

An overview of the TurboFan compiler

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Overview

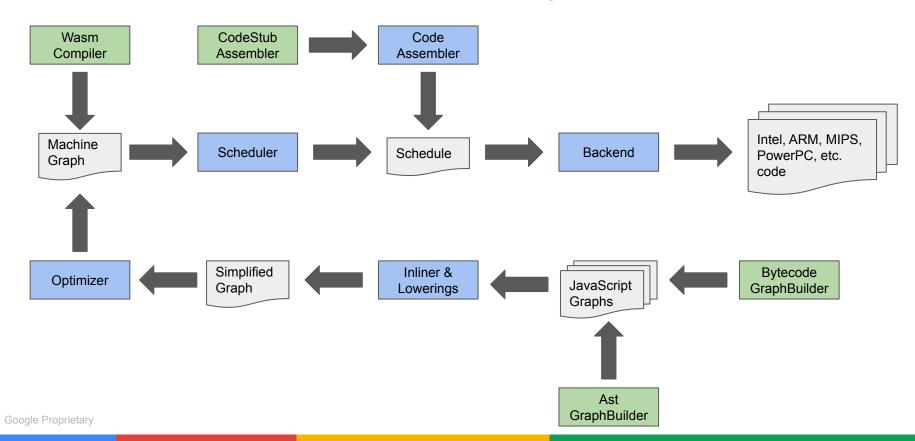
What is TurboFan?

- the new JavaScript compiler to replace the aging Crankshaft
- a high-level, portable assembler

Why a portable assembler?

- Interpreter bytecode handlers
- Code stubs and builtins
- asm.js & WebAssembly

The TurboFan architecture / entry points



TurboFan as a portable assembler

Motivation

Traditionally V8 contains a lot of hand-written native code:

- Fast paths for EcmaScript builtins (i.e. Object.prototype.toString)
- IC handlers and dispatchers
- Code stubs (i.e. FastNewClosureStub)

Native code on the web efforts:

- asm.js
- WebAssembly

Goals

- Provide a uniform code generation architecture
- Reduce porting / maintenance overhead of V8 (currently already 10 ports!)
- Remove performance cliffs due to slow builtins
- Make experimenting with new features easier (i.e. changes to load/store ICs, bootstrapping an interpreter)

CodeAssembler vs. CodeStubAssembler

- CodeAssembler as shim to the TurboFan machine-level IR; wraps RawMachineAssembler, which generates a Schedule with machine instructions
- CodeStubAssembler adds high-level V8/JavaScript-related macros on-top
- InterpreterAssembler adds Ignition-related macros on-top of CodeStubAssembler
- CodeAssembler is portable assembler entry point for TurboFan

Benefits

- Fewer performance cliffs
- Better baseline performance (real world benefit, i.e. ForInFilter, Object.prototype.toString, Function.prototype.bind, etc.)
- Fewer bugs (no more manual register allocation, etc.)
- Reduce porting overhead
- Refactoring viable

TurboFan as JavaScript compiler

Motivation

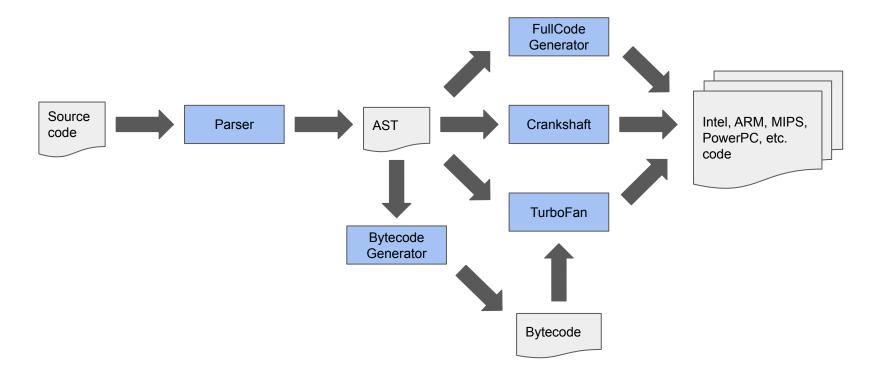
Crankshaft served us well, but has various shortcomings:

