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```

1 .PRECIOUS: ./p%
3 %: p%
  ulimit -s unlimited && ./p%
5 p%: p%.cpp
  g++ -o $@ $< -std=c++17 -Wall -Wextra -Wshadow \
7   -fsanitize=address,undefined

```

1.2. How Did We Get Here?

1.2.1. Macros

Use vectorizations and math optimizations at your own peril.
For gcc \geq 9, there are `[[likely]]` and `[[unlikely]]` attributes.

Call gcc with `-fopt-info-optimized-missed-optall` for optimization info.

```
1 #define _GLIBCXX_DEBUG 1 // for debug mode
2 #define _GLIBCXX_SANITIZE_VECTOR 1 // for asan on vectors
3 #pragma GCC optimize("O3", "unroll-loops")
4 #pragma GCC optimize("fast-math")
5 #pragma GCC target("avx,avx2,abm,bmi,bmi2") // tip: `lscpu`
6 // before a loop
7 #pragma GCC unroll 16 // 0 or 1 -> no unrolling
8 #pragma GCC ivdep
```

1.2.2. Fast I/O

```
1 struct scanner {
2     static constexpr size_t LEN = 32 << 20;
3     char *buf, *buf_ptr, *buf_end;
4     scanner()
5         : buf(new char[LEN]), buf_ptr(buf + LEN),
6           buf_end(buf + LEN) {}
7     ~scanner() { delete[] buf; }
8     char getc() {
9         if (buf_ptr == buf_end) [[unlikely]]
10            buf_end = buf + fread_unlocked(buf, 1, LEN, stdin),
11            buf_ptr = buf;
12        return *(buf_ptr++);
13    }
14    char seek(char del) {
15        char c;
16        while ((c = getc()) < del) {}
17        return c;
18    }
19    void read(int &t) {
20        bool neg = false;
21        char c = seek('-');
22        if (c == '-') neg = true, t = 0;
23        else t = c ^ '0';
24        while ((c = getc()) >= '0') t = t * 10 + (c ^ '0');
25        if (neg) t = -t;
26    }
27 };
28 struct printer {
29     static constexpr size_t CPI = 21, LEN = 32 << 20;
30     char *buf, *buf_ptr, *buf_end, *tbuf;
31     char *int_buf, *int_buf_end;
32     printer()
33         : buf(new char[LEN]), buf_ptr(buf),
34           buf_end(buf + LEN), int_buf(new char[CPI + 1]()),
35           int_buf_end(int_buf + CPI - 1) {}
36     ~printer() {
37         flush();
38         delete[] buf, delete[] int_buf;
39     }
40     void flush() {
41         fwrite_unlocked(buf, 1, buf_ptr - buf, stdout);
42         buf_ptr = buf;
43     }
44     void write_(const char &c) {
45         *buf_ptr = c;
46         if (++buf_ptr == buf_end) [[unlikely]]
47             flush();
48     }
49     void write_(const char *s) {
50         for (; *s != '\0'; ++s) write_(*s);
51     }
52     void write(int x) {
53         if (x < 0) write_('-'), x = -x;
54         if (x == 0) [[unlikely]]
55             return write_('0');
56         for (tbuf = int_buf_end; x != 0; --tbuf, x /= 10)
57             *tbuf = '0' + char(x % 10);
58         write_(++tbuf);
59     }
60 };
```

Kotlin

```
1 import java.io.*
2 import java.util.*
3
4 @JvmField val cin = System.`in`.bufferedReader()
5 @JvmField val cout = PrintWriter(System.out, false)
6 @JvmField var tokenizer: StringTokenizer = StringTokenizer("")
7 fun nextLine() = cin.readLine()!!
8 fun read(): String {
9     while(!tokenizer.hasMoreTokens())
10         tokenizer = StringTokenizer(nextLine())
11     return tokenizer.nextToken()
12 }
13
14 // example
15 fun main() {
```

```
val n = read().toInt()
val a = DoubleArray(n) { read().toDouble() }
cout.println("omg hi")
cout.flush()
}
```

1.2.3. constexpr

Some default limits in gcc (7.x - trunk):

- constexpr recursion depth: 512
- constexpr loop iteration per function: 262144
- constexpr operation count per function: 33554432
- template recursion depth: 900 (gcc *might* segfault first)

```
1 constexpr array<int, 10> fibonacci{[] {
2     array<int, 10> a{};
3     a[0] = a[1] = 1;
4     for (int i = 2; i < 10; i++) a[i] = a[i - 1] + a[i - 2];
5     return a;
6 }}();
7 static_assert(fibonacci[9] == 55, "CE");
8
9 template <typename F, typename INT, INT... S>
10 constexpr void for_constexpr(integer_sequence<INT, S...>,
11                             F &&func) {
12     int _[] = {(func(integral_constant<INT, S>{}), 0)...};
13 }
14 // example
15 template <typename... T> void print_tuple(tuple<T...> t) {
16     for_constexpr(make_index_sequence<sizeof...(T)>{}),
17                 [&](auto i) { cout << get<i>(t) << '\n'; });
18 }
```

1.2.4. Bump Allocator

```
1
2
3 // global bump allocator
4 char mem[256 << 20]; // 256 MB
5 size_t rsp = sizeof mem;
6 void *operator new(size_t s) {
7     assert(s < rsp); // MLE
8     return (void *)&mem[rsp -= s];
9 }
10 void operator delete(void *) {}
11
12 // bump allocator for STL / pbds containers
13 char mem[256 << 20];
14 size_t rsp = sizeof mem;
15 template <typename T> struct bump {
16     typedef T value_type;
17     bump() {}
18     template <typename U> bump(U, ...) {}
19     T *allocate(size_t n) {
20         rsp -= n * sizeof(T);
21         rsp &= 0 - alignof(T);
22         return (T *)(&mem + rsp);
23     }
24     void deallocate(T *, size_t n) {}
25 };
```

1.3. Tools

1.3.1. Floating Point Binary Search

```
1 union di {
2     double d;
3     ull i;
4 };
5 bool check(double);
6 // binary search in [L, R) with relative error 2^-eps
7 double binary_search(double L, double R, int eps) {
8     di l = {L}, r = {R}, m;
9     while (r.i - l.i > 1LL << (52 - eps)) {
10         m.i = (l.i + r.i) >> 1;
11         if (check(m.d)) r = m;
12         else l = m;
13     }
14     return l.d;
15 }
```

1.3.2. SplitMix64

```
1 using ull = unsigned long long;
2 inline ull splitmix64(ull x) {
3     // change to `static ull x = SEED;` for DRBG
4     ull z = (x += 0x9E3779B97F4A7C15);
5     z = (z ^ (z >> 30)) * 0xBF58476D1CE4E5B9;
6     z = (z ^ (z >> 27)) * 0x94D049BB133111EB;
7     return z ^ (z >> 31);
8 }
```

1.3.3. <random>

```

1 #ifdef __unix__
  random_device rd;
  mt19937_64 RNG(rd());
2 #else
3 const auto SEED = chrono::high_resolution_clock::now()
  .time_since_epoch()
  .count();
4 mt19937_64 RNG(SEED);
5 #endif
6 // random uint_fast64_t: RNG();
7 // uniform random of type T (int, double, ...) in [l, r]:
8 // uniform_int_distribution<T> dist(l, r); dist(RNG);

```

1.3.4. x86 Stack Hack

```

1 constexpr size_t size = 200 << 20; // 200MiB
2 int main() {
3   register long rsp asm("rsp");
4   char *buf = new char[size];
5   asm("movq %0, %%rsp\n" :: "r"(buf + size));
6   // do stuff
7   asm("movq %0, %%rsp\n" :: "r"(rsp));
8   delete[] buf;
9 }

```

1.4. Algorithms

1.4.1. Bit Hacks

```

1 // next permutation of x as a bit sequence
2 ull next_bits_permutation(ull x) {
3   ull c = __builtin_ctzll(x), r = x + (1ULL << c);
4   return (r ^ x) >> (c + 2) | r;
5 }
6 // iterate over all (proper) subsets of bitset s
7 void subsets(ull s) {
8   for (ull x = s; x;) { --x &= s; /* do stuff */ }
9 }

```

1.4.2. Aliens Trick

```

1 // min dp[i] value and its i (smallest one)
2 pll get_dp(int cost);
3 ll aliens(int k, int l, int r) {
4   while (l != r) {
5     int m = (l + r) / 2;
6     auto [f, s] = get_dp(m);
7     if (s == k) return f - m * k;
8     if (s < k) r = m;
9     else l = m + 1;
10  }
11  return get_dp(l).first - l * k;
12 }

```

1.4.3. Hilbert Curve

```

1 ll hilbert(ll n, int x, int y) {
2   ll res = 0;
3   for (ll s = n; s /= 2;) {
4     int rx = !(x & s), ry = !(y & s);
5     res += s * s * ((3 * rx) ^ ry);
6     if (ry == 0) {
7       if (rx == 1) x = s - 1 - x, y = s - 1 - y;
8       swap(x, y);
9     }
10  }
11  return res;
12 }

```

1.4.4. Infinite Grid Knight Distance

```

1 ll get_dist(ll dx, ll dy) {
2   if (++dx == abs(dx)) > ++(dy = abs(dy)) swap(dx, dy);
3   if (dx == 1 && dy == 2) return 3;
4   if (dx == 3 && dy == 3) return 4;
5   ll lb = max(dy / 2, (dx + dy) / 3);
6   return ((dx ^ dy ^ lb) & 1) ? ++lb : lb;
7 }

```

1.4.5. Poker Hand

```

1
2
3
4
5
6
7 using namespace std;
8
9 struct hand {
10   static constexpr auto rk = [] {

```

```

11   array<int, 256> x{};
12   auto s = "23456789TJQKACDHS";
13   for (int i = 0; i < 17; i++) x[s[i]] = i % 13;
14   return x;
15 }();
16 vector<pair<int, int>> v;
17 vector<int> cnt, vf, vs;
18 int type;
19 hand() : cnt(4), type(0) {}
20 void add_card(char suit, char rank) {
21   ++cnt[rk[suit]];
22   for (auto &[f, s] : v)
23     if (s == rk[rank]) return ++f, void();
24   v.emplace_back(1, rk[rank]);
25 }
26 void process() {
27   sort(v.rbegin(), v.rend());
28   for (auto [f, s] : v) vf.push_back(f), vs.push_back(s);
29   bool str = 0, flu = find(all(cnt), 5) != cnt.end();
30   if ((str = v.size() == 5))
31     for (int i = 1; i < 5; i++)
32       if (vs[i] != vs[i - 1] + 1) str = 0;
33   if (vs == vector<int>{12, 3, 2, 1, 0})
34     str = 1, vs = {3, 2, 1, 0, -1};
35   if (str && flu) type = 9;
36   else if (vf[0] == 4) type = 8;
37   else if (vf[0] == 3 && vf[1] == 2) type = 7;
38   else if (str || flu) type = 5 + flu;
39   else if (vf[0] == 3) type = 4;
40   else if (vf[0] == 2) type = 2 + (vf[1] == 2);
41   else type = 1;
42 }
43 bool operator<(const hand &b) const {
44   return make_tuple(type, vf, vs) <
45     make_tuple(b.type, b.vf, b.vs);
46 }
47 };

```

1.4.6. Longest Increasing Subsequence

```

1
2
3 template <class I> vi lis(const vector<I> &S) {
4   if (S.empty()) return {};
5   vi prev(sz(S));
6   typedef pair<I, int> p;
7   vector<p> res;
8   rep(i, 0, sz(S)) {
9     // change 0 -> i for longest non-decreasing subsequence
10    auto it = lower_bound(all(res), p{S[i], 0});
11    if (it == res.end())
12      res.emplace_back(), it = res.end() - 1;
13    *it = {S[i], i};
14    prev[i] = it == res.begin() ? 0 : (it - 1)->second;
15  }
16  int L = sz(res), cur = res.back().second;
17  vi ans(L);
18  while (L--) ans[L] = cur, cur = prev[cur];
19  return ans;
20 }

```

1.4.7. Mo's Algorithm on Tree

```

1 void MoAlgoOnTree() {
2   Dfs(0, -1);
3   vector<int> euler(tk);
4   for (int i = 0; i < n; ++i) {
5     euler[tin[i]] = i;
6     euler[tout[i]] = i;
7   }
8   vector<int> l(q), r(q), qr(q), sp(q, -1);
9   for (int i = 0; i < q; ++i) {
10    if (tin[u[i]] > tin[v[i]]) swap(u[i], v[i]);
11    int z = GetLCA(u[i], v[i]);
12    sp[i] = z[i];
13    if (z == u) l[i] = tin[u[i]], r[i] = tin[v[i]];
14    else l[i] = tout[u[i]], r[i] = tin[v[i]];
15    qr[i] = i;
16  }
17  sort(qr.begin(), qr.end(), [&](int i, int j) {
18    if (l[i] / kB == l[j] / kB) return r[i] < r[j];
19    return l[i] / kB < l[j] / kB;
20  });
21  vector<bool> used(n);
22  // Add(v): add/remove v to/from the path based on used[v]
23  for (int i = 0, tl = 0, tr = -1; i < q; ++i) {
24    while (tl < l[qr[i]]) Add(euler[tl++]);
25    while (tl > l[qr[i]]) Add(euler[--tl]);
26    while (tr > r[qr[i]]) Add(euler[tr--]);
27    while (tr < r[qr[i]]) Add(euler[++tr]);
28    // add/remove LCA(u, v) if necessary
29  }
30 }

```

2. Data Structures

2.1. GNU PBDS

```

1 #include <ext/pb_ds/assoc_container.hpp>
2 #include <ext/pb_ds/priority_queue.hpp>
3 #include <ext/pb_ds/tree_policy.hpp>
4 using namespace __gnu_pbds;
5
6 // most std::map + order_of_key, find_by_order, split, join
7 template <typename T, typename U = null_type>
8 using ordered_map = tree<T, U, std::less<>, rb_tree_tag,
9                        tree_order_statistics_node_update>;
10 // useful tags: rb_tree_tag, splay_tree_tag
11
12 template <typename T> struct myhash {
13     size_t operator()(T x) const; // splitmix, bswap(x*R), ...
14 };
15 // most of std::unordered_map, but faster (needs good hash)
16 template <typename T, typename U = null_type>
17 using hash_table = gp_hash_table<T, U, myhash<T>>;
18
19 // most std::priority_queue + modify, erase, split, join
20 using heap = priority_queue<int, std::less<>>;
21 // useful tags: pairing_heap_tag, binary_heap_tag,
22 //              (rc_)?binomial_heap_tag, thin_heap_tag

```

2.2. Segment Tree (ZKW)

```

1 struct segtree {
2     using T = int;
3     T f(T a, T b) { return a + b; } // any monoid operation
4     static constexpr T ID = 0; // identity element
5     int n;
6     vector<T> v;
7     segtree(int n_) : n(n_), v(2 * n, ID) {}
8     segtree(vector<T> &a) : n(a.size()), v(2 * n, ID) {
9         copy_n(a.begin(), n, v.begin() + n);
10        for (int i = n - 1; i > 0; i--)
11            v[i] = f(v[i * 2], v[i * 2 + 1]);
12    }
13    void update(int i, T x) {
14        for (v[i += n] = x; i /= 2;)
15            v[i] = f(v[i * 2], v[i * 2 + 1]);
16    }
17    T query(int l, int r) {
18        T tl = ID, tr = ID;
19        for (l += n, r += n; l < r; l /= 2, r /= 2) {
20            if (l & 1) tl = f(tl, v[l++]);
21            if (r & 1) tr = f(v[--r], tr);
22        }
23        return f(tl, tr);
24    }
25 };

```

2.3. Line Container

```

1
2
3 struct Line {
4     mutable ll k, m, p;
5     bool operator<(const Line &o) const { return k < o.k; }
6     bool operator<(ll x) const { return p < x; }
7 };
8 // add: line y=kx+m, query: maximum y of given x
9 struct LineContainer : multiset<Line, less<>> {
10    // (for doubles, use inf = 1/.0, div(a,b) = a/b)
11    static const ll inf = LLONG_MAX;
12    ll div(ll a, ll b) { // floored division
13        return a / b - ((a ^ b) < 0 && a % b);
14    }
15    bool isect(iterator x, iterator y) {
16        if (y == end()) return x->p = inf, 0;
17        if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
18        else x->p = div(y->m - x->m, x->k - y->k);
19        return x->p >= y->p;
20    }
21    void add(ll k, ll m) {
22        auto z = insert({k, m, 0}); y = z++, x = y;
23        while (isect(y, z)) z = erase(z);
24        if (x != begin() && isect(--x, y))
25            isect(x, y = erase(y));
26        while ((y = x) != begin() && (--x)->p >= y->p)
27            isect(x, erase(y));
28    }
29    ll query(ll x) {
30        assert(!empty());
31        auto l = *lower_bound(x);
32        return l.k * x + l.m;
33    }
34 };

```

2.4. Li-Chao Tree

```

1 public class LiChaoTree {
2
3     // Represents a line y = mx + c
4     static class Line {
5         long m, c;
6
7         public Line(long m, long c) {
8             this.m = m;
9             this.c = c;
10        }
11
12        // Evaluates the line at a given x-coordinate
13        long eval(long x) {
14            return m * x + c;
15        }
16    }
17
18    // Node of the Li-Chao Tree
19    static class Node {
20        Line line;
21        Node left, right;
22
23        public Node(Line line) {
24            this.line = line;
25        }
26    }
27
28    private Node root;
29    private final long minCoord;
30    private final long maxCoord;
31    private final Line identityLine; // Represents "no line" or infinity
32
33    // Constructor for the Li-Chao Tree
34    // minCoord and maxCoord define the range of x-values the tree covers
35    // identityLine should return a very large value for min queries
36    public LiChaoTree(long minCoord, long maxCoord) {
37        this.minCoord = minCoord;
38        this.maxCoord = maxCoord;
39        // For minimum queries, an identity line should return a very large value
40        // Using Long.MAX_VALUE for 'c' and 0 for 'm' ensures it's the worst line
41        this.identityLine = new Line(0, Long.MAX_VALUE);
42        this.root = new Node(identityLine);
43    }
44
45    // Adds a new line to the tree
46    public void addLine(Line newLine) {
47        addLine(root, minCoord, maxCoord, newLine);
48    }
49
50    private void addLine(Node node, long currentMin, long currentMax, Line newLine) {
51        long mid = currentMin + (currentMax - currentMin) / 2;
52        boolean leftBetter = newLine.eval(currentMin) < node.line.eval(currentMin);
53        boolean midBetter = newLine.eval(mid) < node.line.eval(mid);
54
55        if (midBetter) {
56            // If the new line is better at the midpoint, swap it with the current line
57            Line temp = node.line;
58            node.line = newLine;
59            newLine = temp; // The old line now becomes the 'new' line to be added
60        }
61
62        // If the interval is a single point, we are done
63        if (currentMin == currentMax) {
64            return;
65        }
66
67        // Decide which child to push the 'worse' line to
68        if (leftBetter != midBetter) { // Intersection point is in the interval
69            if (node.left == null) {
70                node.left = new Node(identityLine);
71            }
72            addLine(node.left, currentMin, mid, newLine);
73        } else if (leftBetter == midBetter && leftBetter == false) {
74            if (node.right == null) {
75                node.right = new Node(identityLine);
76            }
77            addLine(node.right, mid + 1, currentMax, newLine);
78        }
79        // If leftBetter == midBetter == true, it means the new line is completely dominated,
80        // so no need to add it
81    }
82
83    // Queries the minimum value at a given x-coordinate
84    public long query(long x) {
85        return query(root, minCoord, maxCoord, x);
86    }
87
88    private long query(Node node, long currentMin, long currentMax, long x) {
89        if (node == null) {
90            return identityLine.eval(x); // No line in this path, return infinity
91        }

```



