

Multiple Constraints and Non-regular Solution in Deep Declarative Network

Suikei Wang

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Except where otherwise indicated, this thesis is my own original work.

Suikei Wang
17 October 2020

to my parents

Acknowledgments

The past two years at the Australian National University have been an invaluable experience for me. When I started my Master of Machine Learning and Computer Vision at the beginning of 2019, I could barely understand lectures, knew little about the country, and had never heard of the term "convex optimization". It is unbelievable that I have been doing a research project on this topic for a whole year. ANU has the top tier research group in this realm and how honored I am to be a postgraduate student here.

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Abstract

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Contents

Acknowledgments	vii
Abstract	ix
1 Introduction	1
1.1 Motivation	1
1.2 Thesis Outline	1
1.3 Contribution	1
I Deep Declarative Network: Multiple Constrained Declarative Nodes	3
2 An Overview of Numerical Optimization	5
2.1 Theory of Optimization	5
2.1.1 Existence of Optimizers	5
2.1.2 Optimality Conditions for Unconstrained Problems	7
2.1.3 Optimality Conditions for Constrained Problems	9
2.2 Solution of Unconstrained and Constrained Optimization Problems . .	12
2.2.1 Unconstrained Optimization	13
2.2.2 Equality Constrained Optimization	14
2.2.3 Inequality Constrained Optimization	16
2.3 Differentiable Neural Network	17
2.4 Summary	18
3 Deep Declarative Network	19
3.1 An Overview of Deep Declarative Network	19
3.1.1 Structure	19
3.1.2 Declarative Nodes	19
3.2 Learning	19
3.3 Back-propagation Through Declarative Nodes	19
3.3.1 Unconstrained	19
3.3.2 Equality Constrained	19
3.3.3 Inequality Constrained	19
3.4 Examples of Declarative Nodes	19
3.4.1 Unconstrained	19
3.4.2 Equality Constrained	19
3.4.3 Inequality Constrained	19

3.5	Summary	19
4	The Future of Declarative Nodes	21

II	Deep Declarative Network: Non-regular Solution	23
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5	An Overview of Regular and Non-regular Solution	25
5.1	Problems in Regular Deep Declarative Nodes	25
5.2	Related Work in Non-regular Solution	25
5.2.1	Overdetermined System	25
5.2.2	Rand Deficiency	25
5.2.3	Non-convex Problems	25
6	Solutions of Non-regular Point	27
6.1	Overdetermined System	27
6.1.1	Least-Squared Method	27
6.2	Conjugate Gradient and Preconditioning	27
6.3	Rank Deficiency	27
6.4	Non-convex Problems	27
6.5	Summary	27
7	Conclusion	29
7.1	Future Work	29
Appendix		33
A	An Overview of Numerical Optimization	33
A.1	Theory of Optimization	33
A.1.1	Proof of Theorem 2.3	33
A.1.2	Proof of Theorem 2.4	33
A.1.3	Proof of Theorem 2.6	34
A.1.4	Proof of Theorem 2.8	34
A.1.5	Proof of Lemma 2.10	35
A.2	Solution of Unconstrained and Constrained Optimization Problems	36
A.2.1	Proof of Equation 2.3	36
A.2.2	Proof of Equation 2.4 [Gould et al., 2019]	36

List of Figures

2.1	Contour Graph of $f(x, y) = x^4 - 4xy + y^4$	8
2.2	Function $f(x, y)$ at $x = 0$, $x = -1$ and $x = 1$	9
2.3	Feasible set of constraints c_1 , c_2 and c_3	11
2.4	Constrained function $c(x) = \ x\ _2^2 - 1$	15
5.1	Hello world in Java and C.	26
6.1	The cost of zero initialization	28

List of Tables

5.1	Processors used in our evaluation.	25
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Introduction

1.1 Motivation

Deep learning models composed with multiple parametrized processing layers can learn different levels of features and representations of data through the directed graph structure.

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thesis!!!

1.2 Thesis Outline

How many chapters you have? You may have Chapter 2, Chapter ??, Chapter ??, Chapter 6, and Chapter 7.

1.3 Contribution

Part I

Deep Declarative Network: Multiple Constrained Declarative Nodes

An Overview of Numerical Optimization

In this chapter, we aim to provide readers with an overview of numerical optimization. We begin with the theory of optimization (Section 2.1), from the existence of optimizers, to the optimality conditions for both unconstrained and constrained problems with duality. As the theoretical background of optimization, this field provides a solid solution for the algorithm.

We then formally define the optimization of unconstrained and constrained problems in Section 2.2 and describe the general regular solution for these problems based on the gradient calculation.

Next, we discuss briefly the differentiable optimization in neural network, which is a novel end-to-end network structure involving optimization problem in each layer with and without constraints. (Section 2.3). Finally, we give a summary of the numerical optimization for solving unconstrained and constrained problems in Section 2.4.

2.1 Theory of Optimization

2.1.1 Existence of Optimizers

In optimization, a basic question is to determine the existence of a global minimizer for a given function f . There are several sufficient conditions on f to guarantee the existence, and the optimizer falls in the feasible set of solutions. For a feasible set, some related definitions are following:

Definition 2.1. A subset $\Omega \in \mathbb{R}^n$ is called

- *bounded* if there is a constant $R > 0$ such that $\|x\| \leq R$ for all $x \in \Omega$
- *closed* if the limit point of any convergent sequence in Ω always lies in Ω
- *compact* if any sequence $\{x_k\}$ in Ω contains a subsequence that converges to a point in Ω

The following result gives a characterization of compact sets in \mathbb{R} . When we find the minimum or maximum solution for the problem, there exists a lower bound

or upper bound but not necessarily an optimal solution. Therefore, we have some additional requirements.

Firstly, we give the definition of compact sets in Lemma 2.2. [Oman, 2017] gives a brief proof.

Lemma 2.2 (Bolzano-Weierstrass theorem). A subset Ω in \mathbb{R}^n is *compact* if and only if it is bounded and closed.

We also assume that the function f is continuous and " $+\infty$ at infinity". More precisely, $f(x) \rightarrow +\infty$ if $|x| \rightarrow +\infty$. Such a function is called *inf-compact* or *coercive*. [Nocedal and Wright, 2006] Then the problem can be restricted to a bounded set and existence of a global minimum x^* is guaranteed: a continuous function has a minimum on a compact set. This theorem is defined as follows and the proof is given in Appendix A.1.1.

Theorem 2.3. [Nocedal and Wright, 2006] If f is a continuous function defined on a compact set Ω in \mathbb{R} , then f has a global minimizer x^* on Ω i.e. there exists $x^* \in \Omega$ such that $f(x^*) \leq f(x)$ for all $x \in \Omega$

More general, based on the definition of coercive function f , we can give following theorem. Proof is given in Appendix A.1.2.

Theorem 2.4. [Nocedal and Wright, 2006] If $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is a continuous coercive function, then f has at least one global minimizer.

Theorem 2.4 requires the continuity of f which is somewhat restrictive for applications. However, we can replace it by the lower semi-continuity of f which is a rather weaker condition.

