



MECS1023 ADSA Assignment

“BST-Based Solution of Pharmacy Inventory Control System”

Group Name	: Group 1	
Group Members	: Lau Su Hui (Abby)	MEC245045
	Nurul ‘Aini Binti Hamdan	MEC245046
	Winnie Ngu	MEC245065
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BST-based Medicine Management System

By: Abby

System Overview

Objective: To build an efficient console-based system for managing medicine batches using a Binary Search Tree (BST) implemented in C++.

Key Features:

- **Add:** Captures Batch ID, Expiry, Location, & Quantity.
- **Search:** Instant lookup via Batch ID.
- **Display:** Auto-sorted alphabetical inventory in a table view.
- **Delete:** Removes expired or empty batches.

Why BST?

- **Efficiency:** Faster Search/Insert ($O(\log n)$) than linear arrays ($O(n)$).
- **Scalability:** Dynamic memory allocation (pointers) allows unlimited inventory growth.



Data Structure Design

The Data Object (**struct Medicine**):

- **Key: batchID** (String) - Used for sorting.
- **Details: medicineName, quantity.**
- **Logistics: expiryDate, location** (e.g., "Shelf A01").

```
struct Medicine {  
    string batchID;  
    string medicineName;  
    int quantity;  
    string expiryDate;  
    string location;  
};
```

Data Structure Design

 **The Core Node (`struct Node`):** The building block of the tree contains:

- **Data:** A **Medicine** object.
- **Left Pointer:** Points to a node with a smaller alphabetical Batch ID.
- **Right Pointer:** Points to a node with a larger alphabetical Batch ID.

Sorting Logic:

- The tree is organised strictly by **Batch ID**.
- Left Child < Parent/Root < Right Child.
- For example:
 1. Parent/Root: "B005"
 2. Left Child: "B001" (Smaller suffix)
 3. Right Child: "B009" (Larger suffix)

```
struct Node {  
    Medicine data;  
    Node *left;  
    Node *right;  
  
    Node(Medicine med) {  
        data = med;  
        left = nullptr;  
        right = nullptr;  
    }  
};
```


Key Functions - Insert

Insert Function (**insertInternal**):

- **Logic:** Recursive comparison.
- **Flow:** 1. If the tree is empty, create the Root.
2. If the new Batch ID is **smaller** than the current node, recurse **Left**.
3. If the new Batch ID is **larger** than the current node, recurse **Right**.
4. Duplicate Batch IDs are flagged as errors to prevent data corruption.

```
Node *insertInternal(Node *current, Medicine med) {  
    if (current == nullptr) {  
        return new Node(med);  
    }  
    if (med.batchID < current->data.batchID) {  
        current->left = insertInternal(current->left, med);  
    } else if (med.batchID > current->data.batchID) {  
        current->right = insertInternal(current->right, med);  
    } else {  
        cout << "Error: Medicine with Batch ID " << med.batchID  
              << " already exists.🚫" << endl;  
    }  
    return current;  
}
```

Key Functions - Search

Search Function (**searchInternal**):

- **Logic:** Binary Search principle.
- **Flow:** Checks if the current node matches the Batch ID. If not, it eliminates half the tree instantly by choosing to go only Left or only Right.

```
Node *searchInternal(Node *current, string batchID) {  
    if (current == nullptr || current->data.batchID == batchID) {  
        return current;  
    }  
    if (batchID < current->data.batchID) {  
        return searchInternal(current->left, batchID);  
    } else {  
        return searchInternal(current->right, batchID);  
    }  
}
```