```

93     long res = node.line.eval(x);
95     if (currentMin == currentMax) {
96         return res; // Reached a leaf node
97     }
99     long mid = currentMin + (currentMax - currentMin) / 2;
101    if (x <= mid) {
102        res = Math.min(res, query(node.left, currentMin, mid, x));
103    } else {
104        res = Math.min(res, query(node.right, mid + 1, currentMax, x));
105    }
106    return res;
107 }
109 public static void main(String[] args) {
110     // Example Usage:
111     // Create a Li-Chao Tree for x-coordinates from -1000 to 1000
112     LiChaoTree lct = new LiChaoTree(-1000, 1000);
114     // Add some lines: y = mx + c
115     lct.addLine(new Line(1, 5)); // y = x + 5
116     lct.addLine(new Line(-1, 10)); // y = -x + 10
117     lct.addLine(new Line(0, 7)); // y = 7
119     // Query minimum values at different x-coordinates
120     System.out.println("Min at x = 0: " + lct.query(0));
121     System.out.println("Min at x = 2: " + lct.query(2));
122     System.out.println("Min at x = 6: " + lct.query(6));
123     System.out.println("Min at x = -5: " + lct.query(-5));
125     lct.addLine(new Line(2, -3)); // y = 2x - 3
126     System.out.println("Min at x = 0 (after new line): " + lct.query(0));
127     System.out.println("Min at x = 2 (after new line): " + lct.query(2));
128 }

```

2.5. Heavy-Light Decomposition

```

1  template <bool VALS_EDGES> struct HLD {
2      int N, tim = 0;
3      vector<vi> adj;
4      vi par, siz, depth, rt, pos;
5      Node *tree;
6      HLD(vector<vi> adj_)
7          : N(sz(adj_)), adj(adj_), par(N, -1), siz(N, 1),
8            depth(N), rt(N), pos(N), tree(new Node(0, N)) {
9          dfsSz(0);
10         dfsHld(0);
11     }
12     void dfsSz(int v) {
13         if (par[v] != -1)
14             adj[v].erase(find(all(adj[v]), par[v]));
15         for (int &u : adj[v]) {
16             par[u] = v, depth[u] = depth[v] + 1;
17             dfsSz(u);
18             siz[v] += siz[u];
19             if (siz[u] > siz[adj[v][0]]) swap(u, adj[v][0]);
20         }
21     }
22     void dfsHld(int v) {
23         pos[v] = tim++;
24         for (int u : adj[v]) {
25             rt[u] = (u == adj[v][0] ? rt[v] : u);
26             dfsHld(u);
27         }
28     }
29     template <class B> void process(int u, int v, B op) {
30         for (; rt[u] != rt[v]; v = par[rt[v]]) {
31             if (depth[rt[u]] > depth[rt[v]]) swap(u, v);
32             op(pos[rt[v]], pos[v] + 1);
33         }
34         if (depth[u] > depth[v]) swap(u, v);
35         op(pos[u] + VALS_EDGES, pos[v] + 1);
36     }
37     void modifyPath(int u, int v, int val) {
38         process(u, v, [&](int l, int r) { tree->add(l, r, val); });
39     }
40     int queryPath(int u, int v) { // Modify depending on problem
41         int res = -1e9;
42         process(u, v, [&](int l, int r) {
43             res = max(res, tree->query(l, r));
44         });
45         return res;
46     }
47     int querySubtree(int v) { // modifySubtree is similar
48         return tree->query(pos[v] + VALS_EDGES, pos[v] + siz[v]);
49     }
50 };

```

2.6. Wavelet Matrix

```

1  #pragma GCC target("popcnt,bmi2")
2  #include <immintrin.h>
3  // T is unsigned. You might want to compress values first
4  template <typename T> struct wavelet_matrix {
5      static_assert(is_unsigned_v<T>, "only unsigned T");
6      struct bit_vector {
7          static constexpr uint W = 64;
8          uint n, cnt0;
9          vector<ull> bits;
10         vector<uint> sum;
11         bit_vector(uint n_)
12             : n(n_), bits(n / W + 1), sum(n / W + 1) {}
13         void build() {
14             for (uint j = 0; j != n / W; ++j)
15                 sum[j + 1] = sum[j] + _mm_popcnt_u64(bits[j]);
16             cnt0 = rank0(n);
17         }
18         void set_bit(uint i) { bits[i / W] |= 1ULL << i % W; }
19         bool operator[](uint i) const {
20             return !(bits[i / W] & 1ULL << i % W);
21         }
22         // Expected: min(5, 10, 7) = 5
23         // Expected: rank1(uint i) = 5
24         // Expected: rank1(uint i) = 4
25         // Expected: min(sum[l, r]) = 4
26         // Expected: min(_mm_popcnt_u64(_bshi_u64(bits[i / W], i % W)))
27         uint rank1(uint i) const {
28             return sum[i / W];
29         }
30         uint rank0(uint i) const { return i - rank1(i); }
31         // lct.query(0); // Expected: min(5, 10, 7, -3) = -3
32         // lct.query(2); // Expected: min(7, 8, 7, 1) = 1
33         vector<bit_vector> b;
34         wavelet_matrix(const vector<T> &a) : n(a.size()) {
35             lg =
36                 lg(max(*max_element(a.begin(), a.end()), T(1))) + 1;
37             b.assign(lg, n);
38             vector<T> cur = a, nxt(n);
39             for (int h = lg; h--;) {
40                 for (uint i = 0; i < n; ++i)
41                     if (cur[i] & (T(1) << h)) b[h].set_bit(i);
42                 b[h].build();
43                 int il = 0, ir = b[h].cnt0;
44                 for (uint i = 0; i < n; ++i)
45                     nxt[(b[h][i] ? ir : il)++] = cur[i];
46                 swap(cur, nxt);
47             }
48         }
49         T operator[](uint i) const {
50             T res = 0;
51             for (int h = lg; h--;)
52                 if (b[h][i])
53                     i += b[h].cnt0 - b[h].rank0(i), res |= T(1) << h;
54             return res;
55         }
56         // query k-th smallest (0-based) in a[l, r]
57         T kth(uint l, uint r, uint k) const {
58             T res = 0;
59             for (int h = lg; h--;) {
60                 uint tl = b[h].rank0(l), tr = b[h].rank0(r);
61                 if (k >= tr - tl) {
62                     k -= tr - tl;
63                     l += b[h].cnt0 - tl;
64                     r += b[h].cnt0 - tr;
65                     res |= T(1) << h;
66                 } else l = tl, r = tr;
67             }
68             return res;
69         }
70         // count of i in [l, r] with a[i] < u
71         uint count(uint l, uint r, T u) const {
72             if (u >= T(1) << lg) return r - l;
73             uint res = 0;
74             for (int h = lg; h--;) {
75                 uint tl = b[h].rank0(l), tr = b[h].rank0(r);
76                 if (u & (T(1) << h)) {
77                     l += b[h].cnt0 - tl;
78                     r += b[h].cnt0 - tr;
79                     res += tr - tl;
80                 } else l = tl, r = tr;
81             }
82             return res;
83         }
84     };
85 };

```

2.7. Link-Cut Tree

```

1
3  const int MXN = 100005;
4  const int MEM = 100005;

```

```

5 struct Splay {
6     static Splay nil, mem[MEM], *pmem;
7     Splay *ch[2], *f;
8     int val, rev, size;
9     Splay() : val(-1), rev(0), size(0) {
10         f = ch[0] = ch[1] = &nil;
11     }
12     Splay(int _val) : val(_val), rev(0), size(1) {
13         f = ch[0] = ch[1] = &nil;
14     }
15     bool isr() {
16         return f->ch[0] != this && f->ch[1] != this;
17     }
18     int dir() { return f->ch[0] == this ? 0 : 1; }
19     void setCh(Splay *c, int d) {
20         ch[d] = c;
21         if (c != &nil) c->f = this;
22         pull();
23     }
24     void push() {
25         if (rev) {
26             swap(ch[0], ch[1]);
27             if (ch[0] != &nil) ch[0]->rev ^= 1;
28             if (ch[1] != &nil) ch[1]->rev ^= 1;
29             rev = 0;
30         }
31     }
32     void pull() {
33         size = ch[0]->size + ch[1]->size + 1;
34         if (ch[0] != &nil) ch[0]->f = this;
35         if (ch[1] != &nil) ch[1]->f = this;
36     }
37 } Splay::nil, Splay::mem[MEM], *Splay::pmem = Splay::mem;
38 Splay *nil = &Splay::nil;
39
40 void rotate(Splay *x) {
41     Splay *p = x->f;
42     int d = x->dir();
43     if (!p->isr()) p->f->setCh(x, p->dir());
44     else x->f = p->f;
45     p->setCh(x->ch[!d], d);
46     x->setCh(p, !d);
47     p->pull();
48     x->pull();
49 }
50
51 vector<Splay*> splayVec;
52 void splay(Splay *x) {
53     splayVec.clear();
54     for (Splay *q = x;; q = q->f) {
55         splayVec.push_back(q);
56         if (q->isr()) break;
57     }
58     reverse(begin(splayVec), end(splayVec));
59     for (auto it : splayVec) it->push();
60     while (!x->isr()) {
61         if (x->f->isr()) rotate(x);
62         else if (x->dir() == x->f->dir())
63             rotate(x->f), rotate(x);
64         else rotate(x), rotate(x);
65     }
66 }
67
68 Splay *access(Splay *x) {
69     Splay *q = nil;
70     for (; x != nil; x = x->f) {
71         splay(x);
72         x->setCh(q, 1);
73         q = x;
74     }
75     return q;
76 }
77
78 void evert(Splay *x) {
79     access(x);
80     splay(x);
81     x->rev ^= 1;
82     x->push();
83     x->pull();
84 }
85
86 void link(Splay *x, Splay *y) {
87     // evert(x);
88     access(x);
89     splay(x);
90     evert(y);
91     x->setCh(y, 1);
92 }
93
94 void cut(Splay *x, Splay *y) {
95     // evert(x);
96     access(y);
97     splay(y);
98     y->push();
99     y->ch[0] = y->ch[0]->f = nil;
100 }

```

```

99 int N, Q;
100 Splay *vt[MXN];
101
102 int ask(Splay *x, Splay *y) {
103     access(x);
104     access(y);
105     splay(x);
106     int res = x->f->val;
107     if (res == -1) res = x->val;
108     return res;
109 }
110
111 int main(int argc, char **argv) {
112     scanf("%d%d", &N, &Q);
113     for (int i = 1; i <= N; i++)
114         vt[i] = new (Splay::pmem++) Splay(i);
115     while (Q--) {
116         char cmd[105];
117         int u, v;
118         scanf("%s", cmd);
119         if (cmd[1] == 'i') {
120             scanf("%d%d", &u, &v);
121             link(vt[u], vt[v]);
122         } else if (cmd[0] == 'c') {
123             scanf("%d", &v);
124             cut(vt[1], vt[v]);
125         } else {
126             scanf("%d%d", &u, &v);
127             int res = ask(vt[u], vt[v]);
128             printf("%d\n", res);
129         }
130     }
131 }

```

3. Graph

3.1. Modeling

- Maximum/Minimum flow with lower bound / Circulation problem
 - Construct super source S and sink T .
 - For each edge (x, y, l, u) , connect $x \rightarrow y$ with capacity $u - l$.
 - For each vertex v , denote by $in(v)$ the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - If $in(v) > 0$, connect $S \rightarrow v$ with capacity $in(v)$, otherwise, connect $v \rightarrow T$ with capacity $-in(v)$.
 - To maximize, connect $t \rightarrow s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T . If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T . Connect $t \rightarrow s$ with capacity ∞ and let the flow from S to T be f' . If $f + f' \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X, Y)
 - Redirect every edge: $y \rightarrow x$ if $(x, y) \in M$, $x \rightarrow y$ otherwise.
 - DFS from unmatched vertices in X .
 - $x \in X$ is chosen iff x is unvisited.
 - $y \in Y$ is chosen iff y is visited.
- Minimum cost cyclic flow
 - Construct super source S and sink T .
 - For each edge (x, y, c) , connect $x \rightarrow y$ with $(cost, cap) = (c, 1)$ if $c > 0$, otherwise connect $y \rightarrow x$ with $(cost, cap) = (-c, 1)$.
 - For each edge with $c < 0$, sum these cost as K , then increase $d(y)$ by 1, decrease $d(x)$ by 1.
 - For each vertex v with $d(v) > 0$, connect $S \rightarrow v$ with $(cost, cap) = (0, d(v))$.
 - For each vertex v with $d(v) < 0$, connect $v \rightarrow T$ with $(cost, cap) = (0, -d(v))$.
 - Flow from S to T , the answer is the cost of the flow $C + K$.
- Maximum density induced subgraph
 - Binary search on answer, suppose we're checking answer T .
 - Construct a max flow model, let K be the sum of all weights.
 - Connect source $s \rightarrow v$, $v \in G$ with capacity K .
 - For each edge (u, v, w) in G , connect $u \rightarrow v$ and $v \rightarrow u$ with capacity w .
 - For $v \in G$, connect it with sink $v \rightarrow t$ with capacity $K + 2T - (\sum_{e \in E(v)} w(e)) - 2w(v)$.
 - T is a valid answer if the maximum flow $f < K|V|$.
- Minimum weight edge cover
 - For each $v \in V$ create a copy v' , and connect $u' \rightarrow v'$ with weight $w(u, v)$.
 - Connect $v \rightarrow v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v .
 - Find the minimum weight perfect matching on G' .
- Project selection problem
 - If $p_v > 0$, create edge (s, v) with capacity p_v ; otherwise, create edge (v, t) with capacity $-p_v$.

2. Create edge (u, v) with capacity w with w being the cost of choosing u without choosing v .
 3. The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

$$\sum_x c_x x + \sum_y c_y \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x \bar{y} + x' \bar{y}')$$

can be minimized by the mincut of the following graph:

1. Create edge (x, t) with capacity c_x and create edge (s, y) with capacity c_y .
2. Create edge (x, y) with capacity c_{xy} .
3. Create edge (x, y) and edge (x', y') with capacity $c_{xyx'y'}$.

3.2. Matching/Flows

3.2.1. Dinic's Algorithm

```

1 struct Dinic {
2     struct edge {
3         int to, cap, flow, rev;
4     };
5     static constexpr int MAXN = 1000, MAXF = 1e9;
6     vector<edge> v[MAXN];
7     int top[MAXN], deep[MAXN], side[MAXN], s, t;
8     void make_edge(int s, int t, int cap) {
9         v[s].push_back({t, cap, 0, (int)v[t].size()});
10        v[t].push_back({s, 0, 0, (int)v[s].size() - 1});
11    }
12    int dfs(int a, int flow) {
13        if (a == t || !flow) return flow;
14        for (int &i = top[a]; i < v[a].size(); i++) {
15            edge &e = v[a][i];
16            if (deep[a] + 1 == deep[e.to] && e.cap - e.flow) {
17                int x = dfs(e.to, min(e.cap - e.flow, flow));
18                if (x) {
19                    e.flow += x, v[e.to][e.rev].flow -= x;
20                    return x;
21                }
22            }
23        }
24        deep[a] = -1;
25        return 0;
26    }
27    bool bfs() {
28        queue<int> q;
29        fill_n(deep, MAXN, 0);
30        q.push(s), deep[s] = 1;
31        int tmp;
32        while (!q.empty()) {
33            tmp = q.front(), q.pop();
34            for (edge e : v[tmp])
35                if (!deep[e.to] && e.cap != e.flow)
36                    deep[e.to] = deep[tmp] + 1, q.push(e.to);
37        }
38        return deep[t];
39    }
40    int max_flow(int _s, int _t) {
41        s = _s, t = _t;
42        int flow = 0, tflow;
43        while (bfs()) {
44            fill_n(top, MAXN, 0);
45            while ((tflow = dfs(s, MAXF))) flow += tflow;
46        }
47        return flow;
48    }
49    void reset() {
50        fill_n(side, MAXN, 0);
51        for (auto &i : v) i.clear();
52    }
53 };

```

3.2.2. Minimum Cost Flow

```

1 struct MCF {
2     struct edge {
3         ll to, from, cap, flow, cost, rev;
4     };
5     *fromE[MAXN];
6     vector<edge> v[MAXN];
7     ll n, s, t, flows[MAXN], dis[MAXN], pi[MAXN], flowlim;
8     void make_edge(int s, int t, ll cap, ll cost) {
9         if (!cap) return;
10        v[s].pb(edge{t, s, cap, 0LL, cost, v[t].size()});
11        v[t].pb(edge{s, t, 0LL, 0LL, -cost, v[s].size() - 1});
12    }
13    bitset<MAXN> vis;
14    void dijkstra() {
15        vis.reset();
16        __gnu_pbds::priority_queue<pair<ll, int>> q;
17        vector<decltype(q)::point_iterator> its(n);
18        q.push({0LL, s});
19        while (!q.empty()) {

```

```

19        int now = q.top().second;
20        q.pop();
21        if (vis[now]) continue;
22        vis[now] = 1;
23        ll ndis = dis[now] + pi[now];
24        for (edge &e : v[now]) {
25            if (e.flow == e.cap || vis[e.to]) continue;
26            if (dis[e.to] > ndis + e.cost - pi[e.to]) {
27                dis[e.to] = ndis + e.cost - pi[e.to];
28                flows[e.to] = min(flows[now], e.cap - e.flow);
29                fromE[e.to] = &e;
30                if (its[e.to] == q.end())
31                    its[e.to] = q.push({-dis[e.to], e.to});
32                else q.modify(its[e.to], {-dis[e.to], e.to});
33            }
34        }
35    }
36    bool AP(ll &flow) {
37        fill_n(dis, n, INF);
38        fromE[s] = 0;
39        dis[s] = 0;
40        flows[s] = flowlim - flow;
41        dijkstra();
42        if (dis[t] == INF) return false;
43        flow += flows[t];
44        for (edge *e = fromE[t]; e; e = fromE[e->from]) {
45            e->flow += flows[t];
46            v[e->to][e->rev].flow -= flows[t];
47        }
48        for (int i = 0; i < n; i++)
49            pi[i] = min(pi[i] + dis[i], INF);
50        return true;
51    }
52    pll solve(int _s, int _t, ll _flowlim = INF) {
53        s = _s, t = _t, flowlim = _flowlim;
54        pll re;
55        while (re.F != flowlim && AP(re.F)) {
56            for (int i = 0; i < n; i++)
57                for (edge &e : v[i])
58                    if (e.flow != 0) re.S += e.flow * e.cost;
59            re.S /= 2;
60            return re;
61        }
62    }
63    void init(int _n) {
64        n = _n;
65        fill_n(pi, n, 0);
66        for (int i = 0; i < n; i++) v[i].clear();
67    }
68    void setpi(int s) {
69        fill_n(pi, n, INF);
70        pi[s] = 0;
71        for (ll it = 0, flag = 1, tdis; flag && it < n; it++) {
72            flag = 0;
73            for (int i = 0; i < n; i++)
74                if (pi[i] != INF)
75                    for (edge &e : v[i])
76                        if (e.cap && (tdis = pi[i] + e.cost) < pi[e.to])
77                            pi[e.to] = tdis, flag = 1;
78        }
79    }
80 };

```

3.2.3. Gomory-Hu Tree

Requires: Dinic's Algorithm

