

Chapter 6A More Lists

CS401

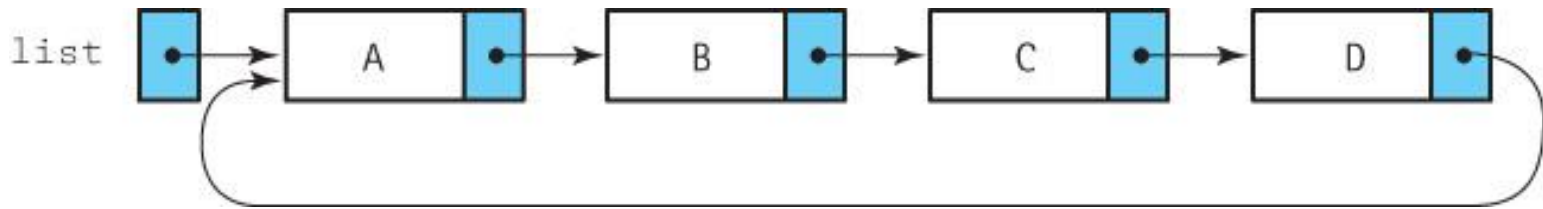
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Revised Nell Dale Presentation

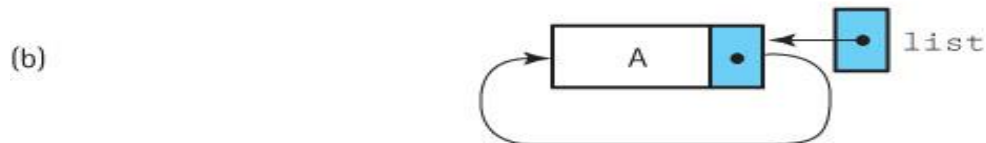
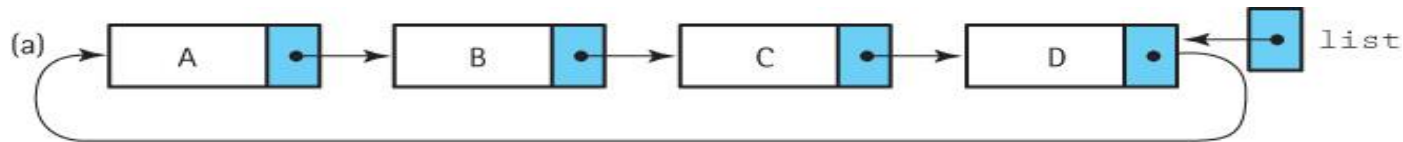
Circular Linked Lists

- **Circular linked list** A list in which every node has a successor; the “last” element is succeeded by the “first” element