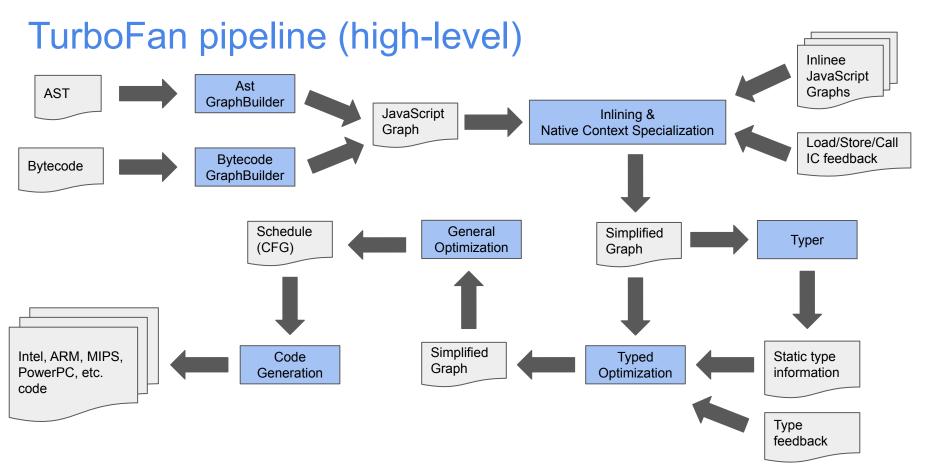
- Doesn't scale to full, modern JavaScript (try-catch, for-of, generators, async/await, ...)
- Defaults to deoptimization (performance cliffs, deoptimization loops)
- Graph construction, inlining and optimization all mixed up
- Tight coupling to fullcodegen / brittle environment tracking
- Limited optimization potential / limited static analysis (i.e. type propagation)
- High porting overhead
- Mixed low-level and high-level semantics of instructions

Goals

- Full ESnext language support (try-catch/-finally, class literals, eval, generators, async functions, modules, destructuring, etc.)
- Utilize and propagate (static) type information
- Separate graph building from optimization / inlining
- No deoptimization loops / deoptimization only when really beneficial
- Sane environment tracking (also for lazy deoptimization)
- Predictable peak performance

V8 compilation overview





Graph Building

- Turns Bytecode or AST into graph with JavaScript nodes
 - JSAdd, JSCallFunction, JSLoadProperty, etc.
 - o Branch, IfTrue, IfFalse, etc.
- Tracks environment for baseline code and inserts appropriate FrameStates
 - CheckPoint with FrameState for eager deoptimization
 - FrameState after nodes that can lead to arbitrary JS execution (i.e. everything except JSTypeOf, JSStrictEqual and JSToBoolean) for lazy deoptimization
- Eliminates dead FrameState uses, i.e. values that baseline cannot see are replaced with optimized_out sentinel
- Optional OSR deconstruction right afterwards to remove the non-OSR entry

Native Context Specialization & Inlining

- Specialize global object property access to the global object for native context
- Specialize to native context based on Load/Store/Call IC feedback
 - JSLoadNamed[x](o) turns into fast access for x on o if baseline collected feedback
 - Specialize function calls to direct calls (i.e. to known JSFunctions) when Call IC provides feedback
- Inline through JSCallFunction and JSCallConstruct sites, even <u>polymorphic</u> <u>sites</u>
 - Supports <u>sane heuristics</u> (not only first come, first serve)
 - Can even inline through Function.prototype.call/.apply
- Eliminates dead nodes in the (control) graph on-the-fly
- Doesn't take into account type feedback for binary/compare operations

