

# Contents

## 1 Discrete structures



# Section outline

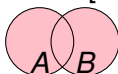
## 1 Discrete structures

- Sets
- Relations
- Lattices
- Lattices (contd.)
- Boolean lattice
- Boolean lattice structure
- Boolean algebra
- Additional Boolean algebra properties

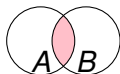


# Sets

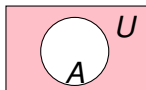
- A set  $A$  of elements:  $A = \{a, b, c\}$
- Natural numbers:  $\mathbb{N} = \{0, 1, 2, 3, \dots\}$  or  $\{1, 2, 3, \dots\} = \mathbb{Z}^+$
- Integers:  $\mathbb{Z} = \{\dots, -3, -2, -1, 0, 1, 2, 3, \dots\}$
- Universal set:  $U$       Empty set:  $\emptyset = \{\}$
- $S = \{X | X \notin X\}$        $S \in S?$  [Russell's paradox]



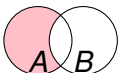
- Set union:  $A \cup B$



- Set intersection:  $A \cap B$



- Complement:  $\bar{S}$

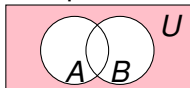


- Set difference:  $A - B = A \cap \bar{B}$

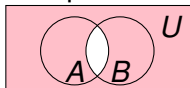


# Sets (contd.)

- Complement of union (De Morgan):  $\overline{A \cup B} = \bar{A} \cap \bar{B}$



- Complement of intersection (De Morgan):  $\overline{A \cap B} = \bar{A} \cup \bar{B}$



- Power set of  $A$ :  $\mathcal{P}(A)$

$$\mathcal{P}(\{a, b\}) = \{\emptyset, \{a\}, \{b\}, \{a, b\}\}$$

- Non-empty  $X_1, \dots, X_k$  is a partition of  $A$  if  $A = X_1 \cup \dots \cup X_k$  and  $X_i \cap X_j = \emptyset \mid i \neq j$

$A \cap \bar{B}$ ,  $B \cap \bar{A}$ ,  $A \cap B$  and  $\overline{A \cup B}$  constitute a partition of  $U$



# Set algebra

Idempotence	$A \cup A = A$	$A \cap A = A$
Associativity	$(A \cup B) \cup C = A \cup (B \cup C)$	$(A \cap B) \cap C = A \cap (B \cap C)$
Commutativity	$A \cup B = B \cup A$	$A \cap B = B \cap A$
Distributivity	$A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$	$A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$
Identity	$A \cup \{\} = A, A \cup U = U$	$A \cap \{\} = \{\}, A \cap U = A$
Involution	$\overline{\overline{A}} = A$	
Complements	$\overline{\overline{U}} = \{\}, A \cup \overline{A} = U$	$\overline{\{\}} = U, A \cap \overline{A} = \{\}$
DeMorgan	$\overline{A \cup B} = \overline{A} \cap \overline{B}$	$\overline{A \cap B} = \overline{A} \cup \overline{B}$



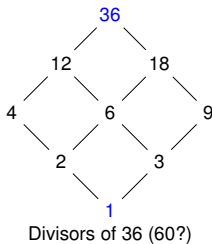
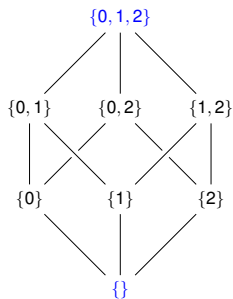
# Relations