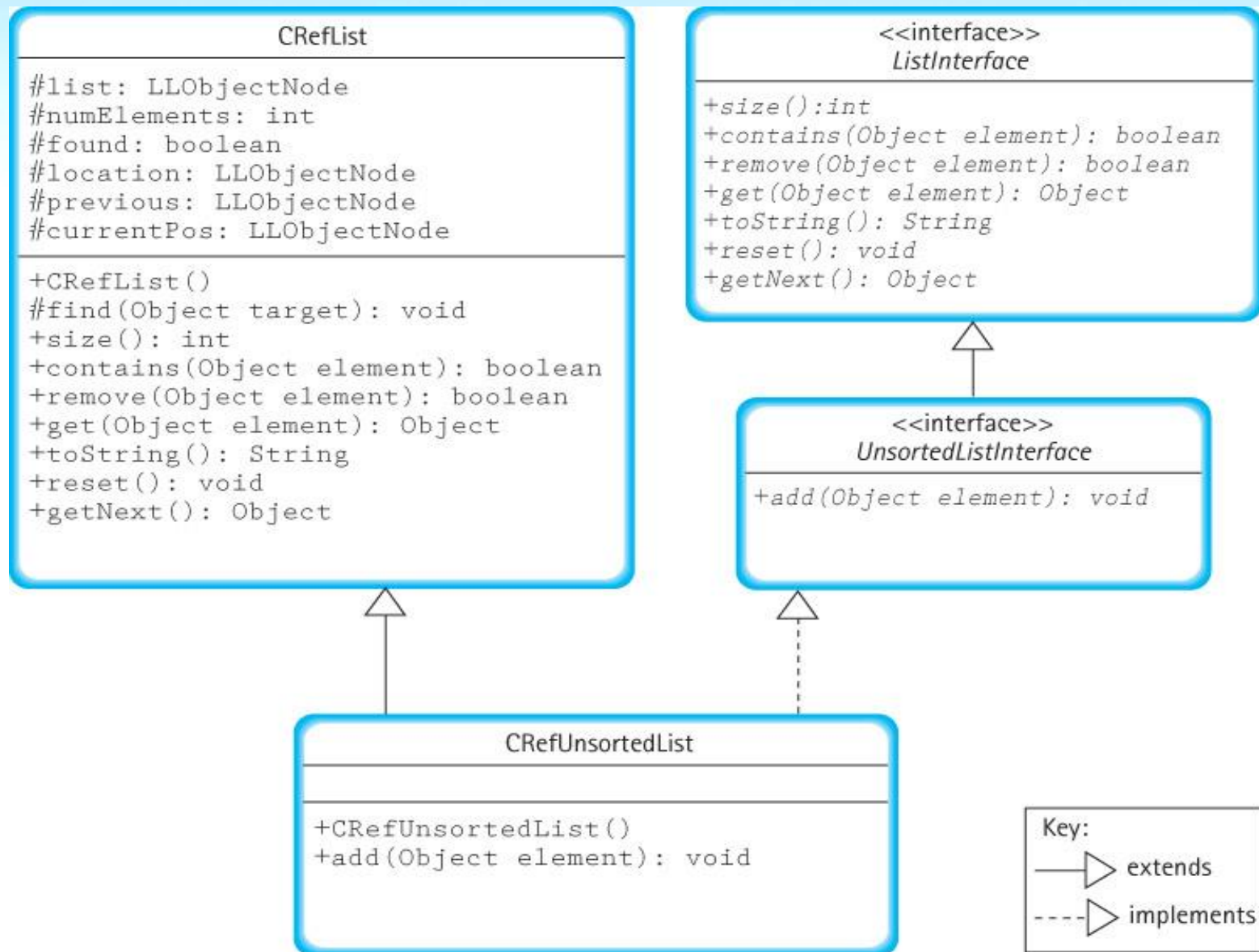


A more efficient approach

- Adding and removing elements at the front of a list might be a common operation for some applications.
- Our linear linked list approach supports these operations very efficiently ← **Why?**
- The previous slide circular linked list approach does not (we need to access the last element also).
- We can fix this problem by letting our list reference point to the last element in the list rather than the first; now we have easy access to both the first and the last elements in the list.



An Unsorted Circular Linked List



The CRefList Class

- The CRefList class can contain the same set of methods (`size`, `contains`, `remove`, `get`, `toString`, `reset`, and `getNext`) as its linear list counterpart RefList.
- The `size` method need not be changed.
- If we provide a revised `find` helper method with functionality that corresponds to the `find` in RefList, we can also reuse both the `contains` and `get` methods.

The Iterator methods

- the `reset` method becomes `reset()` and the `getNext` method becomes `getNext()`

```
public void reset()
{
    if (list != null)
        currentPos = list.getLink();
}

public Object getNext()
{
    Object next = currentPos.getInfo();
    currentPos = currentPos.getLink();

    return next;
}
```

// Linear list

```
Public void reset () {
    currentPos = list;
}
```

// Linear list

```
Public Object getNext () {
    Object next = currentPos.getInfo();
    if (currentPos.getLink() == null)
        currentPos = list;
    currentPos = currentPos.next;
    else
        currentPos = currentPos.getLink();
    return next;
}
```

The toString method

```
public String toString()  
// Returns a nicely formatted String that represents this list.  
{  
    String listString = "List:\n";  
    if (list != null)  
    {  
        LLObjectNode prevNode = list;  
        do  
        {  
            listString = listString + "  " + prevNode.getLink().getInfo() + "\n";  
            prevNode = prevNode.getLink();  
        }  
        while (prevNode != list);  
    }  
    return listString;  
}
```

The find method

```
protected void find(Object target)
{
    boolean moreToSearch;
    location = list;
    found = false;

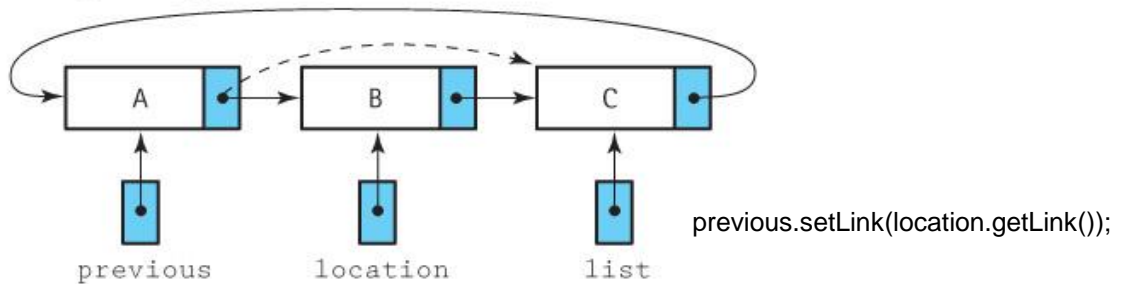
    moreToSearch = (location != null);
    while (moreToSearch && !found)
    {
        // move search to the next node
        previous = location;
        location = location.getLink();

        // check for a match
        if (location.getInfo().equals(target))
            found = true;

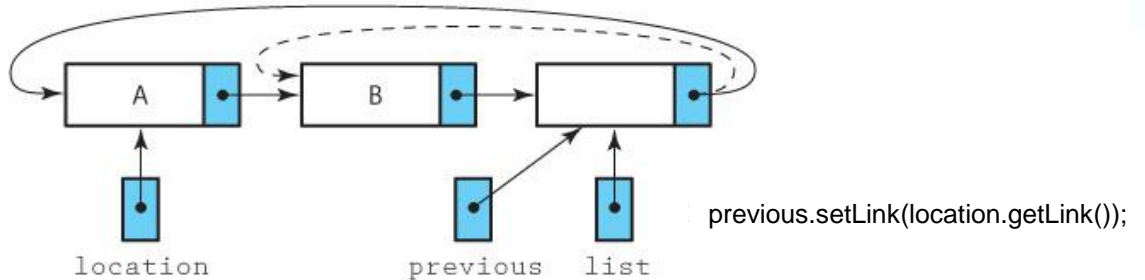
        moreToSearch = (location != list);
    }
}
```


