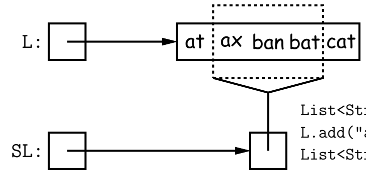
CS61B Lecture #18: Assorted Topics

- Views
- Maps
- More partial implementations
- Array vs. linked: tradeoffs
- Sentinels
- Specialized sequences: stacks, queues, deques
- Circular buffering
- Recursion and stacks
- Adapters

Views

New Concept: A *view* is an alternative presentation of (interface to) an existing object.

- For example, the `sublist` method is supposed to yield a "view of" part of an existing list:




```
List<String> L = new ArrayList<String>();
L.add("at"); L.add("ax"); ...
List<String> SL = L.sublist(1,4);
```

- Example: after `L.set(2, "bag")`, value of `SL.get(1)` is `"bag"`, and after `SL.set(1,"bad")`, value of `L.get(2)` is `"bad"`.
- Example: after `SL.clear()`, `L` will contain only `"at"` and `"cat"`.
- Small challenge: "How do they *do* that?!"

Maps

- A `Map` is a kind of "modifiable function:"

```
package java.util;
public interface Map<Key,Value> {
  Value get(Object key);             // Value at KEY.
  Object put(Key key, Value value);  // Set get(KEY) -> VALUE
  ...
}
-----------------------------------------------------
Map<String,String> f = new TreeMap<String,String>();
f.put("Paul", "George"); f.put("George", "Martin");
f.put("Dana", "John");
// Now f.get("Paul").equals("George")
//     f.get("Dana").equals("John")
//     f.get("Tom") == null
```


Map Views

```
public interface Map<Key,Value> { // Continuation

           /* Views of Maps */

  /** The set of all keys. */
  Set<Key> keySet();

  /** The multiset of all values that can be returned by get.
   *  (A multiset is a collection that may have duplicates). */
  Collection<Value> values();

  /** The set of all(key, value) pairs */
  Set<Map.Entry<Key,Value>> entrySet();
}
```


View Examples

Using example from a previous slide:

```
Map<String,String> f = new TreeMap<String,String>();
f.put("Paul", "George"); f.put("George", "Martin");
f.put("Dana", "John");
```

we can take various views of `f`:

```
for (Iterator<String> i = f.keySet().iterator(); i.hasNext();)
   i.next()  ===>  Dana, George, Paul
// or, more succinctly:
for (String name : f.keySet())
   name  ===>  Dana, George, Paul

for (String parent : f.values())
   parent  ===>   John, Martin, George

for (Map.Entry<String,String> pair : f.entrySet())
   pair   ===>   (Dana,John), (George,Martin), (Paul,George)

f.keySet().remove("Dana");   // Now f.get("Dana") == null
```


Simple Banking I: Accounts

Problem: Want a simple banking system. Can look up accounts by name or number, deposit or withdraw, print.

Account Structure

```
class Account {
  Account(String name, String number, int init) {
     this.name = name; this.number = number;
     this.balance = init;
  }
  /** Account-holder's name */
  final String name;
  /** Account number */
  final String number;
  /** Current balance */
  int balance;

  /** Print THIS on STR in some useful format. */
  void print(PrintStream str) { ... }
}
```

Simple Banking II: Banks

```
class Bank {
  /* These variables maintain mappings of String -> Account.  They keep
   * the set of keys (Strings) in "compareTo" order, and the set of
   * values (Accounts) is ordered according to the corresponding keys. */
  SortedMap<String,Account> accounts = new TreeMap<String,Account>();
  SortedMap<String,Account> names = new TreeMap<String,Account>();

  void openAccount(String name, int initBalance) {
     Account acc =
       new Account(name, chooseNumber(), initBalance);
     accounts.put(acc.number, acc);
     names.put(name, acc);
  }

  void deposit(String number, int amount) {
    Account acc = accounts.get(number);
    if (acc == null) ERROR(...);
    acc.balance += amount;
  }
  // Likewise for withdraw.
```