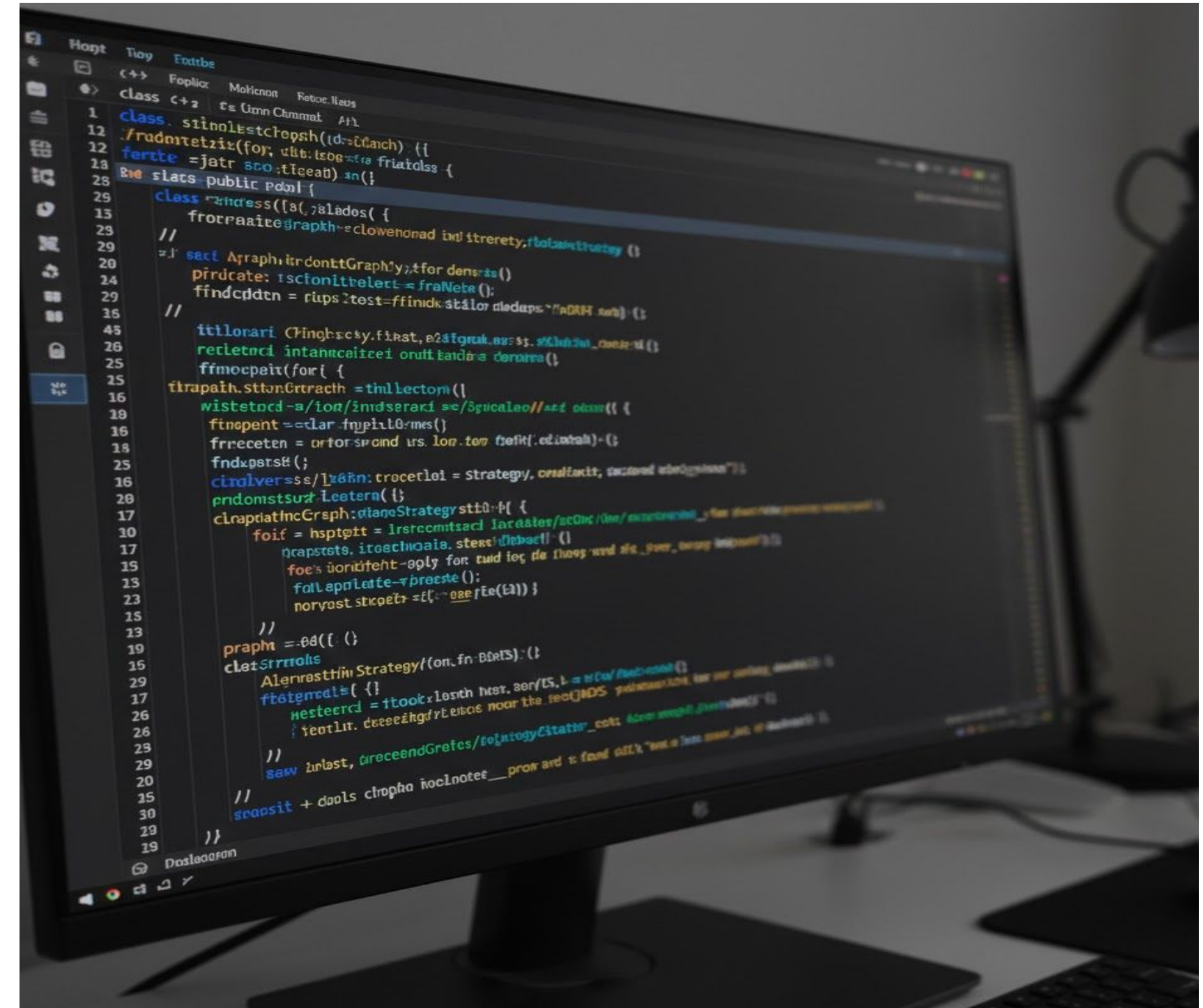
Key Functions - Delete

 **Logic:** **removeInternal** handles 3 cases to keep the tree sorted.

- **Leaf Node:** Delete directly (no children).
- **One Child:** Bypass node; link Parent → Child.
- **Two Children:** Replace with **In-Order Successor** (Right Min), then delete the successor.

```
Node *removeInternal(Node *current, string batchID) {
    if (current == nullptr)
        return current;
    if (batchID < current->data.batchID) {
        current->left = removeInternal(current->left, batchID);
    } else if (batchID > current->data.batchID) {
        current->right = removeInternal(current->right, batchID);
    } else {
        if (current->left == nullptr) {
            Node *temp = current->right;
            delete current;
            return temp;
        } else if (current->right == nullptr) {
            Node *temp = current->left;
            delete current;
            return temp;
        }
        Node *temp = findMinNode(current->right);
        current->data = temp->data;
        current->right = removeInternal(current->right, temp->data.batchID);
    }
    return current;
}
```


Demo of Working Functions



Improvements for Next Draft

- 🖥️ • **Expiry Warning:** Utilise the expiryDate field to list batches that are expiring in < 30 days on a dashboard.
- **File Persistence:** Implement `saveToFile()` and `loadFromFile()` so data is not lost when the console is closed.
- **Category Grouping:** Sort medicines by therapeutic class (e.g., Antibiotics).



