## Problem 7.2 Equivalence of CSP and SAT

## Answer:

1. According to the problem description,

SAT instance P = (V, A), where V is a set of propositional variables and A is propositional formula CSP instance P' = (V', D', C')

## SAT to CSP:

- Variables in CSP are the propositional variables of SAT, thus V' = V
- Domains in CSP are the universe of SAT, thus  $D' \in \{T, F\}$
- Constraints in CSP are the propositional formula of SAT, thus C' = A
- Assignments in CSP are the model of SAT  $M := \langle U, I \rangle$ , where U is universe and I interpretation function
- And two bijections are as follows,
  - Solution of P' will be  $f: I_{\varphi}(A) = T$  for some assignment  $\varphi$
  - $f': I_{\varphi}(A) = F$  for some assignment  $\varphi$
- 2. According to the problem description,

CSP instance P = (V, D, C) and SAT instance P' = (V', A')

## CSP to SAT:

- Propositional variables in SAT are the variables in CSP, thus V' = V
- Propositional formula in SAT is the constrains in CSP, thus A' = C
- And two bijections are as follows,
  - Solution of P' will be the consistent total assignment of CSP, thus  $f\colon a, iff\ \forall \mathcal{C}_{uv}\in\mathcal{C}, \left(a(u),a(v)\right)\in\mathcal{C}_{uv} \text{ where } a \text{ is a consistent total variable assignment}$
  - $f': a, iff \ \forall C_{uv} \in C, (a(u), a(v)) \notin C_{uv}$  where a is an inconsistent variable assignment