```

1
2
3 int e[MAXN][MAXN];
4 int p[MAXN];
5 Dinic D; // original graph
6 void gomory_hu() {
7     fill(p, p + n, 0);
8     fill(e[0], e[n], INF);
9     for (int s = 1; s < n; s++) {
10        int t = p[s];
11        Dinic F = D;
12        int tmp = F.max_flow(s, t);
13        for (int i = 1; i < s; i++)
14            e[s][i] = e[i][s] = min(tmp, e[t][i]);
15        for (int i = s + 1; i <= n; i++)
16            if (p[i] == t && F.side[i]) p[i] = s;
17    }
18 }

```

3.2.4. Global Minimum Cut

```

1
2 // weights is an adjacency matrix, undirected
3 pair<int, vi> getMinCut(vector<vi> &weights) {
4     int N = sz(weights);
5     vi used(N), cut, best_cut;

```

```

7   int best_weight = -1;
9   for (int phase = N - 1; phase >= 0; phase--) {
11      vi w = weights[0], added = used;
13      int prev, k = 0;
15      rep(i, 0, phase) {
17         prev = k;
18         k = -1;
19         rep(j, 1, N) if (!added[j] &&
21             (k == -1 || w[j] > w[k])) k = j;
23         if (i == phase - 1) {
25             rep(j, 0, N) weights[prev][j] += weights[k][j];
27             rep(j, 0, N) weights[j][prev] = weights[prev][j];
29             used[k] = true;
31             cut.push_back(k);
33             if (best_weight == -1 || w[k] < best_weight) {
35                 best_cut = cut;
37                 best_weight = w[k];
39             }
41         } else {
43             rep(j, 0, N) w[j] += weights[k][j];
45             added[k] = true;
47         }
49     }
51     return {best_weight, best_cut};
53 }

```

3.2.5. Bipartite Minimum Cover

Requires: Dinic's Algorithm

```

1   // maximum independent set = all vertices not covered
2   // x : [0, n), y : [0, m]
3   struct Bipartite_vertex_cover {
4       Dinic D;
5       int n, m, s, t, x[maxn], y[maxn];
6       void make_edge(int x, int y) { D.make_edge(x, y + n, 1); }
7       int matching() {
8           int re = D.max_flow(s, t);
9           for (int i = 0; i < n; i++)
10              for (Dinic::edge &e : D.v[i])
11                 if (e.to != s && e.flow == 1) {
13                     x[i] = e.to - n, y[e.to - n] = i;
15                     break;
17                 }
19             return re;
21         }
22         // init() and matching() before use
23         void solve(vector<int> &vx, vector<int> &vy) {
24             bitset<maxn * 2 + 10> vis;
25             queue<int> q;
26             for (int i = 0; i < n; i++)
27                 if (x[i] == -1) q.push(i), vis[i] = 1;
28             while (!q.empty()) {
29                 int now = q.front();
30                 q.pop();
31                 if (now < n) {
32                     for (Dinic::edge &e : D.v[now])
33                         if (e.to != s && e.to - n != x[now] && !vis[e.to])
34                             vis[e.to] = 1, q.push(e.to);
35                 } else {
36                     if (!vis[y[now - n]])
37                         vis[y[now - n]] = 1, q.push(y[now - n]);
38                 }
39             }
40             for (int i = 0; i < n; i++)
41                 if (!vis[i]) vx.pb(i);
42             for (int i = 0; i < m; i++)
43                 if (vis[i + n]) vy.pb(i);
44         }
45         void init(int _n, int _m) {
46             n = _n, m = _m, s = n + m, t = s + 1;
47             for (int i = 0; i < n; i++)
48                 x[i] = -1, D.make_edge(s, i, 1);
49             for (int i = 0; i < m; i++)
50                 y[i] = -1, D.make_edge(i + n, t, 1);
51         }
52     };

```

3.2.6. Edmonds' Algorithm

```

1   struct Edmonds {
2       int n, T;
3       vector<vector<int>> g;
4       vector<int> pa, p, used, base;
5       Edmonds(int n)
6           : n(n), T(0), g(n), pa(n, -1), p(n), used(n),
7             base(n) {}
8       void add(int a, int b) {

```

```

11      g[a].push_back(b);
12      g[b].push_back(a);
13  }
14  int getBase(int i) {
15      while (i != base[i])
16          base[i] = base[base[i]], i = base[i];
17      return i;
18  }
19  vector<int> toJoin;
20  void mark_path(int v, int x, int b, vector<int> &path) {
21      for (; getBase(v) != b; v = p[x]) {
22          p[v] = x, x = pa[v];
23          toJoin.push_back(v);
24          toJoin.push_back(x);
25          if (!used[x]) used[x] = ++T, path.push_back(x);
26      }
27  }
28  bool go(int v) {
29      for (int x : g[v]) {
30          int b, bv = getBase(v), bx = getBase(x);
31          if (bv == bx) continue;
32          if (used[x]) {
33              vector<int> path;
34              toJoin.clear();
35              if (used[bx] < used[bv])
36                  mark_path(v, x, b = bx, path);
37              else mark_path(x, v, b = bv, path);
38              for (int z : toJoin) base[getBase(z)] = b;
39              for (int z : path)
40                  if (go(z)) return 1;
41          } else if (p[x] == -1) {
42              p[x] = v;
43              if (pa[x] == -1) {
44                  for (int y; x != -1; x = v)
45                      y = p[x], v = pa[y], pa[x] = y, pa[y] = x;
46                  return 1;
47              }
48              if (!used[pa[x]]) {
49                  used[pa[x]] = ++T;
50                  if (go(pa[x])) return 1;
51              }
52          }
53      }
54      return 0;
55  }
56  void init_dfs() {
57      for (int i = 0; i < n; i++)
58          used[i] = 0, p[i] = -1, base[i] = i;
59  }
60  bool dfs(int root) {
61      used[root] = ++T;
62      return go(root);
63  }
64  void match() {
65      int ans = 0;
66      for (int v = 0; v < n; v++)
67          for (int x : g[v])
68              if (pa[v] == -1 && pa[x] == -1) {
69                  pa[v] = x, pa[x] = v, ans++;
70                  break;
71              }
72      init_dfs();
73      for (int i = 0; i < n; i++)
74          if (pa[i] == -1 && dfs(i)) ans++, init_dfs();
75      cout << ans * 2 << "\n";
76      for (int i = 0; i < n; i++)
77          if (pa[i] > i)
78              cout << i + 1 << " " << pa[i] + 1 << "\n";
79  }
80  };

```

3.2.7. Minimum Weight Matching

```

1   struct Graph {
2       static const int MAXN = 105;
3       int n, e[MAXN][MAXN];
4       int match[MAXN], d[MAXN], onstk[MAXN];
5       vector<int> stk;
6       void init(int _n) {
7           n = _n;
8           for (int i = 0; i < n; i++)
9               for (int j = 0; j < n; j++)
10                  // change to appropriate infinity
11                  // if not complete graph
12                  e[i][j] = 0;
13       }
14       void add_edge(int u, int v, int w) {
15           e[u][v] = e[v][u] = w;
16       }
17       bool SPFA(int u) {
18           if (onstk[u]) return true;
19           stk.push_back(u);
20           onstk[u] = 1;

```



```

21     for (int v = 0; v < n; v++) {
22         if (u != v && match[u] != v && !onstk[v]) {
23             int m = match[v];
24             if (d[m] > d[u] - e[v][m] + e[u][v]) {
25                 d[m] = d[u] - e[v][m] + e[u][v];
26                 onstk[v] = 1;
27                 stk.push_back(v);
28                 if (SPFA(m)) return true;
29                 stk.pop_back();
30                 onstk[v] = 0;
31             }
32         }
33     }
34     onstk[u] = 0;
35     stk.pop_back();
36     return false;
37 }
38 int solve() {
39     for (int i = 0; i < n; i += 2) {
40         match[i] = i + 1;
41         match[i + 1] = i;
42     }
43     while (true) {
44         int found = 0;
45         for (int i = 0; i < n; i++) onstk[i] = d[i] = 0;
46         for (int i = 0; i < n; i++) {
47             stk.clear();
48             if (!onstk[i] && SPFA(i)) {
49                 found = 1;
50                 while (stk.size() >= 2) {
51                     int u = stk.back();
52                     stk.pop_back();
53                     int v = stk.back();
54                     stk.pop_back();
55                     match[u] = v;
56                     match[v] = u;
57                 }
58             }
59             if (!found) break;
60         }
61         int ret = 0;
62         for (int i = 0; i < n; i++) ret += e[i][match[i]];
63         ret /= 2;
64         return ret;
65     }
66 } graph;

```

```

45     if (order[girl_id][boy_id] <
46         order[girl_id][girl_current[girl_id]]) {
47         if (girl_current[girl_id] < n)
48             que.push(girl_current[girl_id]);
49         girl_current[girl_id] = boy_id;
50     } else {
51         que.push(boy_id);
52     }
53 }
54 }
55 }
56
57 int main() {
58     cin >> n;
59
60     for (int i = 0; i < n; i++) {
61         string p, t;
62         cin >> p;
63         male[p] = i;
64         bname[i] = p;
65         for (int j = 0; j < n; j++) {
66             cin >> t;
67             if (!female.count(t)) {
68                 gname[fit] = t;
69                 female[t] = fit++;
70             }
71             favor[i][j] = female[t];
72         }
73     }
74
75     for (int i = 0; i < n; i++) {
76         string p, t;
77         cin >> p;
78         for (int j = 0; j < n; j++) {
79             cin >> t;
80             order[female[p]][male[t]] = j;
81         }
82     }
83
84     initialize();
85     stable_marriage();
86
87     for (int i = 0; i < n; i++) {
88         cout << bname[i] << " "
89              << gname[favor[i][current[i] - 1]] << endl;
90     }
91 }

```

3.2.8. Stable Marriage

```

1 // normal stable marriage problem
2 /* input:
3 3
4 Albert Laura Nancy Marcy
5 Brad Marcy Nancy Laura
6 Chuck Laura Marcy Nancy
7 Laura Chuck Albert Brad
8 Marcy Albert Chuck Brad
9 Nancy Brad Albert Chuck
10 */
11
12 using namespace std;
13 const int MAXN = 505;
14
15 int n;
16 int favor[MAXN][MAXN]; // favor[boy_id][rank] = girl_id;
17 int order[MAXN][MAXN]; // order[girl_id][boy_id] = rank;
18 int current[MAXN]; // current[boy_id] = rank;
19 // boy_id will pursue current[boy_id] girl.
20 int girl_current[MAXN]; // girl[girl_id] = boy_id;
21
22 void initialize() {
23     for (int i = 0; i < n; i++) {
24         current[i] = 0;
25         girl_current[i] = n;
26         order[i][n] = n;
27     }
28 }
29
30 map<string, int> male, female;
31 string bname[MAXN], gname[MAXN];
32 int fit = 0;
33
34 void stable_marriage() {
35     queue<int> que;
36     for (int i = 0; i < n; i++) que.push(i);
37     while (!que.empty()) {
38         int boy_id = que.front();
39         que.pop();
40
41         int girl_id = favor[boy_id][current[boy_id]];
42         current[boy_id]++;

```

3.2.9. Kuhn-Munkres algorithm

```

1 // Maximum Weight Perfect Bipartite Matching
2 // Detect non-perfect-matching:
3 // 1. set all edge[i][j] as INF
4 // 2. if solve() >= INF, it is not perfect matching.
5
6 typedef long long ll;
7 struct KM {
8     static const int MAXN = 1050;
9     static const ll INF = 1LL << 60;
10     int n, match[MAXN], vx[MAXN], vy[MAXN];
11     ll edge[MAXN][MAXN], lx[MAXN], ly[MAXN], slack[MAXN];
12     void init(int _n) {
13         n = _n;
14         for (int i = 0; i < n; i++)
15             for (int j = 0; j < n; j++) edge[i][j] = 0;
16     }
17     void add_edge(int x, int y, ll w) { edge[x][y] = w; }
18     bool DFS(int x) {
19         vx[x] = 1;
20         for (int y = 0; y < n; y++) {
21             if (vy[y]) continue;
22             if (lx[x] + ly[y] > edge[x][y]) {
23                 slack[y] =
24                     min(slack[y], lx[x] + ly[y] - edge[x][y]);
25             } else {
26                 vy[y] = 1;
27                 if (match[y] == -1 || DFS(match[y])) {
28                     match[y] = x;
29                     return true;
30                 }
31             }
32         }
33         return false;
34     }
35     ll solve() {
36         fill(match, match + n, -1);
37         fill(lx, lx + n, -INF);
38         fill(ly, ly + n, 0);
39         for (int i = 0; i < n; i++)
40             for (int j = 0; j < n; j++)
41                 lx[i] = max(lx[i], edge[i][j]);
42         for (int i = 0; i < n; i++) {
43             fill(slack, slack + n, INF);
44             while (true) {

```

```

45     fill(vx, vx + n, 0);
46     fill(vy, vy + n, 0);
47     if (DFS(i)) break;
48     ll d = INF;
49     for (int j = 0; j < n; j++)
50         if (!vy[j]) d = min(d, slack[j]);
51     for (int j = 0; j < n; j++) {
52         if (vx[j]) lx[j] -= d;
53         if (vy[j]) ly[j] += d;
54         else slack[j] -= d;
55     }
56 }
57 }
58 ll res = 0;
59 for (int i = 0; i < n; i++) {
60     res += edge[match[i]][i];
61 }
62 return res;
63 }
64 } graph;

```

3.3. Shortest Path Faster Algorithm

```

1 struct SPFA {
2     static const int maxn = 1010, INF = 1e9;
3     int dis[maxn];
4     bitset<maxn> inq, inneg;
5     queue<int> q, tq;
6     vector<pii> v[maxn];
7     void make_edge(int s, int t, int w) {
8         v[s].emplace_back(t, w);
9     }
10    void dfs(int a) {
11        inneg[a] = 1;
12        for (pii i : v[a])
13            if (!inneg[i.F]) dfs(i.F);
14    }
15    bool solve(int n, int s) { // true if have neg-cycle
16        for (int i = 0; i <= n; i++) dis[i] = INF;
17        dis[s] = 0, q.push(s);
18        for (int i = 0; i < n; i++) {
19            inq.reset();
20            int now;
21            while (!q.empty()) {
22                now = q.front(), q.pop();
23                for (pii &i : v[now]) {
24                    if (dis[i.F] > dis[now] + i.S) {
25                        dis[i.F] = dis[now] + i.S;
26                        if (!inq[i.F]) tq.push(i.F), inq[i.F] = 1;
27                    }
28                }
29            }
30            q.swap(tq);
31        }
32        bool re = !q.empty();
33        inneg.reset();
34        while (!q.empty()) {
35            if (!inneg[q.front()]) dfs(q.front());
36            q.pop();
37        }
38        return re;
39    }
40    void reset(int n) {
41        for (int i = 0; i <= n; i++) v[i].clear();
42    }
43 };

```

3.4. Strongly Connected Components

```

1 struct TarjanScc {
2     int n, step;
3     vector<int> time, low, instk, stk;
4     vector<vector<int>> e, scc;
5     TarjanScc(int n)
6         : n(n), step(0), time(n), low(n), instk(n), e(n) {}
7     void add_edge(int u, int v) { e[u].push_back(v); }
8     void dfs(int x) {
9         time[x] = low[x] = ++step;
10        stk.push_back(x);
11        instk[x] = 1;
12        for (int y : e[x])
13            if (!time[y]) {
14                dfs(y);
15                low[x] = min(low[x], low[y]);
16            } else if (instk[y]) {
17                low[x] = min(low[x], time[y]);
18            }
19        if (time[x] == low[x]) {
20            scc.emplace_back();
21            for (int y = -1; y != x; ) {
22                y = stk.back();
23                stk.pop_back();
24                instk[y] = 0;

```

```

25        scc.back().push_back(y);
26    }
27 }
28 }
29 void solve() {
30     for (int i = 0; i < n; i++)
31         if (!time[i]) dfs(i);
32     reverse(scc.begin(), scc.end());
33     // scc in topological order
34 }
35 };

```

3.4.1. 2-Satisfiability

Requires: Strongly Connected Components

```

1
2
3 // 1 based, vertex in SCC = MAXN * 2
4 // (not i) is i + n
5 struct two_SAT {
6     int n, ans[MAXN];
7     SCC S;
8     void imply(int a, int b) { S.make_edge(a, b); }
9     bool solve(int _n) {
10         n = _n;
11         S.solve(n * 2);
12         for (int i = 1; i <= n; i++) {
13             if (S.scc[i] == S.scc[i + n]) return false;
14             ans[i] = (S.scc[i] < S.scc[i + n]);
15         }
16         return true;
17     }
18     void init(int _n) {
19         n = _n;
20         fill_n(ans, n + 1, 0);
21         S.init(n * 2);
22     }
23 } SAT;

```

3.5. Biconnected Components

3.5.1. Articulation Points

```

1 void dfs(int x, int p) {
2     tin[x] = low[x] = ++t;
3     int ch = 0;
4     for (auto u : g[x])
5         if (u.first != p) {
6             if (!ins[u.second])
7                 st.push(u.second), ins[u.second] = true;
8             if (tin[u.first]) {
9                 low[x] = min(low[x], tin[u.first]);
10                continue;
11            }
12            ++ch;
13            dfs(u.first, x);
14            low[x] = min(low[x], low[u.first]);
15            if (low[u.first] >= tin[x]) {
16                cut[x] = true;
17                ++sz;
18                while (true) {
19                    int e = st.top();
20                    st.pop();
21                    bcc[e] = sz;
22                    if (e == u.second) break;
23                }
24            }
25        }
26        if (ch == 1 && p == -1) cut[x] = false;
27 }

```

3.5.2. Bridges

```

1 // if there are multi-edges, then they are not bridges
2 void dfs(int x, int p) {
3     tin[x] = low[x] = ++t;
4     st.push(x);
5     for (auto u : g[x])
6         if (u.first != p) {
7             if (tin[u.first]) {
8                 low[x] = min(low[x], tin[u.first]);
9                 continue;
10            }
11            dfs(u.first, x);
12            low[x] = min(low[x], low[u.first]);
13            if (low[u.first] == tin[u.first]) br[u.second] = true;
14        }
15    if (tin[x] == low[x]) {
16        ++sz;
17        while (st.size()) {
18            int u = st.top();
19            st.pop();
20            bcc[u] = sz;

```