The remove method

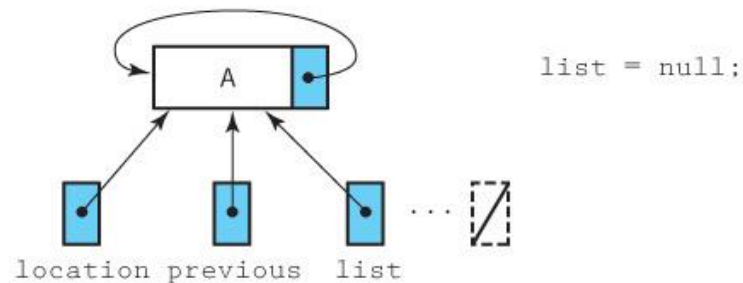
(a) The general case (remove B)



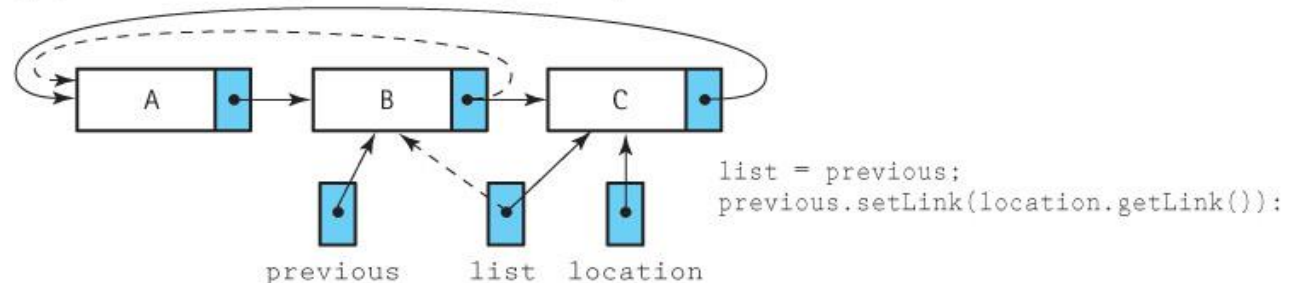
(b) Special case (?): Removing the first element (remove A)



(c) Special case: Removing the only element (remove A)



(d) Special case: Removing the last element (remove C)



The remove method

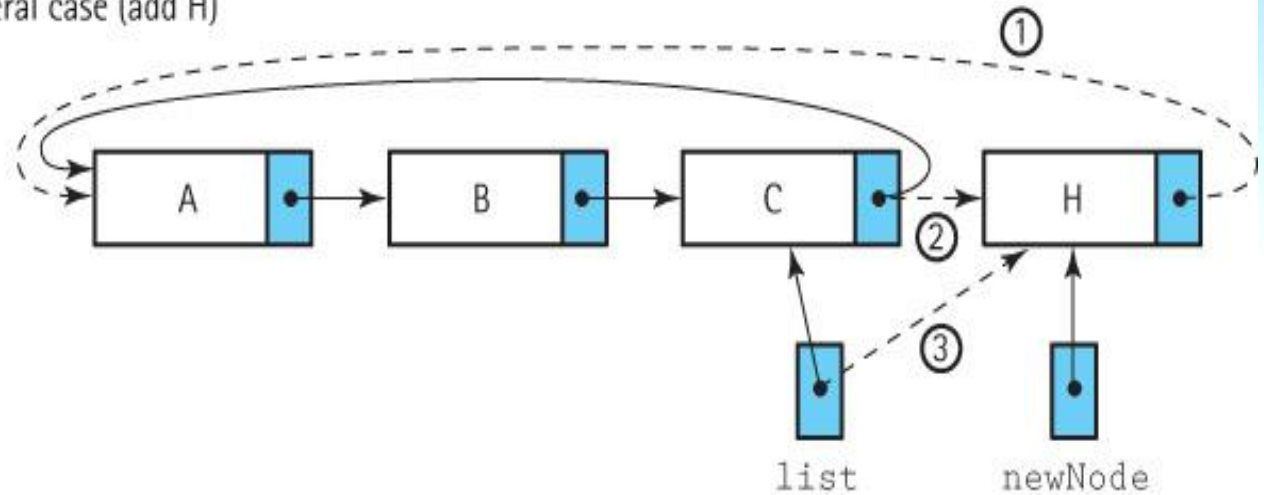
```
public boolean remove (Object element)
// Removes an element e from this list such that e.equals(element)
// and returns true; if no such element exists returns false.
{
    find(element);
    if (found)
    {
        if (list == list.getLink())        // if single element list
            list = null;
        else
            if (previous.getLink() == list) // if removing last node
                list = previous;
            previous.setLink(location.getLink()); // remove node
            numElements--;
    }
    return found;
}
```