- Tuple:  $\langle a, b \rangle, \langle 4, b, \alpha \rangle$
- Cartesian product:  $A \times B = \{\langle a, b \rangle \mid a \in A, b \in B\}$   
 $\{a, b, c\} \times \{\alpha, \beta\} = \{\langle a, \alpha \rangle, \langle b, \alpha \rangle, \langle c, \alpha \rangle, \langle a, \beta \rangle, \langle b, \beta \rangle, \langle c, \beta \rangle\}$   
 $\mathbb{N} \times \mathbb{N} = \{\langle i, j \rangle \mid i, j \geq 1\}$
- Binary relation  $\mathcal{R}$  on sets  $A$  and  $B$ :  $R \subseteq A \times B$
- Characteristic function of  $\mathcal{R}$ :  $\chi_{\mathcal{R}}(a, b) = \begin{cases} 1 & \text{if } \langle a, b \rangle \in \mathcal{R} \\ 0 & \text{if } \langle a, b \rangle \notin \mathcal{R} \end{cases}$
- $\mathcal{R} \subseteq A \times A$  is reflexive if  $\forall x \in A. x\mathcal{R}x$
- $\mathcal{R} \subseteq A \times A$  is symmetric if  $\forall x, y \in A. x\mathcal{R}y \Rightarrow y\mathcal{R}x$
- $\mathcal{R} \subseteq A \times A$  is transitive if  $\forall x, y, z \in A. x\mathcal{R}y \wedge y\mathcal{R}z \Rightarrow x\mathcal{R}z$
- $\mathcal{R} \subseteq A \times A$  is antisymmetric if  $\forall x, y \in A. x\mathcal{R}y \wedge y\mathcal{R}x \Rightarrow x = y$
- Equivalence relation:  $\mathcal{R}$  is reflexive, symmetric and transitive
- An equivalence relation induces a partition and vice versa
- Partial order:  $\mathcal{R}$  is reflexive, antisymmetric and transitive



# Relations (contd.)

- Connected relation:  $\forall x, y \in A$ , either  $x\mathcal{R}y$  or  $y\mathcal{R}x$
- Total order: Connected partial order (eg  $\leq$  on  $\mathbb{R}$ )
- Irreflexive relation:  $\forall x \in A, \langle x, x \rangle \notin \mathcal{R}$
- Asymmetric relation:  $\langle x, y \rangle \in \mathcal{R} \Rightarrow \langle y, x \rangle \notin \mathcal{R}$
- Strict order:  $\mathcal{R}$  is irreflexive and transitive ( $\therefore$  asymmetric)
- If  $\preceq$  is a PO on  $A$ , then  $<: x < y \equiv x \preceq y \wedge x \neq y$  is a SO on  $A$
- If  $<$  is a SO on  $A$ , then  $\preceq: x \preceq y \equiv x < y \vee x = y$  is a PO on  $A$



Suppose  $\langle A, \preceq \rangle$  is a poset,  
 $M \in A$  ( $m \in A$ ),  $S \subseteq A$   


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 $M$  ( $m$ ) is a maximal (minimal)  
 element of  $S$  iff  $M \in S$  ( $m \in S$ )  
 and  $\nexists x \in S$  st  $M < x$  ( $x < m$ )  


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 $M$  ( $m$ ) is a maximum (minimum)

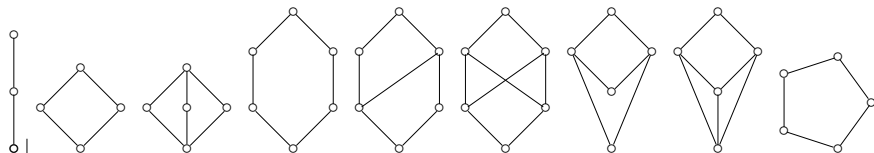


# Lattices

Let  $\langle A, \preceq \rangle$  be a poset, let  $x, y \in A$

- The *meet* of  $x$  and  $y$  ( $x \wedge y$ ), is the maximum of all lower bounds for  $x$  and  $y$ :  $x \wedge y = \max \{w \in A : w \preceq x, w \preceq y\}$ , *glb* for  $x$  and  $y$
- The *join* of  $x$  and  $y$  ( $x \vee y$ ), is the minimum of all upper bounds for  $x$  and  $y$ :  $x \vee y = \min \{z \in A : x \preceq z, y \preceq z\}$ , *lub* for  $x$  and  $y$

A poset  $\langle A, \preceq \rangle$  is a lattice iff every pair of elements in  $A$  have both a meet and a join





# Lattices (contd.)

Basic order properties of meet and join

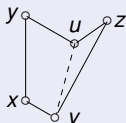
- $x \wedge y \preceq \{x, y\} \preceq x \vee y$
- $x \preceq y$  iff  $x \wedge y = x$
- $x \preceq y$  iff  $x \vee y = y$
- If  $x \preceq y$ , then  $x \wedge z \preceq y \wedge z$  and  $x \vee z \preceq y \vee z$
- If  $x \preceq y$  and  $z \preceq w$ , then  $x \wedge z \preceq y \wedge w$  and  $x \vee z \preceq y \vee w$

## Theorem

*If  $x \preceq y$ , then  $x \wedge z \preceq y \wedge z$  and  $x \vee z \preceq y \vee z$*