Banks (continued): Iterating

Printing out Account Data

```
/** Print out all accounts sorted by number on STR. */
void printByAccount(PrintStream str) {
   // accounts.values() is the set of mapped-to values.  Its
   // iterator produces elements in order of the corresponding keys.
   for (Account account : accounts.values())
     account.print(str);
}

/** Print out all bank accounts sorted by name on STR. */
void printByName(PrintStream str) {
   for (Account account : names.values())
     account.print(str);
}
```

A Design Question: What would be an appropriate representation for keeping a record of all transactions (deposits and withdrawals) against each account?

Partial Implementations

- Besides interfaces (like `List`) and concrete types (like `LinkedList`), Java library provides abstract classes such as `AbstractList`.
- Idea is to take advantage of the fact that operations are related to each other.
- Example: once you know how to do `get(k)` and `size()` for an implementation of `List`, you can implement all the other methods needed for a *read-only* list (and its iterators).
- Now throw in `add(k,x)` and you have all you need for the additional operations of a growable list.
- Add `set(k,x)` and `remove(k)` and you can implement everything else.

Example: The java.util.AbstractList helper class

```
public abstract class AbstractList<Item> implements List<Item>
{
   /** Inherited from List */
   // public abstract int size();
   // public abstract Item get(int k);
   public boolean contains(Object x) {
      for (int i = 0; i < size(); i += 1) {
        if ((x == null && get(i) == null) ||
            (x != null && x.equals(get(i))))
          return true;
      }
      return false;
   }
   /* OPTIONAL: Throws exception; override to do more. */
   void add(int k, Item x) {
     throw new UnsupportedOperationException();
   }
   Likewise for remove, set
```


Example, continued: AListIterator

```
// Continuing abstract class AbstractList<Item>:
public Iterator<Item> iterator() { return listIterator(); }
public ListIterator<Item> listIterator() {
   return new AListIterator(this);
}

private static class AListIterator implements ListIterator<Item> {
   AbstractList<Item> myList;
   AListIterator(AbstractList<Item> L) { myList = L; }
   /** Current position in our list. */
   int where = 0;

   public boolean hasNext() { return where < myList.size(); }
   public Item next() { where += 1; return myList.get(where-1); }
   public void add(Item x) { myList.add(where, x); where += 1; }
   ... previous, remove, set, etc.
}
...
```


Aside: Another way to do AListIterator

It's also possible to make the nested class non-static:

```
public Iterator<Item> iterator() { return listIterator(); }
public ListIterator<Item> listIterator() { return this.new AListIterator(); }

private class AListIterator implements ListIterator<Item> {
  /** Current position in our list. */
  int where = 0;

  public boolean hasNext() { return where < AbstractList.this.size(); }
  public Item next() { where += 1; return AbstractList.this.get(where-1); }
  public void add(Item x) { AbstractList.this.add(where, x); where += 1; }
  ... previous, remove, set, etc.
}
...
```

- Here, `AbstractList.this` means "the `AbstractList` I am attached to" and X`.new AListIterator` means "create a new `AListIterator` that is attached to X."
- In this case you can abbreviate `this.new` as `new` and can leave off some `AbstractList.this` parts, since meaning is unambiguous.

