{6}C \lambda^n &= 0 \\ C \lambda^n \left[\lambda^2 - \frac{5}{6}\lambda + \frac{1}{6} \right] &= 0 \\ y_n \left[\lambda^2 - \frac{5}{6}\lambda + \frac{1}{6} \right] &= 0 \end{aligned}$$

So either $y_n = 0$ or, more interestingly, $\lambda^2 - \frac{5}{6}\lambda + \frac{1}{6} = 0$, whose roots are $\lambda_1 = \frac{1}{2}$; $\lambda_2 = \frac{1}{3}$. So the general solution to equation (5) is

$$y_n = C_1 \left(\frac{1}{2} \right)^n + C_2 \left(\frac{1}{3} \right)^n \quad (6)$$

Now we determine C_1 and C_2 from the I.C.s

$$\begin{array}{lcl} 0 & = & C_1 + C_2 \\ 1 & = & C_1 \left(\frac{1}{2} \right) + C_2 \left(\frac{1}{3} \right) \end{array} \quad \begin{array}{c} \text{solve} \\ \Rightarrow \end{array} \quad \begin{array}{l} C_1 = 6 \\ C_2 = -6 \end{array}$$

The solution to the difference equation (5) is, therefore,

$$y_n = 6 \left(\frac{1}{2} \right)^n - 6 \left(\frac{1}{3} \right)^n \quad (7)$$

Compare equation (7) with the difference equation (5)

$$y_n = 6 \left(\frac{1}{2} \right)^n - 6 \left(\frac{1}{3} \right)^n \quad y_{n+2} - \frac{5}{6} y_{n+1} + \frac{1}{6} y_n = 0$$

$$\begin{array}{ll} y_0 = 0 & y_0 = 0 \\ y_1 = 1 & y_1 = 1 \\ y_2 = \frac{5}{6} & y_2 = \frac{5}{6} \\ y_3 = \frac{19}{36} & y_3 = \frac{19}{36} \\ y_4 = \frac{65}{216} & y_4 = \frac{65}{216} \\ \vdots & \vdots \end{array}$$

* same *

For our purposes, we do not need to get solutions to specific initial conditions, only the general behavior of the difference equation given by equation (6).

SOLVING DIFFERENCE EQUATION

$$E_n \propto \lambda^n \quad [E_n = C\lambda^n]$$

$$E_n = \frac{10}{3} E_{n-1} - E_{n-2}$$

$$\lambda^n = \frac{10}{3} \lambda^{n-1} - \lambda^{n-2}$$

DIVIDE λ^{n-2}

$$\lambda^2 - \frac{10}{3}\lambda + 1 = 0 \rightarrow \text{CHARACTERISTIC EQUATION}$$

$$\lambda = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

$$\Rightarrow \lambda_1 = \frac{1}{3} \text{ & } \lambda_2 = 3$$

$$E_n = C_1 \left(\frac{1}{3}\right)^n + C_2 (3)^n$$

$$\lim_{n \rightarrow \infty} \epsilon_n = c_1 \left(\frac{1}{3}\right)^{n^0} + c_2 \left(\frac{2}{3}\right)^{n^0}$$

ONCE ERRORS INTRODUCE THEY
GROW EXPONENTIALLY

WHY NOT $\frac{10}{3} = \frac{10}{3} - \delta$

BECAUSE ASSUMING $\frac{10}{3}$ BEING

EXACT, WE ALREADY GOT

EXPONENTIAL ERROR

BEHAVIOR //

FURTHER ON ERRORS:

$p_n \rightarrow$ SOLUTION AT n^{th} STEP

$p \rightarrow$ EXACT SOLUTION.