Typer & Typed Optimization

- Typer assigns types to value producing nodes, based on JS typing rules
 - Already considers types for fields known from native context specialization
 - Propagates type information through phis and sigmas (loop exits)
 - Range analysis and constant folding (some) included
- Lowering JS graph to simplified graph based on types
 - Take into account static type information first, but also utilize type feedback (binary/compare)
 - Inline builtin calls based on parameter types (i.e. Math.abs, Array.prototype.push, etc.)
- Lower JS object creation nodes to inline allocations
 - JSCreate becomes an inline allocation when new.target is known
 - Similar for JSCreateArray, JSCreateLiteralObject, etc.
 - Makes those eligible for escape analysis and load/store/check elimination

General Optimizations

- Loop peeling
 - No more deoptimization loops because of aggressive hoisting
- Load/check elimination
 - Kills redundant loads and checks
 - Includes global value numbering pass
- Escape analysis
 - Eliminate non escaping allocations
 - Scalar replacement of aggregates
 - Not yet partial; allocation sinking planned
- Representation selection
 - Truncation analysis
 - Type feedback propagation

General Optimizations

- Effect/control linearization
 - Wires potentially deoptimizing nodes into effect/control chain, assigns check points
 - Expands macro operations with internal control flow (i.e. ChangeTaggedToFloat64)
 - Optimizes branches that depend on phis (branch cloning)
- Redundant store elimination
- Control flow optimization (turns certain branch chains into switches)
- Allocation folding and write barrier elimination
 - Fragmentation-free folding of allocations
 - Even across branches; not yet across merges
- Late optimization pass
 - Another global value numbering pass
 - Strength reduction, dead code elimination, redundant branch elimination

Code Generation

- Sophisticated instruction selection
 - Pattern matching on subgraphs for often optimal cover
 - Architecture-dependent matchings
 - Dead instruction elimination.
- Register allocation
 - Essentially improved version of the lithium allocator
- Jump threading
 - Remove jumps between blocks
- Code emission

Code generation example

Generate code for add

```
function add(x, y) { return x + y; }
Run with --turbo --ignition-staging
add(1,2);
add(2,3);
%OptimizeFunctionOnNextCall(add);
add(3,4);
```

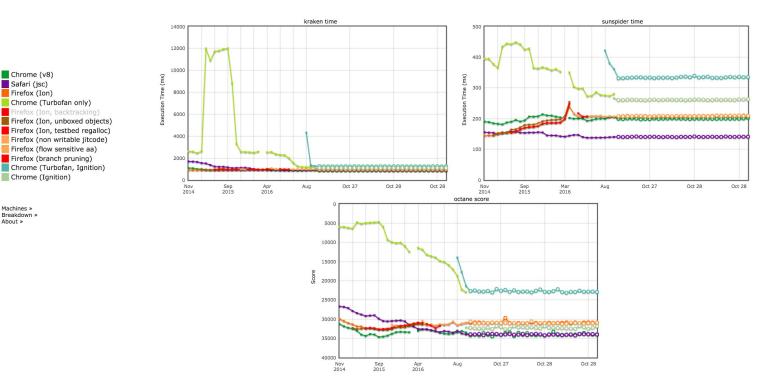
Almost optimal code sequence already (Smi representation missing)

