



Trees Problem Solving 1

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Problem 1



Kuro and Walking Route:

<https://codeforces.com/contest/979/problem/C>



Main Idea

- Consider the node **y** as the root of the tree.
- Now all the paths that start from the subtree of node **x** (including **x**) and do not end at any ancestor of **x** (except **y**) are not safe.
- Subtract all these paths from the value $n * (n - 1)$.
- This will give you the final answer.

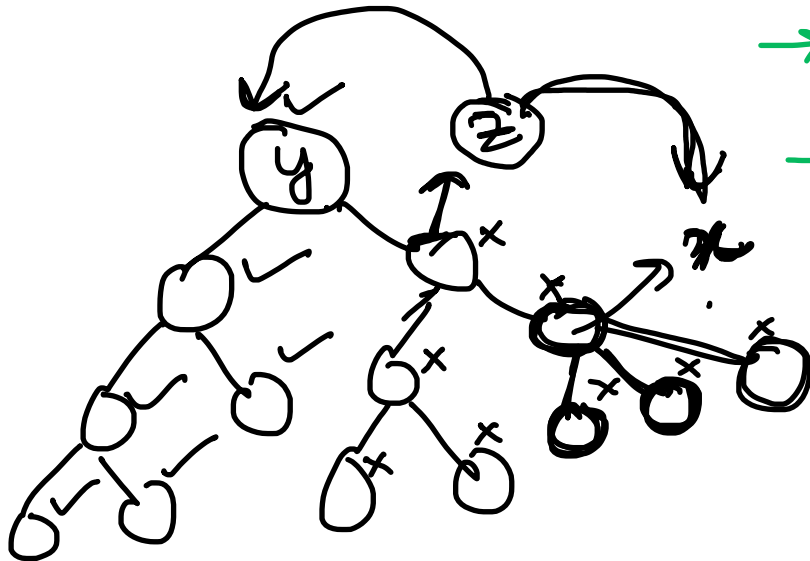
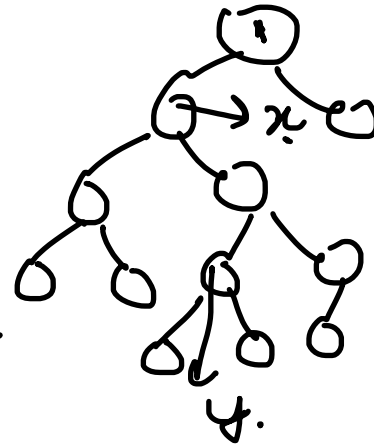
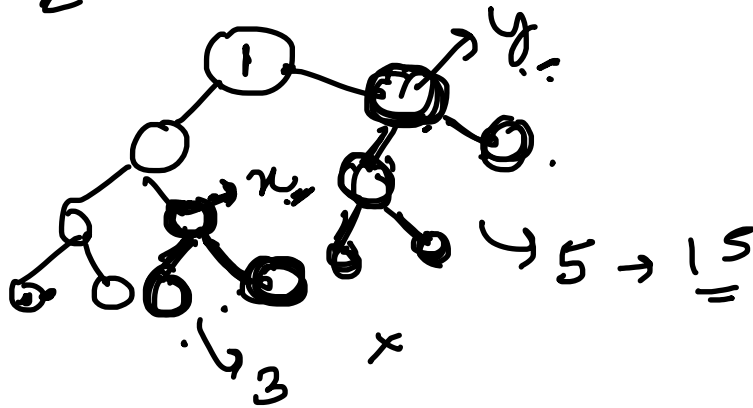
$$(n) \times (n-1)$$

$$(u, v) \rightarrow \begin{matrix} u & \neq & v \\ \swarrow & & \searrow \\ (n) & \times & (n-1) \end{matrix}$$

$$(u, v) \rightarrow \boxed{n_2} \times 2$$

$$(v, w) \rightarrow \frac{n \times (n-1)}{2} \times 2 \rightarrow n \times \underline{(n-1)}$$