IF TRUTH $p_n \neq p$
WE DO NOT KNOW p //

$$\text{RELATIVE ERROR}_n = \frac{|P_n - P|}{|P|}$$

WHAT HAPPENS WHEN YOU
DO NOT KNOW 'P'

$$E_n = \frac{|P_n - P_{n-1}|}{|P_n|} \quad \begin{matrix} \text{BEST GUESS} \\ \text{FOR } P \approx P_n \end{matrix}$$

IF 'P' IS MULTIDIMENSIONAL

$$||\bar{P}_n - \bar{P}_{n-1}||_m$$

$$E_n = \frac{||\bar{P}_n||_m}{||\bar{P}_{n-1}||_m}$$

$|| \cdot ||_m \rightarrow$ IS THE L-m NORM

$$||\bar{x}||_2 = \left(\sum x_i^2 \right)^{1/2}$$

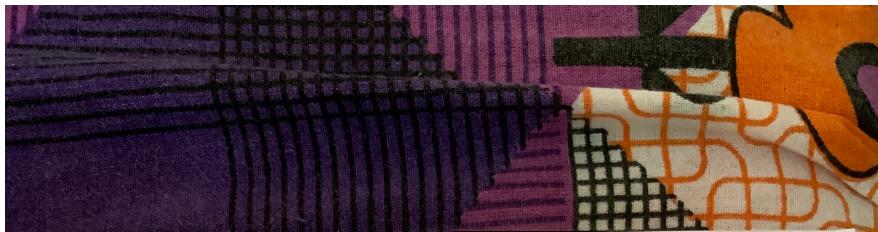
MAGNITUDE

$$||\bar{x}||_\infty = \max(\bar{x})$$

$|P_n - P| \rightarrow$ ABSOLUTE ERROR

$$|P_n - P|$$

WHEN p IS UNKNOWN
 $|P_n - P_{n+1}| \rightarrow$ ABSOLUTE ERROR.



1.2 Round-off Errors and Computer Arithmetic 21

Example 2 Determine the absolute and relative errors when approximating p by p^* when

- (a) $p = 0.3000 \times 10^1$ and $p^* = 0.3100 \times 10^1$; [REDACTED]
- (b) $p = 0.3000 \times 10^{-3}$ and $p^* = 0.3100 \times 10^{-3}$; [REDACTED]
- (c) $p = 0.3000 \times 10^4$ and $p^* = 0.3100 \times 10^4$. [REDACTED]

Solution

- (a) For $p = 0.3000 \times 10^1$ and $p^* = 0.3100 \times 10^1$ the absolute error is 0.1, and the relative error is $0.333\bar{3} \times 10^{-1}$.
- (b) For $p = 0.3000 \times 10^{-3}$ and $p^* = 0.3100 \times 10^{-3}$ the absolute error is 0.1×10^{-4} , and the relative error is $0.333\bar{3} \times 10^{-1}$.
- (c) For $p = 0.3000 \times 10^4$ and $p^* = 0.3100 \times 10^4$, the absolute error is 0.1×10^3 , and the relative error is again $0.333\bar{3} \times 10^{-1}$.

This example shows that the same relative error, $0.333\bar{3} \times 10^{-1}$, occurs for widely varying absolute errors. As a measure of accuracy, the absolute error can be misleading and the relative error more meaningful, because the relative error takes into consideration the size of the value. ■

The following definition uses relative error to give a measure of significant digits of accuracy for an approximation.

Definition 1.16 The number p^* is said to approximate p to **t significant digits** (or figures) if t is the largest nonnegative integer for which

$$\frac{|p - p^*|}{|p|} \leq 5 \times 10^{-t}.$$

Table 1.1 illustrates the continuous nature of significant digits by listing, for the various values of p , the least upper bound of $|p - p^*|$, denoted $\max |p - p^*|$, when p^* agrees with p to four significant digits.

Table 1.1	p	0.1	0.5	100	1000	5000	9990	10000
max $ p - p^* $	0.00005	0.00025	0.05	0.5	2.5	4.995	5.	

RELATIVE ERRORS ARE
MORE MEANINGFUL.

CASES WHERE WE HAVE TO
USE ABSOLUTE ERROR

EX:- ROOT FINDING
FIND P $\exists f(p) = 0$

ERROR GROWTH.
 $E_n = C(K)^n E_0$ EXPONENTIAL
 $|K| > 1$

SO THE CRITICAL PART.

ANALYZE MAGNITUDE OF
ROOTS OF CHARACTERISTIC
EQUATION

ROOTS COULD BE
COMPLEX.