The CRefUnsortedList Class

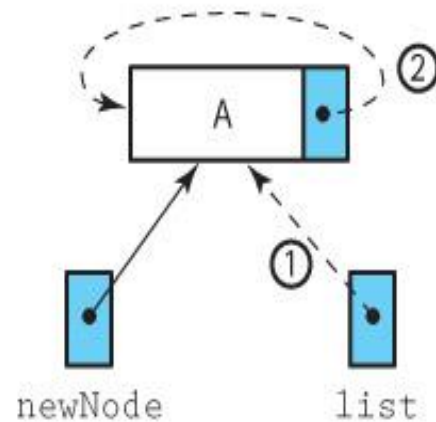
- To create the CRefUnsortedList class and complete our implementation of the unsorted circular list we just need to extend the CRefList class with an add method.
- To implement the add method:
 - create the new node using the LLObjectNode constructor
 - if adding to an empty list set the list variable to reference the new element and link the new element to itself
 - otherwise, add the element to the list in the most convenient place – at the location referenced by list
 - increment numElements.

Adding a node

(a) The general case (add H)



(b) Special case: The empty list (add A)



The add method

```
public void add(Object element)
// Adds element to this list.
{
    LLObjectNode newNode = new LLObjectNode(element);
    if (list == null)
    {
        // add element to an empty list
        list = newNode;
        newNode.setLink(list);
    }
    else
    {
        // add element to a non-empty list
        newNode.setLink(list.getLink());
        list.setLink(newNode);
        list = newNode;
    }
    numElements++;
}
```

Doubly Linked Lists

- **Doubly linked list** A linked list in which each node is linked to both its successor and its predecessor



