

Binomial Coefficient: A Dynamic Programming Approach

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Abstract

The Binomial Coefficient, denoted $C(n, k)$, represents the number of ways to choose k items from a set of n items without regard to order. This document explores a dynamic programming solution to compute $C(n, k)$, addressing its combinatorial significance, recurrence relation, algorithm, complexity analysis, practical applications, and limitations. Additional insights into space optimization and handling large values are included for a comprehensive understanding.

1 Introduction

The Binomial Coefficient $C(n, k)$, also written as $\binom{n}{k}$, is a fundamental concept in combinatorics, representing the number of ways to select k elements from a set of n elements without considering the order of selection. It appears in various fields, including probability, statistics, and algorithm design, and is a building block of Pascal's Triangle.

A naive recursive approach to compute $C(n, k)$ suffers from redundant calculations, leading to exponential time complexity. **Dynamic programming** offers an efficient solution by storing intermediate results, leveraging the problem's overlapping subproblems and optimal substructure.

2 Problem Formulation

Given two non-negative integers n and k (where $k \leq n$), compute $C(n, k)$, the number of ways to choose k items from n items.

Mathematically:

$$C(n, k) = \binom{n}{k} = \frac{n!}{k!(n-k)!}$$

However, direct factorial computation is inefficient for large n . Instead, we use a recursive relation suitable for dynamic programming.

2.1 Recurrence Relation

The Binomial Coefficient satisfies the following recurrence:

$$C(n, k) = C(n-1, k-1) + C(n-1, k)$$

With base cases:

$$C(n, 0) = C(n, n) = 1$$

$$C(n, k) = 0 \quad \text{if } k > n$$

This recurrence reflects the combinatorial interpretation: to choose k items from n , either include the n -th item (requiring $k - 1$ more items from $n - 1$) or exclude it (requiring k items from $n - 1$).

3 Dynamic Programming Algorithm

The bottom-up dynamic programming approach constructs a 2D table dp of size $(n + 1) \times (k + 1)$, where $dp[i][j]$ stores $C(i, j)$. Pseudocode is as follows:

Input: n, k

Output: $C(n, k)$

```
1. Initialize  $dp[n+1][k+1]$ 
2. For  $i$  from 0 to  $n$ :
3.      $dp[i][0] = 1$  // Base case:  $C(i, 0) = 1$ 
4.     If  $i \leq k$ :
5.          $dp[i][i] = 1$  // Base case:  $C(i, i) = 1$ 
6. For  $i$  from 1 to  $n$ :
7.     For  $j$  from 1 to  $\min(i, k)$ :
8.          $dp[i][j] = dp[i-1][j-1] + dp[i-1][j]$ 
9. Return  $dp[n][k]$ 
```

4 Complexity Analysis

- **Time Complexity:**

- Naive recursion: $O(2^n)$, due to overlapping subproblems.
- Dynamic programming: $O(n \cdot k)$, as the algorithm fills a table of size $(n + 1) \times (k + 1)$.

- **Space Complexity:**

- $O(n \cdot k)$ for the 2D dp table.
- Can be optimized to $O(k)$ using a 1D rolling array (see Section 5).

5 Space Optimization

Since $dp[i][j]$ depends only on $dp[i-1][j-1]$ and $dp[i-1][j]$, we can use a 1D array $dp[k+1]$ to reduce space complexity to $O(k)$. The array is updated iteratively, computing each row from right to left to avoid overwriting values:

Input: n, k

Output: $C(n, k)$

```

1. Initialize dp[k+1] with dp[0] = 1
2. For i from 1 to n:
3.     For j from min(i, k) down to 1:
4.         dp[j] = dp[j-1] + dp[j]
5. Return dp[k]

```

This optimization is particularly useful for large n and moderate k .

6 Handling Large Values

For large n , $C(n, k)$ can exceed the range of standard integer types. To prevent overflow:

- Use `long long` or arbitrary-precision arithmetic libraries.
- In competitive programming, apply modular arithmetic (e.g., modulo $10^9 + 7$) to keep intermediate results manageable:

$$C(n, k) \bmod m = (C(n-1, k-1) \bmod m + C(n-1, k) \bmod m) \bmod m$$

7 Practical Applications

The Binomial Coefficient is widely used in:

- **Probability and Statistics:** Computing probabilities in binomial distributions.
- **Pascal's Triangle:** Each entry is a binomial coefficient.
- **Combinatorial Algorithms:** Solving problems involving subsets or permutations.
- **Coding Contests:** Common in dynamic programming and math-based challenges.

8 Example

Compute $C(5, 2)$.

Using the DP table:

$i \backslash j$	0	1	2
0	1	0	0
1	1	1	0
2	1	2	1
3	1	3	3
4	1	4	6
5	1	5	10

Result: $C(5, 2) = 10$, meaning there are 10 ways to choose 2 items from 5.

9 Limitations and Extensions

- **Limitations:**
 - Large n and k may cause overflow without modular arithmetic.
 - Space complexity can be significant for large k .
- **Extensions:**
 - **Lucas' Theorem:** Computes $C(n, k) \bmod p$ for prime p .
 - **Multi-choose:** Computes combinations with repetition.
 - **Fast Fourier Transform:** Used for convolution-based binomial calculations in advanced applications.

10 References

- [GeeksForGeeks: Binomial Coefficient](#)
- [Wikipedia: Binomial Coefficient](#)
- Cormen, T. H., et al., *Introduction to Algorithms*, Chapter on Dynamic Programming.