Example: Using AbstractList

Problem: Want to create a *reversed view* of an existing `List` (same elements in reverse order). Operations on the original list affect the view, and vice-versa.

```
public class ReverseList<Item> extends AbstractList<Item> {
  private final List<Item> L;

  public ReverseList(List<Item> L) { this.L = L; }

  public int size() { return L.size(); }

  public Item get(int k) { return L.get(L.size()-k-1); }

  public void add(int k, Item x) { L.add(L.size()-k, x); }

  public Item set(int k, Item x) { return L.set(L.size()-k-1, x); }

  public Item remove(int k) { return L.remove(L.size() - k - 1); }
}
```


Getting a View: Sublists

Problem: `L.sublist(start, end)` is a `List` that gives a view of part of an existing list. Changes in one must affect the other. How?

```
// Continuation of class AbstractList.  Error checks not shown.
List<Item> sublist(int start, int end) {
  return this.new Sublist(start, end);
}

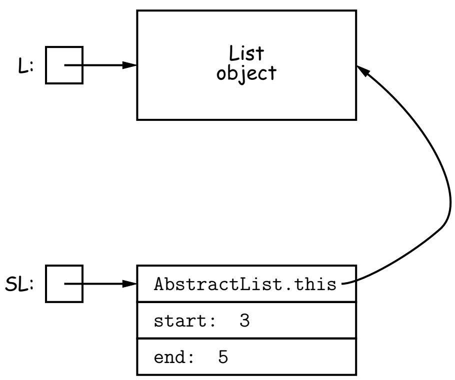
private class Sublist extends AbstractList<Item> {
  private int start, end;
  Sublist(int start, int end) { obvious }

  public int size() { return end-start; }
  public Item get(int k) { return AbstractList.this.get(start+k); }

  public void add(int k, Item x)
    { AbstractList.this.add(start+k, x); end += 1; }
  ...
}
```


What Does a Sublist Look Like?

- Consider `SL = L.sublist(3, 5);`



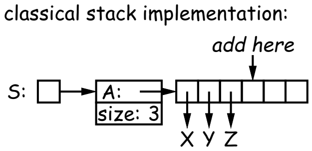
Arrays and Links

- Two main ways to represent a sequence: array and linked list
- In Java Library: `ArrayList` and `Vector` vs. `LinkedList`.
- Array:
  - Advantages: compact, fast ($\Theta(1)$) *random access* (indexing).
  - Disadvantages: insertion, deletion can be slow ($\Theta(N)$)
- Linked list:
  - Advantages: insertion, deletion fast once position found.
  - Disadvantages: space (link overhead), random access slow.

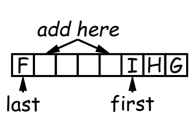
Implementing with Arrays

- Biggest problem using arrays is insertion/deletion in the *middle* of a list (must shove things over).
- Adding/deleting from ends can be made fast:
  - Double array size to grow; amortized cost constant (Lecture #15).
  - Growth at one end really easy; classical stack implementation:

```
S.push("X");
S.push("Y");
S.push("Z");
```



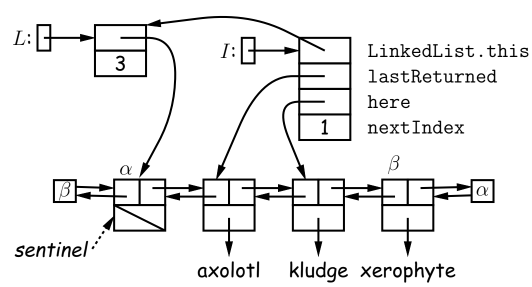
  - To allow growth at either end, use *circular buffering:*



  - Random access still fast.

Linking

- Essentials of linking should now be familiar
- Used in Java `LinkedList`. One possible representation for linked list and an iterator object over it:

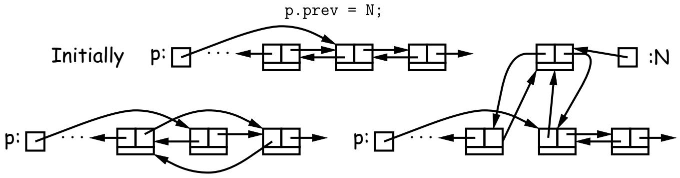



```
L = new LinkedList<String>();
L.add("axolotl");
L.add("kludge");
L.add("xerophyte");
```

```
I = L.listIterator();
I.next();
```