DLLObjectNode Class

```
package support;

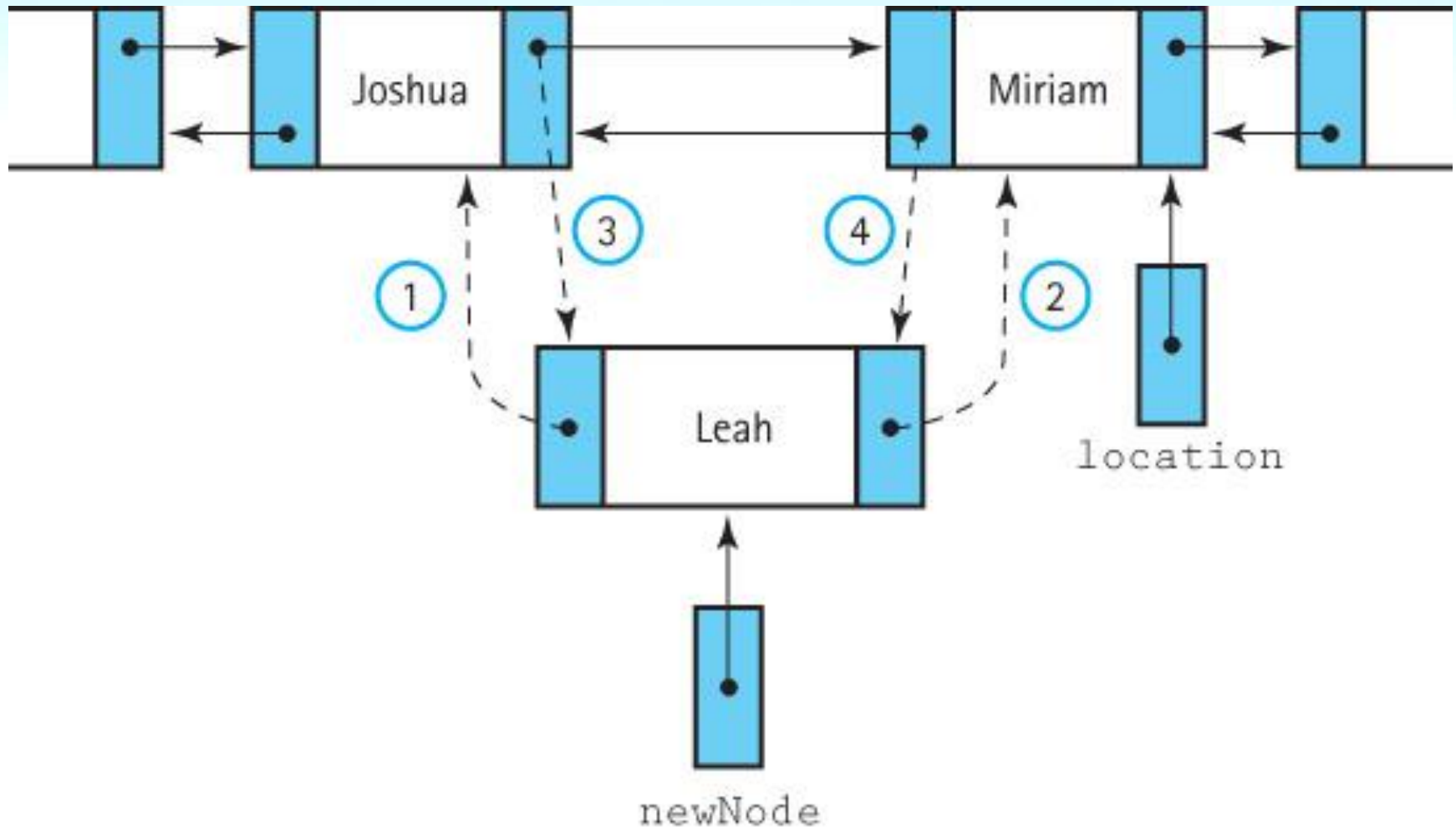
public class DLLObjectNode extends LLObjectNode
{
    private DLLObjectNode back;

    public DLLObjectNode(Object info)
    {
        super(info);
        back = null;
    }

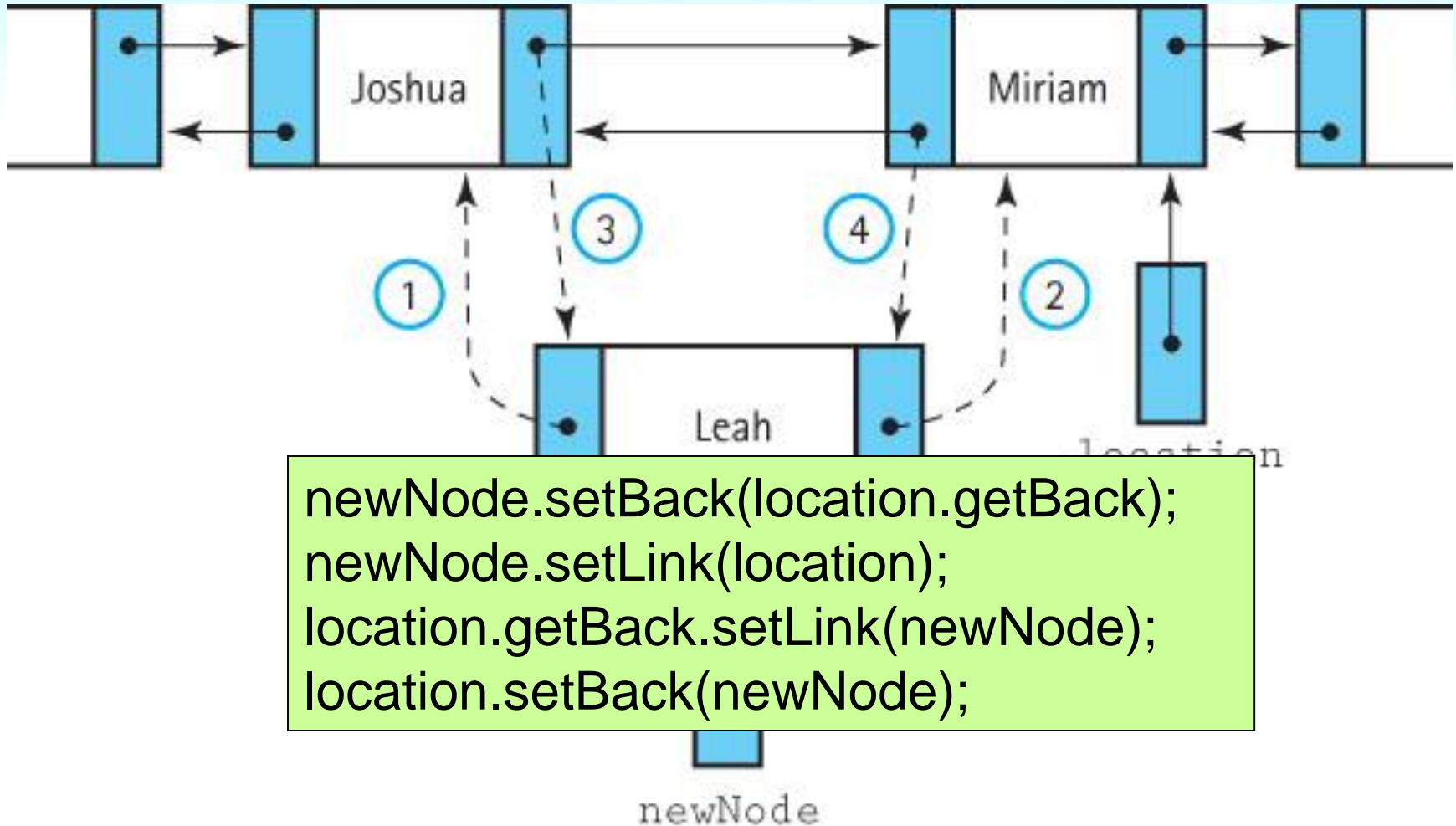
    public void setBack(DLLObjectNode back)
    // Sets back link of this DLLObjectNode.
    {
        this.back = back;
    }

    public DLLObjectNode getBack()
    // Returns back link of this DLLObjectNode.
    {
        return back;
    }
}
```