— END OF CHAPTER 1 —

CHAPTER-2 :

SOLNS TO

EQUATIONS IN ONE VARIABLE

"ROOT FINDING" PROBLEM

FIND $P \ni f(P) = 0$

P IS THE ROOT OR ZERO OF $f(x)$

$$\text{Ex: } x^2 e^x + e^{-2x} = 4$$

$$\Rightarrow f(x) = x^2 e^x + e^{-2x} - 4$$



ITERATIVE METHODS

- SUCCESSIVELY IMPROVING ESTIMATION OF 'P'

(0) SCAN x UNTIL $f(x) = 0$

$$x_0 \rightarrow x_1 = x_0 + s \quad x \in [a, b]$$

STEP SIZE

$$f(x_0) \neq 0$$

$$f(x_1) \neq 0 \rightarrow x_2 = x_1 + s$$

WHAT STEP SIZE?

s TOO SMALL \rightarrow GO STEPS

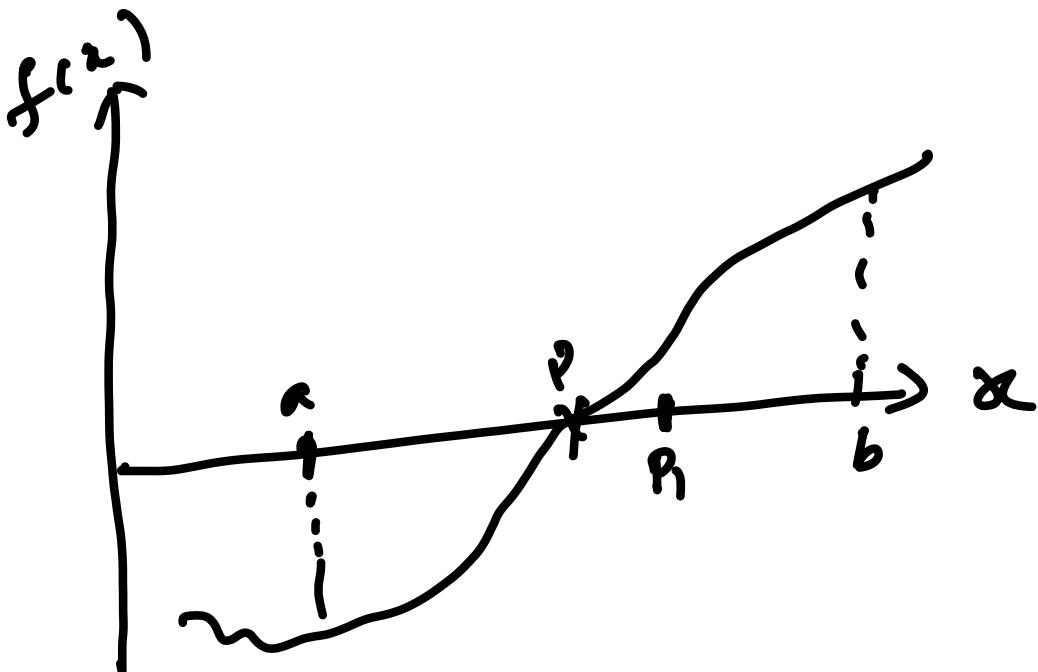
TOO BIG \rightarrow MISS 'P'.

"NO GUARANTEES"

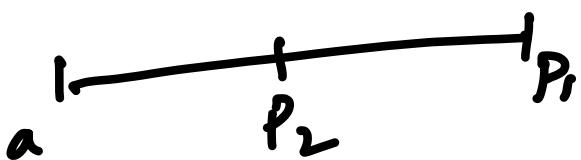
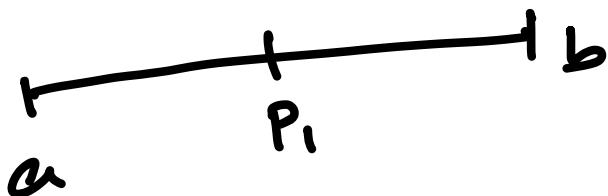
i) BISECTION METHOD
(§ 2.1)