Definition 2.5. Let $f : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{\pm\infty\}$. Then f is called *lower semi-continuous* at a point $x_0 \in \mathbb{R}^n$ if for any sequence $f(x_k)$ converging to x_0 here holds $f(x_0) \leq \lim_{k \rightarrow \infty} f(x_k)$. f is called *lower semi-continuous* if f is lower semi-continuous at every point.

Recall our assumptions on function f , it is a continuous function, which is always lower semi-continuous. However, lower semi-continuous functions are not necessarily continuous. For instance, a binary function equals to 0 when $x \leq 0$ and equals to 1 when $x > 0$ is not continuous at $x_0 = 0$. However, since it is greater than 0 for all x and $f(0) = 0$, we have $f(0) = 0 \leq \liminf_{x \rightarrow 0} f(x)$ and it is lower semi-continuous at $x_0 = 0$.

The theorem of the existence of the optimizer of lower semi-continuous function is given as follows and the proof is given in Appendix A.1.3

Theorem 2.6. [Nocedal and Wright, 2006] Let $f : \mathbb{R}^n \rightarrow \mathbb{R}$ be a lower semi-continuous function. If f has a nonempty, compact sublevel set $D := \{x \in \mathbb{R}^n : f(x) \leq \alpha\}$, then f achieves a global minimizer on \mathbb{R}

Also, we introduce the definition of convex function and convex set which are important in regular optimization problems.

Definition 2.7. A function f is convex when

$$f(\alpha x + (1 - \alpha)y) \leq \alpha f(x) + (1 - \alpha)f(y) \quad \text{for all } x, y, \text{ and } \alpha \in]0, 1[$$

A set $C \subset \mathbb{R}^n$ is convex when

$$\alpha x + (1 - \alpha)y \in C \quad \text{for all } x, y \text{ in } C, \text{ and } \alpha \in]0, 1[$$

The problem we are going to discuss in this part is convex and regular, which means its gradient can be computed and the solution exists. However, although the existence of the optimizer is sufficient, for different problems, the optimality conditions are different. In the next two sections, we will give necessary and sufficient conditions for both unconstrained and constrained problems.

2.1.2 Optimality Conditions for Unconstrained Problems

Firstly, we consider the unconstrained minimization problem

$$\min_{x \in \mathbb{R}^n} f(x), \tag{2.1}$$

where f is given function on \mathbb{R}^n .

In order to determine the minimizer, it is important to understand what can happen at a minimizer, and at what condition a point must be a minimizer. Now we have to recognize the optimum point. There are two necessary conditions and one sufficient condition given below [Nocedal and Wright, 2006]. The proof is given in Appendix A.1.4.

Theorem 2.8. Necessary and Sufficient Conditions. Let $f : \Omega \rightarrow \mathbb{R}$ be a function defined on a set $\Omega \subset \mathbb{R}^n$ and let x^* be an interior point of Ω that is a local minimizer of f .

Necessary conditions:

- (NC1) If f is differentiable at x^* , then x^* is a critical point of f , i.e. $\nabla f(x^*) = 0$.
- (NC2) If f is twice continuous differentiable on Ω , then the Hessian $\nabla^2 f(x^*)$ is positive semidefinite.

Sufficient condition (SC1): if x^* is such that $\nabla f(x^*) = 0$ and $\nabla^2 f(x^*)$ is positive definite, then x^* is a local minimum. (i.e. $f(x) \geq f(x^*)$ for x close to x^*)

Any point satisfying (NC1) as the minimizer of f is called a *critical* or *stationary* point of f . In the objective function f is convex, (NC1) is also the sufficient condition for the global minimum of the solution.

Let us see an example of unconstrained minimization problem. Supposed we have to determine the minimization of function

$$f(x, y) = x^4 - 4xy + y^4$$

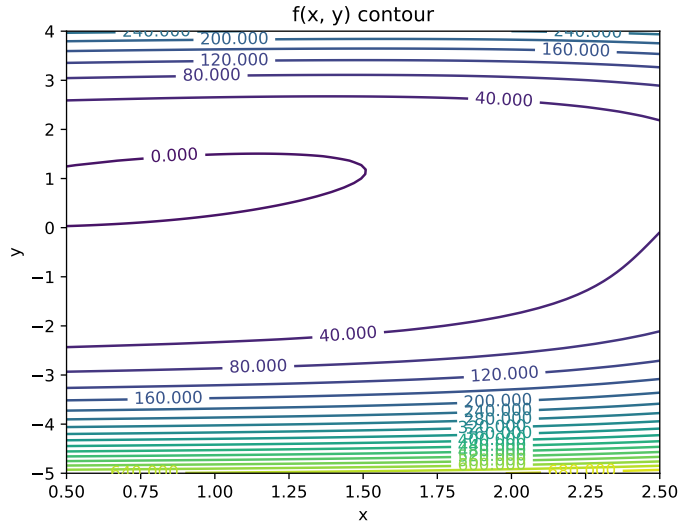


Figure 2.1: Contour Graph of $f(x, y) = x^4 - 4xy + y^4$

From the definition of function f , it is clear that f is continuous. Then we can expand f by writing

$$f(x, y) = (x^4 + y^4) \left(1 - \frac{4xy}{x^4 + y^4} \right)$$

we can see f is coercive. Also, we give the contour graph of function f in Figure 2.1. Therefore f has global minimizers which are critical points. Now according to (NC1), we can find the global minimizer through solving the derivative of f equaling to zero:

$$0 = \nabla f(x, y) = \begin{pmatrix} 4x^3 - 4y \\ -4x + 4y^3 \end{pmatrix}$$

Thus, $y = x^3$ and $x = y^3$. Consequently $y = y^9$, i.e.

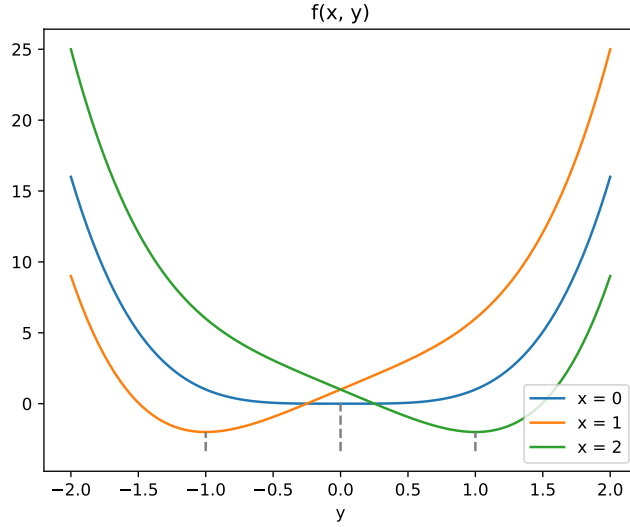
$$0 = y - y^9 = y(1 - y^8) = y(1 - y^4)(1 + y^4) = y(1 - y)(1 + y)(1 + y^2)(1 + y^4)$$

This implies $y = 0, 1, -1$. Thus f has three critical points $(0, 0), (1, 1), (-1, -1)$. Then we can evaluate f as these points since they may be local minimizer:

$$f(0, 0) = 0, \quad f(1, 1) = -2, \quad f(-1, -1) = -2$$

It achieves the same global minimum value on $(1, 1)$ and $(-1, -1)$. Therefore, they are both global minimizers of f . Figure 2.2 shows the function $f(x, y)$ at these two optimal points.

