

Idris

A Programming Language with Dependent Types

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Introduction

What is Idris?

"Idris is a general purpose pure functional programming language with dependent types."

The Idris Website

- **Version 0.1.3** of Idris was released in December of 2009.
- **Version 1.2.0** is the latest stable release and was released on January 9, 2018.
- Idris was named after the singing dragon in the 1970s UK children's television program *Ivor the Engine*.
- Idris development is led by Edwin Brady at the University of St. Andrews.

The Obligatory Picture of This Madman



Properties of Idris

- Idris can be **interpreted, transpiled, or compiled**
- Idris is **statically typed**
- Idris is **strongly typed**
- Idris is **purely functional** (much like Haskell)
- Idris has **first class types** (types can be treated as data)
- Idris has **dependent types** (the types are all high on something)

Idris Features

Idris is a general purpose language, and thus it has a lot of features. We will focus on the following aspects of the language.

- Haskell-like Syntax
- Dependent Types
- Proof Assistant

Syntax

Function Signatures

The Idris function signature syntax is *very* similar to the Haskell function signature syntax. Here are a few examples of Idris function signatures:

```
even : Nat -> Bool  
add  : Nat -> Nat -> Nat  
foo  : (a:Nat) -> (b:Nat) -> a = b  
bar  : (a:Nat) -> (b:Nat) -> LTE a b
```

If you are familiar with Haskell, you will note the use of `:` rather than `::`. This makes it look a bit more like a mathematical function definition:

$$f : \mathbb{N} \rightarrow \mathbb{N}.$$

You will also note that instead of the `(Type x) => x` syntax, it uses a more concise `(x:Type)` syntax.

Currying and Pattern Matching

Because of its foundation in Lambda Calculus, all functions only take a single argument. We can still handle multiple arguments using *currying*. For example, the `plus` operator is defined as follows:

```
plus : Nat -> Nat -> Nat
plus  Z      y  = y
plus  (S k)  y  = S (plus k y)
```

Like Haskell, functions are implemented using *pattern matching*.

Type Definition Syntax

Idris defines several primitives including `Int`, `Integer`, `Double`, `Char`, `String`, and `Ptr`.

There are a bunch of other data types defined in the standard library including `Nat` and `Bool`.

Idris allows programmers to define their own data types. Again, the syntax is similar to Haskell.

```
data Nat      = Z    | S Nat  
data List a = Nil | (::) a (List a)
```

Holes

Idris allows you to leave some of your code unfinished. For example, if we write the following code in a file called `even.idr`:

```
even : Nat -> Bool
even Z = True
even (S k) = ?even_rhs
```

And then load it into Idris:

```
:Idris> :l even
Holes: even_rhs
even> :t even_rhs
      k : Nat
```

```
-----
even_rhs : Bool
Holes: even_rhs
```

Dependent Types

Types

Idris has familiar, Haskell-ish types:

- `Nat` - A natural number
- `Bool` - A boolean
- `Char` - A single character
- `List Int` - A list of integers
- `Nat -> Bool` - A function that takes a natural number and produces a boolean
- `(Nat, Nat)` - A tuple of two natural numbers
- `Int -> Int -> Int` - A function that takes two arguments

Types as Data

Unlike Haskell, data types can be stored, passed, and constructed like data:

```
an_int : Int  
an_int = 5
```

```
a_type : Type  
a_type = Int
```

We could write a function to choose between an `Int` and a `Nat`:

```
PickInt : Bool -> Type  
PickInt True = Int  
PickInt False = Nat
```

This is called a **type constructor**.

Dependent Types

Any expression that returns Type can be used as type itself:

```
foo : PickInt (True && False)  
foo = 5
```

```
bar : case False of  
      True => List Char  
      False => String  
bar = "Hello, World!"
```

These are called **dependent types**, since they *depend* on data.

Useful Dependent Types

List and Vect are examples of type constructors:

- List Int is a dynamically sized list of integers.
- Vect 10 Int is a list of exactly 10 integers.

Since type constructors are simply functions, they support things like currying:

```
TwoOf : Type -> Type
TwoOf = Vect 2
```


The Equality Type Constructor

The basis for proofs in Idris is the $(=)$ function. It takes two inputs, and the return type is a proof that the two inputs have the same value.

- Any **Nat** is a natural number.
- Any **Vect** 2 **Nat** is a list of two natural numbers.
- Any $(=)$ $(2 + 2)$ 4 is a proof that $2+2$ and 4 have the same value.
- Any $1 = 3$ is a proof that 1 and 3 have the same value.
- Any even $x =$ **True** is a proof that x is even

Using Idris as a Proof Assistant

What is a Proof Assistant?

A proof assistant is a software tool to assist with the development of formal proofs by human-machine collaboration.

The Idris type system is robust enough that it can be used as a proof assistant.

How can Idris be a Proof Assistant?

Recall from above that equality is a type constructor. This means that we can pass equalities in and out of functions. This is the basis for all proofs in Idris.

Take this example function declaration:

```
plusReduces : (n:Nat) -> plus Z n = n
```

This is a function which takes any $n \in \mathbb{N}$, and returns a proof that $0 + n = n$. Any successful implementation of this function will prove that $0 + n = n$.

Warning

LIVE DEMO AHEAD

We are not responsible for any harm done to your brain by viewing the following code.

Quotes From Our Exploration

"The concept of a programming language in which the possibility of inline assembly is an entirely foreign concept hurts my brain."

"Where do I put it? Do I put it in the type?"

"When your Rust program compiles, you know it won't seg-fault, or give you any undefined behavior at runtime. When your Idris program compiles, you throw away your executable, and publish your dissertation."

Questions?