## **Artificial Intelligence 1**

# **Assignment7: Propositional Logic**

- Given Dec 8, Due Dec 18 -

### **Problem 7.1 (PL Concepts)**

30 pt

Which of the following statements are true? In each case, give an informal argument why it is true or a counter-example.

- 1. Every satisfiable formula is valid.
- 2. Every valid formula is satisfiable.
- 3. If *A* is satisfiable, then  $\neg A$  is unsatisfiable.
- 4. If  $A \models B$ , then  $A \land C \models B \land C$ .
- 5. Every admissible inference rule is derivable.
- 6. If  $\vdash$  is sound for  $\models$  and  $\{A, B\} \vdash C$ , then C is satisfiable if A and B are.

#### Problem 7.2 (Equivalence of CSP and SAT)

30 pt

We consider

- CSPs (V, D, C) with finite domains as before
- SAT problems (V, A) where V is a set of propositional variables and A is a propositional formula over V.

We will show that these problem classes are equivalent by reducing their instances to each other.

- 1. Given a SAT instance P = (V, A), define a CSP instance P' = (V', D', C') and two bijections
  - f mapping satisfying assignments of P to solutions of P'
  - f' the inverse of f

We already know that binary CSPs are equivalent to higher-order CSPs. Therefore, it is sufficient to give a higher-order CSP.

2. Given a CSP instance (V, D, C), define a SAT instance (V', A') and bijections as above

### Problem 7.3 (Calculi Comparison)

60 pt

Prove (or disprove) the validity of the following formulae in i) Natural Deduction ii) Tableau and iii) Resolution.

- 1.  $(P \land Q) \Rightarrow (P \lor Q)$  (to be done in the tutorial, not part of grading)
- 2.  $((A \lor B) \land (A \Rightarrow C) \land (B \Rightarrow C)) \Rightarrow C$
- 3.  $((P \Rightarrow Q) \Rightarrow P) \Rightarrow P$