## Proof.

- Let  $v = x \wedge z$  and  $u = y \wedge z$
- By transitivity,  $v$  is a lb for  $y$  and  $z$
- By definition of  $\wedge$ ,  $v \preceq u$  (as  $u$  is the maximum among all lbs)



Similarly, the other clause may be proven



# Lattices (contd.)

**Commutativity**  $x \wedge y = y \wedge x$ ,  $x \vee y = y \vee x$

**Associativity**  $(x \wedge y) \wedge z = x \wedge (y \wedge z)$ ,  $(x \vee y) \vee z = x \vee (y \vee z)$

**Absorption**  $x \wedge (x \vee y) = x$ ,  $x \vee (x \wedge y) = x$

**Idempotence**  $x \wedge x = x$ ,  $x \vee x = x$

## Associativity of meet.

- $(x \wedge y) \wedge z \preceq x \wedge y \preceq x$  [ $x \wedge y \preceq \{x, y\}$  applied twice]
- $(x \wedge y) \wedge z \preceq x$  [transitivity of  $\preceq$ ]
- $(x \wedge y) \preceq y$  [ $x \wedge y \preceq \{x, y\}$ ]
- $(x \wedge y) \wedge z \preceq y \wedge z$  [If  $x \preceq y$ , then  $x \wedge z \preceq y \wedge z$ ]
- Thus  $(x \wedge y) \wedge z$  is a lb of both  $x$  and  $y \wedge z$
- $\therefore (x \wedge y) \wedge z \preceq x \wedge (y \wedge z)$  [glb of  $x$  and  $y \wedge z$ ]
- Also,  $x \wedge (y \wedge z) \preceq (x \wedge y) \wedge z$  [on similar lines]
- $\therefore (x \wedge y) \wedge z = x \wedge (y \wedge z)$  [if  $a \preceq b$  and  $b \preceq a$  then  $a = b$ ]



# Lattices (contd.)

## Absorbtion.

- $x \preceq x \vee y$  [ $\{x, y\} \preceq x \vee y$ ]

- $\therefore x \wedge (x \vee y) = x$  [ $x \preceq y$  iff  $x \wedge y = x$ ]



## Idempotence.

- $x \wedge x = x \wedge (x \vee (x \wedge y)) = x$  [Absorbtion, applied twice]



## Principle of Duality

The dual of any theorem in a lattice is also a theorem.



## Lattices (contd.)

**Bounded lattice:** It has a maximum element (1) and a minimum element (0), in which case identity properties hold:

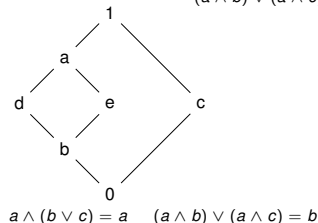
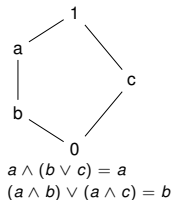
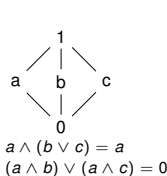
- $0 \vee x = x = x \vee 0, \quad 1 \wedge x = x = x \wedge 1$
- $0 \wedge x = 0 = x \wedge 0, \quad 1 \vee x = 1 = x \vee 1$

Every finite lattice is bounded

**Distributive lattice:** If  $\forall x, y, z \in A$ ,

- $x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z)$  and
- $x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z)$

Are these lattices distributive?



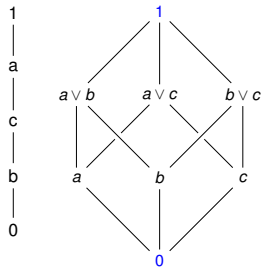
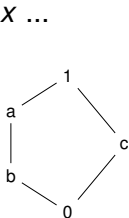
**Is  $\mathcal{P}(A)$  for set  $A$  distributive?**

- ?  $x \cap (y \cup z) = (x \cap y) \cup (x \cap z)$  and
- ?  $x \cup (y \cap z) = (x \cup y) \cap (x \cup z)$

# Complemented lattice

- *Complement in a bounded lattice:*  $z$  is the complement of  $x$  iff
  - $x \wedge z = 0$  and
  - $x \vee z = 1$
- **Bounded complemented lattice:** every element has a complement
- In a bounded distributive lattice with minimum 0 and maximum 1, the complements of elements are unique, provided they exist – let  $\bar{x}$  and  $z$  be complements of  $x$  ...

