

AI Assignment

AI 1 - Assignment OG

[Constraint Propagation]

Problem 6.1 \Rightarrow Scheduling CS Classes as a CSP

01. Since, there will be one var. per class,

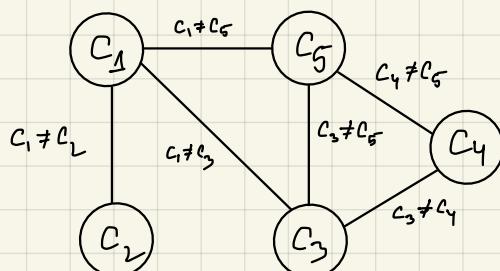
Time	Class	Var	Domain	$f_{\text{var}} \Rightarrow$ each classes
8:30 - 9:30	Class 1	C_1	A	
8:00 - 9:00	Class 2	C_2	A, C	Domain \Rightarrow professors for that class
9:00 - 10:00	Class 3	C_3	A, B, C	
9:30 - 10:30	Class 4	C_4	A, B, C	Constraints \Rightarrow $\sum C_2 = \frac{5!}{2!(5-2)!} = \frac{5!}{2!3!}$
9:00 - 10:00	Class 5	C_5	A, C	$= \frac{5 \times 4 \times 3!}{2 \times 3!} = 10$

Possible Constraints:

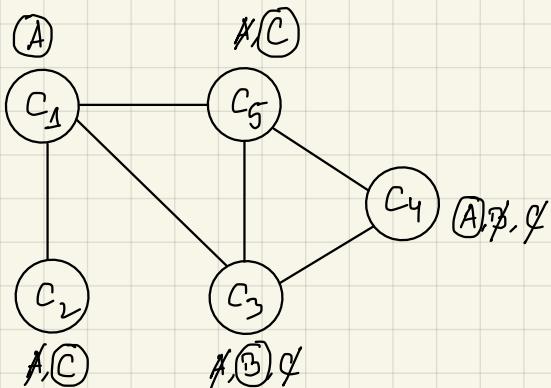
$C_1 \neq C_2$ ✓	$C_2 \neq C_3$ ✗	$C_3 \neq C_4$ ✓	$C_4 \neq C_5$ ✓
$C_1 \neq C_3$ ✗	$C_2 \neq C_4$ ✗	$C_3 \neq C_5$ ✗	
$C_1 \neq C_4$ ✗	$C_2 \neq C_5$ ✗		
$C_1 \neq C_5$ ✓			

a	b	c	d	e
x	x	x	x	/
ab	bc	cd	de	
ac	bd	ce		
ad	be			
ae				

02.



03.



Solⁿ:

$$C_1 = A$$

$$C_2 = C$$

$$C_3 = B$$

$$C_4 = A$$

$$C_5 = C$$

a total inconsistent assignment:

