# A Polyhedral Study of Triplet Formulation for Single Row Facility Layout Problem

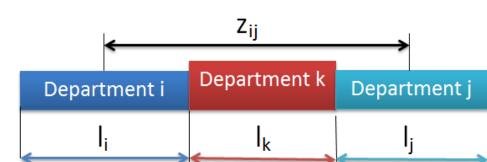
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## SINGLE ROW FACILITY LAYOUT PROBLEM

Single Row Facility Layout Problem (SRFLP): linear arrangement problem with the objective of minimizing the total weighted sum of distances between department pairs.



- $ightharpoonup l_i$  Length of department i
- $ightharpoonup c_{ij}$  Average daily traffic between departments i and j
- $ightharpoonup z_{ii}^{\pi}$  Distance between centroids of departments i and j in permutation  $\pi$

Objective of SRFLP: 
$$\min_{\pi} \sum_{i=1}^{n-1} \sum_{j=i+1}^{n} c_{ij} z_{ij}^{\pi}$$

- ► Amaral's formulation of SRFLP triplet formulation [1].
- ► Two projections of triplet formulation and its valid inequalities presented in [1].
- Linear program solved over valid inequalities in [1] yields optimal solution for several classical instances of sizes n = 5 to n = 30.
- ► This suggests that the valid inequalities are quite strong.

## RESEARCH CONTRIBUTIONS

- ▶ Dimension of the triplet polytope is n'' = n(n-1)(n-2)/3.
- ► Almost all valid inequalities defined for the triplet polytope by Amaral in [1] are facet-defining.
- ▶ The above results are also true for the other projections of the triplet polytope defined in [1].

### TRIPLET POLYTOPE

► Decision variable:

$$\lambda_{ijk} = \begin{cases} 1 & \text{if department } k \text{ lies between departments } i \text{ and } j, i < j \\ 0 & \text{otherwise.} \end{cases}$$

- ► Set of departments:  $N = \{1, ..., n\}$
- ▶ Decision variable vector:  $\lambda = \{\lambda_{ijk} : i, j, k \in N, i < j\}$
- ▶ Number of elements of  $\lambda$ : n' = n(n-1)(n-2)/2
- $ightharpoonup P^1 = \{\lambda \in \{0,1\}^{n'}: \lambda \text{ represents a permutation of } \{1,...,n\}\}$
- ▶ Triplet polytope: convex hull of  $P^1$ .

Objective function of SRFLP: 
$$\min \sum_{i=1}^{n-1} \sum_{j=i+1}^{n} c_{ij} \left( \frac{1}{2} (l_i + l_j) + \sum_{k \neq i, k \neq j}^{n} l_k \lambda_{ijk} \right)$$

Valid inequalities for SRFLP presented in [1]:

$$0 \leq \lambda_{ijk} \leq 1 \qquad i, j, k \in \mathbb{N}, i < j \tag{1}$$

$$\lambda_{ijk} + \lambda_{ikj} + \lambda_{jki} = 1 \qquad i, j, k \in \mathbb{N}, i < j < k$$

$$-\lambda_{ijd} + \lambda_{jkd} + \lambda_{ikd} \geqslant 0 \qquad i, j, k, d \in N, i < j < k \tag{}$$

$$\lambda_{ijd} + \lambda_{jkd} - \lambda_{ikd} \ge 0$$
  $i, j, k, d \in N, i < j < k$   
 $\lambda_{ijd} - \lambda_{jkd} + \lambda_{ikd} \ge 0$   $i, j, k, d \in N, i < j < k$ 

$$\lambda_{ijd} - \lambda_{jkd} + \lambda_{ikd} \ge 0 \qquad i, j, k, d \in N, i < j < k$$

$$\lambda_{ijd} + \lambda_{jkd} + \lambda_{ikd} \le 2 \qquad i, j, k, d \in N, i < j < k$$

$$(5)$$

For a positive even integer  $\beta \le n$ , let  $S = \{i_t : t = 1, ..., \beta\} \subseteq N$  be a set of departments and  $d \in S$ . Let  $(S_1, S_2)$  be a partition of  $S \setminus \{d\}$  such that  $|S_1| = \beta/2$ . Then the following inequality is valid for  $conv(P^1)$  [1]:

$$\sum_{p,q \in S_1: p < q} \lambda_{pqd} + \sum_{p,q \in S_2: p < q} \lambda_{pqd} \leq \sum_{p \in S_h, q \in S_{\{1,2\} \setminus h}: h = 1, 2, p < q} \lambda_{pqd}$$
 (7)