Code generation example

| | B0 start (construct frame) | | 0x3caafc1040b0 | 80 | 48bbd012abb3397f0000 REX.W movq rbx,0x7f39b3ab12d0 | | x7f39b3ab12d0 |
|----------------|--|------------------------------|----------------|-----|--|------------------------|-------------------|
| 0x3caafc104060 | 0 55 | push rbp | 0x3caafc1040ba | 90 | 33c0 | xorl rax, rax | |
| 0x3caafc104061 | 1 4889e5 | REX.W movq rbp,rsp | 0x3caafc1040bc | 92 | 488b75f8 | REX.W movq rsi,[rbp-0x | 8] |
| 0x3caafc104064 | 4 56 | push rsi | | | real.js:229:12 | | |
| 0x3caafc104065 | 5 57 | push rdi | 0x3caafc1040c0 | 96 | e87b02f0ff | call 0x3caafc004340 | ;; code: STUB |
| 0x3caafc104066 | 6 493ba5600c0000 | 0 REX.W cmpq rsp,[r13+0xc60] | 0x3caafc1040c5 | 101 | ebac | jmp 19 (0x3caafc10407 | 3) |
| 0x3caafc10406d | 13 0f863d000000 | jna 80 (0x3caafc1040b0) | 0x3caafc1040c7 | 103 | e834ffd7ff | call 0x3caafbe84000 | ;; deoptimization |
| | B2 start | | bailout 0 | | | | |
| | B3 start (deconstruct frame) | | 0x3caafc1040cc | 108 | e839ffd7ff | call 0x3caafbe8400a | ;; deoptimization |
| 0x3caafc104073 | 19 488b4518 | REX.W movq rax,[rbp+0x18] | bailout 1 | | | | |
| 0x3caafc104077 | 23 a801 | test al,0x1 | 0x3caafc1040d1 | 113 | e83effd7ff | call 0x3caafbe84014 | ;; deoptimization |
| 0x3caafc104079 | 25 0f8548000000 | jnz 103 (0x3caafc1040c7) | bailout 2 | | | | |
| 0x3caafc10407f | 31 488b5d10 | REX.W movq rbx,[rbp+0x10] | 0x3caafc1040d6 | 118 | 90 | nop | |
| 0x3caafc104083 | 35 f6c301 | testb rbx,0x1 | 0x3caafc1040d7 | 119 | 90 | nop | |
| 0x3caafc104086 | 38 0f8540000000 | jnz 108 (0x3caafc1040cc) | 0x3caafc1040d8 | 120 | 90 | nop | |
| 0x3caafc10408c | 44 488bd3 | REX.W movq rdx,rbx | 0x3caafc1040d9 | 121 | 90 | nop | |
| 0x3caafc10408f | 47 48c1ea20 | REX.W shrq rdx, 32 | 0x3caafc1040da | 122 | 90 | nop | |
| 0x3caafc104093 | 51 488bc8 | REX.W movq rcx, rax | 0x3caafc1040db | 123 | 90 | nop | |
| 0x3caafc104096 | 54 48cle920 | REX.W shrq rcx, 32 | 0x3caafc1040dc | 124 | 90 | nop | |
| 0x3caafc10409a | 58 03d1 | addl rdx,rcx | 0x3caafc1040dd | 125 | 90 | nop | |
| 0x3caafc10409c | 60 0f802f000000 | jo 113 (0x3caafc1040d1) | 0x3caafc1040de | 126 | 90 | nop | |
| 0x3caafc1040a2 | 66 48c1e220 | REX.W shlq rdx, 32 | 0x3caafc1040df | 127 | 90 | nop | |
| 0x3caafc1040a6 | 70 488bc2 | REX.W movq rax, rdx | 0x3caafc1040e0 | 128 | 90 | nop | |
| 0x3caafc1040a9 | 73 488be5 | REX.W movq rsp,rbp | 0x3caafc1040e1 | 129 | 90 | nop | |
| 0x3caafc1040ac | 76 5d | pop rbp | 0x3caafc1040e2 | 130 | 90 | nop | |
| 0x3caafc1040ad | 77 c21800 | ret 0x18 | 0x3caafc1040e3 | 131 | 90 | nop | |
| | B4 start (no frame) ;;; Safepoint table. | | | | | | |
| | B1 start (deferred) | | | | | | |

Where do we stand?

Peak-performance benchmarks



What's coming?

Future

- Retire AstGraphBuilder (asm.js only soon)
 - WebAssembly based asm validator coming
- Optimize ESnext features via Ignition+TurboFan
- Add missing optimizations to TurboFan (i.e. arguments object, Smi representation, etc.)
- Eventually route more (traditional) JavaScript through Ignition+TurboFan only

Upcoming talks

- <u>TurboFan IR overview</u> (<u>jarin@chromium.orq</u>)
 - Different IR layers (JavaScript, simplified, machine)
 - Sea-of-nodes overview
- TurboFan deoptimization support (<u>jarin@chromium.org</u>)
 - Check point chain
 - Lazy bailout
 - Deoptimizer support
 - Liveness analysis

Questions?