$$C_1 = A$$

$$C_2 = C$$

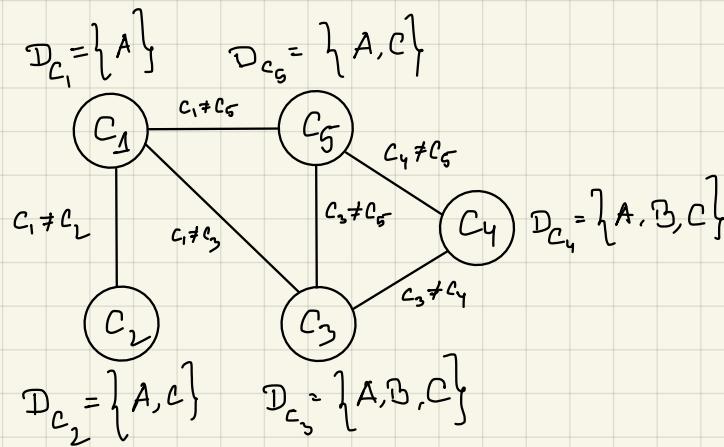
$$C_3 = B$$

$$C_4 = A$$

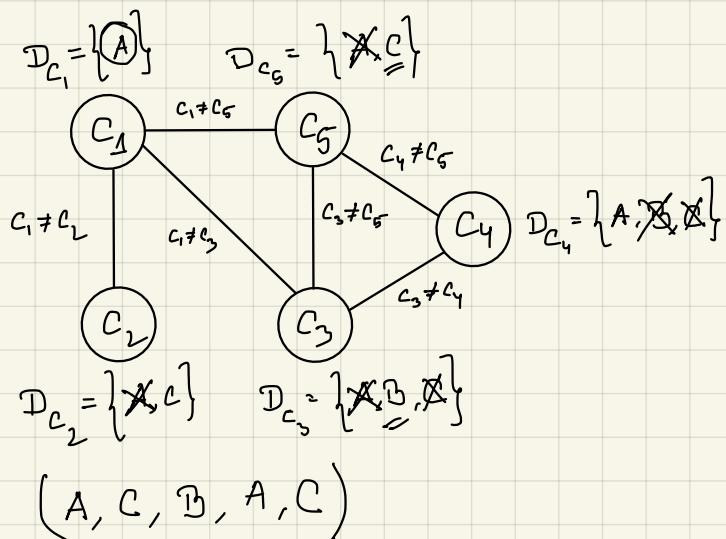
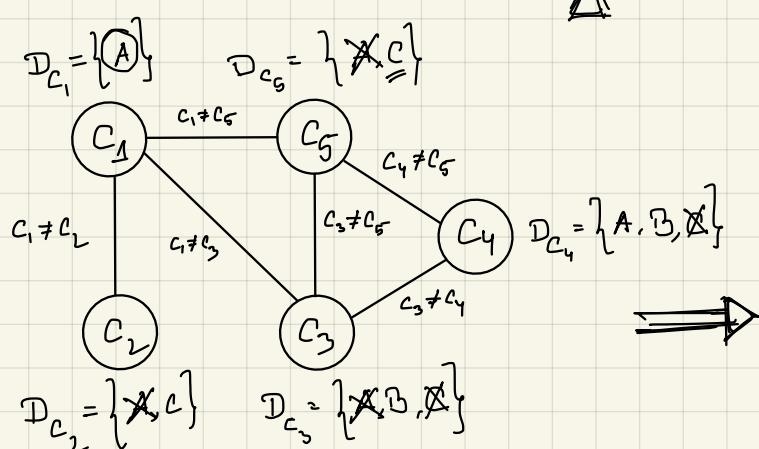
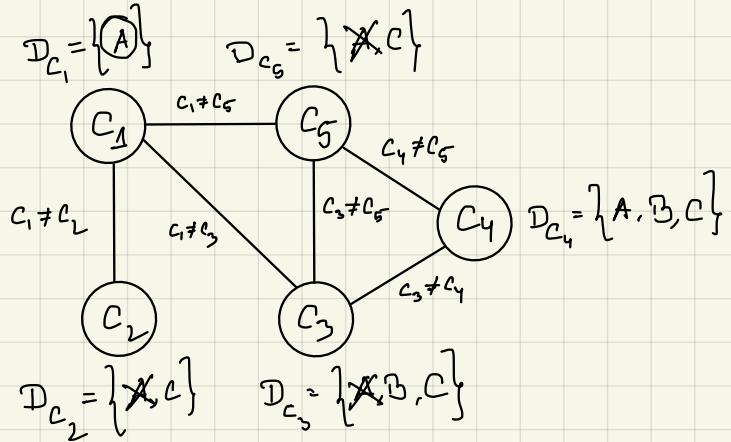
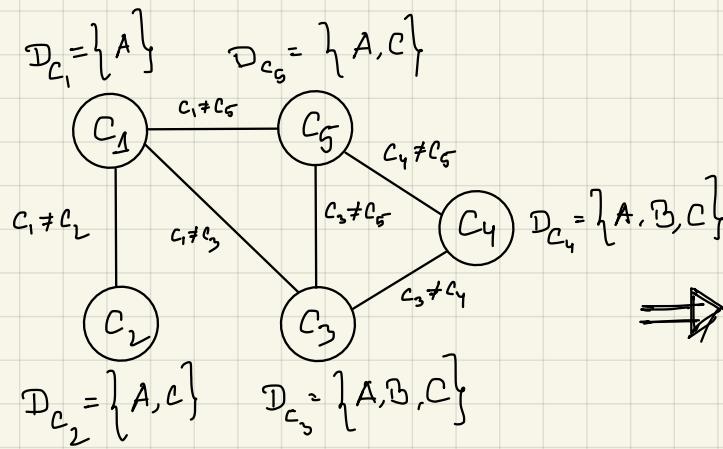
$$C_5 = A \quad X$$