From this example, we verify that through (NC1), we can find the global minimizer. However, not all continuous functions with critical points have any maximizer

Figure 2.2: Function $f(x, y)$ at $x = 0$, $x = -1$ and $x = 1$

or minimizer. If the function goes to infinity along its axes or a line, it does not have any maximizer or minimizer although it has a critical point. The condition of the minimizer as the critical point is that the function f should be a convex function with continuous first partial derivatives.

Let us move to the sufficient condition (SC1). The result obtained under this theorem is best possible for general functions. Specifically, for a convex function f is defined on a convex set $\Omega \subset \mathbb{R}^n$, any local minimizer of f is also a global minimizer. Moreover, if a function f is strictly convex, it has at most one global minimizer.

2.1.3 Optimality Conditions for Constrained Problems

A general formulation for constrained optimization problems is as follows:

$$\begin{aligned} & \text{minimize } f(x) \\ & \text{subject to } \begin{cases} c_i(x) = 0 & \text{for } i = 1, \dots, m_e, \\ c_i(x) \leq 0 & \text{for } i = m_e + 1, \dots, m \end{cases} \end{aligned} \quad (2.2)$$

where f and c_i are smooth real-valued functions on \mathbb{R}^n , and m_e and m are nonnegative integers with $m_e < m$. We set

$$\mathcal{E} := \{1, \dots, m_e\} \quad \text{and} \quad \mathcal{I} := \{m_e + 1, \dots, m\}$$

as index sets of equality constraints and inequality constraints, respectively.

Here, f is so-called the objective function, and $c_i, i \in \mathcal{E}$ and \mathcal{I} are equality constraints and inequality constraints respectively.

To solve the optimization problem (2.2), we define the feasible set of it to be

$$\mathcal{F} := \{x \in \mathbb{R}^n : c_i(x) = 0 \text{ for } i \in \mathcal{E} \text{ and } c_i(x) \leq 0 \text{ for } i \in \mathcal{I}\}$$

Any point $x \in \mathcal{F}$ is called a feasible point of (2.2) and we call (2.2) infeasible if $\mathcal{F} = \emptyset$. Also, in this feasible set, a feasible point $x^* \in \mathcal{F}$ is called a local minimizer of (2.2) if it is the minimum solution in a neighborhood (strict local minimizer if it is the only one minimum solution). The definition of the global minimizer and strict global minimizer is similar, whose neighborhood is the whole feasible set.

Let us move to the constraints in this problem. For equality constraints, they are strictly equivalent. However, for inequality constraints, there are some exceptions. Let x^* be a local minimizer of (2.2). If there is an index $i \in \mathcal{I}$ such that $c_i(x^*) < 0$, then, x^* is still the local minimizer of the problem obtained by deleting i -th constraint. In this situation, we say that the i -th constraint is inactive at x^* since it does not have any effect on the solution. A general definition of active and inactive inequality constraints is as follows:

Definition 2.9. At a feasible point $x \in \mathcal{F}$, the index $i \in \mathcal{I}$ is said to be *active* if $c_i(x) = 0$ and *inactive* if $c_i(x) < 0$.

In the next chapter, we will give different processes for different cases of active or inactive inequality constraints in the deep declarative nodes. In this chapter, we only focus on the necessary and sufficient conditions for a feasible point x to be a local minimizer of (2.2). These conditions will be derived by considering the change of f on the feasible set along with certain directions. We give the lemma for the condition of local minimizer $x^* \in \mathcal{F}$ as follows, which can be proved through Taylor's formula in Appendix A.1.5.

Lemma 2.10. If $x^* \in \mathcal{F}$ is a local minimizer of (2.2), then

$$d^T \nabla f(x^*) \geq 0 \quad \text{for all } d \in T_{x^*} \mathcal{F}$$

where $T_{x^*} \mathcal{F}$ is the set of all vectors tangent to \mathcal{F} .

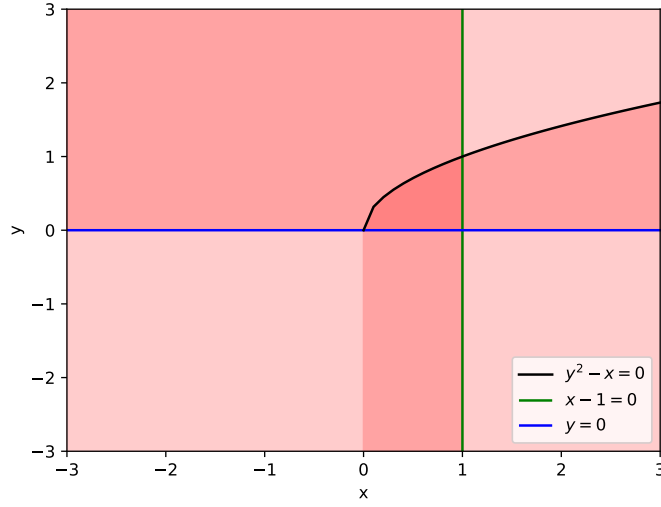
However, we may not be able to extract useful results from this lemma, since $T_{x^*} \mathcal{F}$ depends only on the geometry of \mathcal{F} but not on the constraints functions c_i . Not all local minimum falls on the boundary of the constraint function, which is a part of $T_{x^*} \mathcal{F}$. Therefore, it is necessary to introduce linearized feasible directions to give a characterization of $T_{x^*} \mathcal{F}$ in terms of c_i .

Definition 2.11. Given $x \in \mathcal{F}$, we define

$$\text{LFD}(x) := \left\{ d \in \mathbb{R}^n : d^T \nabla c_i(x) = 0 \text{ for } i \in \mathcal{E}; d^T \nabla c_i(x) \leq 0 \text{ for } i \in \mathcal{I} \cap \mathcal{A}(x) \right\}$$

and call it the set of linearized feasible directions of \mathcal{F} at x .

