

Anagram → Two strings are called Anagrams of each other if they are permutation of each other.

LISTEN , SILENT  
↓  
SILENT

ABC ↔ BAC

Q<sub>⇒</sub> Given two strings, write a program that detects whether they are anagram of each other or not ??

Ex RACE, CARE → True

SANKET, SAMYAK → FALSE

Brute force  $\Rightarrow$  If two string are permutation of each other they're called anagram, so, we can generate all the permutations of the first string and check whether any one of them matches the second string.

(A B C , B A C)  $\Rightarrow$

Yes

ABC  
ACB  
BAC  
BCA  
CAB  
CBA

} 3!

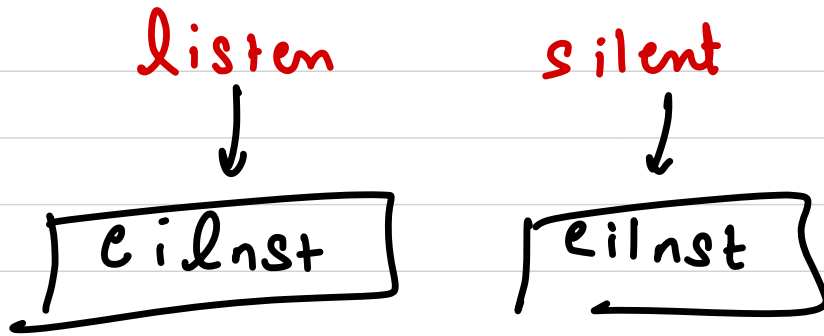
$O(n!)$

How can we optimise ??



The permutation in which the string is sorted -

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We know that if two strings are anagrams, they will have common permutations. And one of the permutations is going to be the sorted arrangement of characters.

Care, race  $\rightarrow$  time  $\rightarrow O(n \log n)$

The diagram illustrates the sorting process for two strings, 'Care' and 'race'. Red arrows point from each string to a red oval containing the sorted string 'acer'. This visualizes that both strings, when sorted by character, result in the same sequence: 'a', 'c', 'e', 'r'.

$\rightarrow$  Let's sort both the strings by chars,  
and check whether the sorted strings are  
same or not.

Can we optimise more on time ?? //

→ If two strings are anagram of each other then →

1) they will have same length

2) they will have exactly same set of

characters i.e. if one string has 2 d's then other should also have 2 d's.

Can we try to detect the fact that characters  
are same in both strings ??

$s_1 = K N \Sigma \Sigma$

$s_2 = K \Sigma \Sigma N$

empty

<del>K</del>	<del>1</del> 0
<del>N</del>	<del>1</del> 0
<del><math>\Sigma</math></del>	<del>2</del> 0



key-value

$S_1 = K N E E$

$S_2 = K E E L$

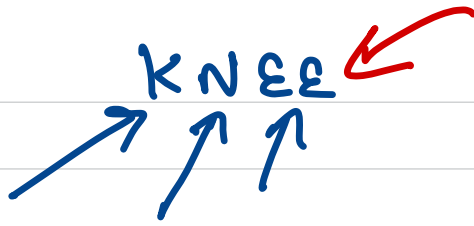
}

<del>K</del>	<del>0</del>
N-1	
<del>E-2</del>	<del>0</del>

Char  $\longleftrightarrow$  frequency

→ we will prepare a frequency mapping of  $S_1$   
→ then for each character in  $S_2$ , check if we have that char in the mapping of  $S_1$ . if not return false, else reduce the frequency as we got the char.

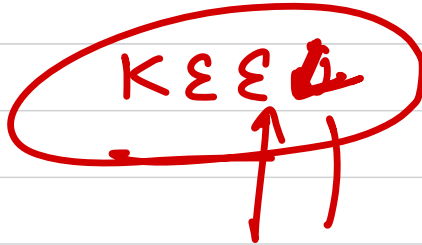
K N E E



mp = {  
  n: 1;  
}

→ ptr

K E E ~~E~~



0 1 2 3 4 5  
K N E A E D

$i = 0, 1, 2, 3, 4$  ✓

{

K: 1

N: 1

E: 1, 2

A: 1

}