```

21     if (u == x) break;
22 }
23 }
}

```

3.6. Triconnected Components

```

1 // requires a union-find data structure
2
3 struct ThreeEdgeCC {
4     int V, ind;
5     vector<int> id, pre, post, low, deg, path;
6     vector<vector<int>> components;
7     UnionFind uf;
8     template <class Graph>
9     void dfs(const Graph &G, int v, int prev) {
10         pre[v] = ++ind;
11         for (int w : G[v])
12             if (w != v) {
13                 if (w == prev) {
14                     prev = -1;
15                     continue;
16                 }
17                 if (pre[w] != -1) {
18                     if (pre[w] < pre[v]) {
19                         deg[v]++;
20                         low[v] = min(low[v], pre[w]);
21                     } else {
22                         deg[v]--;
23                         int &u = path[v];
24                         for (; u != -1 && pre[u] <= pre[w] &&
25                             pre[w] <= post[u];) {
26                             uf.join(v, u);
27                             deg[v] += deg[u];
28                             u = path[u];
29                         }
30                     }
31                     continue;
32                 }
33                 dfs(G, w, v);
34                 if (path[w] == -1 && deg[w] <= 1) {
35                     deg[v] += deg[w];
36                     low[v] = min(low[v], low[w]);
37                     continue;
38                 }
39                 if (deg[w] == 0) w = path[w];
40                 if (low[v] > low[w]) {
41                     low[v] = min(low[v], low[w]);
42                     swap(w, path[v]);
43                 }
44                 for (; w != -1; w = path[w]) {
45                     uf.join(v, w);
46                     deg[v] += deg[w];
47                 }
48             }
49         post[v] = ind;
50     }
51     template <class Graph>
52     ThreeEdgeCC(const Graph &G)
53         : V(G.size()), ind(-1), id(V, -1), pre(V, -1),
54           post(V), low(V, INT_MAX), deg(V, 0), path(V, -1),
55           uf(V) {
56         for (int v = 0; v < V; v++)
57             if (pre[v] == -1) dfs(G, v, -1);
58         components.reserve(uf.cnt);
59         for (int v = 0; v < V; v++)
60             if (uf.find(v) == v) {
61                 id[v] = components.size();
62                 components.emplace_back(1, v);
63                 components.back().reserve(uf.getSize(v));
64             }
65         for (int v = 0; v < V; v++)
66             if (id[v] == -1)
67                 components[id[v] = id[uf.find(v)]] .push_back(v);
68     }
69 };

```

3.7. Centroid Decomposition

```

1 void get_center(int now) {
2     v[now] = true;
3     vtx.push_back(now);
4     sz[now] = 1;
5     mx[now] = 0;
6     for (int u : G[now])
7         if (!v[u]) {
8             get_center(u);
9             mx[now] = max(mx[now], sz[u]);
10            sz[now] += sz[u];
11        }
12 }

```

```

13 void get_dis(int now, int d, int len) {
14     dis[d][now] = cnt;
15     v[now] = true;
16     for (auto u : G[now])
17         if (!v[u.first]) { get_dis(u, d, len + u.second); }
18 }
19 void dfs(int now, int fa, int d) {
20     get_center(now);
21     int c = -1;
22     for (int i : vtx) {
23         if (max(mx[i], (int)vtx.size() - sz[i]) <=
24             (int)vtx.size() / 2)
25             c = i;
26         v[i] = false;
27     }
28     get_dis(c, d, 0);
29     for (int i : vtx) v[i] = false;
30     v[c] = true;
31     vtx.clear();
32     dep[c] = d;
33     p[c] = fa;
34     for (auto u : G[c])
35         if (u.first != fa && !v[u.first]) {
36             dfs(u.first, c, d + 1);
37         }
38 }

```

3.8. Minimum Mean Cycle

```

1
2 // d[i][j] == 0 if {i,j} !in E
3 long long d[1003][1003], dp[1003][1003];
4
5 pair<long long, long long> MMWC() {
6     memset(dp, 0x3f, sizeof(dp));
7     for (int i = 1; i <= n; ++i) dp[0][i] = 0;
8     for (int i = 1; i <= n; ++i) {
9         for (int j = 1; j <= n; ++j) {
10             for (int k = 1; k <= n; ++k) {
11                 dp[i][k] = min(dp[i - 1][j] + d[j][k], dp[i][k]);
12             }
13         }
14     }
15     long long au = 1ll << 31, ad = 1;
16     for (int i = 1; i <= n; ++i) {
17         if (dp[n][i] == 0x3f3f3f3f3f3f3f3f) continue;
18         long long u = 0, d = 1;
19         for (int j = n - 1; j >= 0; --j) {
20             if ((dp[n][i] - dp[j][i]) * d > u * (n - j)) {
21                 u = dp[n][i] - dp[j][i];
22                 d = n - j;
23             }
24         }
25         if (u * ad < au * d) au = u, ad = d;
26     }
27     long long g = __gcd(au, ad);
28     return make_pair(au / g, ad / g);
29 }

```

3.9. Directed MST

```

1 template <typename T> struct DMST {
2     T g[maxn][maxn], fw[maxn];
3     int n, fr[maxn];
4     bool vis[maxn], inc[maxn];
5     void clear() {
6         for (int i = 0; i < maxn; ++i) {
7             for (int j = 0; j < maxn; ++j) g[i][j] = inf;
8             vis[i] = inc[i] = false;
9         }
10    }
11    void addedge(int u, int v, T w) {
12        g[u][v] = min(g[u][v], w);
13    }
14    T operator()(int root, int _n) {
15        n = _n;
16        if (dfs(root) != n) return -1;
17        T ans = 0;
18        while (true) {
19            for (int i = 1; i <= n; ++i) fw[i] = inf, fr[i] = i;
20            for (int i = 1; i <= n; ++i)
21                if (!inc[i]) {
22                    for (int j = 1; j <= n; ++j) {
23                        if (!inc[j] && i != j && g[j][i] < fw[i]) {
24                            fw[i] = g[j][i];
25                            fr[i] = j;
26                        }
27                    }
28                }
29            int x = -1;
30            for (int i = 1; i <= n; ++i)
31                if (i != root && !inc[i]) {

```

```

33     int j = i, c = 0;
34     while (j != root && fr[j] != i && c <= n)
35         ++c, j = fr[j];
36     if (j == root || c > n) continue;
37     else {
38         x = i;
39         break;
40     }
41     if (!x) {
42         for (int i = 1; i <= n; ++i)
43             if (i != root && !inc[i]) ans += fw[i];
44         return ans;
45     }
46     int y = x;
47     for (int i = 1; i <= n; ++i) vis[i] = false;
48     do {
49         ans += fw[y];
50         y = fr[y];
51         vis[y] = inc[y] = true;
52     } while (y != x);
53     inc[x] = false;
54     for (int k = 1; k <= n; ++k)
55         if (vis[k]) {
56             for (int j = 1; j <= n; ++j)
57                 if (!vis[j]) {
58                     if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
59                     if (g[j][k] < inf &&
60                         g[j][k] - fw[k] < g[j][x])
61                         g[j][x] = g[j][k] - fw[k];
62                 }
63         }
64     }
65     return ans;
66 }
67 int dfs(int now) {
68     int r = 1;
69     vis[now] = true;
70     for (int i = 1; i <= n; ++i)
71         if (g[now][i] < inf && !vis[i]) r += dfs(i);
72     return r;
73 }
};

```

3.10. Maximum Clique

```

1 // source: KACTL
2
3 typedef vector<bitset<200>> vb;
4 struct Maxclique {
5     double limit = 0.025, pk = 0;
6     struct Vertex {
7         int i, d = 0;
8     };
9     typedef vector<Vertex> vv;
10    vb e;
11    vv V;
12    vector<vi> C;
13    vi qmax, q, S, old;
14    void init(vv &r) {
15        for (auto &v : r) v.d = 0;
16        for (auto &v : r)
17            for (auto j : r) v.d += e[v.i][j.i];
18        sort(all(r), [](auto a, auto b) { return a.d > b.d; });
19        int mxD = r[0].d;
20        rep(i, 0, sz(r)) r[i].d = min(i, mxD) + 1;
21    }
22    void expand(vv &R, int lev = 1) {
23        S[lev] += S[lev - 1] - old[lev];
24        old[lev] = S[lev - 1];
25        while (sz(R)) {
26            if (sz(q) + R.back().d <= sz(qmax)) return;
27            q.push_back(R.back().i);
28            vv T;
29            for (auto v : R)
30                if (e[R.back().i][v.i]) T.push_back({v.i});
31            if (sz(T)) {
32                if (S[lev]++ / ++pk < limit) init(T);
33                int j = 0, mxk = 1,
34                    mnk = max(sz(qmax) - sz(q) + 1, 1);
35                C[1].clear(), C[2].clear();
36                for (auto v : T) {
37                    int k = 1;
38                    auto f = [&](int i) { return e[v.i][i]; };
39                    while (any_of(all(C[k]), f)) k++;
40                    if (k > mxk) mxk = k, C[mxk + 1].clear();
41                    if (k < mnk) T[j++].i = v.i;
42                    C[k].push_back(v.i);
43                }
44                if (j > 0) T[j - 1].d = 0;
45                rep(k, mnk, mxk + 1) for (int i : C[k]) T[j].i = i,
46                    T[j++].d = k;
47                expand(T, lev + 1);
48            }
49        }
50    }
51    } else if (sz(q) > sz(qmax)) qmax = q;
52    q.pop_back(), R.pop_back();
53    }
54    }
55    vi maxClique() {
56        init(V), expand(V);
57        return qmax;
58    }
59    Maxclique(vb conn)
60        : e(conn), C(sz(e) + 1), S(sz(C)), old(S) {
61        rep(i, 0, sz(e)) V.push_back({i});
62    };

```

3.11. Dominator Tree

```

1 // idom[n] is the unique node that strictly dominates n but
2 // does not strictly dominate any other node that strictly
3 // dominates n. idom[n] = 0 if n is entry or the entry
4 // cannot reach n.
5 struct DominatorTree {
6     static const int MAXN = 200010;
7     int n, s;
8     vector<int> g[MAXN], pred[MAXN];
9     vector<int> cov[MAXN];
10    int dfn[MAXN], nfd[MAXN], ts;
11    int par[MAXN];
12    int sdom[MAXN], idom[MAXN];
13    int mom[MAXN], mn[MAXN];
14
15    inline bool cmp(int u, int v) { return dfn[u] < dfn[v]; }
16
17    int eval(int u) {
18        if (mom[u] == u) return u;
19        int res = eval(mom[u]);
20        if (cmp(sdom[mn[mom[u]]], sdom[mn[u]]))
21            mn[u] = mn[mom[u]];
22        return mom[u] = res;
23    }
24
25    void init(int _n, int _s) {
26        n = _n;
27        s = _s;
28        REP1(i, 1, n) {
29            g[i].clear();
30            pred[i].clear();
31            idom[i] = 0;
32        }
33    }
34    void add_edge(int u, int v) {
35        g[u].push_back(v);
36        pred[v].push_back(u);
37    }
38    void DFS(int u) {
39        ts++;
40        dfn[u] = ts;
41        nfd[ts] = u;
42        for (int v : g[u])
43            if (dfn[v] == 0) {
44                par[v] = u;
45                DFS(v);
46            }
47    }
48    void build() {
49        ts = 0;
50        REP1(i, 1, n) {
51            dfn[i] = nfd[i] = 0;
52            cov[i].clear();
53            mom[i] = mn[i] = sdom[i] = i;
54        }
55        DFS(s);
56        for (int i = ts; i >= 2; i--) {
57            int u = nfd[i];
58            if (u == 0) continue;
59            for (int v : pred[u])
60                if (dfn[v]) {
61                    eval(v);
62                    if (cmp(sdom[mn[v]], sdom[u]))
63                        sdom[u] = sdom[mn[v]];
64                }
65            cov[sdom[u]].push_back(u);
66            mom[u] = par[u];
67            for (int w : cov[par[u]]) {
68                eval(w);
69                if (cmp(sdom[mn[w]], par[u])) idom[w] = mn[w];
70                else idom[w] = par[u];
71            }
72            cov[par[u]].clear();
73        }
74        REP1(i, 2, ts) {
75            int u = nfd[i];
76            if (u == 0) continue;
77            if (idom[u] != sdom[u]) idom[u] = idom[idom[u]];
78        }
79    }

```

```

79 }
} dom;

```

3.12. Manhattan Distance MST

```

1
3 // returns [(dist, from, to), ...]
  // then do normal mst afterwards
5 typedef Point<int> P;
  vector<array<int, 3>> manhattanMST(vector<P> ps) {
7     vi id(sz(ps));
      iota(all(id), 0);
9     vector<array<int, 3>> edges;
      rep(k, 0, 4) {
11         sort(all(id), [&](int i, int j) {
13             return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;
14         });
15         map<int, int> sweep;
16         for (int i : id) {
17             for (auto it = sweep.lower_bound(-ps[i].y);
18                  it != sweep.end(); sweep.erase(it++)) {
19                 int j = it->second;
20                 P d = ps[i] - ps[j];
21                 if (d.y > d.x) break;
22                 edges.push_back({d.y + d.x, i, j});
23             }
24             sweep[-ps[i].y] = i;
25         }
26         for (P &p : ps)
27             if (k & 1) p.x = -p.x;
28             else swap(p.x, p.y);
29     }
    return edges;
}

```

4. Math

4.1. Number Theory

4.1.1. Mod Struct

A list of safe primes: 26003, 27767, 28319, 28979, 29243, 29759, 30467, 910927547, 919012223, 947326223, 990669467, 1007939579, 1019126699, 929760389146037459, 975500632317046523, 989312547895528379

NTT prime p	$p - 1$	primitive root
65537	$1 \ll 16$	3
998244353	$119 \ll 23$	3
2748779069441	$5 \ll 39$	3
1945555039024054273	$27 \ll 56$	5

Requires: Extended GCD

```

1
3 template <typename T> struct M {
4     static T MOD; // change to constexpr if already known
5     T v;
6     M(T x = 0) {
7         v = (-MOD <= x && x < MOD) ? x : x % MOD;
8         if (v < 0) v += MOD;
9     }
10    explicit operator T() const { return v; }
11    bool operator==(const M &b) const { return v == b.v; }
12    bool operator!=(const M &b) const { return v != b.v; }
13    M operator-() { return M(-v); }
14    M operator+(M b) { return M(v + b.v); }
15    M operator-(M b) { return M(v - b.v); }
16    M operator*(M b) { return M((__int128)v * b.v % MOD); }
17    M operator/(M b) { return *this * (b ^ (MOD - 2)); }
18    // change above implementation to this if MOD is not prime
19    M inv() {
20        auto [p, _, g] = extgcd(v, MOD);
21        return assert(g == 1), p;
22    }
23    friend M operator^(M a, ll b) {
24        M ans(1);
25        for (; b >= 1, a *= a)
26            if (b & 1) ans *= a;
27        return ans;
28    }
29    friend M &operator+=(M &a, M b) { return a = a + b; }
30    friend M &operator-=(M &a, M b) { return a = a - b; }
31    friend M &operator*=(M &a, M b) { return a = a * b; }
32    friend M &operator/=(M &a, M b) { return a = a / b; }
33 };
34 using Mod = M<int>;
35 template <> int Mod::MOD = 1'000'000'007;
36 int &MOD = Mod::MOD;

```

4.1.2. Miller-Rabin

Requires: Mod Struct

```

1
3 // checks if Mod::MOD is prime
4 bool is_prime() {
5     if (MOD < 2 || MOD % 2 == 0) return MOD == 2;
6     Mod A[] = {2, 7, 61}; // for int values (< 2^31)
7     // ll: 2, 325, 9375, 28178, 450775, 9780504, 1795265022
8     int s = __builtin_ctzll(MOD - 1), i;
9     for (Mod a : A) {
10         Mod x = a ^ (MOD >> s);
11         for (i = 0; i < s && (x + 1).v > 2; i++) x *= x;
12         if (i && x != -1) return 0;
13     }
14     return 1;
15 }

```

4.1.3. Linear Sieve

```

1 constexpr ll MAXN = 1000000;
2 bitset<MAXN> is_prime;
3 vector<ll> primes;
4 ll mpf[MAXN], phi[MAXN], mu[MAXN];
5
6 void sieve() {
7     is_prime.set();
8     is_prime[1] = 0;
9     mu[1] = phi[1] = 1;
10    for (ll i = 2; i < MAXN; i++) {
11        if (is_prime[i]) {
12            mpf[i] = i;
13            primes.push_back(i);
14            phi[i] = i - 1;
15            mu[i] = -1;
16        }
17        for (ll p : primes) {
18            if (p > mpf[i] || i * p >= MAXN) break;
19            is_prime[i * p] = 0;
20            mpf[i * p] = p;
21            mu[i * p] = -mu[i];
22            if (i % p == 0)
23                phi[i * p] = phi[i] * p, mu[i * p] = 0;
24            else phi[i * p] = phi[i] * (p - 1);
25        }
26    }
27 }

```

4.1.4. Get Factors

Requires: Linear Sieve

```

1
3 vector<ll> all_factors(ll n) {
4     vector<ll> fac = {1};
5     while (n > 1) {
6         const ll p = mpf[n];
7         vector<ll> cur = {1};
8         while (n % p == 0) {
9             n /= p;
10            cur.push_back(cur.back() * p);
11        }
12        vector<ll> tmp;
13        for (auto x : fac)
14            for (auto y : cur) tmp.push_back(x * y);
15        tmp.swap(fac);
16    }
17    return fac;
}

```

4.1.5. Binary GCD

```

1 // returns the gcd of non-negative a, b
2 ull bin_gcd(ull a, ull b) {
3     if (!a || !b) return a + b;
4     int s = __builtin_ctzll(a | b);
5     a >>= __builtin_ctzll(a);
6     while (b) {
7         if ((b >>= __builtin_ctzll(b)) < a) swap(a, b);
8         b -= a;
9     }
10    return a << s;
11 }

```


4.1.6. Extended GCD

```

1 // returns (p, q, g): p * a + q * b == g == gcd(a, b)
  // g is not guaranteed to be positive when a < 0 or b < 0
2 tuple<ll, ll, ll> extgcd(ll a, ll b) {
3     ll s = 1, t = 0, u = 0, v = 1;
4     while (b) {
5         ll q = a / b;
6         swap(a -= q * b, b);
7         swap(s -= q * t, t);
8         swap(u -= q * v, v);
9     }
10    return {s, u, a};
11 }

```

4.1.7. Chinese Remainder Theorem

Requires: Extended GCD

```

1 // for 0 <= a < m, 0 <= b < n, returns the smallest x >= 0
2 // such that x % m == a and x % n == b
3 ll crt(ll a, ll m, ll b, ll n) {
4     if (n > m) swap(a, b), swap(m, n);
5     auto [x, y, g] = extgcd(m, n);
6     assert((a - b) % g == 0); // no solution
7     x = ((b - a) / g * x) % (n / g) * m + a;
8     return x < 0 ? x + m / g * n : x;
9 }

```

4.1.8. Baby-Step Giant-Step

Requires: Mod Struct