#### PRELIMINARY RESULTS

#### Notation:

- $ightharpoonup \Pi_N$  set of all permutations of departments in N.
- $\triangleright \lambda^{\pi}$  vector that represents a permutation  $\pi \in \Pi_N$ .
- ▶ To simplify notation, wherever we have  $\lambda_{ijk}$  with i > j, we mean  $\lambda_{ijk}$ .

**Lemma 1.** For given distinct departments  $x, y, z \in N$ , let  $\pi^1, \pi^2, \pi^3, \pi^4$  be four permutations of the departments in N satisfying the following conditions:

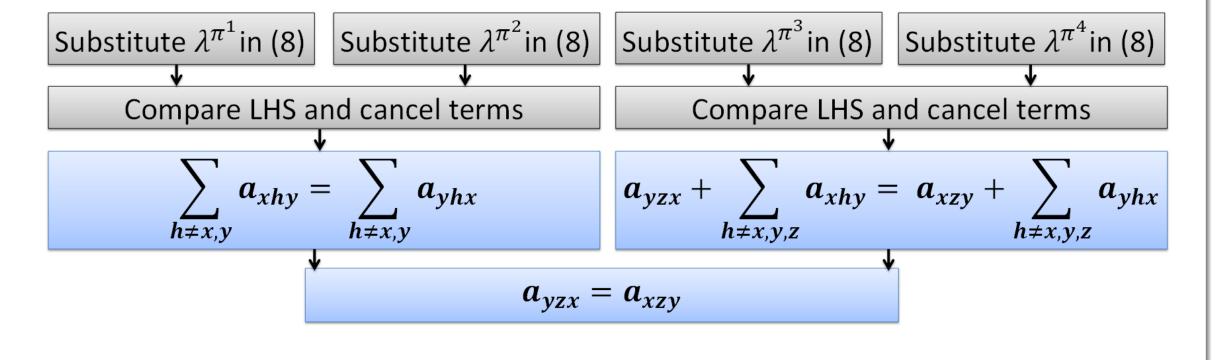
- $\blacktriangleright \pi^1$ : First two departments are x,y
- $\blacktriangleright \pi^2$ : First two departments are y,x, other departments are in the same order as  $\pi^1$
- $\blacktriangleright \pi^3$ : First three departments are z,x,y
- $\blacktriangleright \pi^4$ : First three departments are z,y,x, other departments are in the same order as  $\pi^3$

If the  $\lambda$  vectors corresponding to these permutations lie on the hyperplane

$$\sum_{i,j,k \in N: i < j} a_{ijk} \lambda_{ijk} = b, \tag{}$$

then  $a_{yzx} = a_{xzy}$ .