Heuristically, for $i \in \mathcal{E}$ we should travel along directions d with $d^T \nabla c_i(x) = 0$ in order to stay on the curve $c_i(x) = 0$; for $i \in \mathcal{I}$ we should travel along directions with

Figure 2.3: Feasible set of constraints c_1 , c_2 and c_3

$d^T \nabla c_i(x) \leq 0$ in order to stay in the region $c_i(x) \leq 0$. Let us see an example of the linearized feasible directions and the tangent. Supposed we are considering a set \mathcal{F} with variables $(x, y) \in \mathbb{R}^2$ and three inequality constraints functions:

$$c_1(x, y) = x - 1 \leq 0$$

$$c_2(x, y) = -y \leq 0$$

$$c_3(x, y) = y^2 - x \leq 0$$

We can illustrate the feasible set of constraints c_1 , c_2 and c_3 in Fig 2.3. The active set of $0 = (0, 0)$ is $\{2, 3\}$, since $c_1(0) = -1 < 0$, which is inactive. And we can get the derivative of c_2 and c_3 at 0:

$$\nabla c_2(0) = (0, -1)^T \quad \text{and} \quad \nabla c_3(0) = (-1, 0)^T$$

Then we have the linearized feasible directions on $x = 0$:

$$\begin{aligned} \text{LFD}(0) &= \left\{ d \in \mathbb{R}^2 : d^T \nabla c_2(0) \leq 0 \text{ and } d^T \nabla c_3(0) \leq 0 \right\} \\ &= \{ d \in \mathbb{R}^2 : d \geq 0 \} \end{aligned}$$

which equals to the set of all vectors tangent to the feasible set $T_0 \mathcal{F}$.

Unlike the unconstrained optimization problem, the first order necessary condition of the existence of the optimizer is different since we should consider its linearized feasible directions and constraints feasibility. This is so-called the Karush-Kuhn-Tucker theorem:

Theorem 2.12 (Karush-Kuhn-Tucker Theorem). *Let $x^* \in \mathcal{F}$ be a local minimizer of*

problem (2.2). If

$$T_{x^*}\mathcal{F} = \text{LFD}(x^*),$$

then there exists $\lambda^* = (\lambda_1^*, \dots, \lambda_m^*)^T \in \mathbb{R}^m$ such that

$$\nabla f(x^*) + \sum_{i \in \mathcal{E} \cup \mathcal{J}} \lambda_i^* \nabla c_i(x^*) = 0, \quad (\text{Lagrangian stationary})$$

$$\left. \begin{array}{l} c_i(x^*) = 0 \quad \text{for all } i \in \mathcal{E}, \\ c_i(x^*) \leq 0 \quad \text{for all } i \in \mathcal{J}, \end{array} \right\} \quad (\text{primal feasibility})$$

$$\lambda_i^* \geq 0 \quad \text{for all } i \in \mathcal{J}, \quad (\text{dual feasibility})$$

$$\lambda_i^* c_i(x^*) = 0 \quad \text{for all } i \in \mathcal{E} \cup \mathcal{J}. \quad (\text{complementary slackness})$$

This set of equations are Karush-Kuhn-Tucker (KKT) conditions and a point x^* is called a KKT point if there exists λ^* such that (x^*, λ^*) satisfies the KKT conditions.

For constrained optimization problem, the classic solution is using Lagrange multipliers [Bertsekas, 2014]. This introduces the function

$$\mathcal{L}(x, \lambda) := f(x) + \sum_{i \in \mathcal{E} \cup \mathcal{J}} \lambda_i c_i(x)$$

which is called the Lagrange function. x is the primal variables and $\lambda_i, i = 1, \dots, m$ are the Lagrange multipliers or the dual variables. According to the Lagrange multipliers method, we can solve this problem through the gradient of the Lagrange function:

$$\nabla_x \mathcal{L}(x, \lambda) = \nabla f(x) + \sum_{i \in \mathcal{E} \cup \mathcal{J}} \lambda_i \nabla c_i(x)$$

Therefore, the first equation in KKT conditions can be written as

$$\nabla_x \mathcal{L}(x^*, \lambda^*) = 0$$

2.2 Solution of Unconstrained and Constrained Optimization Problems

According to the sufficient conditions for unconstrained optimization problems, we can easily compute the optimality through the first and second derivative of the objective function. For equality and inequality constrained problems, the introduction of Lagrangian \mathcal{L} is useful for their closed-form solution. Gould et al. [2016] collected both argmin and argmax bi-level optimization results with and without constraints, which also provide insightful examples of these cases. Amos and Kolter [2017] also present a solution for exact, constrained optimization within a neural network. In this thesis, we only focus on argmin problems, but the argmax problems have similar results.

In this section, we are going to provide some background for the solution of

both unconstrained and constrained optimization problems, which is based on the gradient of the regular point.

2.2.1 Unconstrained Optimization

For unconstrained optimization problems, the solution is easy to obtain since we only need to focus on the optimality of the objective function. We consider an objective function $f : \mathbb{R}^n \times \mathbb{R}^m \rightarrow \mathbb{R}$:

$$y(x) \in \operatorname{argmin} f(x, y)$$

The derivative of $y(x)$ with respect to x is

$$\frac{dy(x)}{dx} = -[\frac{\partial^2 f}{\partial y(x)^2}]^{-1} \frac{\partial^2 f}{\partial x \partial y(x)} \quad (2.3)$$

which can be proved through differentiating and chain rule. [A.2.1]

A very classic example of the unconstrained minimization problem based on a closed convex nonempty set is the L2 norm $\|\cdot\|_2$. Let $\Omega \in \mathbb{R}^n$ be a closed convex nonempty set. For any $x \in \mathbb{R}^n$, the minimization problem is defined as follows:

$$\min_{y \in \Omega} \|y - x\|_2^2$$

This problem has a unique minimizer, which can be denoted by $P_\Omega(x)$, the Euclidean projection of x onto Ω .

Proof. Let $m := \inf_{y \in \Omega} \|y - x\|_2^2$. Since $\Omega \neq \emptyset$, we have $0 \leq m < \infty$. Let $\{y_k\} \subset \Omega$ be a minimizing sequence such that $\|y_k - x\|_2^2 \rightarrow m$ as $k \rightarrow \infty$. Thus $\|y_k - x\|_2^2 \leq m + 1$ for large k which implies that $\|y_k\|_2 \leq \|x\|_2 + \sqrt{m + 1}$ for large k . Therefore $\{y_k\}$ is a bounded sequence. Consequently $\{y_k\}$ has a convergent subsequence $\{y_{k_l}\}$ with limit y^* . Since Ω is closed, we have $y^* \in \Omega$. Thus

$$m = \lim_{l \rightarrow \infty} \|y_{k_l} - x\|_2^2 = \|y^* - x\|_2^2$$

which means that m is achieved at y^* , i.e. the given minimization problem has a solution.

Next we show that the given minimization problem has a unique solution by contradiction. If the solution is not unique, let y_0 and y_1 be two distinct solutions. Then for $0 < t < 1$ we set $y_t = ty_1 + (1 - t)y_0$. Since Ω is convex, we have $y_t \in \Omega$.