```

1 // returns x such that a ^ x = b where x \in [l, r]
2 ll bsgs(Mod a, Mod b, ll l = 0, ll r = MOD - 1) {
3     int m = sqrt(r - l + 1), i;
4     unordered_map<ll, ll> tb;
5     Mod d = (a ^ l) / b;
6     for (i = 0, d = (a ^ l) / b; i < m; i++, d *= a)
7         if (d == 1) return l + i;
8     else tb[(ll)d] = l + i;
9     Mod c = Mod(1) / (a ^ m);
10    for (i = 0, d = 1; i < m; i++, d *= c)
11        if (auto j = tb.find((ll)d); j != tb.end())
12            return j->second + i * m;
13    return assert(0), -1; // no solution
14 }

```

4.1.9. Pollard's Rho

```

1 ll f(ll x, ll mod) { return (x * x + 1) % mod; }
2 // n should be composite
3 ll pollard_rho(ll n) {
4     if (!(n & 1)) return 2;
5     while (1) {
6         ll y = 2, x = RNG() % (n - 1) + 1, res = 1;
7         for (int sz = 2; res == 1; sz *= 2) {
8             for (int i = 0; i < sz && res == 1; i++) {
9                 x = f(x, n);
10                res = __gcd(abs(x - y), n);
11            }
12            y = x;
13        }
14        if (res != 0 && res != n) return res;
15    }
16 }

```

4.1.10. Tonelli-Shanks Algorithm

Requires: Mod Struct

```

1 // returns (p, q, g): p * a + q * b == g == gcd(a, b)
2 // g is not guaranteed to be positive when a < 0 or b < 0
3 int legendre(Mod a) {
4     if (a == 0) return 0;
5     return (a ^ ((MOD - 1) / 2)) == 1 ? 1 : -1;
6 }
7 Mod sqrt(Mod a) {
8     assert(legendre(a) != -1); // no solution
9     ll p = MOD, s = p - 1;
10    if (a == 0) return 0;
11    if (p == 2) return 1;
12    if (p % 4 == 3) return a ^ ((p + 1) / 4);
13    int r, m;
14    for (r = 0; !(s & 1); r++) s >>= 1;
15    Mod n = 2;
16    while (legendre(n) != -1) n += 1;
17    Mod x = a ^ ((s + 1) / 2), b = a ^ s, g = n ^ s;
18    while (b != 1) {
19        Mod t = b;
20        for (m = 0; t != 1; m++) t *= b;
21        Mod gs = g ^ (1LL << (r - m - 1));

```

```

22    }
23    }
24    return x;
25 }
26 // to get sqrt(x) modulo p^k, where p is an odd prime:
27 // c = x^2 (mod p), c = x^2 (mod p^k), q = p^(k-1)
28 // x = x^q * c^((p^k-2q+1)/2) (mod p^k)

```

4.1.11. Chinese Sieve

```

1 const ll N = 1000000;
2 // f, g, h multiplicative, h = f (dirichlet convolution) g
3 ll pre_g(ll n);
4 ll pre_h(ll n);
5 // preprocessed prefix sum of f
6 ll pre_f[N];
7 // prefix sum of multiplicative function f
8 ll solve_f(ll n) {
9     static unordered_map<ll, ll> m;
10    if (n < N) return pre_f[n];
11    if (m.count(n)) return m[n];
12    ll ans = pre_h(n);
13    for (ll l = 2, r; l <= n; l = r + 1) {
14        r = n / (n / l);
15        ans -= (pre_g(r) - pre_g(l - 1)) * djs_f(n / l);
16    }
17    return m[n] = ans;
18 }

```

4.1.12. Rational Number Binary Search

```

1 struct QQ {
2     ll p, q;
3     QQ go(QQ b, ll d) { return {p + b.p * d, q + b.q * d}; }
4 };
5 bool pred(QQ);
6 // returns smallest p/q in [lo, hi] such that
7 // pred(p/q) is true, and 0 <= p, q <= N
8 QQ frac_bs(ll N) {
9     QQ lo{0, 1}, hi{1, 0};
10    if (pred(lo)) return lo;
11    assert(pred(hi));
12    bool dir = 1, L = 1, H = 1;
13    for (; L || H; dir = !dir) {
14        ll len = 0, step = 1;
15        for (int t = 0; t < 2 && (t ? step /= 2 : step *= 2);)
16            if (QQ mid = hi.go(lo, len + step);
17                mid.p > N || mid.q > N || dir ^ pred(mid))
18                t++;
19        else len += step;
20        swap(lo, hi = hi.go(lo, len));
21        (dir ? L : H) = !len;
22    }
23    return dir ? hi : lo;
24 }

```

4.1.13. Farey Sequence

```

1 // returns (e/f), where (a/b, c/d, e/f) are
2 // three consecutive terms in the order n farey sequence
3 // to start, call next_farey(n, 0, 1, 1, n)
4 pll next_farey(ll n, ll a, ll b, ll c, ll d) {
5     ll p = (n + b) / d;
6     return pll(p * c - a, p * d - b);
7 }

```

4.2. Combinatorics

4.2.1. Matroid Intersection

This template assumes 2 weighted matroids of the same type, and that removing an element is much more expensive than checking if one can be added. **Remember to change the implementation details.**

The ground set is $0, 1, \dots, n - 1$, where element i has weight $w[i]$. For the unweighted version, remove weights and change BF/SPFA to BFS.

```

1 constexpr int N = 100;
2 constexpr int INF = 1e9;
3 struct Matroid {
4     // represents an independent set
5     Matroid(bitset<N> S); // initialize from an independent set
6     bool can_add(int); // if adding will break independence
7     Matroid remove(int); // removing from the set
8 };
9 auto matroid_intersection(int n, const vector<int> &w) {
10    bitset<N> S;
11    for (int sz = 1; sz <= n; sz++) {
12        Matroid M1(S), M2(S);
13        vector<vector<pii>> e(n + 2);
14        for (int j = 0; j < n; j++)

```

```

17     if (!S[j]) {
18         if (M1.can_add(j)) e[n].emplace_back(j, -w[j]);
19         if (M2.can_add(j)) e[j].emplace_back(n + 1, 0);
20     }
21     for (int i = 0; i < n; i++)
22         if (S[i]) {
23             Matroid T1 = M1.remove(i), T2 = M2.remove(i);
24             for (int j = 0; j < n; j++)
25                 if (!S[j]) {
26                     if (T1.can_add(j)) e[i].emplace_back(j, -w[j]);
27                     if (T2.can_add(j)) e[j].emplace_back(i, w[i]);
28                 }
29         }
30
31     vector<pii> dis(n + 2, {INF, 0});
32     vector<int> prev(n + 2, -1);
33     dis[n] = {0, 0};
34     // change to SPFA for more speed, if necessary
35     bool upd = 1;
36     while (upd) {
37         upd = 0;
38         for (int u = 0; u < n + 2; u++)
39             for (auto [v, c] : e[u]) {
40                 pii x(dis[u].first + c, dis[u].second + 1);
41                 if (x < dis[v]) dis[v] = x, prev[v] = u, upd = 1;
42             }
43     }
44
45     if (dis[n + 1].first < INF)
46         for (int x = prev[n + 1]; x != n; x = prev[x])
47             S.flip(x);
48     else break;
49
50     // S is the max-weighted independent set with size sz
51 }
52 return S;
53 }

```

4.2.2. De Bruijn Sequence

```

1 int res[kN], aux[kN], a[kN], sz;
2 void Rec(int t, int p, int n, int k) {
3     if (t > n) {
4         if (n % p == 0)
5             for (int i = 1; i <= p; ++i) res[sz++] = aux[i];
6     } else {
7         aux[t] = aux[t - p];
8         Rec(t + 1, p, n, k);
9         for (aux[t] = aux[t - p] + 1; aux[t] < k; ++aux[t])
10             Rec(t + 1, t, n, k);
11     }
12 }
13 int DeBruijn(int k, int n) {
14     // return cyclic string of length k^n such that every
15     // string of length n using k character appears as a
16     // substring.
17     if (k == 1) return res[0] = 0, 1;
18     fill(aux, aux + k * n, 0);
19     return sz = 0, Rec(1, 1, n, k), sz;
20 }

```

4.2.3. Multinomial

```

1 // ways to permute v[i]
2 ll multinomial(vi &v) {
3     ll c = 1, m = v.empty() ? 1 : v[0];
4     for (int i = 1; i < v.size(); i++)
5         for (int j = 0; j < v[i]; j++) c = c * ++m / (j + 1);
6     return c;
7 }

```

4.3. Algebra

4.3.1. Formal Power Series

```

1
2
3 template <typename mint>
4 struct FormalPowerSeries : vector<mint> {
5     using vector<mint>::vector;
6     using FPS = FormalPowerSeries;
7
8     FPS &operator+=(const FPS &r) {
9         if (r.size() > this->size()) this->resize(r.size());
10        for (int i = 0; i < (int)r.size(); i++)
11            (*this)[i] += r[i];
12        return *this;
13    }
14
15    FPS &operator+=(const mint &r) {

```

```

16        if (this->empty()) this->resize(1);
17        (*this)[0] += r;
18        return *this;
19    }
20
21    FPS &operator-=(const FPS &r) {
22        if (r.size() > this->size()) this->resize(r.size());
23        for (int i = 0; i < (int)r.size(); i++)
24            (*this)[i] -= r[i];
25        return *this;
26    }
27
28    FPS &operator-=(const mint &r) {
29        if (this->empty()) this->resize(1);
30        (*this)[0] -= r;
31        return *this;
32    }
33
34    FPS &operator*=(const mint &v) {
35        for (int k = 0; k < (int)this->size(); k++)
36            (*this)[k] *= v;
37        return *this;
38    }
39
40    FPS &operator/=(const FPS &r) {
41        if (this->size() < r.size()) {
42            this->clear();
43            return *this;
44        }
45        int n = this->size() - r.size() + 1;
46        if ((int)r.size() <= 64) {
47            FPS f(*this), g(r);
48            g.shrink();
49            mint coeff = g.back().inverse();
50            for (auto &x : g) x *= coeff;
51            int deg = (int)f.size() - (int)g.size() + 1;
52            int gs = g.size();
53            FPS quo(deg);
54            for (int i = deg - 1; i >= 0; i--) {
55                quo[i] = f[i + gs - 1];
56                for (int j = 0; j < gs; j++)
57                    f[i + j] -= quo[i] * g[j];
58            }
59            *this = quo * coeff;
60            this->resize(n, mint(0));
61            return *this;
62        }
63        return *this = ((*this).rev().pre(n) * r.rev().inv(n))
64            .pre(n)
65            .rev();
66    }
67
68    FPS &operator%=(const FPS &r) {
69        *this -= *this / r * r;
70        shrink();
71        return *this;
72    }
73
74    FPS operator+(const FPS &r) const {
75        return FPS(*this) += r;
76    }
77
78    FPS operator+(const mint &v) const {
79        return FPS(*this) += v;
80    }
81
82    FPS operator-(const FPS &r) const {
83        return FPS(*this) -= r;
84    }
85
86    FPS operator-(const mint &v) const {
87        return FPS(*this) -= v;
88    }
89
90    FPS operator*(const FPS &r) const {
91        return FPS(*this) *= r;
92    }
93
94    FPS operator*(const mint &v) const {
95        return FPS(*this) *= v;
96    }
97
98    FPS operator/(const FPS &r) const {
99        return FPS(*this) /= r;
100    }
101
102    FPS operator/(const FPS &r) const {
103        return FPS(*this) /= r;
104    }
105
106    FPS operator%(const FPS &r) const {
107        return FPS(*this) %= r;
108    }
109
110    FPS operator-() const {
111        FPS ret(this->size());
112        for (int i = 0; i < (int)this->size(); i++)
113            ret[i] = -(*this)[i];
114        return ret;
115    }
116
117    void shrink() {
118        while (this->size() && this->back() == mint(0))
119            this->pop_back();
120    }

```

```

111 FPS rev() const {
112     FPS ret(*this);
113     reverse(begin(ret), end(ret));
114     return ret;
115 }
117 FPS dot(FPS r) const {
118     FPS ret(min(this->size(), r.size()));
119     for (int i = 0; i < (int)ret.size(); i++)
120         ret[i] = (*this)[i] * r[i];
121     return ret;
122 }
123 FPS pre(int sz) const {
124     return FPS(begin(*this),
125               begin(*this) + min((int)this->size(), sz));
126 }
127 }
129 FPS operator>>(int sz) const {
130     if ((int)this->size() <= sz) return {};
131     FPS ret(*this);
132     ret.erase(ret.begin(), ret.begin() + sz);
133     return ret;
134 }
135 FPS operator<<(int sz) const {
136     FPS ret(*this);
137     ret.insert(ret.begin(), sz, mint(0));
138     return ret;
139 }
141 FPS diff() const {
142     const int n = (int)this->size();
143     FPS ret(max(0, n - 1));
144     mint one(1), coeff(1);
145     for (int i = 1; i < n; i++) {
146         ret[i - 1] = (*this)[i] * coeff;
147         coeff += one;
148     }
149     return ret;
150 }
151 }
153 FPS integral() const {
154     const int n = (int)this->size();
155     FPS ret(n + 1);
156     ret[0] = mint(0);
157     if (n > 0) ret[1] = mint(1);
158     auto mod = mint::get_mod();
159     for (int i = 2; i <= n; i++)
160         ret[i] = (-ret[mod % i]) * (mod / i);
161     for (int i = 1; i < n; i++) ret[i + 1] *= (*this)[i];
162     return ret;
163 }
165 mint eval(mint x) const {
166     mint r = 0, w = 1;
167     for (auto &v : *this) r += w * v, w *= x;
168     return r;
169 }
171 FPS log(int deg = -1) const {
172     assert((*this)[0] == mint(1));
173     if (deg == -1) deg = (int)this->size();
174     return (this->diff() * this->inv(deg))
175         .pre(deg - 1)
176         .integral();
177 }
179 FPS pow(int64_t k, int deg = -1) const {
180     const int n = (int)this->size();
181     if (deg == -1) deg = n;
182     for (int i = 0; i < n; i++) {
183         if ((*this)[i] != mint(0)) {
184             if (i * k > deg) return FPS(deg, mint(0));
185             mint rev = mint(1) / (*this)[i];
186             FPS ret =
187                 (((*this * rev) >> i).log(deg) * k).exp(deg) *
188                 ((*this)[i].pow(k));
189             ret = (ret << (i * k)).pre(deg);
190             if ((int)ret.size() < deg) ret.resize(deg, mint(0));
191             return ret;
192         }
193     }
194     return FPS(deg, mint(0));
195 }
197 static void *ntt_ptr;
198 static void set_fft();
199 FPS &operator*=(const FPS &r);
200 void ntt();
201 void intt();
202 void ntt_doubling();
203 static int ntt_pr();
204 FPS inv(int deg = -1) const;

```

```

205 FPS exp(int deg = -1) const;
206 };
207 template <typename mint>
void *FormalPowerSeries<mint>::ntt_ptr = nullptr;

```

4.4. Theorems

4.4.1. Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G , where $L_{ii} = d(i)$, $L_{ij} = -c$ where c is the number of edge (i, j) in G .

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

4.4.2. Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if $i < j$ and $(i, j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{\text{rank}(D)}{2}$ is the maximum matching on G .

4.4.3. Cayley's Formula

- Given a degree sequence d_1, d_2, \dots, d_n for each *labeled* vertices, there are

$$\frac{(n-2)!}{(d_1-1)!(d_2-1)! \cdots (d_n-1)!}$$

spanning trees.

- Let $T_{n,k}$ be the number of *labeled* forests on n vertices with k components, such that vertex $1, 2, \dots, k$ belong to different components. Then $T_{n,k} = kn^{n-k-1}$.

4.4.4. Erdős–Gallai Theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \dots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1 + d_2 + \dots + d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$$

holds for all $1 \leq k \leq n$.

4.4.5. Burnside's Lemma

Let X be a set and G be a group that acts on X . For $g \in G$, denote by X^g the elements fixed by g :

$$X^g = \{x \in X \mid gx \in X\}$$

Then

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

5. Numeric

5.1. Barrett Reduction

```

1 using ull = unsigned long long;
2 using ul = uint128_t;
3 // very fast calculation of a % m
4 struct reduction {
5     const ull m, d;
6     explicit reduction(ull m) : m(m), d(((ull)1 << 64) / m) {}
7     inline ull operator()(ull a) const {
8         ull q = (ull)((ull)d * a) >> 64;
9         return (a - q * m) >= m ? a - m : a;
10    }
11 };

```

5.2. Long Long Multiplication

```

1 using ull = unsigned long long;
2 using ll = long long;
3 using ld = long double;
4 // returns a * b % M where a, b < M < 2**63
5 ull mult(ull a, ull b, ull M) {
6     ll ret = a * b - M * ull(ld(a) * ld(b) / ld(M));
7     return ret + M * (ret < 0) - M * (ret >= (ll)M);
8 }

```

5.3. Fast Fourier Transform

```

1 template <typename T>
2 void fft(int n, vector<T> &a, vector<T> &rt, bool inv) {
3     vector<int> br(n);
4     for (int i = 1; i < n; i++) {
5         br[i] = (i & 1) ? br[i - 1] + n / 2 : br[i / 2] / 2;
6         if (br[i] > i) swap(a[i], a[br[i]]);
7     }
8     for (int len = 2; len <= n; len *= 2)
9         for (int i = 0; i < n; i += len)
10             for (int j = 0; j < len / 2; j++) {
11                 int pos = n / len * (inv ? len - j : j);
12                 T u = a[i + j], v = a[i + j + len / 2] * rt[pos];
13                 a[i + j] = u + v, a[i + j + len / 2] = u - v;
14             }
15     if (T minv = T(1) / T(n); inv)
16         for (T &x : a) x *= minv;
17 }

```

Requires: Mod Struct

```

1
2
3 void ntt(vector<Mod> &a, bool inv, Mod primitive_root) {
4     int n = a.size();
5     Mod root = primitive_root ^ (MOD - 1) / n;
6     vector<Mod> rt(n + 1, 1);
7     for (int i = 0; i < n; i++) rt[i + 1] = rt[i] * root;
8     fft(n, a, rt, inv);
9 }
10
11 void fft(vector<complex<double>> &a, bool inv) {
12     int n = a.size();
13     vector<complex<double>> rt(n + 1);
14     double arg = acos(-1) * 2 / n;
15     for (int i = 0; i <= n; i++)
16         rt[i] = {cos(arg * i), sin(arg * i)};
17     fft_(n, a, rt, inv);
18 }

```

5.4. Fast Walsh-Hadamard Transform

Requires: Mod Struct

```

1
2
3 void fwht(vector<Mod> &a, bool inv) {
4     int n = a.size();
5     for (int d = 1; d < n; d <= 1)
6         for (int m = 0; m < n; m++)
7             if (!(m & d)) {
8                 inv ? a[m] -= a[m | d] : a[m] += a[m | d]; // AND
9                 inv ? a[m | d] -= a[m] : a[m | d] += a[m]; // OR
10                 Mod x = a[m], y = a[m | d]; // XOR
11                 a[m] = x + y, a[m | d] = x - y; // XOR
12             }
13     if (Mod iv = Mod(1) / n; inv) // XOR
14         for (Mod &i : a) i *= iv; // XOR
15 }

```

5.5. Subset Convolution

Requires: Mod Struct