BST-based Pharmacy Stock System

By: Winnie

System Overview

Objective: To build an efficient console application for managing stock in pharmacy.

Key Features:

- Add new medicines (ID, Name, Qty, Price).
- Search instantly by ID.
- Update stock details.
- Delete records dynamically.
- Display sorted inventory.

Why BST?

- Faster Search/Insert than Arrays ($O(\log n)$ vs $O(n)$).
- Dynamic memory (no fixed size limit).



Data Structure Design

 Structure **Medicine** contains:

- Medicine unique ID, name
- Quantity
- Price

```
struct Medicine {  
    string id;  
    string name;  
    int quantity{};  
    double price{};  
};
```

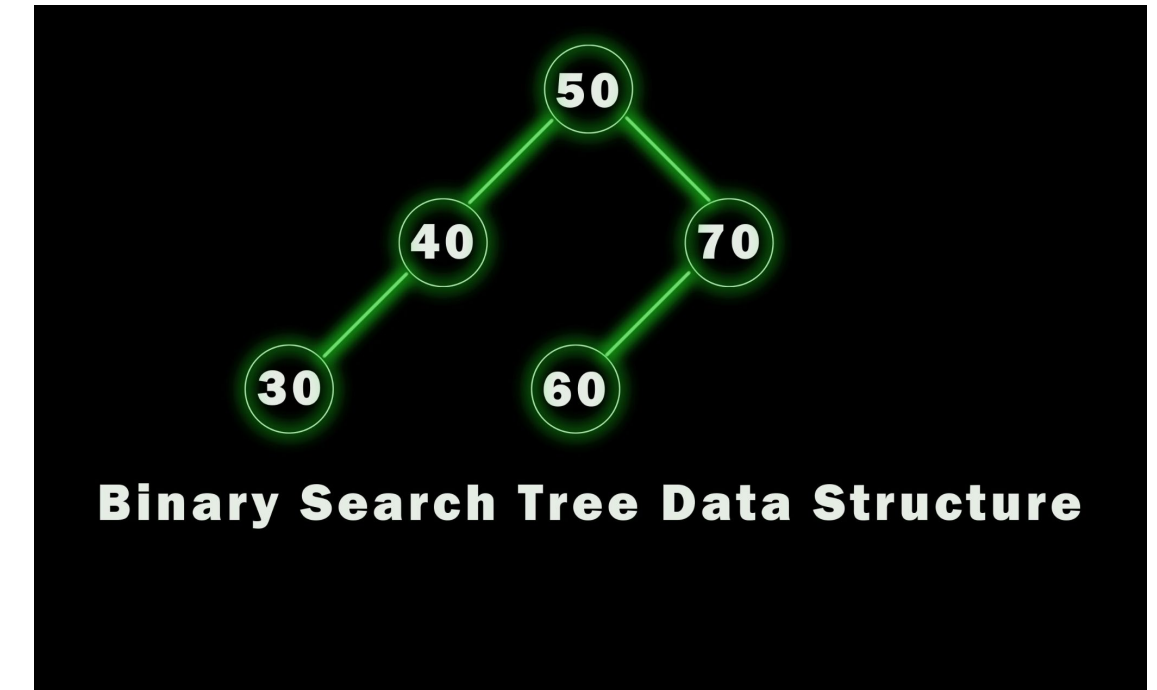
Data Structure Design

The Core Node:

- Structure **Node** contains:
 - **Medicine data** (Struct)
 - **Node* left** (Pointer to smaller IDs)
 - **Node* right** (Pointer to larger IDs)

Sorting Logic:

- The tree is organized by **Medicine ID** (String comparison).
- Left Child < Parent < Right Child.



```
struct Node {  
    Medicine data;  
    Node* left;  
    Node* right;  
  
    Node(const Medicine& m) : data(m), left(nullptr), right(nullptr) {}  
};
```