Thus

$$\begin{aligned}
\|y_0 - x\|_2^2 &= \|y_1 - x\|_2^2 \leq \|y_t - x\|_2^2 = \|t(y_1 - x) + (1-t)(y_0 - x)\|_2^2 \\
&= t^2 \|y_1 - x\|_2^2 + (1-t)^2 \|y_0 - x\|_2^2 + 2t(1-t) \langle y_1 - x, y_0 - x \rangle \\
&= t \|y_1 - x\|_2^2 + (1-t) \|y_0 - x\|_2^2 - (t-t^2) \|y_1 - x\|_2^2 \\
&\quad - (1-t - (1-t)^2) \|y_0 - x\|_2^2 + 2t(1-t) \langle y_1 - x, y_0 - x \rangle \\
&= t \|y_1 - x\|_2^2 + (1-t) \|y_0 - x\|_2^2 \\
&\quad - t(1-t) \left(\|y_1 - x\|_2^2 + \|y_0 - x\|_2^2 - 2 \langle y_1 - x, y_0 - x \rangle \right) \\
&= \|y_0 - x\|_2^2 - t(1-t) \|y_1 - y_0\|_2^2
\end{aligned}$$

where $\langle \cdot, \cdot \rangle$ denotes the inner product on \mathbb{R}^n . Therefore $t(1-t) \|y_1 - y_0\|_2^2 \leq 0$ for $0 < t < 1$ and thus $\|y_1 - y_0\|_2^2 \leq 0$. So $y_1 = y_0$ which is a contradiction.

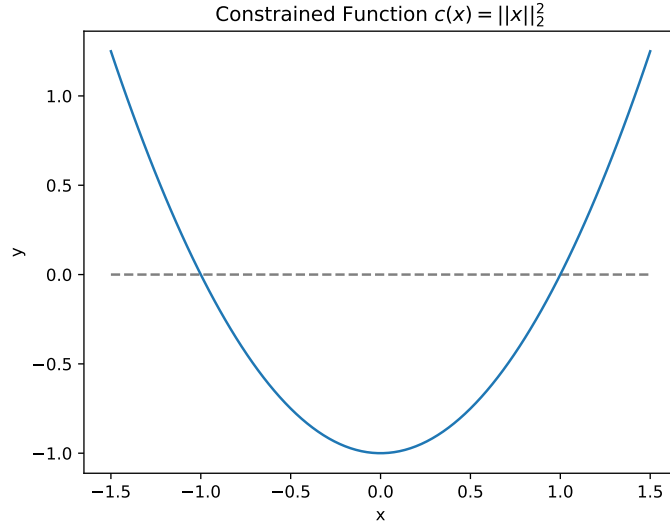
Overall, the minimization problem defined above has a unique minimizer. \square

There are many different methods to solve the unconstrained optimization problem since generally, we treat this kind of problem as basic one. There are two most classical methods, Newton method [Newton and Colson, 1736] and the Method of Steepest Descent [Debye, 1909]. The former one, Newton method starts from an initial guess x_0 and defines a sequence $\{x_k\}$ iteratively according to some rule. It uses the tangent line of the objective function f at x_k to replace f and uses the root of $L(x) = 0$, where $L(x)$ is the updated $f(x)$ as the next iterate x_{k+1} . Finally, the iteration is terminated as long as the difference between x_k and x_{k+1} less than a preassigned small number. The later one, steepest descent is a basic gradient method, which decreases the value of the objective function in a direction of most rapid change. The change rate of a function f at x in the direction u , a unit vector in \mathbb{R} is determined by the directional derivative. Therefore, at x the value of f decrease fastest in the direction $u = -\nabla f(x) / \|\nabla f(x)\|$, which leads to the gradient method: we update the x through the direction with the step length.

2.2.2 Equality Constrained Optimization

Constrained problems are usually more complicated since the solution is restricted on a boundary or in a feasible region. For equality constraints, the basic case is the linear equality constraints $Ay = b$. Again, we consider an objective function $f : \mathbb{R}^m \times \mathbb{R}^n \rightarrow \mathbb{R}$. Let $A \in \mathbb{R}^{p \times m}$ and $b \in \mathbb{R}^p$. A is a set of p linear equations as constraints $Ay = b$. The problem is defined as follows:

$$\begin{aligned}
y(x) &\in \arg \min_{y \in \mathbb{R}^m} f(x, y) \\
&\text{subject to } Ay = b
\end{aligned}$$

Figure 2.4: Constrained function $c(x) = \|x\|_2^2 - 1$

The derivative of $y(x)$ with respect to x is

$$\frac{dy(x)}{dx} = \left(H^{-1} A^T \left(A H^{-1} A^T \right)^{-1} A H^{-1} - H^{-1} \right) B \quad (2.4)$$

where $H = \partial^2 f(x, y) / \partial y(x)^2$ and $B = \partial^2 f(x, y) / \partial x \partial y(x)$.

The solution in 2.4 can be proved through the Lagrange multipliers [Bertsekas, 2014] in A.2.2. More generally, constraints can be non-linear. That means we cannot use A as a weight matrix for constrained parameters anymore. Therefore, we define the equality constraints problem using a set of m constraints functions $c(x, y)$:

$$\begin{aligned} y(x) \in \arg \min_{y \in \mathbb{R}^m} & f(x, y) \\ \text{subject to} & c_i(x, y) = 0, \quad i = 1, \dots, m \end{aligned}$$

Solution for general multiple non-linear equality constraints is discussed in the chapter of deep declarative network nodes. Here, we are giving a simple example of non-linear equality constrained optimization problem.

For any given nonzero vector $y \in \mathbb{R}^n$, we define the minimization problem as follows:

$$\begin{aligned} \text{minimize} & -x^T y \\ \text{subject to} & \|x\|_2^2 = 1 \end{aligned}$$

From the constraint defined above, we can write the constraint function as $c(x) = \|x\|_2^2 - 1$ and illustrate it in Fig 2.4. Differentiating $c(x)$ with respect to x , we get $\nabla c(x) = x \neq 0$. Therefore, it follows the definition of LFD 2.11 and the theorem of KKT 2.12, which means that every local minimizer of this problem is a KKT point.

Now we can write the Lagrangian function:

$$\mathcal{L}(x, \lambda) = -x^T y + \lambda (\|x\|_2^2 - 1)$$

and the KKT conditions are:

$$\nabla_x \mathcal{L} = -y + 2\lambda x = 0, \quad \|x\|_2^2 = 1$$

From $-y + 2\lambda x = 0$ and $y \neq 0$ defined in the question, we must have $\lambda \neq 0$ and $x = y/2\lambda$. Combined with $\|x\|_2^2 = 1$, we get

$$4\lambda^2 = \|y\|_2^2 \Leftrightarrow \lambda = \pm \frac{\|y\|_2}{2}$$

Therefore, consequently, we have $x = \pm \frac{y}{\|y\|_2}$. For each x , we can compute its corresponding value of the objective function:

$$x = \frac{y}{\|y\|_2}, \quad -x^T y = -\|y\|_2$$

$$x = -\frac{y}{\|y\|_2}, \quad -x^T y = \|y\|_2$$

Obviously, the minimum is achieved $-\|y\|_2$ at $x = \frac{y}{\|y\|_2}$.

Algorithms for solving constrained problems are various. For basic linear programming, which means that all functions involved are linear, we can transform it into standard form with matrix A , then solve the problem using Lagrangian function based on the KKT condition.

