1st Internal Suggestions

1. Define finite automata. Why it is called finite?

Ans. Finite automata (Singular: automaton) are the machine format of regular expression, which is

the language format of type 3 grammar. A finite automata is defined as $M = \{Q, \Sigma, \delta, q0, F\}$.

Where *Q*: Finite non-empty set of states.

 Σ : Finite non-empty set of input symbols.

 δ : Transitional function.

*q*0: Beginning state.

F : Finite non-empty set of final states.

Finite automata are a type of finite state machines. It has finite number of states.

Finite automata can be

thought of a finite state machine without output.

- 2. Construction of DFA, NFA, Moore and Mealy Machine. (As discussed in class)
- 3. Problem on Minimization of DFA (Myhill-Nerode Theorem)
- 4. NFA to DFA conversion
- 5. Moore to Mealy, Mealy to Moore Conversion. (Practice problems from the class)

6. Arden's Theorem:

Statement: Let P and Q be two Regular Expression s over Σ . If P does not contain Λ , then for the equation R = Q + RP has a unique (one and only one) solution $R = QP^*$.

Proof:

If $R = QP^*$ is a solution of the equation R = Q + RP then by putting the value of R in the RHS in the equation we get –

$$R = Q + RP$$

$$R = Q + QP*P$$

 $R = Q (\Lambda + P*P)$ [By taking Q as common from both terms]

$$R = QP^*$$
 [By identity rule we know, $\Lambda + R^*R = R^*$]

So, from here it is proved that $R = QP^*$ is a solution of the equation R = Q + RP.

**Now we have to prove that $R = QP^*$ is the one and only one solution of the equation R = Q + RP.

As R = Q + RP, so put the value of R again and again in the RHS of the equation.

$$R = Q + RP$$

$$= Q + (Q + RP)P$$

$$= Q + QP + RPP$$

$$= Q(\Lambda + P) + RPP$$

$$= Q(\Lambda + P) + (Q + RP)PP$$

$$= Q(\Lambda + P) + QPP + RPPP$$

$$= Q(\Lambda + P + PP) + RPPP$$

After several steps we shall get

$$R = Q(\Lambda + P + PP + PPP + \dots + Pn) + RPn + 1.$$

$$R = Q(\Lambda + P + PP + PPP + ... + Pn)$$
 [By neglecting RPn+1]**

$$R = QP^*$$
 [As we know, $\Lambda + P + PP + PPP + ... + Pn = P^*$]

That means $R = QP^*$ is the unique solution of the equation R = Q + RP.

**(This part only contains explanation. You can exclude this.

Now let a string w belongs to R. If w is a Λ string then in which part the Λ will belong?

This string will belong to either in $Q(\Lambda + P + P2 + P3 + ---- Pn)$ part or in RPn + 1 part.

However, according to point number (ii) P does not contain Λ , so the string w does not belong to RPn + 1.

So, it will obviously belong to $Q(\Lambda + P + P2 + P3 + ---- Pn)$ which is nothing but QP^* . $[(\Lambda + P + P2 + P3 + ---- Pn)$ is any combination of P including $\Lambda]$.)