Clever trick: Sentinels

- A *sentinel* is a dummy object containing no useful data except links.
- Used to eliminate special cases and to provide a fixed object to point to in order to access a data structure.
- Avoids special cases ('**if**' statements) by ensuring that the first and last item of a list always have (non-null) nodes—possibly sentinels—before and after them:
-
```
// To delete list node at p:
p.next.prev = p.prev;
p.prev.next = p.next;
```
```
// To add new node N before p:
N.prev = p.prev; N.next = p;
p.prev.next = N;
p.prev = N;
```



Specialization

- Traditional special cases of general list:
  - **Stack**: Add and delete from one end (LIFO).
  - **Queue**: Add at end, delete from front (FIFO).
  - **Dequeue**: Add or delete at either end.
- All of these easily representable by either array (with circular buffering for queue or deque) or linked list.
- Java has the `List` types, which can act like any of these (although with non-traditional names for some of the operations).
- Also has `java.util.Stack`, a subtype of `List`, which gives traditional names ("push", "pop") to its operations. There is, however, no "stack" interface.

Stacks and Recursion

- Stacks related to recursion. In fact, can convert any recursive algorithm to stack-based (however, generally no great performance benefit):
  - Calls become "push current variables and parameters, set parameters to new values, and loop."
  - Return becomes "pop to restore variables and parameters."

```
findExit(start):
  if isExit(start)
    FOUND
  else if (!isCrumb(start))
    leave crumb at start;
    for each square, x,
      adjacent to start:
        if legal(start,x) && !isCrumb(x)
          findExit(x)
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Call: findExit((0,0))
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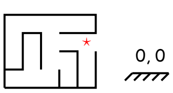

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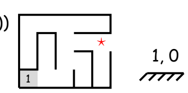

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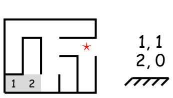

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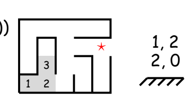

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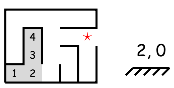

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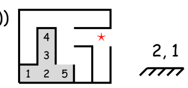

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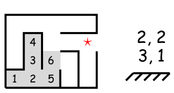

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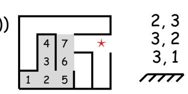

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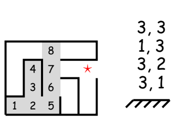

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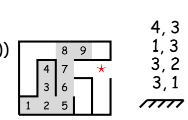

Stacks and Recursion

- Stacks related to recursion. In fact, can convert any recursive algorithm to stack-based (however, generally no great performance benefit):
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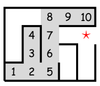

1, 3
3, 2
3, 1

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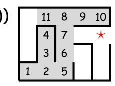

0, 3
3, 2
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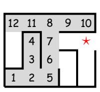

0, 2
3, 2
3, 1

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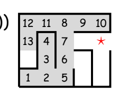

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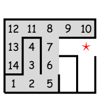

3, 2
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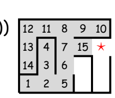

4, 2
3, 1

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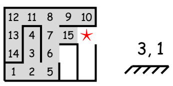

Design Choices: Extension, Delegation, Adaptation

- The standard `java.util.Stack` type *extends* `Vector`:

```
class Stack<Item> extends Vector<Item> { void push(Item x) { add(x); } ... }
```

- Could instead have *delegated* to a field:

```
class ArrayStack<Item> {
   private ArrayList<Item> repl = new ArrayList<Item>();
   void push(Item x) { repl.add(x); } ...
}
```

- Or, could generalize, and define an *adapter:* a class used to make objects of one kind behave as another:

```
public class StackAdapter<Item> {
   private List repl;
   /** A stack that uses REPL for its storage. */
   public StackAdapter(List<Item> repl) { this.repl = repl; }
   public void push(Item x) { repl.add(x); } ...
}

class ArrayStack<Item> extends StackAdapter<Item> {
  ArrayStack() { super(new ArrayList<Item>()); }
}
```