

## Unit 1

### PART-A

- 1) benefits of separating lexical Analysis from Syntax Analysis:-
- 1) simple design
    - Separation allows the simplification of one or the other
    - e.g:- A parser with comments or white spaces is more complex
  - 2) compiler efficiency is improved
    - Optimization of lexical Analysis because large Amount of time is spent on reading the source program and partitioning into tokens.

2) Lexeme	Token	Patterns
int	keyword	
diff	identifier	
,	symbol	
int1	Identifier	
int2	identifier	
;	SP Symbol/semicolon	
if	keyword	
>	operator	
=	operator	
-	operator	
else	keyword	

3) The error recovery for the given program is:-

\* Transforming two Adjacent characters

ex:- `esle`  $\rightarrow$  `else`

4) 1) Structure editor

2) Pretty printers

3) Static checkers

4) Interpreters.

5) \* It is non-trivial test of language

\* This technique makes the compiler self-hosting compiler

\* compiler development can be done in higher language is also compiled.

\* compiler developer only needs to know that language in which target code is compiled.

6) Components of compiler

1) preprocessor

2) Assembler

3) Loader and Link editor.

7) The two parts of compilation is

1) Analysis

2) Synthesis

1) Analysis:- This include lexical and Syntactic Analysis and creation of symbol table, semantic Analysis and Intermediate code generation. and Includes error handling on all these phases

7) Synthesis :- It include code optimization and code generation along with necessary error handling operations.

8) Structure editor :-

- \* Takes as input a sequence of commands to build a source program

- \* It not only performs the text creation and modification functions of an ordinary text editor, but it also analyze the program text, putting an appropriate hierarchical structure of the source program.

ex:- It can supply keywords automatically:-  
while... do and begin... end.

9) i) lexical

ii) syntax

iii) semantic

iv) Intermediate code

v) code optimization

vi) code generator.

10) Pretty printers

- \* A pretty printer analyze a program and prints it in such a way that the structure of the program becomes clearly visible.

ex:- comments may appear in special fonts

Static checkers :-

- \* A static checker reads a program, analyze it and attempts





to discover potential bugs without running the program.  
ex:- A static checker may detect the parts of the source program can never be executed.

$$\begin{aligned}
 11) \quad L &= \{w \in \{a, b\}^* / w \text{ in } abb\} \\
 &= a, b \text{ (or) } a|b \\
 &= \cancel{a+bb}^* (a+b)^* abb \\
 &= \{abb, aabb, aabbb\}
 \end{aligned}$$

12) Sentinel:- The sentinel is a special character that cannot be part of the source program, and natural choice is the character is eof.

### Usage

eof retains its use as a marker for the end of the entire input. Any eof that appears other than at the end of buffer means that the input is at ~~than~~ end.

13) Rational preprocessor:-

These processors change older languages with more modern flow-of-control and data structure facilities.

- 14)
- 1) macro processing
  - 2) file inclusion
  - 3) Rational pre processors
  - 4) language extension



- 15)
- 1) Deleting an extraneous character
  - 2) Inserting a missing character
  - 3) Replacing Incorrect character with correct character
  - 4) Transforming two adjacent characters
  - 5) panic mode recovery:- deletion of successive char from the token until error is resolved
- 16) one buffer  
two - buffer scheme.

### PART-B & C

17) cousins of compiler

- 1) Preprocessor
- 2) Assembler
- 3) Linker and loader

1) preprocessor:- A preprocessor is a program that process the input data <sup>to</sup> produce output that is used as input to another program

\* The output is said to be preprocessed form of the input data, which is often used by some subsequent programs like compilers.

They may perform some function like:-

- \* macro procession
- \* file inclusion
- \* Rational preprocessors
- \* language extension.



- 1) macro processing:- A macro is a pattern which certain sequence of input is mapped to output sequence according to defined procedure.
- 2) file inclusion:- preprocessor includes header files into program text files. When preprocessor finds #include directive it replaces it by the entire content of the specified file.
- 3) Rational preprocessors:- These preprocessors change older languages with modern flow-of-control and data structuring facilities.
- 4) Language extension:- These processors attempt to add capabilities to the language by what amounts to built-in macros.

### 2) Assembler:-

Assembler creates object code by translating assembly instruction mnemonics into machine code.

Two types of Assemblers:-

- \* One-pass assembler goes through the source code once and assumes that all symbols will be defined before any instruction references them.
- \* Two-pass assembler creates a table with all symbols and their values in the first pass and then uses the table in the second pass to generate code.

### 3) Linker & Loader

A linker or link editor is a program that takes one or more objects generated by a compiler and combines them into a single executable program.

Three tasks of Linker

1. Searches the program to find library routines used by a program.
2. Determine the memory locations and reallocate the instructions.
3. Resolves references among files.





A loader is part of OS that is responsible for loading programs in memory.

## 13) Compiler construction Tools

### 1) Parser Generators

\* These produce Syntax Analyzers, normally from output that is based on CFG.

\* It consumes a large fraction of the running time of a compiler.

\* ex:- Yacc.

### 2) Scanner Generators:-

\* These generate lexical Analyzers, normally from a specification based on RE

\* The basic organization of lexical analyzer is based on finite Automata.

### 3) Syntax - Directed Translation

→ These produces routines that walk the parse tree and generate intermediate code as result

→ Each Translation is defined in terms of translation as its neighbour nodes in the tree.