```

1 #pragma GCC target("popcnt")
2 #include <immintrin.h>
3
4 void fwht(int n, vector<vector<Mod>> &a, bool inv) {
5     for (int h = 0; h < n; h++)
6         for (int i = 0; i < (1 << n); i++)
7             if (!(i & (1 << h)))
8                 for (int k = 0; k <= n; k++)
9                     inv ? a[i | (1 << h)][k] -= a[i][k] :
10                         : a[i | (1 << h)][k] += a[i][k];
11 }
12
13 // c[k] = sum(popcnt(i & j) == sz && i | j == k) a[i] * b[j]
14 vector<Mod> subset_convolution(int n, int sz,
15                               const vector<Mod> &a_,
16                               const vector<Mod> &b_) {
17     int len = n + sz + 1, N = 1 << n;
18     vector<vector<Mod>> a(1 << n, vector<Mod>(len, 0)), b = a;
19     for (int i = 0; i < N; i++)
20         a[i][_mm_popcnt_u64(i)] = a_[i],
21         b[i][_mm_popcnt_u64(i)] = b_[i];
22     fwht(n, a, 0), fwht(n, b, 0);
23     for (int i = 0; i < N; i++) {
24         vector<Mod> tmp(len);
25         for (int j = 0; j < len; j++)
26             for (int k = 0; k <= j; k++)
27                 tmp[j] += a[i][k] * b[i][j - k];
28         a[i] = tmp;
29     }
30     fwht(n, a, 1);
31     vector<Mod> c(N);
32     for (int i = 0; i < N; i++)
33         c[i] = a[i][_mm_popcnt_u64(i) + sz];
34     return c;
35 }

```

5.6. Linear Recurrences

5.6.1. Berlekamp-Massey Algorithm

```

1 template <typename T>
2 vector<T> berlekamp_massey(const vector<T> &s) {
3     int n = s.size(), l = 0, m = 1;
4     vector<T> r(n), p(n);
5     r[0] = p[0] = 1;
6     T b = 1, d = 0;
7     for (int i = 0; i < n; i++, m++, d = 0) {
8         for (int j = 0; j <= l; j++) d += r[j] * s[i - j];
9         if ((d /= b) == 0) continue; // change if T is float
10         auto t = r;
11         for (int j = m; j < n; j++) r[j] -= d * p[j - m];
12         if (l * 2 <= i) l = i + 1 - l, b *= d, m = 0, p = t;
13     }
14     return r.resize(l + 1), reverse(r.begin(), r.end()), r;
15 }

```

5.6.2. Linear Recurrence Calculation

```

1 template <typename T> struct lin_rec {
2     using poly = vector<T>;
3     poly mul(poly a, poly b, poly m) {
4         int n = m.size();
5         poly r(n);
6         for (int i = n - 1; i >= 0; i--) {
7             r.insert(r.begin(), 0), r.pop_back();
8             T c = r[n - 1] + a[n - 1] * b[i];
9             // c /= m[n - 1]; if m is not monic
10             for (int j = 0; j < n; j++)
11                 r[j] += a[j] * b[i] - c * m[j];
12         }
13         return r;
14     }
15     poly pow(poly p, ll k, poly m) {
16         poly r(m.size());
17         r[0] = 1;
18         for (; k >= 1; p = mul(p, p, m))
19             if (k & 1) r = mul(r, p, m);
20         return r;
21     }
22     T calc(poly t, poly r, ll k) {
23         int n = r.size();
24         poly p(n);
25         p[1] = 1;
26         poly q = pow(p, k, r);
27         T ans = 0;
28         for (int i = 0; i < n; i++) ans += t[i] * q[i];
29         return ans;
30     }
31 };

```

5.7. Matrices

5.7.1. Determinant

Requires: Mod Struct

```

1
2
3 Mod det(vector<vector<Mod>> a) {
4     int n = a.size();
5     Mod ans = 1;
6     for (int i = 0; i < n; i++) {
7         int b = i;
8         for (int j = i + 1; j < n; j++)
9             if (a[j][i] != 0) {
10                 b = j;
11                 break;
12             }
13         if (i != b) swap(a[i], a[b]), ans = -ans;
14         ans *= a[i][i];
15         if (ans == 0) return 0;
16         for (int j = i + 1; j < n; j++) {
17             Mod v = a[j][i] / a[i][i];
18             if (v != 0)
19                 for (int k = i + 1; k < n; k++)
20                     a[j][k] -= v * a[i][k];
21         }
22     }
23     return ans;
24 }

```

```

1 double det(vector<vector<double>> a) {
2     int n = a.size();
3     double ans = 1;
4     for (int i = 0; i < n; i++) {
5         int b = i;
6         for (int j = i + 1; j < n; j++)
7             if (fabs(a[j][i]) > fabs(a[b][i])) b = j;
8         if (i != b) swap(a[i], a[b]), ans = -ans;
9         ans *= a[i][i];

```

```

    if (ans == 0) return 0;
    for (int j = i + 1; j < n; j++) {
        double v = a[j][i] / a[i][i];
        if (v != 0)
            for (int k = i + 1; k < n; k++)
                a[j][k] -= v * a[i][k];
    }
    return ans;
}

```

5.7.2. Inverse

```

1 // Returns rank.
2 // Result is stored in A unless singular (rank < n).
3 // For prime powers, repeatedly set
4 //  $A^{-1} = A^{-1} (2I - A A^{-1}) \pmod{p^k}$ 
5 // where  $A^{-1}$  starts as the inverse of A mod p,
6 // and k is doubled in each step.
7
8 int matInv(vector<vector<double>> &A) {
9     int n = sz(A);
10    vi col(n);
11    vector<vector<double>> tmp(n, vector<double>(n));
12    rep(i, 0, n) tmp[i][i] = 1, col[i] = i;
13
14    rep(i, 0, n) {
15        int r = i, c = i;
16        rep(j, i, n)
17            rep(k, i, n) if (fabs(A[j][k]) > fabs(A[r][c])) r = j, c = k;
18
19        if (fabs(A[r][c]) < 1e-12) return i;
20        A[i].swap(A[r]);
21        tmp[i].swap(tmp[r]);
22        rep(j, 0, n) swap(A[j][i], A[j][c]),
23            swap(tmp[j][i], tmp[j][c]);
24        swap(col[i], col[c]);
25        double v = A[i][i];
26        rep(j, i + 1, n) {
27            double f = A[j][i] / v;
28            A[j][i] = 0;
29            rep(k, i + 1, n) A[j][k] -= f * A[i][k];
30            rep(k, 0, n) tmp[j][k] -= f * tmp[i][k];
31        }
32        rep(j, i + 1, n) A[i][j] /= v;
33        rep(j, 0, n) tmp[i][j] /= v;
34        A[i][i] = 1;
35    }
36
37    for (int i = n - 1; i > 0; --i) rep(j, 0, i) {
38        double v = A[j][i];
39        rep(k, 0, n) tmp[j][k] -= v * tmp[i][k];
40    }
41
42    rep(i, 0, n) rep(j, 0, n) A[col[i]][col[j]] = tmp[i][j];
43    return n;
44 }
45
46 int matInv_mod(vector<vector<ll>> &A) {
47     int n = sz(A);
48     vi col(n);
49     vector<vector<ll>> tmp(n, vector<ll>(n));
50     rep(i, 0, n) tmp[i][i] = 1, col[i] = i;
51
52     rep(i, 0, n) {
53         int r = i, c = i;
54         rep(j, i, n) rep(k, i, n) if (A[j][k]) {
55             r = j;
56             c = k;
57             goto found;
58         }
59         return i;
60     }
61 found:
62     A[i].swap(A[r]);
63     tmp[i].swap(tmp[r]);
64     rep(j, 0, n) swap(A[j][i], A[j][c]),
65         swap(tmp[j][i], tmp[j][c]);
66     swap(col[i], col[c]);
67     ll v = modpow(A[i][i], mod - 2);
68     rep(j, i + 1, n) {
69         ll f = A[j][i] * v % mod;
70         A[j][i] = 0;
71         rep(k, i + 1, n) A[j][k] =
72             (A[j][k] - f * A[i][k]) % mod;
73         rep(k, 0, n) tmp[j][k] =
74             (tmp[j][k] - f * tmp[i][k]) % mod;
75     }
76     rep(j, i + 1, n) A[i][j] = A[i][j] * v % mod;
77     rep(j, 0, n) tmp[i][j] = tmp[i][j] * v % mod;
78     A[i][i] = 1;
79 }
80 }

```

```

83 for (int i = n - 1; i > 0; --i) rep(j, 0, i) {
84     ll v = A[j][i];
85     rep(k, 0, n) tmp[j][k] =
86         (tmp[j][k] - v * tmp[i][k]) % mod;
87 }
88
89 rep(i, 0, n) rep(j, 0, n) A[col[i]][col[j]] =
90     tmp[i][j] % mod + (tmp[i][j] < 0 ? mod : 0);
91 return n;
92 }

```

5.7.3. Characteristic Polynomial

```

1 // calculate det(a - xI)
2
3 template <typename T>
4 vector<T> CharacteristicPolynomial(vector<vector<T>> a) {
5     int N = a.size();
6
7     for (int j = 0; j < N - 2; j++) {
8         for (int i = j + 1; i < N; i++) {
9             if (a[i][j] != 0) {
10                swap(a[j + 1], a[i]);
11                for (int k = 0; k < N; k++)
12                    swap(a[k][j + 1], a[k][i]);
13                break;
14            }
15        }
16        if (a[j + 1][j] != 0) {
17            T inv = T(1) / a[j + 1][j];
18            for (int i = j + 2; i < N; i++) {
19                if (a[i][j] == 0) continue;
20                T coe = inv * a[i][j];
21                for (int l = j; l < N; l++)
22                    a[i][l] -= coe * a[j + 1][l];
23                for (int k = 0; k < N; k++)
24                    a[k][j + 1] += coe * a[k][i];
25            }
26        }
27    }
28
29    vector<vector<T>> p(N + 1);
30    p[0] = {T(1)};
31    for (int i = 1; i <= N; i++) {
32        p[i].resize(i + 1);
33        for (int j = 0; j < i; j++) {
34            p[i][j + 1] = p[i - 1][j];
35            p[i][j] += p[i - 1][j] * a[i - 1][i - 1];
36        }
37        T x = 1;
38        for (int m = 1; m < i; m++) {
39            x *= -a[i - m][i - m - 1];
40            T coe = x * a[i - m - 1][i - 1];
41            for (int j = 0; j < i - m; j++)
42                p[i][j] += coe * p[i - m - 1][j];
43        }
44    }
45    return p[N];
46 }

```

5.7.4. Solve Linear Equation

```

1
2 typedef vector<double> vd;
3 const double eps = 1e-12;
4
5 // solves for x: A * x = b
6
7 int solveLinear(vector<vd> &A, vd &b, vd &x) {
8     int n = sz(A), m = sz(x), rank = 0, br, bc;
9     if (n) assert(sz(A[0]) == m);
10    vi col(m);
11    iota(all(col), 0);
12
13    rep(i, 0, n) {
14        double v, bv = 0;
15        rep(r, i, n) rep(c, i, m) if ((v = fabs(A[r][c])) > bv)
16            br = r, bc = c, bv = v;
17        if (bv <= eps) {
18            rep(j, i, n) if (fabs(b[j]) > eps) return -1;
19            break;
20        }
21        swap(A[i], A[br]);
22        swap(b[i], b[br]);
23        swap(col[i], col[bc]);
24        rep(j, 0, n) swap(A[j][i], A[j][bc]);
25        bv = 1 / A[i][i];
26        rep(j, i + 1, n) {
27            double fac = A[j][i] * bv;

```



```

29     b[j] -= fac * b[i];
    rep(k, i + 1, m) A[j][k] -= fac * A[i][k];
31 }
    rank++;
33 }

35 x.assign(m, 0);
for (int i = rank; i--;) {
37     b[i] /= A[i][i];
    x[col[i]] = b[i];
39     rep(j, 0, i) b[j] -= A[j][i] * b[i];
41 }
return rank; // (multiple solutions if rank < m)
}

```

5.8. Polynomial Interpolation

```

1 // returns a, such that a[0]x^0 + a[1]x^1 + a[2]x^2 + ...
// passes through the given points
5 typedef vector<double> vd;
vd interpolate(vd x, vd y, int n) {
7     vd res(n), temp(n);
    rep(k, 0, n - 1) rep(i, k + 1, n) y[i] =
9     (y[i] - y[k]) / (x[i] - x[k]);
    double last = 0;
11    temp[0] = 1;
    rep(k, 0, n) rep(i, 0, n) {
13        res[i] += y[k] * temp[i];
        swap(last, temp[i]);
15        temp[i] -= last * x[k];
    }
17    return res;
}

```

5.9. Simplex Algorithm

```

1 // Two-phase simplex algorithm for solving linear programs
// of the form
3 //
//      maximize      c^T x
//      subject to    Ax <= b
//                  x >= 0
7 //
// INPUT: A -- an m x n matrix
//        b -- an m-dimensional vector
//        c -- an n-dimensional vector
11 //        x -- a vector where the optimal solution will be
//            stored
13 //
// OUTPUT: value of the optimal solution (infinity if
15 //        unbounded
//        above, nan if infeasible)
17 //
// To use this code, create an LPSolver object with A, b,
19 // and c as arguments. Then, call Solve(x).

21 typedef long double ld;
typedef vector<ld> vd;
typedef vector<vd> vvd;
typedef vector<int> vi;

25 const ld EPS = 1e-9;

27 struct LPSolver {
29     int m, n;
    vd B, N;
    vvd D;

33     LPSolver(const vvd &A, const vd &b, const vd &c)
        : m(b.size()), n(c.size()), N(n + 1), B(m),
        D(m + 2, vd(n + 2)) {
35         for (int i = 0; i < m; i++)
            for (int j = 0; j < n; j++) D[i][j] = A[i][j];
37         for (int i = 0; i < m; i++) {
            B[i] = n + i;
            D[i][n] = -1;
            D[i][n + 1] = b[i];
41         }
        for (int j = 0; j < n; j++) {
43             N[j] = j;
            D[m][j] = -c[j];
45         }
        N[n] = -1;
        D[m + 1][n] = 1;
49     }

51     void Pivot(int r, int s) {
        double inv = 1.0 / D[r][s];
53         for (int i = 0; i < m + 2; i++)
            if (i != r)
55             for (int j = 0; j < n + 2; j++)

```

```

        if (j != s) D[i][j] -= D[r][j] * D[i][s] * inv;
57     for (int j = 0; j < n + 2; j++)
        if (j != s) D[r][j] *= inv;
59     for (int i = 0; i < m + 2; i++)
        if (i != r) D[i][s] *= -inv;
61     D[r][s] = inv;
    swap(B[r], N[s]);
63 }

65 bool Simplex(int phase) {
    int x = phase == 1 ? m + 1 : m;
67     while (true) {
        int s = -1;
69         for (int j = 0; j <= n; j++) {
            if (phase == 2 && N[j] == -1) continue;
71             if (s == -1 || D[x][j] < D[x][s] ||
                D[x][j] == D[x][s] && N[j] < N[s])
                s = j;
73         }
        if (D[x][s] > -EPS) return true;
        int r = -1;
75         for (int i = 0; i < m; i++) {
            if (D[i][s] < EPS) continue;
77             if (r == -1 ||
                D[i][n + 1] / D[i][s] < D[r][n + 1] / D[r][s] ||
                (D[i][n + 1] / D[i][s]) ==
81                 (D[r][n + 1] / D[r][s]) &&
                B[i] < B[r])
                r = i;
83         }
        if (r == -1) return false;
85         Pivot(r, s);
87     }
89 }

91 ld Solve(vd &x) {
    int r = 0;
93     for (int i = 1; i < m; i++)
        if (D[i][n + 1] < D[r][n + 1]) r = i;
95     if (D[r][n + 1] < -EPS) {
        Pivot(r, n);
97         if (!Simplex(1) || D[m + 1][n + 1] < -EPS)
            return numeric_limits<ld>::infinity();
99         for (int i = 0; i < m; i++)
            if (B[i] == -1) {
101             int s = -1;
            for (int j = 0; j <= n; j++)
103                 if (s == -1 || D[i][j] < D[i][s] ||
                    D[i][j] == D[i][s] && N[j] < N[s])
                    s = j;
105             Pivot(i, s);
107         }
    }
109     if (!Simplex(2)) return numeric_limits<ld>::infinity();
    x = vd(n);
111     for (int i = 0; i < m; i++)
        if (B[i] < n) x[B[i]] = D[i][n + 1];
113     return D[m][n + 1];
115 };

117 int main() {
119     const int m = 4;
    const int n = 3;
121     ld _A[m][n] = {
        {6, -1, 0}, {-1, -5, 0}, {1, 5, 1}, {-1, -5, -1}};
123     ld _b[m] = {10, -4, 5, -5};
    ld _c[n] = {1, -1, 0};

125     vvd A(m);
    vd b(_b, _b + m);
127     vd c(_c, _c + n);
    for (int i = 0; i < m; i++) A[i] = vd(_A[i], _A[i] + n);
129

131     LPSolver solver(A, b, c);
    vd x;
133     ld value = solver.Solve(x);

135     cerr << "VALUE: " << value << endl; // VALUE: 1.29032
    cerr << "SOLUTION: "; // SOLUTION: 1.74194 0.451613 1
137     for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];
    cerr << endl;
139     return 0;
}

```

6. Geometry

6.1. Point

```

1 template <typename T> struct P {
    T x, y;
3     P(T x = 0, T y = 0) : x(x), y(y) {}

```

```

5  bool operator<(const P &p) const {
    return tie(x, y) < tie(p.x, p.y);
}
7  bool operator==(const P &p) const {
    return tie(x, y) == tie(p.x, p.y);
}
9  P operator-(const P &p) const { return {-x, -y}; }
11 P operator+(P p) const { return {x + p.x, y + p.y}; }
13 P operator-(P p) const { return {x - p.x, y - p.y}; }
15 P operator*(T d) const { return {x * d, y * d}; }
17 P operator/(T d) const { return {x / d, y / d}; }
19 T dist2() const { return x * x + y * y; }
21 double len() const { return sqrt(dist2()); }
23 P unit() const { return *this / len(); }
friend T dot(P a, P b) { return a.x * b.x + a.y * b.y; }
friend T cross(P a, P b) { return a.x * b.y - a.y * b.x; }
friend T cross(P a, P b, P o) {
    return cross(a - o, b - o);
}
};
using pt = P<ll>;

```

6.1.1. Quaternion

