

# Contest design with a finite type-space: A unifying approach

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- This paper: contests with a finite type-space





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- **Sisak [2009]:** “The case of asymmetric individuals, where **types are private information but drawn from discrete, identical** or maybe even different distributions, has not been addressed so far. From the results discussed above, especially on asymmetric types with full information, one could **conjecture that multiple prizes might be optimal even with linear costs.**”

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- **Finite types:** Xiao [2018], Liu and Chen [2016], Szech [2011], Konrad [2004], Chen [2021]

# Model

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- Design problem: given a budget, find  $v = (v_1, \dots, v_N)$  to maximize

$$\mathbb{E}[X] = \sum_{k=1}^K p_k \mathbb{E}[X(\theta_k)].$$

## Lemma 1.

*For any contest  $v = (v_1, \dots, v_{N-1}, 0)$ , there is a unique symmetric Bayes-Nash equilibrium and it is such that there exist boundary points  $b_1 < b_2 < \dots < b_K$  so that for any  $\theta_k \in \Theta$ , an agent of type  $\theta_k$  mixes between  $[b_{k-1}, b_k]$  with  $b_0 = 0$ .*

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- More efficient agents (those with lower  $\theta$ ) exert higher effort
- Exhibits both the mixed structure from complete information, and monotonic structure from continuum type-space environments

# Proof sketch

Let  $F_k$  denote the equilibrium cdf's,  $[a_k, b_k]$  denote the support of  $F_k$ , and  $u_k$  denote the equilibrium utility.

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Together, the properties imply the structure in the equilibrium.

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- By the indifference condition, for all  $x_k \in [b_{k-1}, b_k]$ ,

$$\pi_v(P_{k-1} + p_k F_k(x_k)) - \theta_k c(x_k) = u_k,$$

where  $\pi_v(t) = \sum_{m=1}^N v_m H_{N-m}^{N-1}(t)$ , and  $H_{N-m}^{N-1}(t)$  is the probability  $Y \sim \text{Bin}(N-1, t)$  takes the value  $N-m$ .

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- Solve for  $u_k$  and  $b_k$  by using  $u_1 = 0$ ,  $F_k(b_k) = 1$  and  $F_{k+1}(b_k) = 0$ .

## Theorem 2.

*For any  $v$ , the unique symmetric Bayes-Nash equilibrium is such that for any  $\theta_k \in \Theta$ , the distribution function  $F_k : [b_{k-1}, b_k] \rightarrow [0, 1]$  is defined by*

$$\pi_v(P_{k-1} + p_k F_k(x_k)) - \theta_k c(x_k) = u_k \text{ for all } x_k \in [b_{k-1}, b_k],$$

*where the points  $b = (b_1, \dots, b_K)$  and utilities  $u = (u_1, \dots, u_K)$  are*

$$c(b_k) = \sum_{j=1}^k \frac{\pi_v(P_j) - \pi_v(P_{j-1})}{\theta_j} \text{ for any } k \in \{1, 2, \dots, K\},$$

*and*

$$u_k = \theta_k \left[ \sum_{j=1}^{k-1} \pi_v(P_j) \left( \frac{1}{\theta_{j+1}} - \frac{1}{\theta_j} \right) \right] \text{ for any } k \in \{1, 2, \dots, K\}.$$

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- For utilities, observe that for any  $k \in \{1, \dots, K\}$  and any prize  $m \in \{1, \dots, N-1\}$ ,

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- For expected effort, we can write

$$\pi_v(P_{k-1} + p_k F_k(X_k)) - \theta_k c(X_k) = u_k.$$

This gives  $\mathbb{E}[X_k]$ , and since  $\mathbb{E}[X] = \sum_{k=1}^K p_k \mathbb{E}[X_k]$ , we get

$$\mathbb{E}[X] = \sum_{k=1}^K \int_{P_{k-1}}^{P_k} g \left( \frac{\pi_v(t) - u_k}{\theta_k} \right) dt.$$

# Optimal contest: Linear costs

## Lemma 3.

Suppose  $c(x) = x$ . For any contest  $v = (v_1, \dots, v_{N-1}, 0)$ , the expected equilibrium effort of an arbitrary agent is  $\mathbb{E}[X] = \sum_{m=1}^{N-1} \alpha_m v_m$  where

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- Effect of competition: compute  $\alpha_b - \alpha_w$  (maybe  $< 0$  when  $b \neq 1$ )

# Effect of competition: Linear to general costs

## Theorem 4.

*For any pair of prizes  $b, w \in \{1, \dots, N - 1\}$  with  $b < w$  such that, either  $b = 1$  or  $\left( \frac{\partial u_K}{\partial v_b} - \frac{\partial u_K}{\partial v_w} \right) \leq 0$ , the following hold:*

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- If  $K = 1$ , increasing competition encourages effort if costs are concave, and discourages effort if costs are convex.
- For  $K > 1$ , since  $\alpha_1 - \alpha_w > 0$ , the winner-takes-all contest is strictly optimal under concave costs.

## Theorem 5.

*Suppose there are  $N$  agents and consider a fixed contest  $v = (v_1, \dots, v_{N-1}, 0)$ . Let  $G : [\underline{\theta}, \bar{\theta}] \rightarrow [0, 1]$  be a differentiable CDF and let  $G^1, G^2, \dots$ , be any sequence of CDF's, each with a finite support, such that for all  $\theta \in [\underline{\theta}, \bar{\theta}]$ ,*

$$\lim_{n \rightarrow \infty} G^n(\theta) = G(\theta).$$

*Let  $F^n : \mathbb{R} \rightarrow [0, 1]$  denote CDF of the equilibrium effort under the finite type-space distribution  $G^n$ , and let  $F : \mathbb{R} \rightarrow [0, 1]$  denote CDF of the equilibrium under continuum type-space distribution  $G$ . Then, the sequence of CDF's  $F^1, F^2, \dots$ , converges to the CDF  $F$ , i.e., for all  $x \in \mathbb{R}$ ,*

$$\lim_{n \rightarrow \infty} F^n(x) = F(x).$$

- Intuitively, as  $K$  increases,  $[b_{k-1}, b_k]$  shrinks, and converges to the pure-strategy equilibrium effort under the continuum type-space.

# Summary

- Study effect of increasing competition in contests with a finite type-space
- Provide a unifying approach to studying contests simultaneously in complete and incomplete information environments
- Identify effects under linear costs, and find conditions under which they extend to general costs, which pertain to how competition affects the equilibrium utility of the most efficient agent
- Generate insights into what drives some of the differences in the complete and incomplete information environments
- Solve the design problem under linear and concave costs, showing that the winner-takes-all contest is optimal with any limited uncertainty

# Thank you!