### 4) Automatic code generators:-

→ It takes a collection of rules which translate intermediate code to machine code.

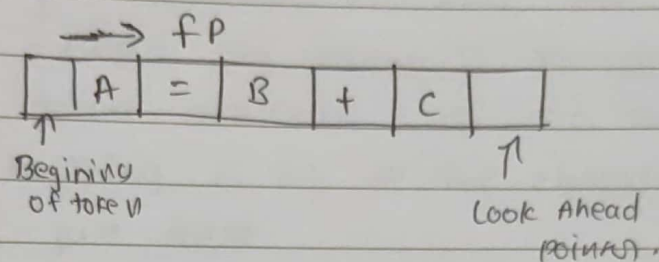
### 5) Data flow engine:-

→ It does code optimization using data-flow analysis, that is the gathering of information about how values are transmitted from one part of program to each other part.



### 19) Input Buffering

We often have to look one or more characters beyond the next lexeme before we can be sure we have right lexem. As we know that as characters are read from left to right, each character is stored in the buffer to form meaning full token.



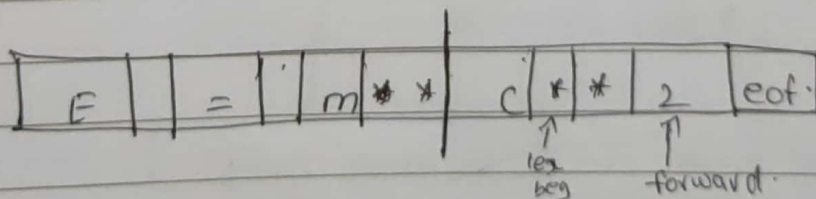
We have two input buffering techniques:-

one-buffer scheme, two-buffer scheme

\* we introduce two buffer scheme that handles large look ahead safely.

We then consider an improvement involving "sentinals" which is used to check end of buffer.

ex:-



- A buffer is divided into two  $N$ -character halves.
- \* Each buffer is of the same size of  $N$
- \* Using one system read command we can read  $N$  char into a buffer
- \* if fewer than  $N$  char remain in input file, then special character, represented by eof, marks the end of source file.





\* two pointers to the input is maintained.

- 1) lexeme beginning :- which marks the beginning of current lexeme
- 2) forward :- scans ahead until a pattern match is found.  
 one the next lexeme is determined forward is set to the character at its right end.

\* The string of characters b/w the two pointers is current lexeme. After the lexeme is recorded as an Attribute value of a token returned to the parser.  
 lexeme-beginning is set to the character immediately after the lexeme is just found.

\* eof is special character which tells the end of buffer or end of input from source file.

20) example  $a = b + c * 2$   
 ↓ source code.

Lexical Analyzer

→ which reads characters of sc and converts in tokens.

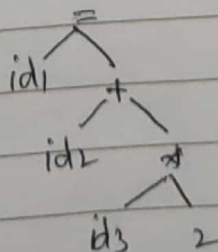
$id1 = id2 + id3 * 2$

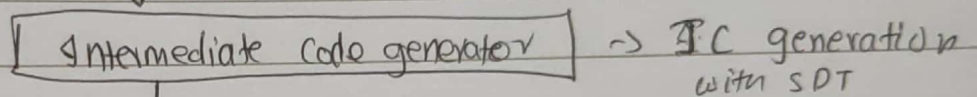
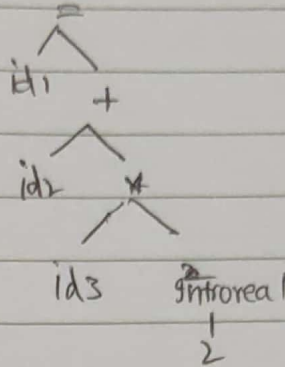
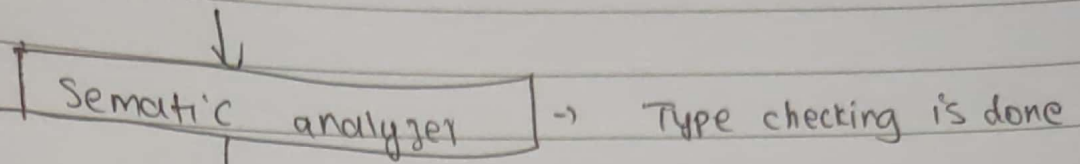
Syntax Analyzer

→ parse Tree and symbol table.

Symbol Table.

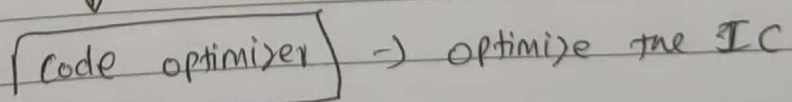
a	id1
b	id2
c	id3





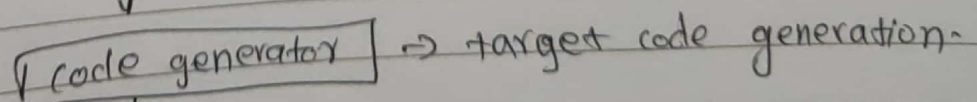
```

temp1 = intoreal(2)
temp2 = id3 * temp1
temp3 = id2 + temp2
id1 = temp3
  
```



```

temp1 = id3 * 2.0
id1 = id2 + temp1
  
```



```

movf id3, R2
MULF #2.0, R2
MOVf id2, R1
ADDf R2, R1
MOVf R1, id1.
  
```