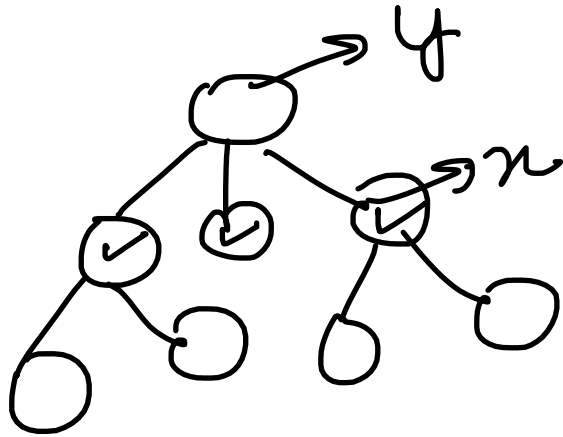
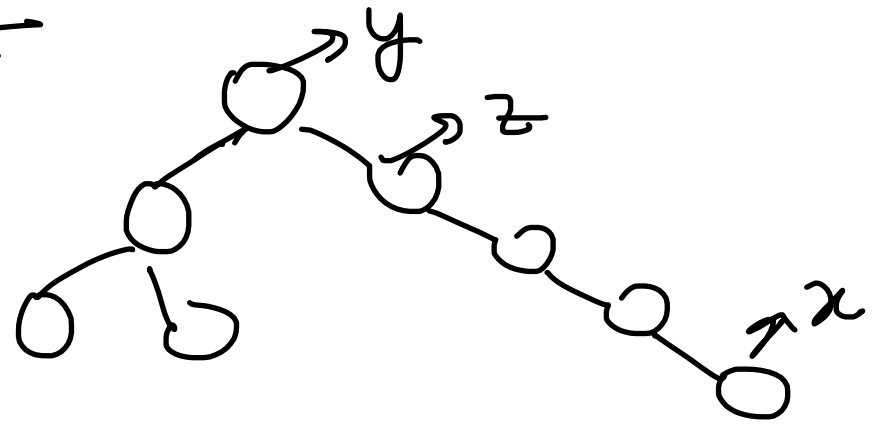
$$\begin{pmatrix} \check{x} \\ \check{y} \end{pmatrix} = \begin{pmatrix} u \\ v \end{pmatrix}$$



- z is direct child of y .
- z is an ancestor of x .

$$\underline{\underline{z, x}}$$

$$\boxed{\text{sub}[x] \times (n - \text{sub}[z])} \leftarrow$$



Problem 2



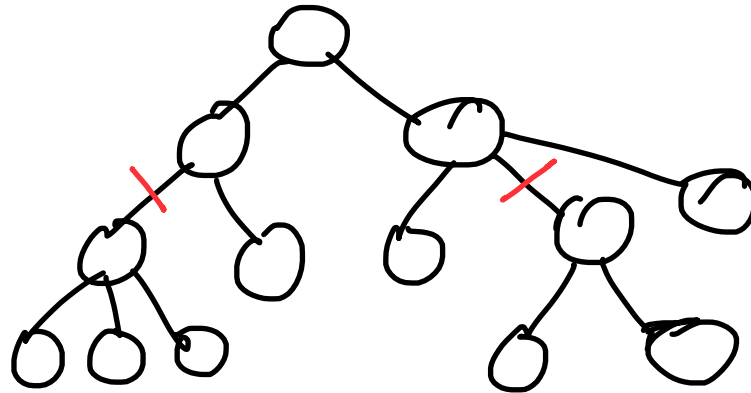
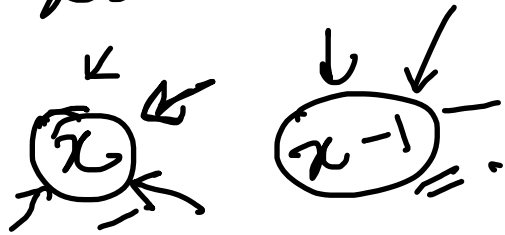
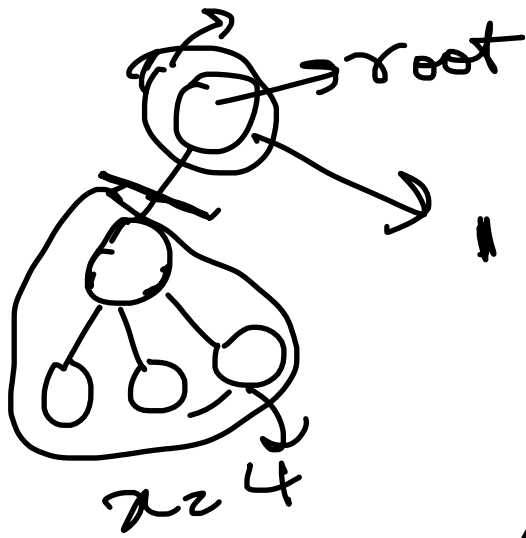
Tree Cutting:

<https://codeforces.com/problemset/problem/1946/C>



Main Idea

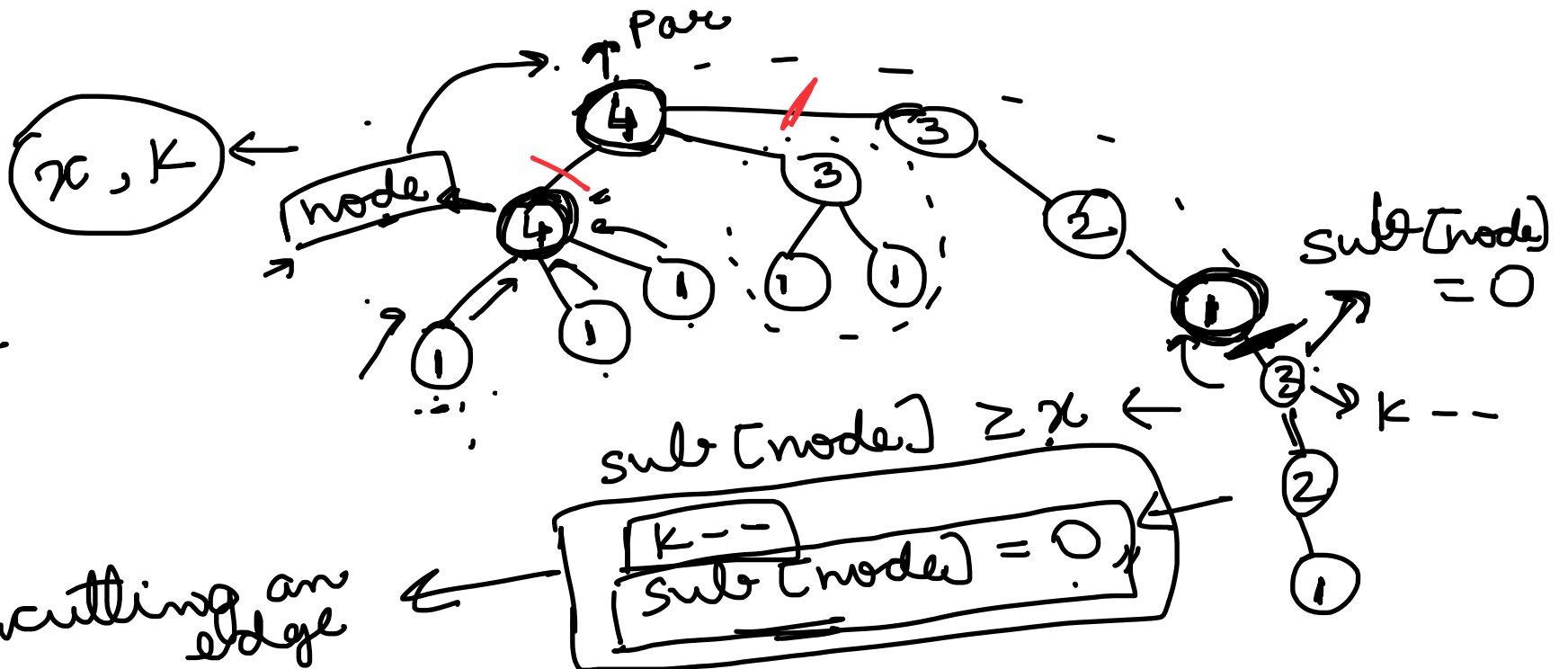
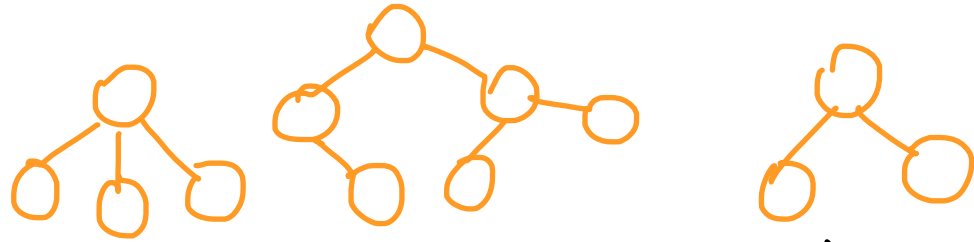
- We will binary search over the value of x .
- For a given value of x , take an arbitrary root node and start calculating the subtree sizes of the nodes from leaf nodes to root.
- At any point of time, if the subtree size of a node is greater than the given x , then we can remove the edge from the given node to its parent.
- Just be careful about the calculations for root node.