Problem 6.2 \Rightarrow Scheduling CS classes with constraint propagation

CSP graph:



* running arc consistency \Rightarrow

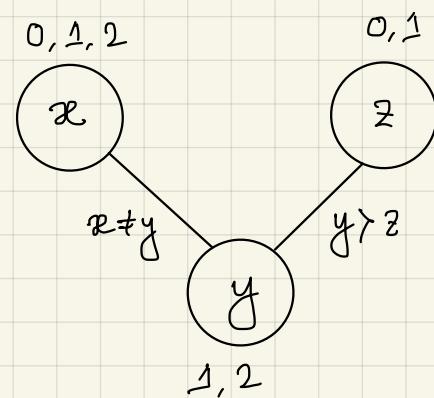


Problem 6.3 \Rightarrow CSP Formalization

Binary CSP $\Pi := (\mathcal{V}, \mathcal{D}, \mathcal{C})$

	\mathcal{D}
x	$D_x = \{0, 1, 2\}$
y	$D_y = \{1, 2\}$
z	$D_z = \{0, 1\}$

$C: x \neq y, y > z$



* Forward checking : Updating domain

(heuristics)

Choose var. $\xrightarrow{\text{MCV}}$ most constrained var.

One-step lookahead

[that has fewest consistent values]



Choose val $\xrightarrow{\text{LCV}}$ least constrained value

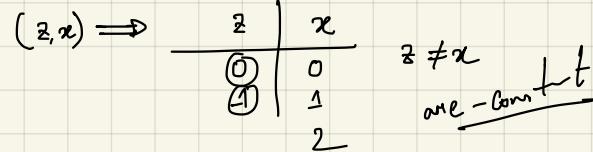
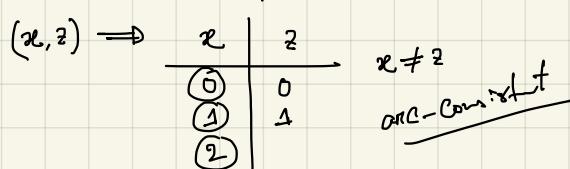
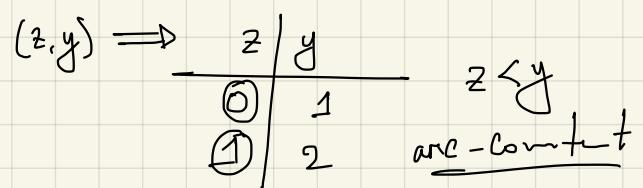
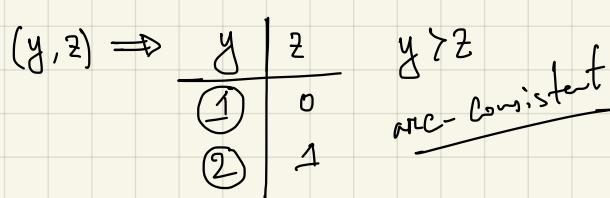
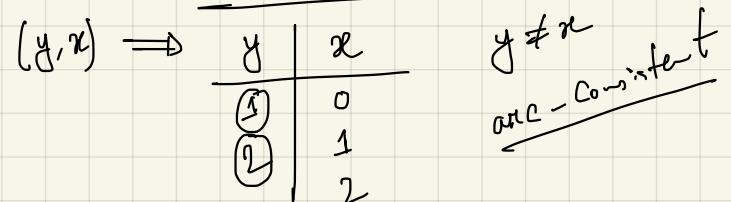
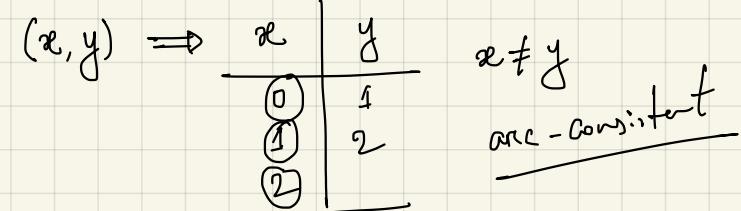
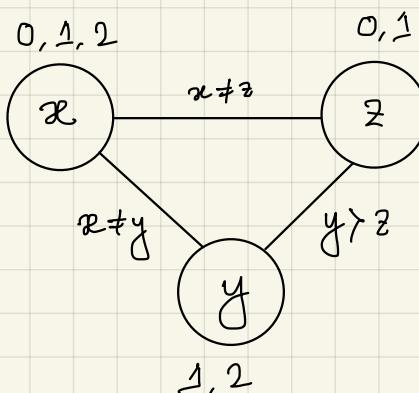
(AC-3)