Penalty method[Yeniay, 2005], a function determining when a point x is feasible or not, is used to replace the constrained problem with an unconstrained one. For a minimization problem $f(x)$, the penalty function $P(x)$ associated with a penalty parameter are introduced to combine with $f(x)$ and now we are going to solve a series of unconstrained problems. These problems have converged solutions of the original constrained problem.

2.2.3 Inequality Constrained Optimization

Similar to equality constrained problems, inequality constrained problem usually defined the solution in a feasible set. In general, the standard form of inequality constrained problem is negative constraints:

$$\begin{aligned} y(x) &\in \arg \min_{y \in \mathbb{R}^m} f(x, y) \\ \text{subject to} \quad &c_i(x, y) \leq 0, \quad i = 1, \dots, m \end{aligned}$$

Solutions for inequality constrained problems are various. According to the properties of inequality constraints, active and inactive constraints have different criteria. We aim to find the gradient of the optimal solution, $f'(x)$, based on the inequality

constrained argmin function. Panier et al. [1988] proposed a globally convergent algorithm for solving the minimization of smooth objective function based on smooth inequality constraints. This algorithm is based on the Quasi-Newton [Dennis and Moré, 1977] iteration for the solution of the first order condition of the optimality in KKT. An updated version of this algorithm, a new QP-free method demonstrated by Qi and Qi [2000], emphasizes the feasibility of all iterates. It reformulates the KKT optimality condition Fischer–Burmeister function [Jiang, 1999] for nonlinear complementarity problems. The classical solution is still based on the Lagrange multipliers. Bertsekas [2014] proposed that for one-sided inequality constrained problems, it cannot be converted to equality constrained problem. Therefore, it introduced the method minimizing the augmented Lagrangian with respect to x for various value of the Lagrange parameters, which is presented by Powell [1969] and Hestenes [1969]:

$$\begin{aligned}\bar{L}_c(x, z, \lambda, \mu) = & f(x) + \lambda' h(x) + \frac{1}{2} c |h(x)|^2 \\ & + \sum_{j=1}^r \left\{ \mu_j [g_j(x) + z_j^2] + \frac{1}{2} c |g_j(x) + z_j^2|^2 \right\}\end{aligned}$$

The minimization of this augmented Lagrangian can be found through computing the first order derivative with respect to z explicitly for each fixed x .

More recently, Gould et al. [2016] introduced a method approximating the gradient of the inequality constrained problem based on ideas from interior-point methods [Boyd et al., 2004]. It gives a demonstration of log-barrier function, which transforms the original constrained problem into a unconstrained minimization problem.

$$\phi(x, y) = \sum_{i=1}^m \log(-c_i(x, y)) \quad (2.5)$$

$$\text{minimize}_y \quad t f(x, y) - \sum_{i=1}^m \log(-c_i(x, y)) \quad (2.6)$$

Equation 2.5 is the log-barrier function, which takes the sum of the logarithm for all constraints. Then subtracting it in the unconstrained minimization problem approximates the original inequality constrained problem. t in Equation 2.6 is a scaling factor for duality gap control if the solution set is convex.

Similar to the solution for unconstrained and equality constrained problem, minimizing Equation 2.6 is based on the gradient and hessian of the log-barrier function.

2.3 Differentiable Neural Network

If a problem is differentiable, that means the solution of this problem can be back-propagated. In neural network,

2.4 Summary

Summary what you discussed in this chapter, and mention the story in next chapter. Readers should roughly understand what your thesis takes about by only reading words at the beginning and the end (Summary) of each chapter.

Deep Declarative Network

Same as the last chapter, introduce the motivation and the high-level picture to readers, and introduce the sections in this chapter.

3.1 An Overview of Deep Declarative Network

3.1.1 Structure

3.1.2 Declarative Nodes

3.2 Learning

3.3 Back-propagation Through Declarative Nodes

3.3.1 Unconstrained

3.3.2 Equality Constrained

3.3.3 Inequality Constrained

3.4 Examples of Declarative Nodes

3.4.1 Unconstrained

3.4.2 Equality Constrained

3.4.3 Inequality Constrained

3.5 Summary

Same as the last chapter, summary what you discussed in this chapter and be the bridge to next chapter. s

The Future of Declarative Nodes

Same as the last chapter, introduce the motivation and the high-level picture to readers, and introduce the sections in this chapter.

Part II

Deep Declarative Network: Non-regular Solution

An Overview of Regular and Non-regular Solution

5.1 Problems in Regular Deep Declarative Nodes

5.2 Related Work in Non-regular Solution

5.2.1 Overdetermined System

5.2.2 Rand Deficiency

5.2.3 Non-convex Problems

Table 5.1 shows how to include tables and Figure 5.1 shows how to include codes.

Architecture	Pentium 4	Atom D510	i7-2600
Model	P4D 820	Atom D510	Core i7-2600
Technology	90nm	45nm	32nm
Clock	2.8GHz	1.66GHz	3.4GHz
Cores \times SMT	2×2	2×2	4×2
L2 Cache	1MB \times 2	512KB \times 2	256KB \times 4
L3 Cache	none	none	8MB
Memory	1GB DDR2-400	2GB DDR2-800	4GB DDR3-1066

Table 5.1: Processors used in our evaluation.

```
1 int main(void)
2 {
3     printf("Hello_World\n");
4     return 0;
5 }
```

(a)

```
1 void main(String[] args)
2 {
3     System.out.println("Hello_World");
4 }
```

(b)

Figure 5.1: Hello world in Java and C.

Solutions of Non-regular Point

6.1 Overdetermined System

6.1.1 Least-Squared Method

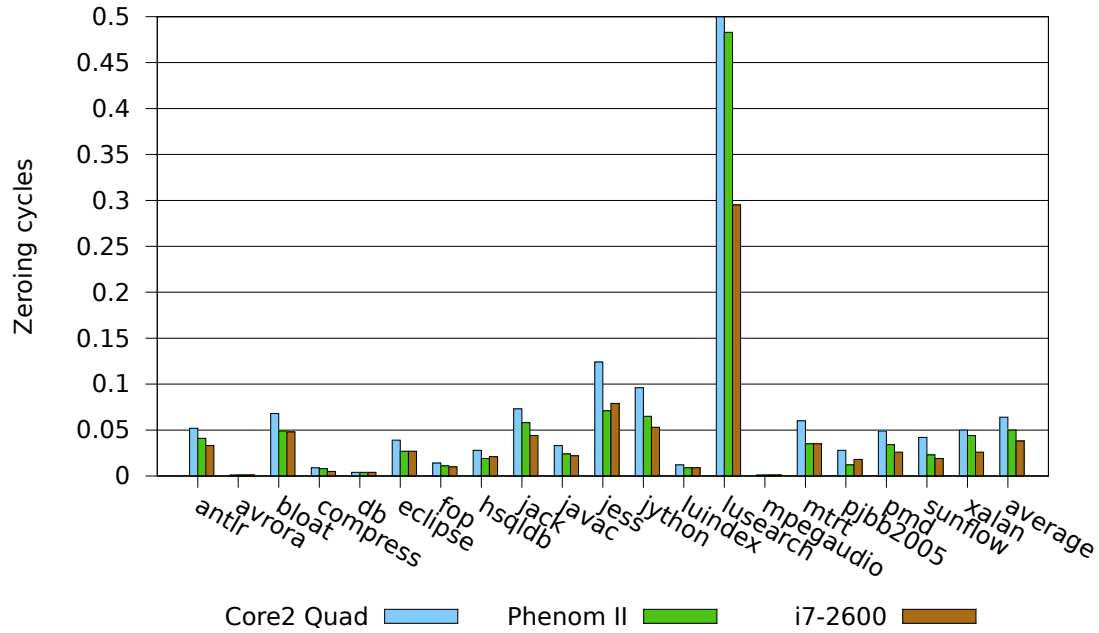
6.2 Conjugate Gradient and Preconditioning