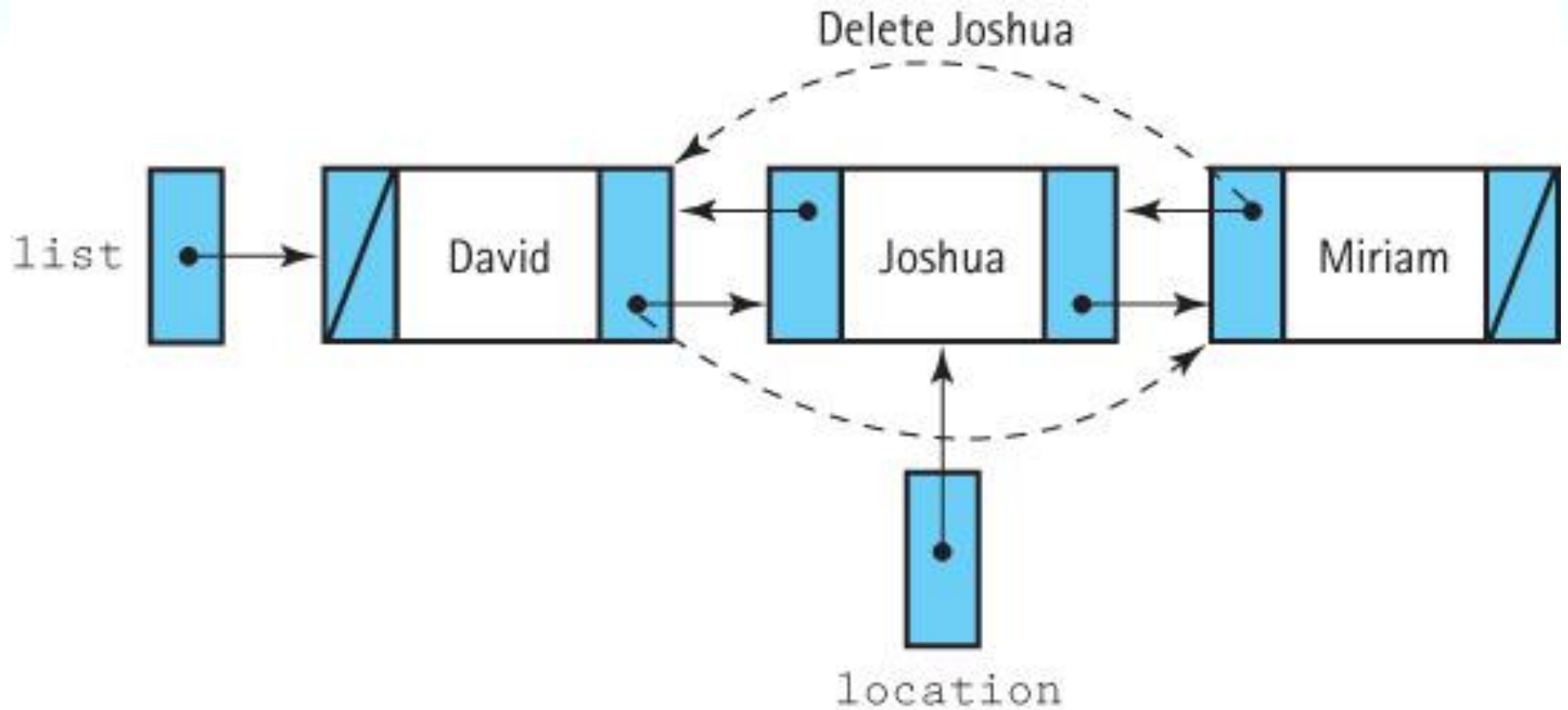
The add operation



The add operation

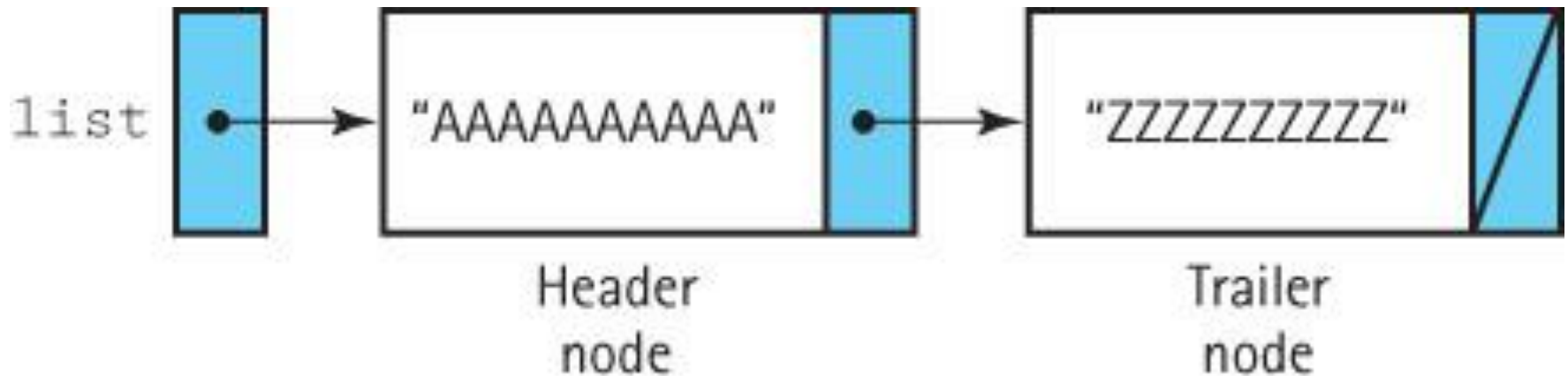


The remove operation



Linked Lists with Headers and Trailers

- **Header node** A placeholder node at the beginning of a list; used to simplify list processing
- **Trailer node** A placeholder node at the end of a list; used to simplify list processing