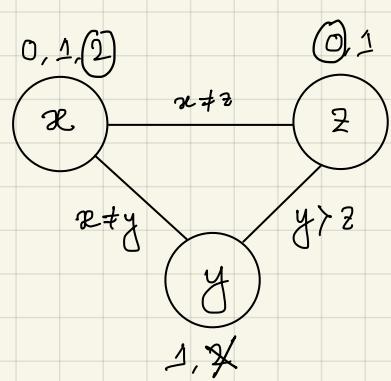
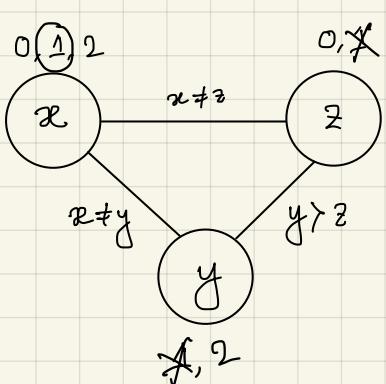
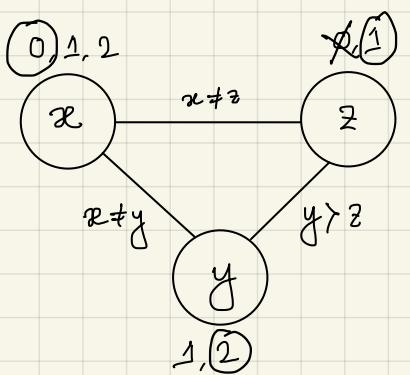
[order values of selected x_i (var.) by decreasing number of consistent values of neighboring var.]

* Arc consistency : Updating domain (Not just neighbors, ~~the whole~~ constraints graph)

$x_i \xrightarrow[\text{consistent}]{\text{arc}} x_j$: remove values from D_i

01. $(f, w) \Rightarrow f$ is arc-consistent w.r.t w





02.

x	y	z	$x \neq y$	$x \neq z$	$y > z$
0	2	1	✓	✓	✓
1	2	0	✓	✓	✓
2	1	0	✓	✓	✓

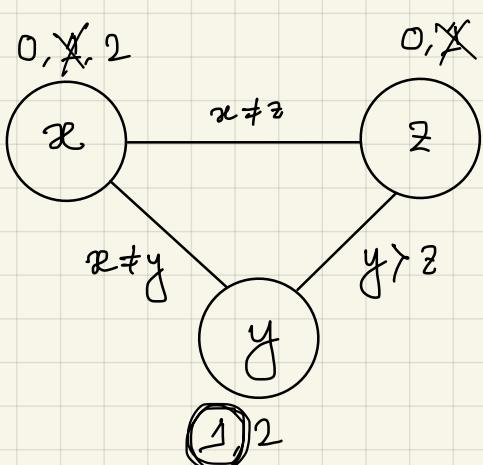
(0, 2, 1)

(1, 2, 0)

(2, 1, 0)



03.



Assign, $y = 1$, then

$$D_x = \{0, 2\}$$

$$D_y = \{1\}$$

$$D_z = \{0\}$$

AI 1 - Assignment 07 (Propositional Logic)

Propositional symbol (atomic formulas): A, B, C

Logical connectives: \neg , \wedge , \vee , \Rightarrow , \Leftrightarrow

$\neg f$, $f \wedge g$, $f \vee g$, $f \Rightarrow g$, $f \Leftrightarrow g$

Syntax	Semantics
Valid Express.	Meaning of Express.

* Model: model ω in propositional logic \Rightarrow an assignment of True val. to propositional symbols

A, B, C: $2^3 = 8$ possible ω

$$\mapsto \{(0,0,0), (0,0,1), \dots, (1,1,1)\}$$

* Interpretation functⁿ:

$I(f, \omega)$ $\begin{cases} \text{returning True (1)} \\ \text{returning False (0)} \end{cases}$

