# Idris A Programming Language with Dependent Types

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#### What is Idris?

"Idris is a general purpose pure functional programming language with dependent types."

The Idris Website

- **Version 0.1.3** of Idris was released in December of 2009.
- Version 1.2.0 is the latest stable release and was released on January 9, 2018.
- Idris was named after the singing dragon in the 1970s UK children's television program Ivor the Engine.
- Idris development is led by Edwin Brady at the University of St. Andrews.

# The Obligatory Picture of This Madman



#### **Properties of Idris**

- Idris can be interpreted, transpiled, or compiled.
- Idris is statically typed.
- Idris is strongly typed.
- Idris has first class functions, much like Haskell.
- Idris has **first class types**. This means that types can be treated as data.

#### **Idris Features**

Idris is a general purpose language, and thus it has a lot of features. We will focus on the following aspects of the language.

- Dependent Types
- Haskell-like Syntax
- Proof Assistant

# Dependent Types

Consider these two arrays:

[1, 2, 3] [1, 2, 3, 4]

What are their types? The first is a 3-dimensional array, the second is 4-dimensional.

# **Idris Syntax: Function Signatures**

The Idris function signature syntax is *very* similar to the Haskell function signature syntax. Here are a few examples of Idris function signatures:

```
even : Nat -> Bool
add : Nat -> Nat -> Nat
foo : (a:Nat) -> (b:Nat) -> a = b
bar : (a:Nat) -> (b:Nat) -> LTE a b
```

If you are familiar with Haskell, you will note the use of : rather than ::. This makes it look a bit more like a mathematical function definition:

$$f: \mathbb{N} \to \mathbb{N}$$
.

You will also note that instead of the (Type x) = x syntax, it uses a more concise (x:Type) syntax.

# Idris Syntax: Currying and Pattern Matching

Because of its foundation in Lambda Calculus, all functions only take a single argument.

We can still handle multiple arguments using *currying*. For example, the plus operator is defined as follows:

```
plus : Nat -> Nat -> Nat
plus Z y = y
plus (S k) y = S (plus k y)
```

Like Haskell, functions are implemented using pattern matching.

# Idris Syntax: Type Definition Syntax

Idris defines several primitives including Int, Integer, Double, Char, String, and Ptr.

There are a bunch of other data types defined in the standard library including Nat and Bool.

Idris allows programmers to define their own data types. Again, the syntax is similar to Haskell.

```
data Nat = Z | S Nat
data List a = Nil | (::) a (List a)
```

### **Idris Syntax: Holes**

Idris allows you to leave some of your code unfinished. For example, if we write the following code in a file called even.idr:

```
even : Nat -> Bool
even Z = True
even (S k) = ?even_rhs
```

And then load it into Idris:

```
Holes: even_rhs
even> :t even_rhs
k : Nat
```

:Idris> :l even

even\_rhs : Bool
Holes: even\_rhs

# Using Idris as a Proof Assistant

A proof assistant is a software tool to assist with the development of formal proofs by human-machine collaboration.

The ldris type system is robust enough that it can be used as a proof assistant.

# The Equality Type

The basis for proofs in Idris is the (=) type constructor. It constructs a type that is dependent on two expressions.

- Any Nat is a natural number.
- Any Vect 2 Nat is a list of two natural numbers.
- Any (=) (2 + 2) 4 is a proof that 2+2 and 4 have the same value.
- Any 1 = 3 is a proof that 1 and 3 have the same value.

It is impossible to create an instance of 1 = 3

### **Quotes From Our Exploration**

"The concept of a programming language in which the possibility of inline assembly is an entirely foreign concept hurts my brain."

"Where do I put it? Do I put it in the type?"

"When your Rust program compiles, you know it won't segfault, or give you any undefined behavior at runtime. When your Idris program compiles, you throw away your executable, and publish your dissertation."