```

1  constexpr double PI = 3.141592653589793;
2  constexpr double EPS = 1e-7;
3  struct Q {
4      using T = double;
5      T x, y, z, r;
6      Q(T r = 0) : x(0), y(0), z(0), r(r) {}
7      Q(T x, T y, T z, T r = 0) : x(x), y(y), z(z), r(r) {}
8      friend bool operator==(const Q &a, const Q &b) {
9          return (a - b).abs2() <= EPS;
10     }
11     friend bool operator!=(const Q &a, const Q &b) {
12         return !(a == b);
13     }
14     Q operator-() const { return {-x, -y, -z, -r}; }
15     Q operator+(const Q &b) const {
16         return Q(x + b.x, y + b.y, z + b.z, r + b.r);
17     }
18     Q operator-(const Q &b) const {
19         return Q(x - b.x, y - b.y, z - b.z, r - b.r);
20     }
21     Q operator*(const T &t) const {
22         return Q(x * t, y * t, z * t, r * t);
23     }
24     Q operator*(const Q &b) const {
25         return Q(
26             r * b.x + x * b.r + y * b.z - z * b.y,
27             r * b.y - x * b.z + y * b.r + z * b.x,
28             r * b.z + x * b.y - y * b.x + z * b.r,
29             r * b.r - x * b.x - y * b.y - z * b.z);
30     }
31     Q operator/(const Q &b) const { return *this * b.inv(); }
32     T abs2() const { return r * r + x * x + y * y + z * z; }
33     T len() const { return sqrt(abs2()); }
34     Q conj() const { return Q(-x, -y, -z, r); }
35     Q unit() const { return *this * (1.0 / len()); }
36     Q inv() const { return conj() * (1.0 / abs2()); }
37     friend T dot(Q a, Q b) {
38         return a.x * b.x + a.y * b.y + a.z * b.z;
39     }
40     friend Q cross(Q a, Q b) {
41         return Q(
42             a.y * b.z - a.z * b.y, a.z * b.x - a.x * b.z,
43             a.x * b.y - a.y * b.x);
44     }
45     friend Q rotation_around(Q axis, T angle) {
46         return axis.unit() * sin(angle / 2) + cos(angle / 2);
47     }
48     Q rotated_around(Q axis, T angle) {
49         Q u = rotation_around(axis, angle);
50         return u * *this / u;
51     }
52     friend Q rotation_between(Q a, Q b) {
53         a = a.unit(), b = b.unit();
54         if (a == -b) {
55             // degenerate case
56             Q ortho = abs(a.y) > EPS ? cross(a, Q(1, 0, 0))
57                 : cross(a, Q(0, 1, 0));
58             return rotation_around(ortho, PI);
59         }
60         return (a * (a + b)).conj();
61     }
62 };

```

6.1.2. Spherical Coordinates

```

1  struct car_p {
2      double x, y, z;
3  };
4  struct sph_p {
5      double r, theta, phi;
6  };

```

```

7  sph_p conv(car_p p) {
8      double r = sqrt(p.x * p.x + p.y * p.y + p.z * p.z);
9      double theta = asin(p.y / r);
10     double phi = atan2(p.y, p.x);
11     return {r, theta, phi};
12 }
13 car_p conv(sph_p p) {
14     double x = p.r * cos(p.theta) * sin(p.phi);
15     double y = p.r * cos(p.theta) * cos(p.phi);
16     double z = p.r * sin(p.theta);
17     return {x, y, z};
18 }

```

6.2. Segments

```

1  // for non-collinear ABCD, if segments AB and CD intersect
2  bool intersects(pt a, pt b, pt c, pt d) {
3      if (cross(b, c, a) * cross(b, d, a) > 0) return false;
4      if (cross(d, a, c) * cross(d, b, c) > 0) return false;
5      return true;
6  }
7  // the intersection point of lines AB and CD
8  pt intersect(pt a, pt b, pt c, pt d) {
9      auto x = cross(b, c, a), y = cross(b, d, a);
10     if (x == y) {
11         // if(abs(x, y) < 1e-8) {
12         // is parallel
13     } else {
14         return d * (x / (x - y)) - c * (y / (x - y));
15     }
16 }

```

6.3. Convex Hull

```

1  // returns a convex hull in counterclockwise order
2  // for a non-strict one, change cross >= to >
3  vector<pt> convex_hull(vector<pt> p) {
4      sort(ALL(p));
5      if (p[0] == p.back()) return {p[0]};
6      int n = p.size(), t = 0;
7      vector<pt> h(n + 1);
8      for (int _ = 2, s = 0; _--; s = --t, reverse(ALL(p)))
9          for (pt i : p) {
10              while (t > s + 1 && cross(i, h[t - 1], h[t - 2]) >= 0)
11                  t--;
12              h[t++] = i;
13          }
14      return h.resize(t), h;
15 }

```

6.3.1. 3D Hull

```

1  typedef Point3D<double> P3;
2
3  struct PR {
4      void ins(int x) { (a == -1 ? a : b) = x; }
5      void rem(int x) { (a == x ? a : b) = -1; }
6      int cnt() { return (a != -1) + (b != -1); }
7      int a, b;
8  };
9
10 struct F {
11     P3 q;
12     int a, b, c;
13 };
14
15 vector<F> hull3d(const vector<P3> &A) {
16     assert(sz(A) >= 4);
17     vector<vector<PR>> E(sz(A), vector<PR>(sz(A), {-1, -1}));
18     #define E(x, y) E[f.x][f.y]
19     vector<F> FS;
20     auto mf = [&](int i, int j, int k, int l) {
21         P3 q = (A[j] - A[i]).cross((A[k] - A[i]));
22         if (q.dot(A[l]) > q.dot(A[i])) q = q * -1;
23         F f{q, i, j, k};
24         E(a, b).ins(k);
25         E(a, c).ins(j);
26         E(b, c).ins(i);
27         FS.push_back(f);
28     };
29     rep(i, 0, 4) rep(j, i + 1, 4) rep(k, j + 1, 4)
30         mf(i, j, k, 6 - i - j - k);
31
32     rep(i, 4, sz(A)) {
33         rep(j, 0, sz(FS)) {
34             F f = FS[j];
35             if (f.q.dot(A[i]) > f.q.dot(A[f.a])) {
36                 E(a, b).rem(f.c);
37                 E(a, c).rem(f.b);
38                 E(b, c).rem(f.a);
39             }
40         }
41     }
42 }

```

```

41     swap(FS[j--], FS.back());
42     FS.pop_back();
43 }
44 }
45 int nw = sz(FS);
46 rep(j, 0, nw) {
47     F f = FS[j];
48     #define C(a, b, c)
49     if (E(a, b).cnt() != 2) mf(f.a, f.b, i, f.c);
50     C(a, b, c);
51     C(a, c, b);
52     C(b, c, a);
53 }
54 }
55 for (F &it : FS)
56     if ((A[it.b] - A[it.a])
57         .cross(A[it.c] - A[it.a])
58         .dot(it.q) <= 0)
59         swap(it.c, it.b);
60 return FS;
61 };

```

6.4. Angular Sort

```

1 auto angle_cmp = [](const pt &a, const pt &b) {
2     auto btm = [] (const pt &a) {
3         return a.y < 0 || (a.y == 0 && a.x < 0);
4     };
5     return make_tuple(btm(a), a.y * b.x, abs2(a)) <
6         make_tuple(btm(b), a.x * b.y, abs2(b));
7 };
8 void angular_sort(vector<pt> &p) {
9     sort(p.begin(), p.end(), angle_cmp);
10 }

```

6.5. Convex Polygon Minkowski Sum

```

1 // O(n) convex polygon minkowski sum
2 // must be sorted and counterclockwise
3 vector<pt> minkowski_sum(vector<pt> p, vector<pt> q) {
4     auto diff = [] (vector<pt> &c) {
5         auto rcmp = [] (pt a, pt b) {
6             return pt{a.y, a.x} < pt{b.y, b.x};
7         };
8         rotate(c.begin(), min_element(ALL(c), rcmp), c.end());
9         c.push_back(c[0]);
10        vector<pt> ret;
11        for (int i = 1; i < c.size(); i++)
12            ret.push_back(c[i] - c[i - 1]);
13        return ret;
14    };
15    auto dp = diff(p), dq = diff(q);
16    pt cur = p[0] + q[0];
17    vector<pt> d(dp.size() + dq.size(), ret = {cur});
18    // include angle_cmp from angular-sort.cpp
19    merge(ALL(dp), ALL(dq), d.begin(), angle_cmp);
20    // optional: make ret strictly convex (UB if degenerate)
21    int now = 0;
22    for (int i = 1; i < d.size(); i++) {
23        if (cross(d[i], d[now]) == 0) d[now] = d[now] + d[i];
24        else d[++now] = d[i];
25    }
26    d.resize(now + 1);
27    // end optional part
28    for (pt v : d) ret.push_back(cur = cur + v);
29    return ret.pop_back(), ret;
30 }

```

6.6. Point In Polygon

```

1 bool on_segment(pt a, pt b, pt p) {
2     return cross(a, b, p) == 0 && dot((p - a), (p - b)) <= 0;
3 }
4 // p can be any polygon, but this is O(n)
5 bool inside(const vector<pt> &p, pt a) {
6     int cnt = 0, n = p.size();
7     for (int i = 0; i < n; i++) {
8         pt l = p[i], r = p[(i + 1) % n];
9         // change to return 0; for strict version
10        if (on_segment(l, r, a)) return 1;
11        cnt ^= ((a.y < l.y) - (a.y < r.y)) * cross(l, r, a) > 0;
12    }
13    return cnt;
14 }

```

6.6.1. Convex Version

```

1 // no preprocessing version
2 // p must be a strict convex hull, counterclockwise
3 // if point is inside or on border
4 bool is_inside(const vector<pt> &c, pt p) {
5     int n = c.size(), l = 1, r = n - 1;
6     if (cross(c[0], c[1], p) < 0) return false;

```

```

7     if (cross(c[n - 1], c[0], p) < 0) return false;
8     while (l < r - 1) {
9         int m = (l + r) / 2;
10        T a = cross(c[0], c[m], p);
11        if (a > 0) l = m;
12        else if (a < 0) r = m;
13        else return dot(c[0] - p, c[m] - p) <= 0;
14    }
15    if (l == r) return dot(c[0] - p, c[l] - p) <= 0;
16    else return cross(c[l], c[r], p) >= 0;
17 }
18
19 // with preprocessing version
20 vector<pt> vecs;
21 pt center;
22 // p must be a strict convex hull, counterclockwise
23 // BEWARE OF OVERFLOWS!!
24 void preprocess(vector<pt> p) {
25     for (auto &v : p) v = v * 3;
26     center = p[0] + p[1] + p[2];
27     center.x /= 3, center.y /= 3;
28     for (auto &v : p) v = v - center;
29     vecs = (angular_sort(p), p);
30 }
31 bool intersect_strict(pt a, pt b, pt c, pt d) {
32     if (cross(b, c, a) * cross(b, d, a) > 0) return false;
33     if (cross(d, a, c) * cross(d, b, c) >= 0) return false;
34     return true;
35 }
36 // if point is inside or on border
37 bool query(pt p) {
38     p = p * 3 - center;
39     auto pr = upper_bound(ALL(vecs), p, angle_cmp);
40     if (pr == vecs.end()) pr = vecs.begin();
41     auto pl = (pr == vecs.begin()) ? vecs.back() : *(pr - 1);
42     return !intersect_strict({0, 0}, p, pl, *pr);
43 }

```

6.6.2. Offline Multiple Points Version

Requires: Point, GNU PBDS

```

1
2
3
4
5 using Double = __float128;
6 using Point = pt<Double, Double>;
7
8
9 int n, m;
10 vector<Point> poly;
11 vector<Point> query;
12 vector<int> ans;
13
14 struct Segment {
15     Point a, b;
16     int id;
17 };
18 vector<Segment> segs;
19
20 Double Xnow;
21 inline Double get_y(const Segment &u, Double xnow = Xnow) {
22     const Point &a = u.a;
23     const Point &b = u.b;
24     return (a.y * (b.x - xnow) + b.y * (xnow - a.x)) /
25         (b.x - a.x);
26 }
27 bool operator<(Segment u, Segment v) {
28     Double yu = get_y(u);
29     Double yv = get_y(v);
30     if (yu != yv) return yu < yv;
31     return u.id < v.id;
32 }
33 ordered_map<Segment> st;
34
35 struct Event {
36     int type; // +1 insert seg, -1 remove seg, 0 query
37     Double x, y;
38     int id;
39 };
40 bool operator<(Event a, Event b) {
41     if (a.x != b.x) return a.x < b.x;
42     if (a.type != b.type) return a.type < b.type;
43     return a.y < b.y;
44 }
45 vector<Event> events;
46
47 void solve() {
48     set<Double> xs;
49     set<Point> ps;
50     for (int i = 0; i < n; i++) {
51         xs.insert(poly[i].x);
52         ps.insert(poly[i]);

```

```

53 }
54 for (int i = 0; i < n; i++) {
55     Segment s{poly[i], poly[(i + 1) % n], i};
56     if (s.a.x > s.b.x ||
57         (s.a.x == s.b.x && s.a.y > s.b.y)) {
58         swap(s.a, s.b);
59     }
60     segs.push_back(s);
61
62     if (s.a.x != s.b.x) {
63         events.push_back({+1, s.a.x + 0.2, s.a.y, i});
64         events.push_back({-1, s.b.x - 0.2, s.b.y, i});
65     }
66 }
67 for (int i = 0; i < m; i++) {
68     events.push_back({0, query[i].x, query[i].y, i});
69 }
70 sort(events.begin(), events.end());
71 int cnt = 0;
72 for (Event e : events) {
73     int i = e.id;
74     Xnow = e.x;
75     if (e.type == 0) {
76         Double x = e.x;
77         Double y = e.y;
78         Segment tmp = {{x - 1, y}, {x + 1, y}, -1};
79         auto it = st.lower_bound(tmp);
80
81         if (ps.count(query[i]) > 0) {
82             ans[i] = 0;
83         } else if (xs.count(x) > 0) {
84             ans[i] = -2;
85         } else if (it != st.end() &&
86             get_y(*it) == get_y(tmp)) {
87             ans[i] = 0;
88         } else if (it != st.begin() &&
89             get_y(*prev(it)) == get_y(tmp)) {
90             ans[i] = 0;
91         } else {
92             int rk = st.order_of_key(tmp);
93             if (rk % 2 == 1) {
94                 ans[i] = 1;
95             } else {
96                 ans[i] = -1;
97             }
98         }
99     } else if (e.type == 1) {
100         st.insert(segs[i]);
101         assert((int)st.size() == ++cnt);
102     } else if (e.type == -1) {
103         st.erase(segs[i]);
104         assert((int)st.size() == --cnt);
105     }
106 }
107 }

```

6.7. Closest Pair

```

1 vector<pll> p; // sort by x first!
2 bool cmpy(const pll &a, const pll &b) const {
3     return a.y < b.y;
4 }
5 ll sq(ll x) { return x * x; }
6 // returns (minimum dist)^2 in [l, r]
7 ll solve(int l, int r) {
8     if (r - l <= 1) return 1e18;
9     int m = (l + r) / 2;
10    ll mid = p[m].x, d = min(solve(l, m), solve(m, r));
11    auto pb = p.begin();
12    inplace_merge(pb + l, pb + m, pb + r, cmpy);
13    vector<pll> s;
14    for (int i = l; i < r; i++)
15        if (sq(p[i].x - mid) < d) s.push_back(p[i]);
16    for (int i = 0; i < s.size(); i++)
17        for (int j = i + 1;
18            j < s.size() && sq(s[j].y - s[i].y) < d; j++)
19            d = min(d, dis(s[i], s[j]));
20    return d;
21 }

```

6.8. Minimum Enclosing Circle

```

1
2
3 typedef Point<double> P;
4 double ccRadius(const P &A, const P &B, const P &C) {
5     return (B - A).dist() * (C - B).dist() * (A - C).dist() /
6         abs((B - A).cross(C - B)) / 2;
7 }
8 P ccCenter(const P &A, const P &B, const P &C) {
9     P b = C - A, c = B - A;
10    return A + (b * c.dist2() - c * b.dist2()).perp() /
11        b.cross(c) / 2;

```

```

}
pair<P, double> mec(vector<P> ps) {
    shuffle(all(ps), mt19937(time(0)));
    P o = ps[0];
    double r = 0, EPS = 1 + 1e-8;
    rep(i, 0, sz(ps)) if ((o - ps[i]).dist() > r * EPS) {
        o = ps[i], r = 0;
        rep(j, 0, i) if ((o - ps[j]).dist() > r * EPS) {
            o = (ps[i] + ps[j]) / 2;
            r = (o - ps[i]).dist();
            rep(k, 0, j) if ((o - ps[k]).dist() > r * EPS) {
                o = ccCenter(ps[i], ps[j], ps[k]);
                r = (o - ps[i]).dist();
            }
        }
    }
    return {o, r};
}

```

6.9. Delaunay Triangulation

```

1
2
3 typedef Point<ll> P;
4 typedef struct Quad *Q;
5 typedef __int128_t lll; // (can be ll if coords are < 2e4)
6 P arb(LLONG_MAX, LLONG_MAX); // not equal to any other point
7
8 struct Quad {
9     bool mark;
10    P o, rot;
11    P p;
12    P F() { return r()->p; }
13    Q r() { return rot->rot; }
14    Q prev() { return rot->o->rot; }
15    Q next() { return r()->prev(); }
16 };
17
18 bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
19     lll p2 = p.dist2(), A = a.dist2() - p2,
20         B = b.dist2() - p2, C = c.dist2() - p2;
21     return p.cross(a, b) * C + p.cross(b, c) * A +
22         p.cross(c, a) * B >
23         0;
24 }
25
26 Q makeEdge(P orig, P dest) {
27     Q q[] = {new Quad{0, 0, 0, orig}, new Quad{0, 0, 0, arb},
28         new Quad{0, 0, 0, dest}, new Quad{0, 0, 0, arb}};
29     rep(i, 0, 4) q[i]->o = q[(i + 1) & 3],
30         q[i]->rot = q[(i + 1) & 3];
31     return *q;
32 }
33
34 void splice(Q a, Q b) {
35     swap(a->o->rot->o, b->o->rot->o);
36     swap(a->o, b->o);
37 }
38
39 Q connect(Q a, Q b) {
40     Q q = makeEdge(a->F(), b->p);
41     splice(q, a->next());
42     splice(q->r(), b);
43     return q;
44 }
45
46 pair<Q, Q> rec(const vector<P> &s) {
47     if (sz(s) <= 3) {
48         Q a = makeEdge(s[0], s[1]),
49             b = makeEdge(s[1], s.back());
50         if (sz(s) == 2) return {a, a->r()};
51         splice(a->r(), b);
52         auto side = s[0].cross(s[1], s[2]);
53         Q c = side ? connect(b, a) : 0;
54         return {side < 0 ? c->r() : a, side < 0 ? c : b->r()};
55     }
56
57     #define H(e) e->F(), e->p
58     #define valid(e) (e->F().cross(H(base)) > 0)
59     Q A, B, ra, rb;
60     int half = sz(s) / 2;
61     tie(ra, A) = rec({all(s) - half});
62     tie(B, rb) = rec({sz(s) - half + all(s)});
63     while ((B->p.cross(H(A)) < 0 && (A = A->next())) ||
64         (A->p.cross(H(B)) > 0 && (B = B->r()->o)));
65     Q base = connect(B->r(), A);
66     if (A->p == ra->p) ra = base->r();
67     if (B->p == rb->p) rb = base;
68
69     #define DEL(e, init, dir)
70     Q e = init->dir;
71     if (valid(e))
72         while (circ(e->dir->F(), H(base), e->F())) {
73             Q t = e->dir;
74             splice(e, e->prev());
75             splice(e->r(), e->r()->prev());
76             e = t;

```

```

    }
75   for (;;) {
77       DEL(LC, base->r(), 0);
77       DEL(RC, base, prev());
77       if (!valid(LC) && !valid(RC)) break;
79       if (!valid(LC) || !valid(RC) && circ(H(RC), H(LC)))
81           base = connect(RC, base->r());
81       else base = connect(base->r(), LC->r());
83   }
83   return {ra, rb};
85 }