True (1) \Rightarrow if ω satisfies f

False (0) \Rightarrow if ω does not satisfy f

f	g	$\neg f$	$f \wedge g$	$f \vee g$	$f \Rightarrow g$	$f \Leftrightarrow g$
0	0	1	0	0	1	1
0	1	1	0	1	1	0
1	0	0	0	1	0	0
1	1	0	1	1	1	1

$$I(f, \omega) \in \{0, 1\}$$

Conditional proposition

* Models: $M(f)$ — set of models ω for which $I(f, \omega) = 1$

Knowledge Base: set of formulas

$$M(KB) = \bigcap_{f \in KB} M(f)$$

* SAT \Rightarrow checking satisfiability in PL

Model checking

propositional symbol — variable
 formula — constraints
 model — assignment

CSP

Assignment 11

Problem 11.1

according to problem description,

- concept : student, course, woman
- roles : has-registered, taught-by, is-student

ALC conversion :

01. $\text{student} \sqcap \exists \text{has-registered}. \text{course}$
02. $(\text{student} \sqcap \forall \text{has-registered}. \text{course}) \sqsubseteq \forall \text{taught-by}. \text{woman}$
03. $\forall \text{has-registered}. T \sqsubseteq \forall \text{is-student}. T$

ALC formula into FOL :

01. $\forall x (\text{student}(x) \wedge \exists y (\text{has-registered}(x,y) \wedge \text{course}(y)))$
02. $\forall x (\text{student}(x) \wedge \forall y (\text{has-registered}(x,y) \rightarrow \text{course}(y)))$
 $\Rightarrow \forall z (\text{taught-by}(x,z) \rightarrow \text{woman}(z))$
03. $\forall x \forall y \text{ has-registered}(x,y) \rightarrow \text{is-student}(x)$

Problem 11.2 (ALC Semantics)

01. ALC Concepts $\forall R. (C \sqcap D) \Leftrightarrow (\forall R.C) \sqcap (\forall R.D)$

$$\begin{aligned} [[\forall R. (C \sqcap D)]] &= \left\{ x \in D \mid \forall y. \text{ if } \langle xy \rangle \in [[R]] \text{ then } y \in [[C \sqcap D]] \right\} \\ &= \left\{ x \in D \mid \forall y. \text{ if } \langle xy \rangle \in [[R]] \text{ then } (y \in [[C]]) \wedge y \in [[D]] \right\} \\ &= \left\{ x \in D \mid \forall y. \text{ if } \langle xy \rangle \in [[R]] \text{ then } y \in [[C]] \right\} \wedge \\ &\quad \left\{ x \in D \mid \forall y. \text{ if } \langle xy \rangle \in [[R]] \text{ then } y \in [[D]] \right\} \\ &= [[\forall R.C]] \wedge [[\forall R.D]] \end{aligned}$$

02. FOL conversion:

$$(i) \forall R. (C \sqcap D) = \forall y R(x,y) \rightarrow (C(y) \wedge D(y))$$

$$(ii) \forall R.C \sqcap \forall R.D = (\forall y R(x,y) \rightarrow C(y)) \wedge (\forall y R(x,y) \rightarrow D(y))$$

$$\forall y R(x,y) \rightarrow (C(y) \wedge D(y)) \Leftrightarrow (\forall y R(x,y) \rightarrow C(y)) \wedge (\forall y R(x,y) \rightarrow D(y))$$

Since, we need to show that two formulas are equivalent, then
Natural Deduction calculus will be best suited for this case.

Problem 11.3 (ALC TBox)

- primitive concepts: woman, man
- rules: has-child, has-parent, has-sibling, has-spouse

Concepts are defined using ALC TBox as follows,

person = woman \sqcup man

parent = person \sqcap \exists has-child. person

mother = woman \sqcap \exists has-child. person

father = man \sqcap \exists has-child. person

grandmother = woman \sqcap \exists has-child. (mother \sqcup father)

aunt = woman \sqcap \exists has-sibling. (mother \sqcup father)

uncle = man \sqcap \exists has-sibling. (mother \sqcup father)

sister = woman \sqcap \exists has-sibling. person

brother = man \sqcap \exists has-sibling. person

only child = person \sqcap \neg \exists has-sibling. person

cousin = person \sqcap \exists has-parent. \exists has-sibling. (mother \sqcup father)

nephew = man \sqcap \exists has-parent. (brother \sqcup sister)

niece = woman \sqcap \exists has-parent. (brother \sqcup sister)

father-in-law = man \sqcap \exists has-child. \exists has-spouse. person

mother-in-law = woman \sqcap \exists has-child. \exists has-spouse. person