Key Functions - Insert & Search


 **Insert:** Recursive logic. Smaller IDs go Left, Larger go Right. Handles duplicates by updating.

Search: Recursive traversal based on ID comparison.

```
Node* insert(Node* node, const Medicine& med) {
    if (node == nullptr) {
        return new Node(med);
    }
    if (med.id < node->data.id) {
        node->left = insert(node->left, med);
    } else if (med.id > node->data.id) {
        node->right = insert(node->right, med);
    } else {
        //Update record if same medicine exists
        cout << "Medicine with ID " << med.id
              << " already exists. Updating record...\n";
        node->data.name = med.name;
        node->data.quantity = med.quantity;
        node->data.price = med.price;
    }
    return node;
}
```

```
//Search
Node* search(Node* node, const string& id) const {
    if (node == nullptr) return nullptr;
    if (id == node->data.id) return node;
    if (id < node->data.id) return search(node->left, id);
    return search(node->right, id);
}
```

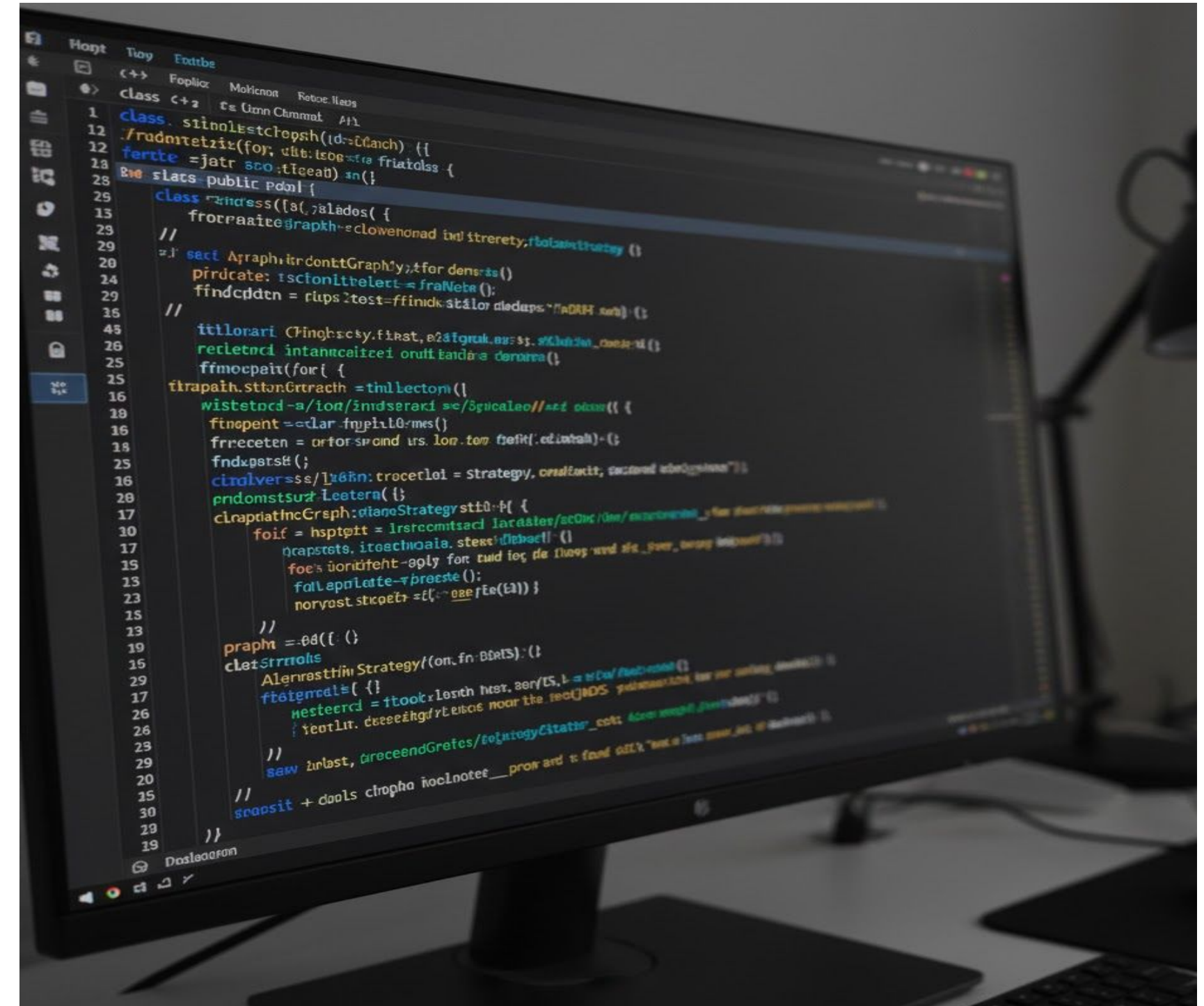
Key Functions - Delete

-  **Case 1:** Leaf Node (Delete).
- Case 2:** One Child (Bypass/Link Parent to Child).
- Case 3:** Two Children (Replace with **In-Order Successor**).

```
//Remove
Node* remove(Node* node, const string& id) {
    if (node == nullptr) return nullptr;

    if (id < node->data.id) {
        node->left = remove(node->left, id);
    } else if (id > node->data.id) {
        node->right = remove(node->right, id);
    } else {
        // Found node to delete
        if (node->left == nullptr && node->right == nullptr) {
            // Leaf node
            delete node;
            return nullptr;
        } else if (node->left == nullptr) {
            Node* temp = node->right;
            delete node;
            return temp;
        } else if (node->right == nullptr) {
            Node* temp = node->left;
            delete node;