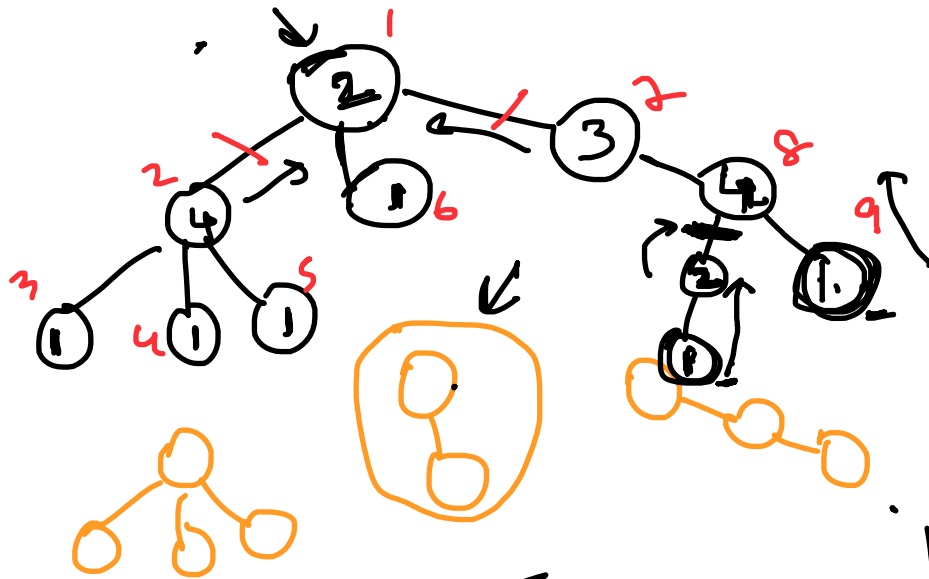


$K=2$

$x=3$

$x=2$





$K=2$
 $K=1$
 $K=0$

$\geq n_{11}$

$\chi=2$

(n)

$\downarrow \chi$
 $2 \leq 1$
 3

$\chi(K)$

Problem 3



13th Labour of Heracles:

<https://codeforces.com/problemset/problem/1466/D>

Main Idea

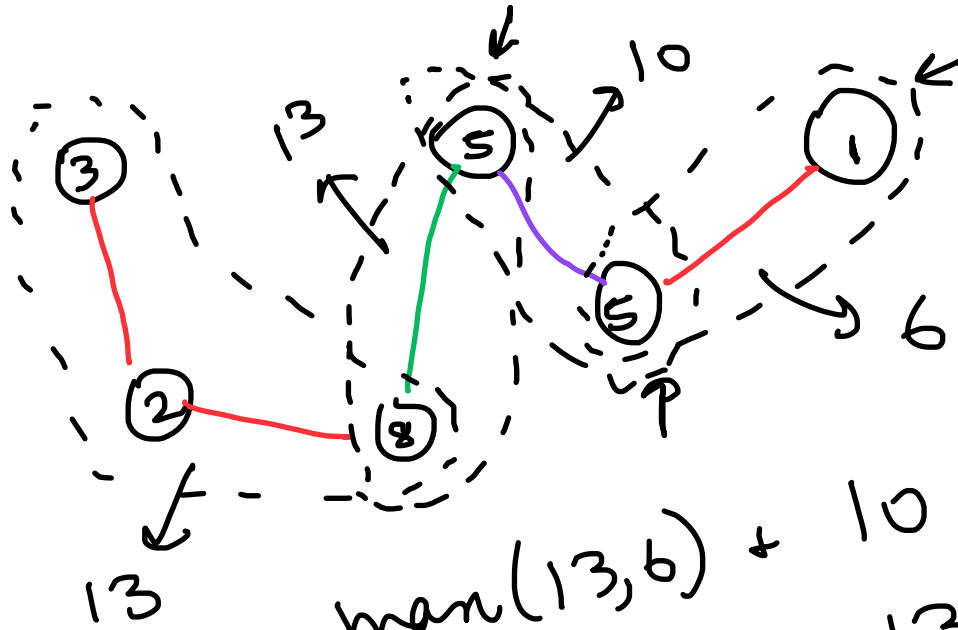


- Each node's weight can be considered in the sum for at max the number of times equal to its degree.
- Also, it is always possible to configure the colour of the edges for any k such that a given node is considered again.
- Hence, at every step, colour the edges in such a way that the node with maximum weight is counted again in the current step.