If $f(a)f(b) < 0$ THEN INT
GUARANTEES A ROOT $p \in [a, b]$

\rightarrow BRACKETING THE
ROOT.



$$P_1 = a + \frac{b-a}{2} = \frac{b+a}{2}$$



EVALUATE $f(a_n), f(b_n), f(P_n)$

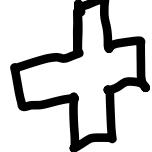
$$a_0 = a ; b_0 = b$$

DETERMINE NEW INTERVAL

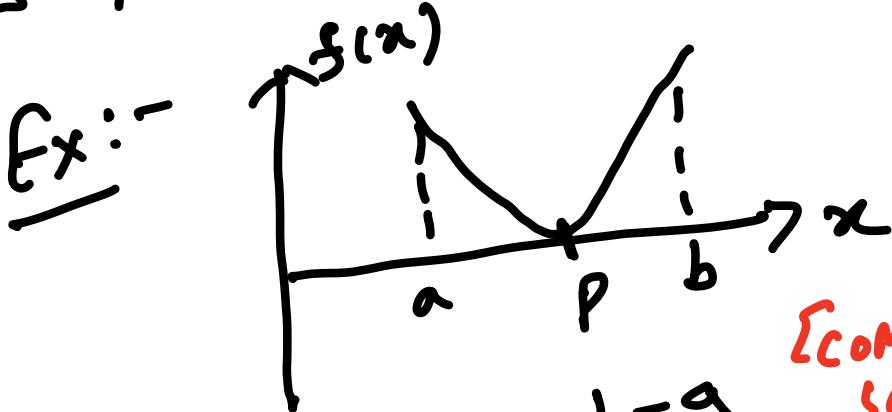
SUCH THAT $f(a_{n+1}) f(b_{n+1}) < 0$

$$a_{n+1} = \begin{cases} a_n & f(a_n) f(p_n) < 0 \\ p_n & f(a_n) f(p_n) > 0 \end{cases}$$

$$b_{n+1} = \begin{cases} b_n & f(b_n) f(p_n) < 0 \\ p_n & f(b_n) f(p_n) > 0 \end{cases}$$

 SIMPLE, ALWAYS FINDS ROOT.

 NEED $[a, b]$ $\Rightarrow f(a) f(b) < 0$



* SLOW $\epsilon_n \leq \frac{b-a}{2^n}$ [COMPUTER SCIENCE:-]

ERROR BEHAVIOR [BINARY SEARCH PROJ]

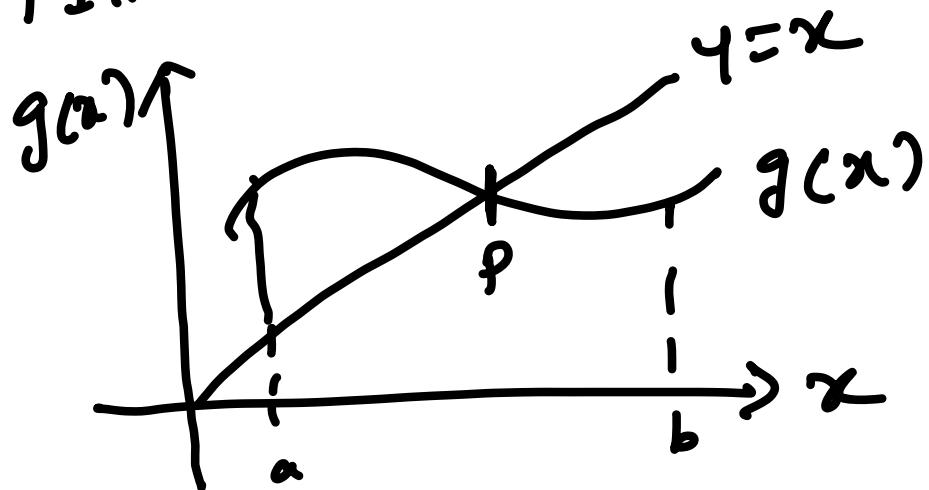
(PROOF IN TEXT THEOREM 2.1)