            return temp;
        } else {
            // Two children- replace with inorder successor
            Node* succ = findMin(node->right);
            node->data = succ->data;
            node->right = remove(node->right, succ->data.id);
        }
    }
}
```


Demo of Working Functions



Improvements for Next Draft

- **Low Stock Alert:** Identifies medicines falling below a specific threshold (e.g., less than 10 units).
- **Range Search (ID):** Finds all medicines within a specific ID range (e.g., from **M002** to **M008**).
- **Total Asset Valuation:** Calculates the total monetary value of the entire stock ($\text{Price} \times \text{Quantity}$).



BST-based Supplier and Vendor Management System

By: Aini

System Overview

Objective: To centralize and streamline the management of supplier relationships to maximize value, minimize risks, and ensure efficient procurement and high-quality supply chain performance

Key Features:

- Add new supplier/vendor
- Search supplier/vendor by ID
- Delete supplier/vendor by ID
- Update supplier/vendor information
- Display supplier/vendor sorted list

```
class Supplier:
    def __init__(self, supplier_id,
                 supplier_name,
                 supplier_contact,
                 supplier_product,
                 supplier_rating):
```

Why BST?

- Provide efficient average-case insertion and search operations with time complexity of $O(\log n)$.
- Returns data sorted by keys without needing extra sorting.
- Enables intuitive hierarchical data organization and recursive operations for tasks



Data Structure Design

 Class **Supplier** contains:

- Supplier ID, name, contact
- Type of products supplied
- Rating of supplier

```
class Supplier:
    def __init__(self, supplier_id
        self.supplier_id = supplie
        self.supplier_name = suppl
        self.supplier_contact = su
        self.supplier_product = su
        self.supplier_rating = sup
```