$(K) \rightarrow$ no. of colours.
(edges).

atmost K colours
in edges.

3 colours,



$$\max(13, 6) + 10 + 13$$

$$13 + 10 + 13$$

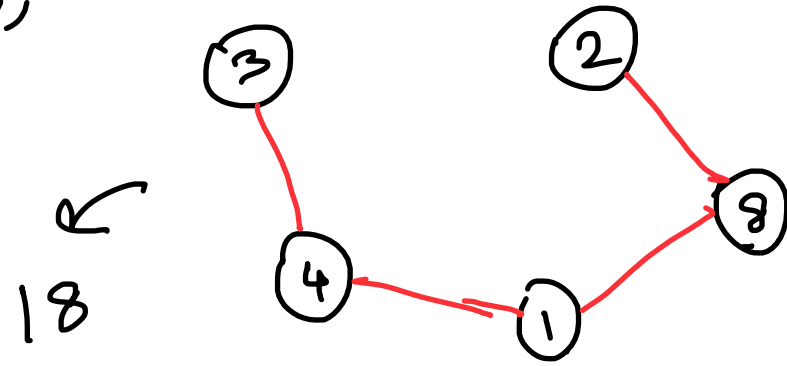
$$= (36)$$

$$K = 5,$$

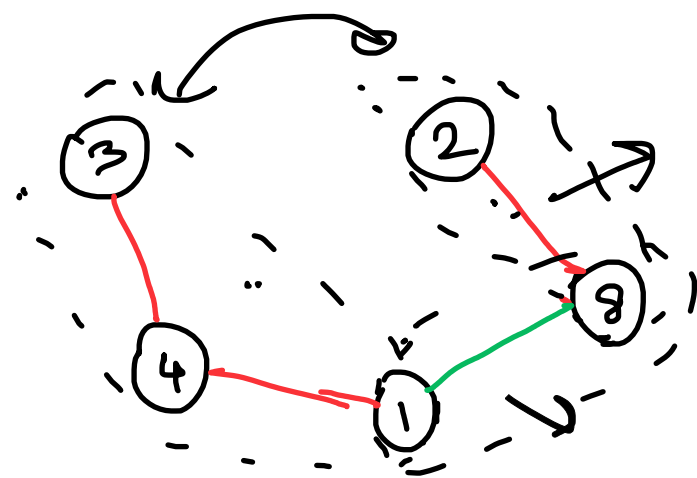
$$K = 3,$$

$$h-1 = 5,$$

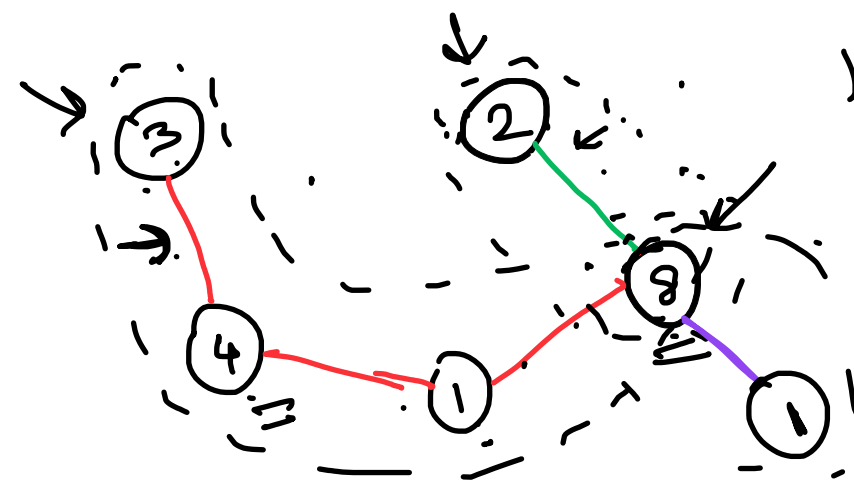
$K \rightarrow 1$ to $n-1,$



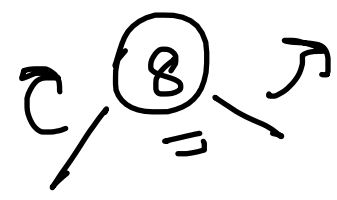
$$K=2,,$$



$$10 + 9 = 19$$



$$16 + 10 = 26 =$$



$$K=1,,$$

$$K=2$$

$$K=3,,$$

$$K=4$$

$$val \rightarrow val + 8,, \quad val + 4 + 8$$