## Outline of proof:



## **Dimension of** $conv(P^1)$

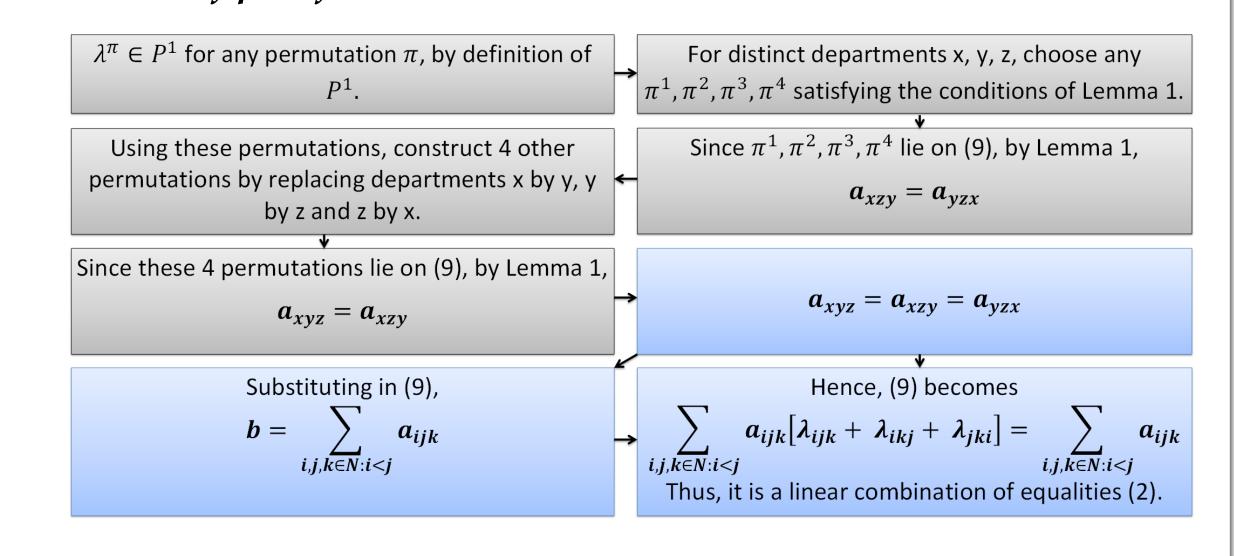
**Theorem 2.**  $conv(P^1)$  is of dimension n'' = n(n-1)(n-2)/3.

 $conv(P^1) \subset \mathbb{R}^{n'}$  and any  $\lambda \in P^1$  satisfies the set of  $\binom{n}{3}$  linearly independent equalities (2). Hence, dim  $\left(conv(P^1)\right) \leqslant n' - \binom{n}{3} = n''$ . To prove that the dimension is actually equal to n'', we just need to show that any other hyperplane like

$$\sum_{i,j,k \in N: i < j} a_{ijk} \lambda_{ijk} = b \tag{9}$$

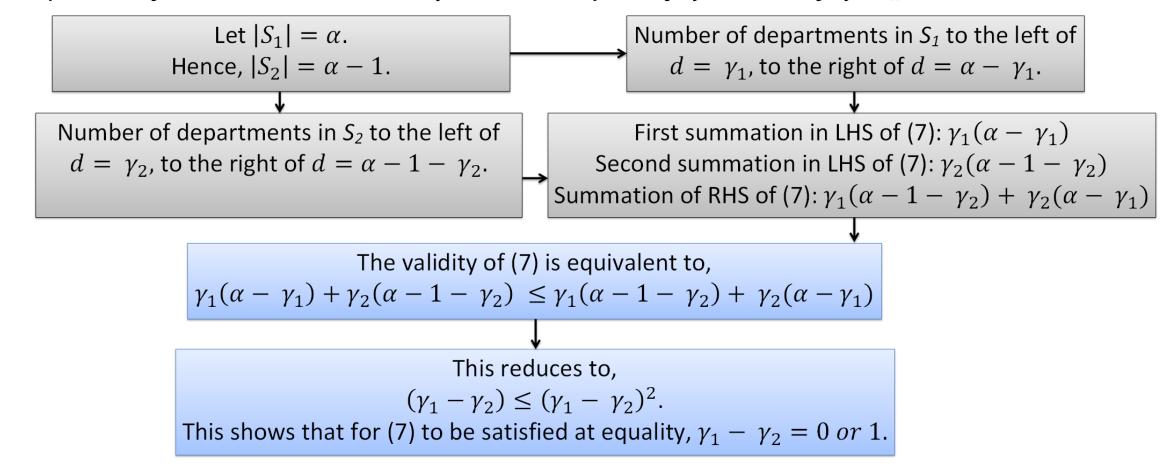
satisfied by all  $\lambda \in P^1$  will be a linear combination of the equalities (2).

