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CSE DS D1

Exp. 3: Strassen's Multiplication

AIM: Multiplying two matrices using Strassen's Multiplication Method (Divide and Conquer Approach)

Theory:

Let us consider two matrices X and Y. We want to calculate the resultant matrix Z by multiplying X and Y.

General Method

First, we will discuss naïve method and its complexity. Here, we are calculating $Z = X \times Y$. Using Naïve method, two matrices (X and Y) can be multiplied if the order of these matrices are $p \times q$ and $q \times r$. Following is the algorithm.

Algorithm: Matrix-Multiplication (X, Y, Z)

for $i = 1$ to p do

for $j = 1$ to r do

$Z[i,j] := 0$

for $k = 1$ to q do

$Z[i,j] := Z[i,j] + X[i,k] \times Y[k,j]$

Complexity

Here, we assume that integer operations take $O(1)$ time. There are three for loops in this algorithm and one is nested in other. Hence, the algorithm takes $O(n^3)$ time to execute.

Strassen's Matrix Multiplication Algorithm

In this context, using Strassen's Matrix multiplication algorithm, the time consumption can be improved a little bit.

Strassen's Matrix multiplication can be performed only on square matrices where n is a power of 2. Order of both of the matrices are $n \times n$.

Divide X , Y and Z into four $(n/2) \times (n/2)$ matrices as represented below –

$$Z = \begin{bmatrix} I & K & J & L \end{bmatrix}$$

$$X = \begin{bmatrix} A & C & B & D \end{bmatrix} \text{ and } Y = \begin{bmatrix} E & G & F & H \end{bmatrix}$$

Using Strassen's Algorithm compute the following –

$$M_1 = (A + C) \times (E + F)$$

$$M_2 = (B + D) \times (G + H)$$

$$M_3 = (A - D) \times (E + H)$$

$$M_4 = A \times (F - H)$$

$$M_5 = (C + D) \times (E)$$

$$M_6 = (A + B) \times (H)$$

$$M_7 = D \times (G - E)$$

Then,

$$I = M_2 + M_3 - M_6 - M_7$$

$$J = M_4 + M_6$$

$$K = M_5 + M_7$$

$$L = M_1 - M_3 - M_4 - M_5$$

Analysis

$$T(n) = \begin{cases} c_7 \times T(n^2) + d \times n^2 & \text{if } n=1 \\ \text{otherwise} \end{cases}$$

where c and d are constants

Using this recurrence relation, we get $T(n) = O(n \log^7)$

Hence, the complexity of Strassen's matrix multiplication algorithm is

$$O(n \log^7)$$

PROGRAM:

```
#include<stdio.h>
```

```
int
```

```
main ()
```

```
{
```

```
int a[2][2], b[2][2], c[2][2], i, j;
```

```
int m1, m2, m3, m4, m5, m6, m7;
```

```
printf ("\nEnter the first matrix\n");
```

```
for (i = 0; i < 2; i++)
```

```
{
```

```
for (j = 0; j < 2; j++)
```

```
{
```

```
printf ("Enter the element %d%d: ", i+1, j+1);
```

```
scanf ("%d", &a[i][j]);
```

```
}
```

```
}
```

```
printf ("\nEnter the second matrix\n");
```

```
for (i = 0; i < 2; i++)
```

```
{
```

```
for (j = 0; j < 2; j++)
```

```
{
```

```
printf ("Enter the element %d%d: ", i+1, j+1);
```

```
scanf ("%d", &b[i][j]);
```

```
}
```

```
}
```

```
printf ("\nThe first matrix is\n\n");  
for (i = 0; i < 2; i++)  
{  
    printf ("\t");  
    for (j = 0; j < 2; j++)  
    {  
        printf ("%d\t", a[i][j]);  
    }  
    printf ("\n");  
}
```

```
printf ("\nThe second matrix is\n\n");  
for (i = 0; i < 2; i++)  
{  
    printf ("\t");  
    for (j = 0; j < 2; j++)  
    {  
        printf ("%d\t", b[i][j]);  
    }  
    printf ("\n");  
}
```

$m1 = (a[0][0] + a[1][1]) * (b[0][0] + b[1][1]);$

$m2 = (a[1][0] + a[1][1]) * b[0][0];$

$m3 = a[0][0] * (b[0][1] - b[1][1]);$

$m4 = a[1][1] * (b[1][0] - b[0][0]);$

$m5 = (a[0][0] + a[0][1]) * b[1][1];$

$m6 = (a[1][0] - a[0][0]) * (b[0][0] + b[0][1]);$

```
m7 = (a[0][1] - a[1][1]) * (b[1][0] + b[1][1]);
```

```
c[0][0] = m1 + m4 - m5 + m7;
```

```
c[0][1] = m3 + m5;
```

```
c[1][0] = m2 + m4;
```

```
c[1][1] = m1 - m2 + m3 + m6;
```

```
printf ("\nAfter multiplication using Strassen's algorithm \n\n");
```

```
for (i = 0; i < 2; i++)
```

```
{
```

```
    printf ("\t");
```

```
    for (j = 0; j < 2; j++)
```

```
    {
```

```
        printf ("%d\t", c[i][j]);
```

```
    }
```

```
    printf ("\n");
```

```
}
```

```
return 0;
```

```
}
```

RESULT:

```
Enter the first matrix
Enter the element 11: 3
Enter the element 12: 4
Enter the element 21: 5
Enter the element 22: 6

Enter the second matrix
Enter the element 11: 3
Enter the element 12: 4
Enter the element 21: 5
Enter the element 22: 6

The first matrix is

    3    4
    5    6

The second matrix is

    3    4
    5    6

After multiplication using Strassen's algorithm

    29    36
    45    56
```

OBSERVATIONS:

Strassen's Multiplication is better than the Normal method of matrix multiplication as its time complexity comes to $n^{2.81}$ compared to the general method which has time complexity of n^3 .

We use only 7 multiplications in Strassen's Multiplication whereas we use 8 multiplications in the general method, so the time complexity is better.

While usual multiplication requires 8 multiplications and 4 additions Strassen's Multiplication requires 7 multiplications and 18 additions.

CONCLUSION:

I implemented Strassen's Matrix Multiplication method in this experiment and solved user-given matrices to find their product.