// returns [A_0, B_0, C_0, A_1, B_1, ...]
87 // where A_i, B_i, C_i are counter-clockwise triangles
vector<P> triangulate(vector<P> pts) {
89     sort(all(pts));
91     assert(unique(all(pts)) == pts.end());
91     if (sz(pts) < 2) return {};
91     Q e = rec(pts).first;
93     vector<Q> q = {e};
95     int qi = 0;
95     while (e->o->F().cross(e->F(), e->p) < 0) e = e->o;
97 #define ADD
97 {
99     Q c = e;
101     do {
101         c->mark = 1;
101         pts.push_back(c->p);
103         q.push_back(c->r());
103         c = c->next();
105     } while (c != e);
105 }
107 ADD;
107 pts.clear();
107 while (qi < sz(q))
109     if (!e = q[qi++])->mark) ADD;
111 return pts;
}

```

6.9.1. Slower Version

```

1
3 template <class P, class F>
4 void delaunay(vector<P> &ps, F trifun) {
5     if (sz(ps) == 3) {
6         int d = (ps[0].cross(ps[1], ps[2]) < 0);
7         trifun(0, 1 + d, 2 - d);
8     }
9     vector<P3> p3;
10    for (P p : ps) p3.emplace_back(p.x, p.y, p.dist2());
11    if (sz(ps) > 3)
12        for (auto t : hull3d(p3))
13            if ((p3[t.b] - p3[t.a])
14                .cross(p3[t.c] - p3[t.a])
15                .dot(P3(0, 0, 1)) < 0)
16                trifun(t.a, t.c, t.b);
17 }

```

6.10. Half Plane Intersection

```

1 struct Line {
2     Point P;
3     Vector v;
4     bool operator<(const Line &b) const {
5         return atan2(v.y, v.x) < atan2(b.v.y, b.v.x);
6     }
7 };
8 bool OnLeft(const Line &L, const Point &p) {
9     return Cross(L.v, p - L.P) > 0;
10 }
11 Point GetIntersection(Line a, Line b) {
12     Vector u = a.P - b.P;
13     Double t = Cross(b.v, u) / Cross(a.v, b.v);
14     return a.P + a.v * t;
15 }
16 int HalfplaneIntersection(Line *L, int n, Point *poly) {
17     sort(L, L + n);
19     int first, last;
19     Point *p = new Point[n];
21     Line *q = new Line[n];
21     q[first = last = 0] = L[0];
23     for (int i = 1; i < n; i++) {
24         while (first < last && !OnLeft(L[i], p[last - 1]))
25             last--;
26         while (first < last && !OnLeft(L[i], p[first])) first++;
27         q[++last] = L[i];
27         if (fabs(Cross(q[last].v, q[last - 1].v)) < EPS) {
28             last--;
29             if (OnLeft(q[last], L[i].P)) q[last] = L[i];
31         }
31         if (first < last)

```

```

33         p[last - 1] = GetIntersection(q[last - 1], q[last]);
35     }
35     while (first < last && !OnLeft(q[first], p[last - 1]))
36         last--;
37     if (last - first <= 1) return 0;
37     p[last] = GetIntersection(q[last], q[first]);
39     int m = 0;
41     for (int i = first; i <= last; i++) poly[m++] = p[i];
41     return m;
43 }

```

7. Strings

7.1. Knuth-Morris-Pratt Algorithm

```

1
3 vector<int> pi(const string &s) {
4     vector<int> p(s.size());
5     for (int i = 1; i < s.size(); i++) {
6         int g = p[i - 1];
7         while (g && s[i] != s[g]) g = p[g - 1];
8         p[i] = g + (s[i] == s[g]);
9     }
9     return p;
11 }
12 vector<int> match(const string &s, const string &pat) {
13     vector<int> p = pi(pat + '\0' + s), res;
14     for (int i = p.size() - s.size(); i < p.size(); i++)
15         if (p[i] == pat.size())
16             res.push_back(i - 2 * pat.size());
17     return res;
18 }

```

7.2. Aho-Corasick Automaton

```

1 struct Aho_Corasick {
2     static const int maxc = 26, maxn = 4e5;
3     struct NODES {
4         int Next[maxc], fail, ans;
5     };
6     NODES T[maxn];
7     int top, qtop, q[maxn];
8     int get_node(const int &fail) {
9         fill_n(T[top].Next, maxc, 0);
10        T[top].fail = fail;
11        T[top].ans = 0;
12        return top++;
13    }
14    int insert(const string &s) {
15        int ptr = 1;
16        for (char c : s) { // change char id
17            c -= 'a';
18            if (!T[ptr].Next[c]) T[ptr].Next[c] = get_node(ptr);
19            ptr = T[ptr].Next[c];
20        }
21        return ptr;
22    } // return ans_last_place
23    void build_fail(int ptr) {
24        int tmp;
25        for (int i = 0; i < maxc; i++)
26            if (T[ptr].Next[i]) {
27                tmp = T[ptr].fail;
28                while (tmp != 1 && !T[tmp].Next[i])
29                    tmp = T[tmp].fail;
30                if (T[tmp].Next[i] != T[ptr].Next[i])
31                    if (T[tmp].Next[i]) tmp = T[tmp].Next[i];
32                T[T[ptr].Next[i]].fail = tmp;
33                q[qtop++] = T[ptr].Next[i];
34            }
35    }
36    void AC_auto(const string &s) {
37        int ptr = 1;
38        for (char c : s) {
39            while (ptr != 1 && !T[ptr].Next[c]) ptr = T[ptr].fail;
40            if (T[ptr].Next[c]) {
41                ptr = T[ptr].Next[c];
42                T[ptr].ans++;
43            }
44        }
45    }
46    void Solve(string &s) {
47        for (char &c : s) // change char id
48            c -= 'a';
49        for (int i = 0; i < qtop; i++) build_fail(q[i]);
50        AC_auto(s);
51        for (int i = qtop - 1; i > -1; i--)
52            T[T[q[i]].fail].ans += T[q[i]].ans;
53    }
54    void reset() {
55        qtop = top = q[0] = 1;

```



```

    get_node(1);
}
} AC;
// usage example
string s, S;
int n, t, ans_place[50000];
int main() {
    Tie cin >> t;
    while (t--) {
        AC.reset();
        cin >> S >> n;
        for (int i = 0; i < n; i++) {
            cin >> s;
            ans_place[i] = AC.insert(s);
        }
        AC.Solve(S);
        for (int i = 0; i < n; i++)
            cout << AC.T[ans_place[i]].ans << '\n';
    }
}

```

7.3. Suffix Array

```

1 // sa[i]: starting index of suffix at rank i
2 // 0-indexed, sa[0] = n (empty string)
3 // lcp[i]: lcp of sa[i] and sa[i - 1], lcp[0] = 0
4 struct SuffixArray {
5     vector<int> sa, lcp;
6     SuffixArray(string &s,
7         int lim = 256) { // or basic_string<int>
8         int n = sz(s) + 1, k = 0, a, b;
9         vector<int> x(all(s) + 1), y(n), ws(max(n, lim)),
10             rank(n);
11         sa = lcp = y, iota(all(sa), 0);
12         for (int j = 0, p = 0; p < n;
13             j = max(1, j * 2), lim = p) {
14             p = j, iota(all(y), n - j);
15             for (int i = 0; i < n; i++)
16                 ws[x[i]]++;
17             for (int i = 1; i < lim; i++) ws[i] += ws[i - 1];
18             for (int i = n; i--;) sa[-ws[x[y[i]]]] = y[i];
19             swap(x, y), p = 1, x[sa[0]] = 0;
20             for (int i = 1; i < n; i++)
21                 a = sa[i - 1], b = sa[i],
22                 x[b] = (y[a] == y[b] && y[a + j] == y[b + j])
23                     ? p - 1 : p++;
24         }
25         for (int i = 1; i < n; i++) rank[sa[i]] = i;
26         for (int i = 0, j; i < n - 1; lcp[rank[i++]] = k)
27             for (k && k--, j = sa[rank[i] - 1];
28                 s[i + k] == s[j + k]; k++);
29     }
30 };

```

7.4. Suffix Tree

```

1 struct SAM {
2     static const int maxc = 26; // char range
3     static const int maxn = 10010; // string len
4     struct Node {
5         Node *green, *edge[maxc];
6         int max_len, in, times;
7     } *root, *last, reg[maxn * 2];
8     int top;
9     Node *get_node(int _max) {
10         Node *re = &reg[top++];
11         re->in = 0, re->times = 1;
12         re->max_len = _max, re->green = 0;
13         for (int i = 0; i < maxc; i++) re->edge[i] = 0;
14         return re;
15     }
16     void insert(const char c) { // c in range [0, maxc)
17         Node *p = last;
18         last = get_node(p->max_len + 1);
19         while (p && !p->edge[c])
20             p->edge[c] = last, p = p->green;
21         if (!p) last->green = root;
22         else {
23             Node *pot_green = p->edge[c];
24             if ((pot_green->max_len == (p->max_len + 1))
25                 last->green = pot_green;
26             else {
27                 Node *wish = get_node(p->max_len + 1);
28                 wish->times = 0;
29                 while (p && p->edge[c] == pot_green)
30                     p->edge[c] = wish, p = p->green;

```

```

31         for (int i = 0; i < maxc; i++)
32             wish->edge[i] = pot_green->edge[i];
33         wish->green = pot_green->green;
34         pot_green->green = wish;
35         last->green = wish;
36     }
37 }
38 Node *q[maxn * 2];
39 int ql, qr;
40 void get_times(Node *p) {
41     ql = 0, qr = -1, reg[0].in = 1;
42     for (int i = 1; i < top; i++) reg[i].green->in++;
43     for (int i = 0; i < top; i++)
44         if (!reg[i].in) q[++qr] = &reg[i];
45     while (ql <= qr) {
46         q[ql]->green->times += q[ql]->times;
47         if (!(--q[ql]->green->in)) q[++qr] = q[ql]->green;
48         ql++;
49     }
50 }
51 void build(const string &s) {
52     top = 0;
53     root = last = get_node(0);
54     for (char c : s) insert(c - 'a'); // change char id
55     get_times(root);
56 }
57 // call build before solve
58 int solve(const string &s) {
59     Node *p = root;
60     for (char c : s)
61         if (!(p = p->edge[c - 'a'])) // change char id
62             return 0;
63     return p->times;
64 }
65 };

```

7.5. Cocke-Younger-Kasami Algorithm

```

1 struct rule {
2     // s -> xy
3     // if y == -1, then s -> x (unit rule)
4     int s, x, y, cost;
5 };
6 int state;
7 // state (id) for each letter (variable)
8 // lowercase letters are terminal symbols
9 map<char, int> rules;
10 vector<rule> cnf;
11 void init() {
12     state = 0;
13     rules.clear();
14     cnf.clear();
15 }
16 // convert a cfg rule to cnf (but with unit rules) and add
17 // it
18 void add_to_cnf(char s, const string &p, int cost) {
19     if (!rules.count(s)) rules[s] = state++;
20     for (char c : p)
21         if (!rules.count(c)) rules[c] = state++;
22     if (p.size() == 1) {
23         cnf.push_back({rules[s], rules[p[0]], -1, cost});
24     } else {
25         // length >= 3 -> split
26         int left = rules[s];
27         int sz = p.size();
28         for (int i = 0; i < sz - 2; i++) {
29             cnf.push_back({left, rules[p[i]], state, 0});
30             left = state++;
31         }
32         cnf.push_back(
33             {left, rules[p[sz - 2]], rules[p[sz - 1]], cost});
34     }
35 }
36 }
37 }
38 constexpr int MAXN = 55;
39 vector<long long> dp[MAXN][MAXN];
40 // unit rules with negative costs can cause negative cycles
41 vector<bool> neg_INF[MAXN][MAXN];
42 void relax(int l, int r, rule c, long long cost,
43     bool neg_c = 0) {
44     if (!neg_INF[l][r][c.s] &&
45         (neg_INF[l][r][c.x] || cost < dp[l][r][c.s])) {
46         if (neg_c || neg_INF[l][r][c.x]) {
47             dp[l][r][c.s] = 0;
48             neg_INF[l][r][c.s] = true;
49         } else {
50             dp[l][r][c.s] = cost;
51         }
52     }
53 }
54 }
55 }

```

```

void bellman(int l, int r, int n) {
    for (int k = 1; k <= state; k++)
        for (rule c : cnf)
            if (c.y == -1)
                relax(l, r, c, dp[l][r][c.x] + c.cost, k == n);
}

void cyk(const string &s) {
    vector<int> tok;
    for (char c : s) tok.push_back(rules[c]);
    for (int i = 0; i < tok.size(); i++) {
        for (int j = 0; j < tok.size(); j++) {
            dp[i][j] = vector<long long>(state + 1, INT_MAX);
            neg_INF[i][j] = vector<bool>(state + 1, false);
        }
        dp[i][i][tok[i]] = 0;
        bellman(i, i, tok.size());
    }
    for (int r = 1; r < tok.size(); r++) {
        for (int l = r - 1; l >= 0; l--) {
            for (int k = l; k < r; k++)
                for (rule c : cnf)
                    if (c.y != -1)
                        relax(l, r, c,
                            dp[l][k][c.x] + dp[k + 1][r][c.y] +
                            c.cost);
            bellman(l, r, tok.size());
        }
    }
}

// usage example
int main() {
    init();
    add_to_cnf('S', "aSc", 1);
    add_to_cnf('S', "BBB", 1);
    add_to_cnf('S', "SB", 1);
    add_to_cnf('B', "b", 1);
    cyk("abbbbc");
    // dp[0][s.size() - 1][rules[start]] = min cost to
    // generate s
    cout << dp[0][5][rules['S']] << '\n'; // 7
    cyk("acbc");
    cout << dp[0][3][rules['S']] << '\n'; // INT_MAX
    add_to_cnf('S', "S", -1);
    cyk("abbbbc");
    cout << neg_INF[0][5][rules['S']] << '\n'; // 1
}

```

7.6. Z Value

```

int z[n];
void zval(string s) {
    // z[i] => longest common prefix of s and s[i:], i > 0
    int n = s.size();
    z[0] = 0;
    for (int b = 0, i = 1; i < n; i++) {
        if (z[b] + b <= i) z[i] = 0;
        else z[i] = min(z[i - b], z[b] + b - i);
        while (s[i + z[i]] == s[z[i]]) z[i]++;
        if (i + z[i] > b + z[b]) b = i;
    }
}

```

7.7. Manacher's Algorithm

```

int z[n];
void manacher(string s) {
    // z[i] => longest odd palindrome centered at i is
    // s[i - z[i] ... i + z[i]]
    // to get all palindromes (including even length),
    // insert a '#' between each s[i] and s[i + 1]
    int n = s.size();
    z[0] = 0;
    for (int b = 0, i = 1; i < n; i++) {
        if (z[b] + b >= i)
            z[i] = min(z[2 * b - i], b + z[b] - i);
        else z[i] = 0;
        while (i + z[i] + 1 < n && i - z[i] - 1 >= 0 &&
            s[i + z[i] + 1] == s[i - z[i] - 1])
            z[i]++;
        if (z[i] + i > z[b] + b) b = i;
    }
}

```

7.8. Minimum Rotation

```

int min_rotation(string s) {
    int a = 0, n = s.size();
    s += s;
    for (int b = 0; b < n; b++) {
        for (int k = 0; k < n; k++) {
            if (a + k == b || s[a + k] < s[b + k]) {
                b += max(0, k - 1);
            }
        }
    }
}

```

```

        break;
    }
    if (s[a + k] > s[b + k]) {
        a = b;
        break;
    }
}
return a;
}

```

7.9. Palindromic Tree

```

1 struct palindromic_tree {
2     struct node {
3         int next[26], fail, len;
4         int cnt,
5         num; // cnt: appear times, num: number of pal. suf.
6         node(int l = 0) : fail(0), len(l), cnt(0), num(0) {
7             for (int i = 0; i < 26; ++i) next[i] = 0;
8         }
9     };
10    vector<node> St;
11    vector<char> s;
12    int last, n;
13    palindromic_tree() : St(2), last(1), n(0) {
14        St[0].fail = 1, St[1].len = -1, s.pb(-1);
15    }
16    inline void clear() {
17        St.clear(), s.clear(), last = 1, n = 0;
18        St.pb(0), St.pb(-1);
19        St[0].fail = 1, s.pb(-1);
20    }
21    inline int get_fail(int x) {
22        while (s[n - St[x].len - 1] != s[n]) x = St[x].fail;
23        return x;
24    }
25    inline void add(int c) {
26        s.push_back(c - 'a'), ++n;
27        int cur = get_fail(last);
28        if (!St[cur].next[c]) {
29            int now = SZ(St);
30            St.pb(St[cur].len + 2);
31            St[now].fail = St[get_fail(St[cur].fail)].next[c];
32            St[cur].next[c] = now;
33            St[now].num = St[St[now].fail].num + 1;
34        }
35        last = St[cur].next[c], ++St[last].cnt;
36    }
37    inline void count() { // counting cnt
38        auto i = St.rbegin();
39        for (; i != St.rend(); ++i) {
40            St[i->fail].cnt += i->cnt;
41        }
42    }
43    inline int size() { // The number of diff. pal.
44        return SZ(St) - 2;
45    }
46 };

```

8. Debug List

- 1 - Pre-submit:
 - 2 - Did you make a typo when copying a template?
 - 3 - Test more cases if unsure.
 - 4 - Write a naive solution and check small cases.
 - 5 - Submit the correct file.
- 7 - General Debugging:
 - 8 - Read the whole problem again.
 - 9 - Have a teammate read the problem.
 - 10 - Have a teammate read your code.
 - 11 - Explain your solution to them (or a rubber duck).
 - 12 - Print the code and its output / debug output.
 - 13 - Go to the toilet.
- 15 - Wrong Answer:
 - 16 - Any possible overflows?
 - 17 - > `__int128`?
 - 18 - Try `-ftrapv` or `#pragma GCC optimize("trapv")`
 - 19 - Floating point errors?
 - 20 - > `long double`?
 - 21 - turn off math optimizations
 - 22 - check for `==`, `>=`, `acos(1.000000001)`, etc.
 - 23 - Did you forget to sort or unique?
 - 24 - Generate large and worst "corner" cases.
 - 25 - Check your `m` / `n`, `i` / `j` and `x` / `y`.
 - 26 - Are everything initialized or reset properly?
 - 27 - Are you sure about the STL thing you are using?
 - 28 - Read cppreference (should be available).

```
29 - Print everything and run it on pen and paper.
31 - Time Limit Exceeded:
33   - Calculate your time complexity again.
33   - Does the program actually end?
35     - Check for while(q.size()) etc.
35   - Test the largest cases locally.
37   - Did you do unnecessary stuff?
37     - e.g. pass vectors by value
39     - e.g. memset for every test case
39   - Is your constant factor reasonable?
41 - Runtime Error:
43   - Check memory usage.
43   - Forget to clear or destroy stuff?
45     - > vector::shrink_to_fit()
45   - Stack overflow?
47   - Bad pointer / array access?
47     - Try -fsanitize=address
47   - Division by zero? NaN's?
```