6.3 Rank Deficiency

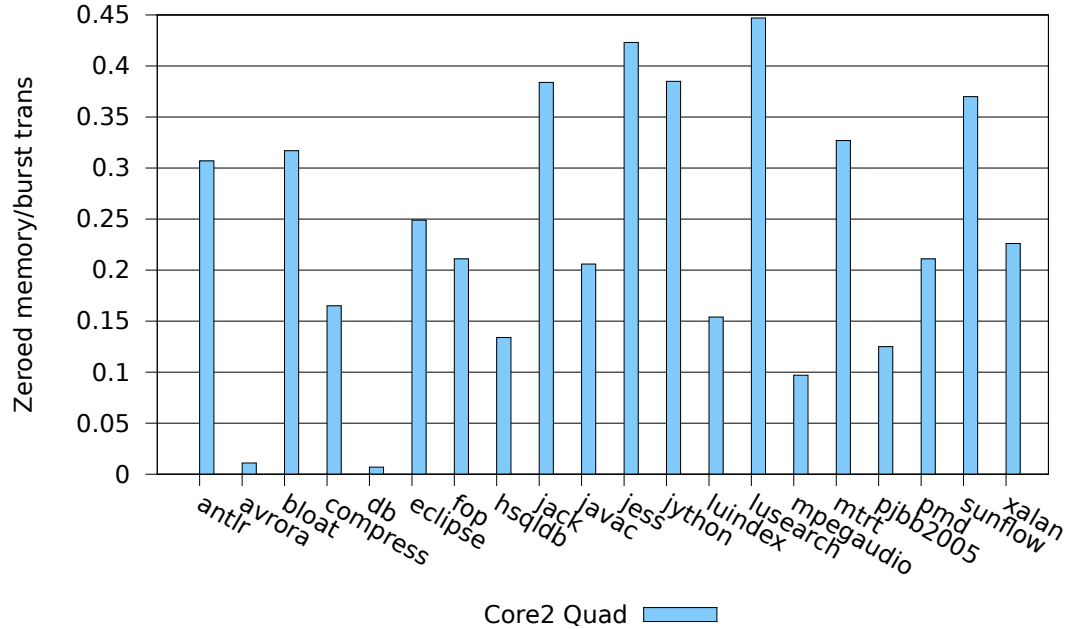
6.4 Non-convex Problems

Here is the example to show how to include a figure. Figure 6.1 includes two subfigures (Figure 6.1(a), and Figure 6.1(b));

6.5 Summary



(a) Fraction of cycles spent on zeroing



(b) BytesZeroed / BytesBurstTransactionsTransferred

Figure 6.1: The cost of zero initialization

Conclusion

Summary your thesis and discuss what you are going to do in the future in Section 7.1.

7.1 Future Work

Good luck.

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Appendix

A An Overview of Numerical Optimization

A.1 Theory of Optimization

A.1.1 Proof of Theorem 2.3

Proof. Let

$$m := \inf\{f(x) : x \in \Omega\}$$

By the definition of m we may pick a sequence $\{x_k\} \subset \Omega$ with $f(x_k) \rightarrow m$ as $k \rightarrow \infty$. Because Ω is compact, we can extract a convergent subsequence $\{x_{k_j}\}$ from $\{x_k\}$. Let $x^* \in \Omega$ denote the limit point of $\{x_{k_j}\}$. Since f is continuous, $f(x^*) = \lim_{j \rightarrow \infty} f(x_{k_j}) = m$. Thus m is finite and x^* is a global minimizer of f on Ω .

When $\Omega = \mathbb{R}^n$, we need to impose conditions on f at infinity to guarantee the existence of a global minimizer. \square

A.1.2 Proof of Theorem 2.4

Proof. Let $m := \inf\{f(x) : x \in \mathbb{R}^n\}$, and take a sequence $\{x_k\}$ such that

$$f(x_k) \rightarrow m \quad \text{as } k \rightarrow \infty.$$

Since f is coercive, $\{x_k\}$ must be bounded; otherwise it has a subsequence $\{x_{k_j}\}$ with $\|x_{k_j}\| \rightarrow \infty$ as $j \rightarrow \infty$, and hence $m = \lim_{j \rightarrow \infty} f(x_{k_j}) = +\infty$, a contradiction.

Thus there is $r > 0$ such that

$$\{x_k\} \subset \{x \in \mathbb{R}^n : \|x\| \leq r\}.$$

Because $\{x \in \mathbb{R}^n : \|x\| \leq r\}$ is compact, $\{x_k\}$ has a convergent subsequence $\{x_{k_j}\}$

with $x_{k_j} \rightarrow x^*$ as $j \rightarrow \infty$. In view of the continuity of f , we have

$$f(x^*) = \lim_{j \rightarrow \infty} f(x_{k_j}) = m$$

Therefore m is finite and f achieves its minimum on \mathbb{R}^n at x^* □

A.1.3 Proof of Theorem 2.6

Proof. We may assume that $\alpha > f_* := \inf \{f(x) : x \in \mathbb{R}^n\}$. Let $\{x_k\}$ be a minimizing sequence for f , i.e.

$$f(x_k) \rightarrow f_* \quad \text{as } k \rightarrow \infty$$

Then there is an N such that $f(x_k) \leq \alpha$ for all $k \geq N$, that is, $x_k \in D$ for all $k \geq N$. Since D is compact, $\{x_k\}_{k=N}^\infty$ has a convergent subsequence $\{x_{k_j}\}$ with $x_{k_j} \rightarrow x_* \in D$ as $j \rightarrow \infty$. In view of the lower semi-continuity of f , we have

$$f(x_*) \leq \lim_{j \rightarrow \infty} f(x_{k_j}) = f_*$$

By the definition of f_* we must have $f(x_*) = f_*$. Therefore f achieves its minimum on \mathbb{R} at x_* . □

A.1.4 Proof of Theorem 2.8

Proof. (NC1): First, recall that for any $v \in \mathbb{R}^n$ there holds

$$v^T \nabla f(x^*) = D_v f(x^*) = \lim_{t \searrow 0} \frac{f(x^* + tv) - f(x^*)}{t}.$$