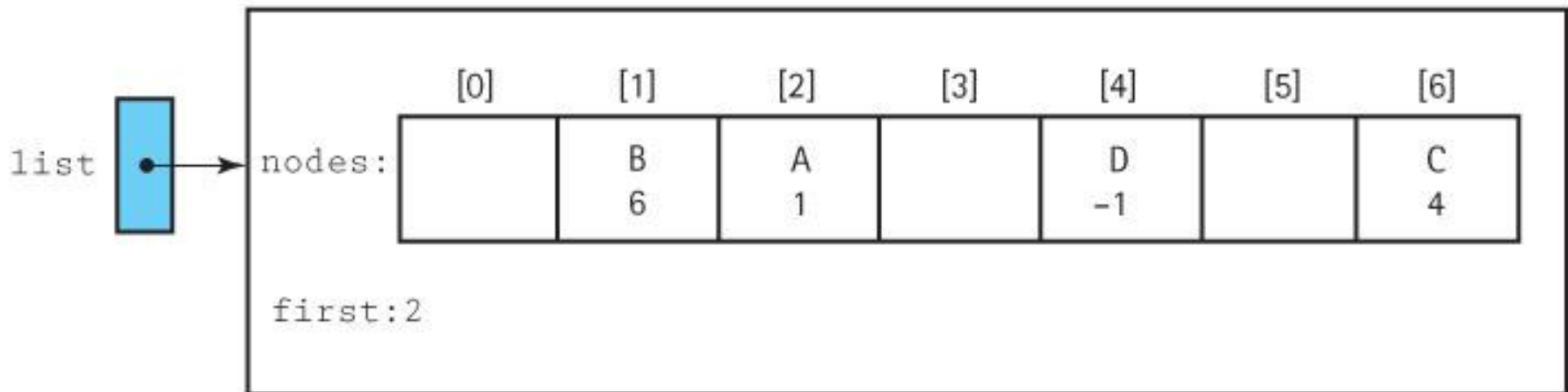


A Linked List as an Array of Nodes

(a) A linked list in dynamic storage



(b) A linked list in static storage



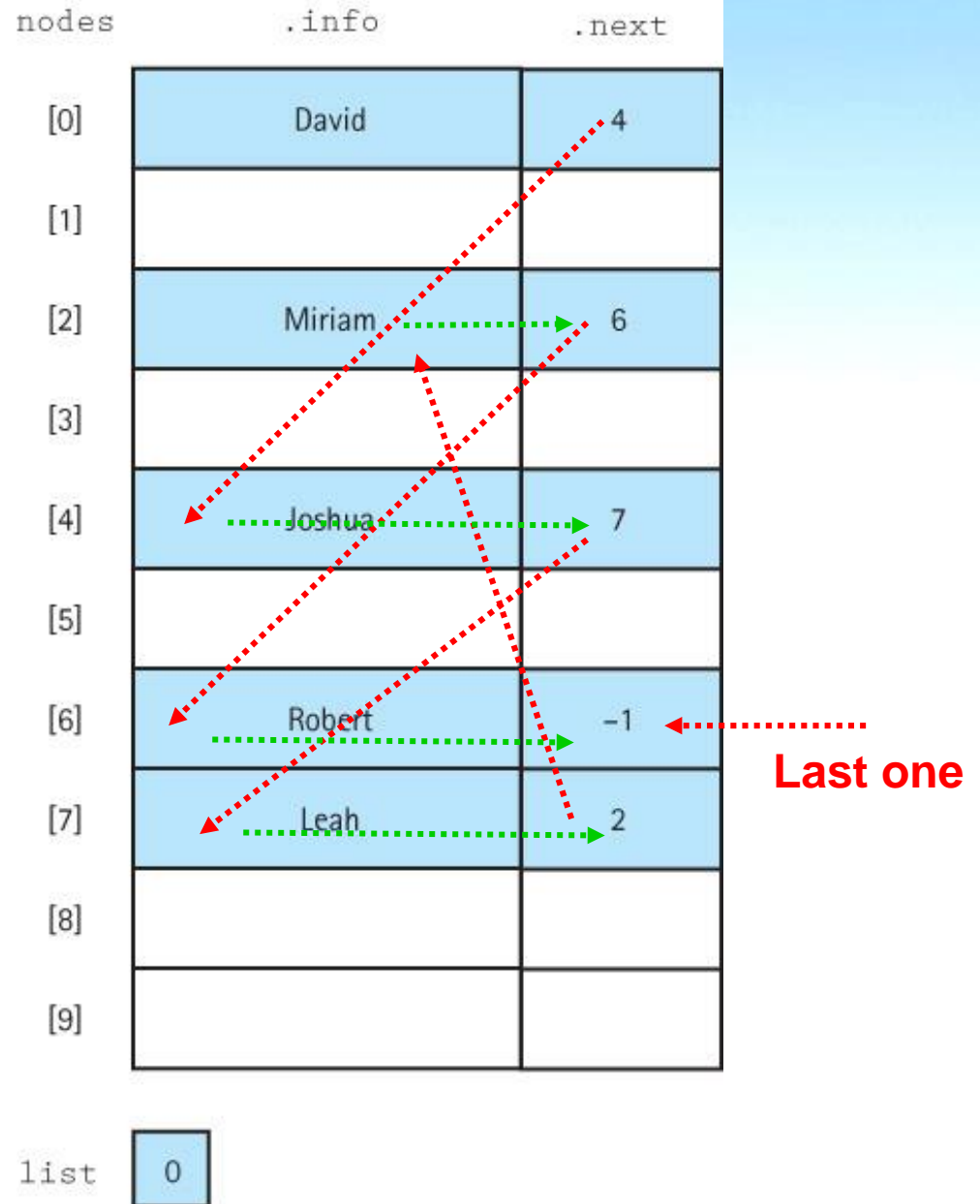
Why Use an Array?