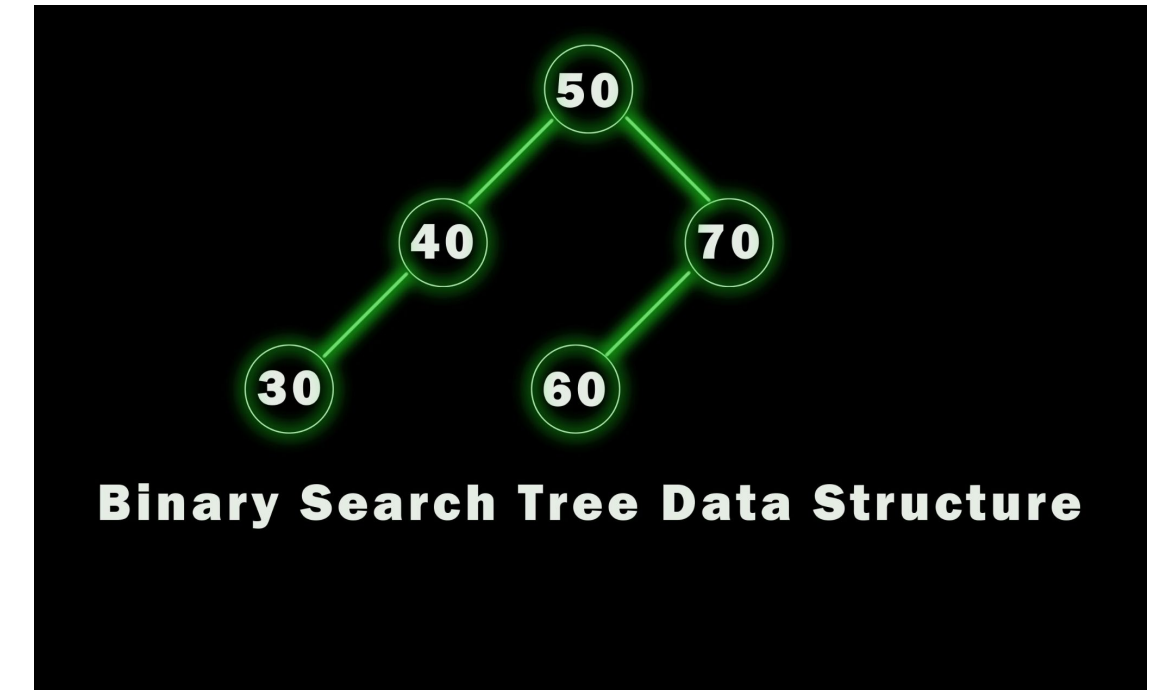
Data Structure Design

The Core Node:

- Class `BSTNode` contains:
 - `self.supplier` (store supplier object)
 - `self.left` (Pointer to smaller IDs)
 - `self.right` (Pointer to larger IDs)

Sorting Logic:

- The tree is organized by **Supplier ID** (Numeric comparison).



```
class BSTNode:
    def __init__(self, supplier):
        self.supplier = supplier # Store supplier object
        self.left = None # Pointer to left child (smaller ID)
        self.right = None # Pointer to right child (larger ID)
```

Key Functions - Insert & Search

 **Insert:** Recursive logic. Smaller IDs go Left, Larger go Right.

Search: Recursive traversal based on ID comparison.

```
# Insert new supplier into BST based on ID
def insert(self, supplier):
    if supplier.supplier_id < self.supplier.supplier_id:
        if self.left is None:
            self.left = BSTNode(supplier)
        else:
            self.left.insert(supplier)
    else:
        if self.right is None:
            self.right = BSTNode(supplier)
        else:
            self.right.insert(supplier)
```

```
# Search for supplier by ID in BST
def search(self, supplier_id):
    if supplier_id == self.supplier.supplier_id:
        return self.supplier
    elif supplier_id < self.supplier.supplier_id and self.left:
        return self.left.search(supplier_id)
    elif supplier_id > self.supplier.supplier_id and self.right:
        return self.right.search(supplier_id)
    else:
        return None
```


Key Functions - Delete



Case 1: Node with No Children (Leaf Node)

Case 2: Node with One Child (bypass node and connect the parent directly to grandchild)

Case 3: Node with Two Children (in-order successor using smallest node in the right subtree).

```
# Delete supplier from BST
def delete(self, supplier_id):
    if supplier_id < self.supplier.supplier_id:
        if self.left:
            self.left = self.left.delete(supplier_id)
    elif supplier_id > self.supplier.supplier_id:
        if self.right:
            self.right = self.right.delete(supplier_id)
    else:
        # Node to be deleted found
        if self.left is None:
            return self.right
        elif self.right is None:
            return self.left

        # Node with two children: get inorder successor
        temp = self.right
        while temp.left:
            temp = temp.left
        self.supplier = temp.supplier
        self.right = self.right.delete(temp.supplier.supplier_id)
```


Demo of Working Functions



Improvements for Next Draft

- **Supplier Category:** Add product categories for better organization.
- **Audit Trail:** Track all changes made to supplier records.
- **Search Enhancements:** Add search by name, product, contact or rating range.





Thank You

Questions?