#### Outline of proof:



### FACET-DEFINING PROPERTY OF VALID INEQUALITIES

**Lemma 3.** Consider inequality (7) for given  $\beta$ , S,  $S_1$ ,  $S_2$  and d. Let  $\pi \in \Pi_N$ , and  $\gamma_1$  and  $\gamma_2$  be the number of departments in  $S_1$  and  $S_2$  which are to the left of d in  $\pi$ , respectively. Then  $\lambda^{\pi} \in P^1$  satisfies (7) at equality if and only if  $\gamma_1 - \gamma_2 = 0$  or 1.



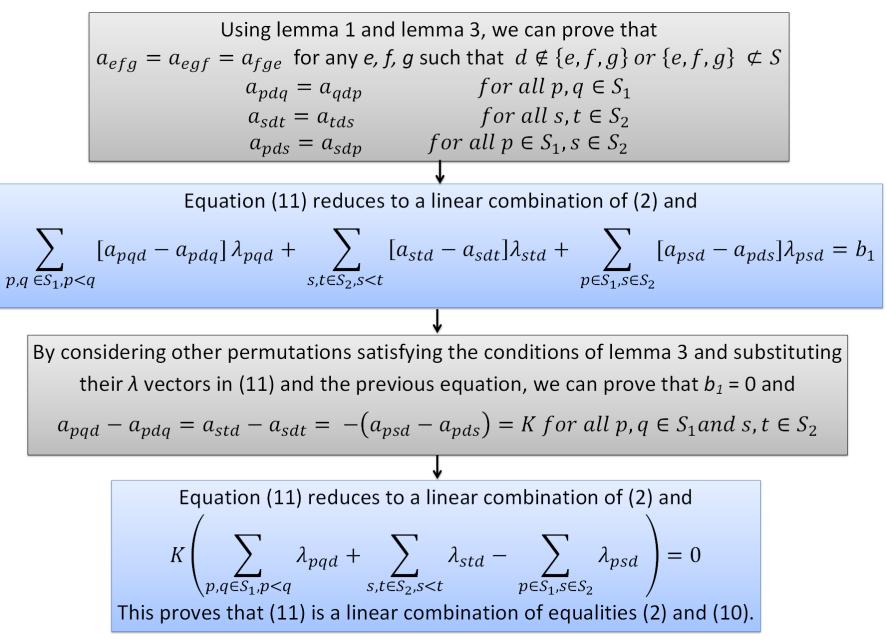
**Theorem 4.** Any of inequalities (7) is facet-defining for  $conv(P^1)$ . Let P' be the face of  $conv(P^1)$  defined by an inequality of the form (7) for given  $\beta$ , S, S<sub>1</sub>, S<sub>2</sub> and d. Hence, for every point in P', (7) is satisfied at equality, i.e.

$$\sum_{p,q \in S_1: p < q} \lambda_{pqd} + \sum_{p,q \in S_2: p < q} \lambda_{pqd} - \sum_{p \in S_1, q \in S_2} \lambda_{pqd} = 0.$$
 (10)

To prove that (10) is a facet, we need to show that any hyperplane like

$$\sum_{e,f,g \in N: e < f} a_{efg} \lambda_{efg} = b \tag{}$$

that passes through P' is a linear combination of hyperplanes (2) and hyperplane (10).



- ▶ (3), (4) and (5) special cases of (7) for  $\beta = 4$ , so they are facet-defining.
- ▶ (1) and (6) not facet-defining in general can be checked numerically.
- ➤ We have also shown that the above results are true for other projections of the triplet polytope in [1] using a simple result related to projection of polyhedra [2].

## Concluding Remarks

Our results provide theoretical support for the fact that the LP solution over these valid inequalities gives the optimal solution for all instances studied in [1].

#### Major References

- 1. Amaral, A. R. S. 2009. A new lower bound for the single row facility layout problem. *Discrete Applied Mathematics* **157(1)**, 183–190.
- 2. Sanjeevi, S. and Kianfar, K. 2010. A polyhedral study of triplet formulation for single row facility layout problem. *Discrete Applied Mathematics*, **in press**, DOI: 10.1016/j.dam.2010.07.005.