- Sometimes managing the free space ourselves gives us greater flexibility
- There are programming languages that do not support dynamic allocation or reference types
- There are times when dynamic allocation of each node, one at a time, is too costly in terms of time

Boundedness

- A desire for static allocation is one of the primary motivations for the array-based linked approach
- We drop our assumption that our lists are of unlimited size in this section - our lists will not grow as needed.
- Applications should not add elements to a full list.
- Our list will export an `isFull` operation, in addition to all the other standard list operations

A sorted list

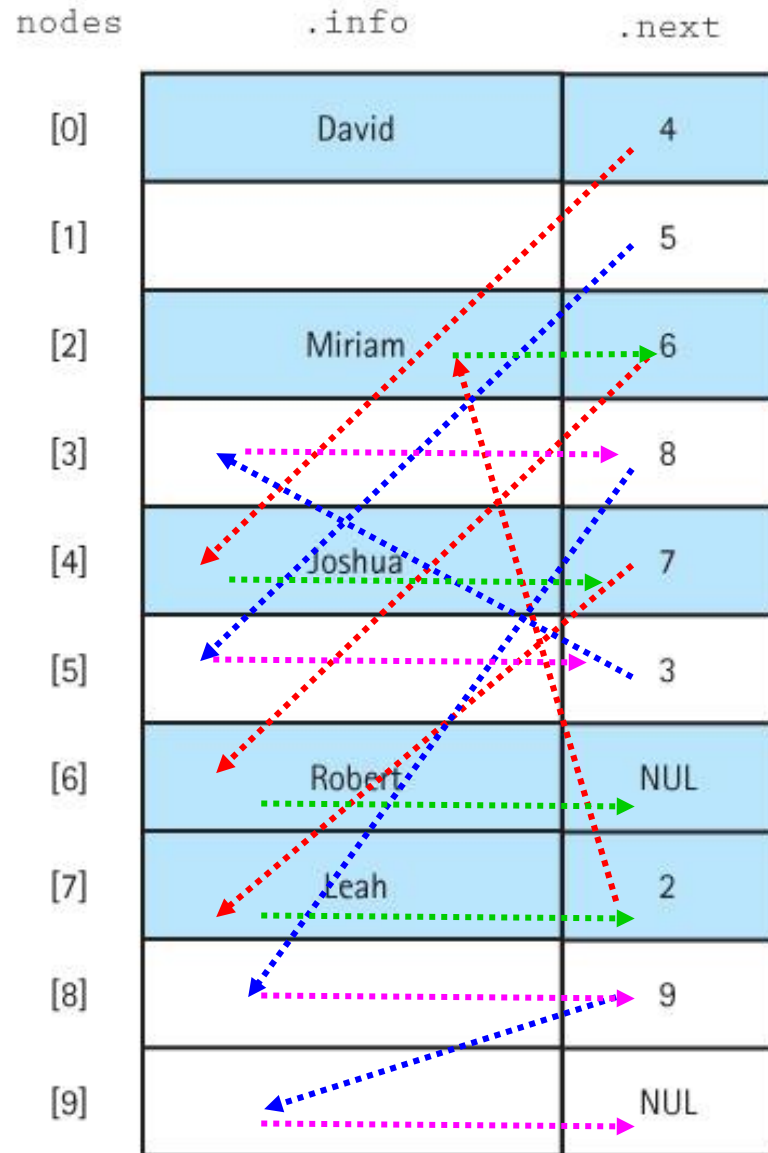


Implementation Issues

- We mark the end of the list with a “null” value
 - the “null” value must be an invalid address for a real list element
 - we use the value `-1`
 - we use the identifier `NUL` and define it to be `-1`

```
private static final int NUL = -1;
```
- We must directly manage the free space available for new list elements.
 - We link the collection of unused array elements together into a linked list of free nodes.
 - We write our own method to allocate nodes from the free space. We call this method `getNode`. We use `getNode` when we add new elements onto the list.
 - We write our own method, `freeNode`, to put a node back into the pool of free space when it is de-allocated.

A linked list and free space



list	0
free	1

More than one list

free 7

nodes	.info	.next
[0]	John	4
[1]	Mark	5
[2]		3
[3]		NUL
[4]	Nell	8
[5]	Naomi	6
[6]	Robert	NUL
[7]		2
[8]	Susan	9
[9]	Susanne	NUL

list1	0
list2	1