AS n INCREASES $p_n \rightarrow p$

FIXED POINT SCHEMES

- USEFUL FOR ROOT FINDING

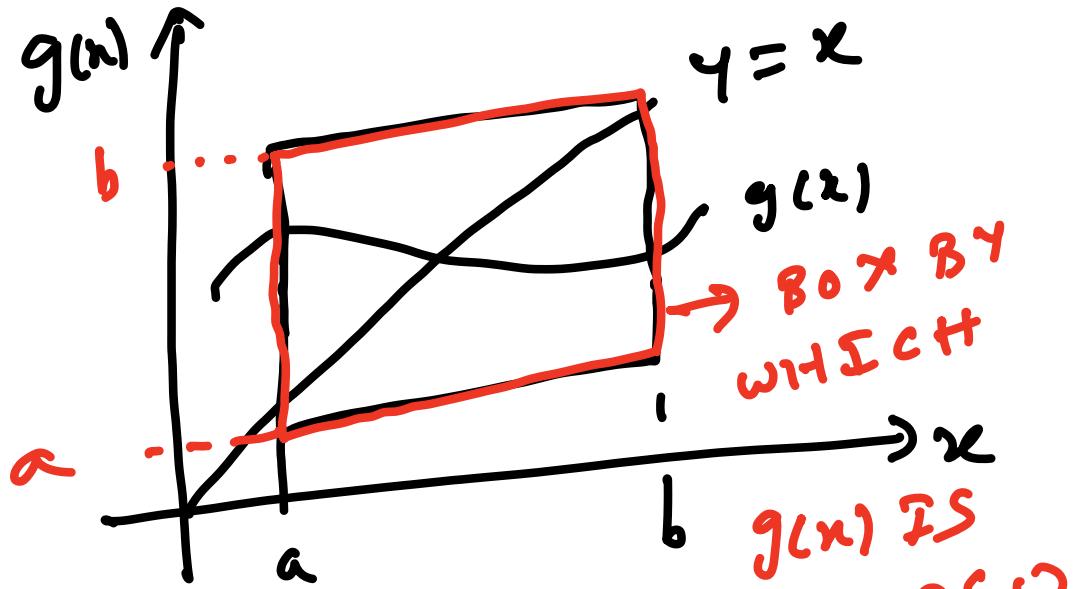
* FIXED POINT ' p ' WHEN $g(p) = p$



p EXISTS ON $[a, b]$ IF

(i) $g(x)$ MUST BE CONTINUOUS
ON $[a, b]$

(ii) $g(x)$ MUST BE "BOUNDED" ON
 $[a, b]$ BY $[a, b]$



(i) & (ii) ARE SUFFICIENT
 BUT NOT NECESSARY
 ALONE DOES NOT GUARANTEE
 UNIQUENESS.

SUFFICIENT FOR EXISTENCE
 OF UNIQUE FIXED POINT
 $\in [a, b]$

(i) $g(x) \in G[a, b]$

(ii) $g(x)$ IS BOUNDED BY $[a, b]$

(iii) $|g'(x)| < 1$ if $x \in [a, b]$

↳ ADDITIONAL CONSTRAINT

RECALL "ROOT FINDING" $f(p)=0$
HOW TO GET $g(x)$ FROM $f(x)$
→ WHEN $g(p)=p$ THEN
 $f(p)=0$

Ex: $g(x) = x - f(x)$
WHEN $g(p)=p$ THEN $f(p)=0$

* FIND FIXED POINT p USING
 $g(x)$ WHICH CORRESPONDS TO
ROOT IN $f(x)$ *

USE FIXED POINT METHODS
BECAUSE THEY HAVE DESIRABLE
& PROVABLE FEATURES SUCH AS
* SIMPLE ALGORITHMS
* CONVERGENCE CONDITIONS
* CONVERGENCE ARE PROVABLE
* ROUND-OFF ERRORS BEING
SMALLER.

$$\text{if } f(p) = 0 \equiv g(p) = p$$

REPRESENTATION IS HARDER
IN COMPUTER