

## 1. REGULAR LANGUAGES AND REGULAR GRAMMARS

### → INTRODUCTION:

- The languages accepted by finite automata are described by regular expressions. So to prove a language is accepted by finite automata it is sufficient to prove the regular expression of that language is accepted by finite automata.
- Regular expressions are used in many programming languages and language tools like lex, vi editor, PHP, PERL etc. They are used as powerful tools in search engines.

### → REGULAR EXPRESSIONS:

- Regular Expressions are shorthand notations to describe a language.
- A Regular expression is a notation to represent certain sets of strings in an algebraic fashion.
- Regular Expression describes the language accepted by FA.

#### Definition

\* Any terminal symbol/element of  $\Sigma$  is R.E.

Eg:  $\phi, \epsilon, a$  in  $\Sigma$

$\phi$  is a regular expression & denotes the empty set.

$\epsilon$  is a regular expression and denotes the set  $\{\epsilon\}$ .

$a$  is a regular expression and denotes the set  $\{a\}$ .

\* ) Union of two regular expressions  $R_1$  and  $R_2$  is regular expression  $R$  ( $R = R_1 + R_2$ ).

Eg: Let "a" be regular expression  $R_1$

"b" be regular expression  $R_2$

$(a+b)$  is also a regular expression & having the elements  $\{a, b\}$ .

\* ) Concatenation of two regular expressions  $R_1$  and  $R_2$  written as  $R_1.R_2$  is also regular expression  $R$  ( $R = R_1 \cdot R_2$ )

Eg: Let "a" be regular expression  $R_1$

"b" be regular expression  $R_2$

$(a.b)$  also a regular expression  $R$  having the element  $ab$ .

\* ) Iteration of a regular expression  $R$  written as  $R^*$  is also a regular expression.  
(closure)

Eg: Let "z" be a regular expression

then  $\epsilon, a, aa, aaa, \dots$  are also regular expression.

If  $L$  is a language represented by the regular expression  $R$  then Kleen closure of  $L$  is denoted as  $L^*$  and given as

$$L^* = \bigcup_{i=0}^{\infty} L^i$$

The positive closure of  $L$ , denoted  $L^+$ , is the set

$$L^+ = \bigcup_{i=1}^{\infty} L^i$$

\* ) If  $R$  is a r.e. then  $R^*$  is also a regular expression.

\* ) R.E over  $\Sigma$  is precisely those obtained recursively by the application of the above rules once or several times.

- The basic operations performed on regular expressions are
  - Union
  - Concatenation
  - Kleene's closure.

Among the three closure has highest priority (precedence)  
 next highest is for concatenation and least is for union.

→ REGULAR SET:

- ✓ Any set represented by a regular expression is called a regular set.

Eg:  $\Sigma = \{0,1\}$ , then 0 denotes a set  $\{0\}$ .

- ✓ If  $a, b$  are the elements of  $\Sigma$  then regular expressions

- $\phi$  is a regular expression which denotes an empty set.
- $\epsilon$  is a regular expression which denotes a set  $\{\epsilon\}$ .
- $a$  denotes the set  $\{a\}$
- $a+b$  denotes the set  $\{a,b\}$
- $ab$  denotes the set  $\{ab\}$
- $a^k$  denotes the set  $\{\epsilon, a, aa, aaa, \dots\}$
- $(a+b)^*$  denotes the set  $\{\epsilon, a, b, aa, ab, ba, bb, \dots\}$

- ✓ Representing the regular set by regular expression

Regular set	Regular expression
$\{101\}$	$101$
$\{\epsilon, a\}$	$\epsilon + a$
$\{\epsilon, a, aa, ab, ba, bb, \dots\}$	$(a+b)^*$
$\{ab, ba\}$	$(ab+ba)$ .

> Describing the following sets by regular expression.

D) All strings of 0's and 1's.

$$(0+1)^*$$

② Set of all strings of 0's and 1's ending with 00

$$(0+1)^*00$$

③ Set of all strings ending with 011

$$(0+1)^*011$$

④ Set of all strings of 0's and 1's beginning with 00.

$$00(0+1)^*$$

⑤ Set of all strings of 0's and 1's beginning with 0 and ending with 1.

$$0(0+1)^*1$$

⑥ Set of all strings of 0's and 1's beginning with 1 and ending with 00.

$$1(0+1)^*00$$

⑦ Strings of 0's and 1's with at least two consecutive 0's

$$(0+1)^*00(0+1)^*$$

⑧ Regular expression that denotes all the words with at least two 0's.

$$(0+1)^*0(0+1)^*0(0+1)^*$$

⑨ Regular expression that denotes exactly two 0's

$$1^*01^*01^*$$

⑩ Set of all strings having even no. of 1's

$$(11)^*$$

⑪ Set of all strings having odd number of 1's

$$1(11)^* \text{ or } (11)^*1$$

⑫ Regular expression to denote any number of 0's followed by any number of 1's followed by any number of 2's

$$0^*1^*2^*$$

- (13) Regular expression to denote any number of 0's followed by any number of 1's followed by any number of 2's with at least one of each symbol.

$$00^*11^*22^* = 0^+1^+2^+$$

- (14) Regular expression for the set of strings, that consists of alternating 0's and 1's

$$(0+1)(01)^*(\epsilon+0)$$

(or).

$$(\epsilon+0)(10^*)(\epsilon+1)$$

- (15) Set of all strings of 0's and 1's whose last two symbols are the same.

$$(0+1)^*(00+11).$$

- (16) Set of strings in which every 0 is immediately followed by at least two 1's  
To find R.E consider two possibilities

(a) String with only 1's

(b) Every 0 preceded by 11. i.e., 011

Hence,  $(1+011)^*$

- (17) All strings of 0's and 1's beginning with 1 or 0 and not having two consecutive 0's.

$$(\epsilon+1)(1+0)^*$$