Since x^* is a local minimizer, we have

$$f(x^* + tv) - f(x^*) \geq 0 \quad \text{for small } |t|.$$

Therefore

$$v^T \nabla f(x^*) \geq 0 \quad \text{for all } v \in \mathbb{R}^n.$$

In particular this implies $(-v)^T \nabla f(x^*) \geq 0$ and thus

$$v^T \nabla f(x^*) \leq 0 \quad \text{for all } v \in \mathbb{R}^n.$$

Therefore $v^T \nabla f(x^*) = 0$ for all $v \in \mathbb{R}^n$. Taking $v = \nabla f(x^*)$ gives $\|\nabla f(x^*)\|^2 = 0$ which shows that $\nabla f(x^*) = 0$ □

Proof. (NC2): Recall that for any $v \in \mathbb{R}^n$ and small $t > 0$ there is $0 < s < 1$ such that

$$f(x^* + tv) = f(x^*) + tv^T \nabla f(x^*) + \frac{1}{2} t^2 v^T \nabla^2 f(x^* + stv) v.$$

Since x^* is a local minimizer of f , we have $f(x^* + tv) \geq f(x^*)$ and $\nabla f(x^*) = 0$ by (NC1). Therefore

$$\frac{1}{2}t^2 v^T \nabla^2 f(x^* + stv) v = f(x^* + tv) - f(x^*) \geq 0.$$

This implies that

$$v^T \nabla^2 f(x^* + stv) v \geq 0.$$

Taking $t \rightarrow 0$ gives

$$v^T \nabla^2 f(x^*) v \geq 0 \quad \text{for all } v \in \mathbb{R}^n$$

i.e. $\nabla^2 f(x^*)$ is semi-definite. □

Proof. (SC1): Since $\nabla^2 f(x)$ is continuous and $\nabla^2 f(x^*) \geq 0$, we can find $r > 0$ such that

$$B_r(x^*) \subset \Omega \quad \text{and} \quad \nabla^2 f(x) > 0 \text{ for all } x \in B_r(x^*).$$

By Taylor's formula we have

$$f(x) = f(x^*) + \nabla f(x^*) \cdot (x - x^*) + \frac{1}{2}(x - x^*)^T \nabla^2 f(\hat{x})(x - x^*)$$

where $\hat{x} := x^* + t(x - x^*)$ for some $0 < t < 1$.

It is clear that $\hat{x} \in B_r(x^*)$ and hence $\nabla^2 f(\hat{x}) > 0$ which implies that

$$(x - x^*)^T \nabla^2 f(\hat{x})(x - x^*) > 0 \quad \text{for } x \neq x^*$$

Consequently

$$f(x) > f(x^*) + \nabla f(x^*) \cdot (x - x^*)$$

for all $x \in B_r(x^*)$ with $x \neq x^*$.

Since $\nabla f(x^*) = 0$, we can obtain $f(x) > f(x^*)$ for all $x \in B_r(x^*)$ with $x \neq x^*$. □

A.1.5 Proof of Lemma 2.10

Proof. For $d \in T_{x^*}\mathcal{F}$, we have $z_k \subset \mathcal{F}$ and t_k such that

$$z_k \rightarrow x^*, \quad 0 < t_k \rightarrow 0 \quad \text{and} \quad \frac{z_k - x^*}{t_k} \rightarrow d$$

as $k \rightarrow \infty$. As $f(x^*) \leq f(z_k)$, by Taylor's formula we have

$$\begin{aligned} f(x^*) &\leq f(z_k) = f(x^* + (z_k - x^*)) \\ &= f(x^*) + (z_k - x^*)^T \nabla f(x^*) + \frac{1}{2}(z_k - x^*)^T \nabla^2 f(\hat{z}_k)(z_k - x^*) \end{aligned}$$

where \hat{z}_k is a point on the line segment joining x^* and z_k . This implies that

$$0 \leq \left(\frac{z_k - x^*}{t_k} \right)^T \nabla f(x^*) + \frac{1}{2} (z_k - x^*)^T \nabla^2 f(\hat{z}_k) \left(\frac{z_k - x^*}{t_k} \right)$$

Letting $k \rightarrow \infty$ gives $d^T \nabla f(x^*) \geq 0$ □

A.2 Solution of Unconstrained and Constrained Optimization Problems

A.2.1 Proof of Equation 2.3

Proof. Firstly, for any optimal y , according to the first-order optimality condition, we have

$$\frac{df(x, y)}{dy} = \mathbf{0} \in \mathbb{R}^{1 \times m}$$

Then from the implicit function theorem, rearranging and differentiating both sides we have

$$\begin{aligned} D\left(\frac{df(x, y)}{dy}\right)^T &= \mathbf{0} \in \mathbb{R}^{m \times n} \\ &= \frac{\partial^2}{\partial x \partial y} f(x, y) + \frac{\partial^2}{\partial y^2} f(x, y) \frac{dy(x)}{dx} \\ \frac{dy(x)}{dx} &= -\left[\left(\frac{\partial^2}{\partial y^2}\right) f(x, y)\right]^{-1} \left(\frac{\partial^2}{\partial x \partial y}\right) f(x, y) \end{aligned}$$

□

A.2.2 Proof of Equation 2.4 [Gould et al., 2019]

Proof. According to the definition of Lagrange multipliers, we can define the Lagrangian:

$$\mathcal{L}(x, y, \lambda) = f(x, y) - \sum_{i=1}^p \lambda_i (A_i y_i - b_i)$$

We are going to find the stationary point (y, λ) for this lagrangian. Therefore, we calculate the derivative of \mathcal{L} with respect to y and λ separately:

$$\frac{\partial}{\partial y} f(x, y) - \sum_{i=1}^p \lambda_i \frac{\partial}{\partial y} (A_i y_i - b_i) = 0 \quad (1)$$

$$A y - b = 0 \quad (2)$$

Since y is the optimal point, we have $\frac{\partial}{\partial y} f(x, y) = 0$, which can be an unconstrained problem or it is orthogonal to the constraint surface. For unconstrained cases, we can

set $\lambda = 0$ directly. For the orthogonal case, from Equation 1, we have

$$\frac{\partial}{\partial y} f(x, y) = \sum_{i=1}^p \lambda_i \frac{\partial}{\partial y} (A_i y_i - b_i) = \lambda^T A$$

Now we are going to calculate the derivative of the Lagrangian with respect to x for both 1 and 2:

$$\frac{\partial^2}{\partial x \partial y} f(x, y) + \frac{\partial^2}{\partial y^2} f(x, y) Dy - \frac{\partial}{\partial y} (A y - b)^T D \lambda = 0 \quad (3)$$

$$\frac{\partial}{\partial x} (A y - b) + \frac{\partial}{\partial y} (A y - b) Dy = 0 \quad (4)$$

Solving 3 and 4, we get:

$$Dy(x) = \left(H^{-1} A^T \left(A H^{-1} A^T \right)^{-1} A H^{-1} - H^{-1} \right) B$$

where

$$H = \frac{\partial^2}{\partial y^2} f(x, y), \quad B = \frac{\partial^2}{\partial x \partial y} f(x, y)$$

□