- $\bar{x} = \bar{x} \wedge 1 = \bar{x} \wedge (x \vee z) =$
- $(\bar{x} \wedge x) \vee (\bar{x} \wedge z) =$
- $0 \vee (\bar{x} \wedge z) =$
- $(x \wedge z) \vee (\bar{x} \wedge z) =$
- $(x \vee \bar{x}) \wedge z = 1 \wedge z = z$



# Boolean lattice

**Boolean lattice:** Bounded complemented distributive lattice

- Extreme elements: Max: 1, Min: 0
- Distributivity holds
- Every element has unique complement
- De Morgan's laws apply

## De Morgan's laws in a Boolean lattice $\langle \mathcal{A}, \preceq, -, 0, 1 \rangle$

$$1 \quad \overline{x \wedge y} = \bar{x} \vee \bar{y}$$

$$2 \quad \overline{x \vee y} = \bar{x} \wedge \bar{y}$$

Meet of complements is 0

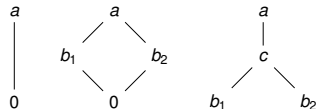
- $(x \wedge y) \wedge (\bar{x} \vee \bar{y}) =$   
 $(x \wedge y \wedge \bar{x}) \vee (x \wedge y \wedge \bar{y})$
- $= 0 \vee 0 = 0$

Join of complements is 1

- $(x \wedge y) \vee (\bar{x} \vee \bar{y}) = ((x \wedge y) \vee \bar{x}) \vee \bar{y}$
- $= ((x \vee \bar{x}) \wedge (y \vee \bar{x})) \vee \bar{y}$
- $= (1 \wedge (y \vee \bar{x})) \vee \bar{y}$
- $= (y \vee \bar{x}) \vee \bar{y}$
- $= \bar{x} \vee (y \vee \bar{y})$
- $= \bar{x} \vee 1 = 1$

# Boolean lattice structure

- Let  $\mathcal{A}$  be a lattice with  $\min 0$
- $a \in \mathcal{A}$  is join irreducible if  $a \neq x \vee y$  for  $x, y \preceq a$ , alternatively  $a = x \vee y$  implies  $a = x$  or  $a = y$
- $0$  is join irreducible



- If  $b_1 \preceq c$  and  $b_2 \preceq c$  (immediate preds) of  $c$  then  $c = b_1 \vee b_2$
- $a \neq 0$  is join irreducible if and only if  $a$  has a unique immediate predecessor
- Elements immediately succeeding  $0$  are atoms (join irreducible)
- Any element  $a$  can be expressed as the join of a set of atoms
- Not unique for non-distributive lattice (diamond lattice)
- For finite lattice  $a = d_1 \vee d_2 \vee \dots \vee d_n$ ,  $d_j$  are join irreducible
- $d_j = d_i \vee d_j$  for  $d_i \preceq d_j$
- Any  $d_i \preceq d_j$  can be dropped to make the join irredundant
- Unique (up to permutation) for distributive lattice



# Boolean lattice representation (contd.)

## Unique irredundant irreducible sum representation

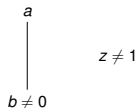
- Let  $a = c_1 \vee c_2 \vee \dots \vee c_m = d_1 \vee d_2 \vee \dots \vee d_n$
- Now,  $c_i, d_j \preceq c_1 \vee c_2 \vee \dots \vee c_m = d_1 \vee d_2 \vee \dots \vee d_n$
- $\therefore c_i = c_i \wedge (d_1 \vee d_2 \vee \dots \vee d_n) = (c_i \wedge d_1) \vee (c_i \wedge d_2) \vee \dots \vee (c_i \wedge d_n)$
- Since  $c_i$  is join irreducible,  $\exists d_j | c_i = c_i \wedge d_j$ , so that  $c_i \preceq d_j$
- But similar working,  $d_j \preceq c_k$ , so that  $c_i \preceq d_j \preceq c_k$
- This requires  $c_i = c_k$ , since these are irredundant
- Thus,  $c_i \preceq d_j$  and  $d_j \preceq c_i$ ,  $c_i = d_j$ ,
- This way, all the  $c_i$ s may be paired off with the  $d_j$ s,  
– making the representation unique (up to permutation)





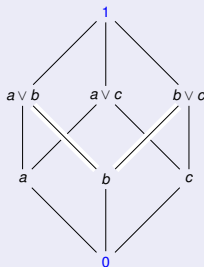
## Boolean lattice structure (contd.)

- Let  $z$  be the complement of  $a$  in a lattice as shown
- So,  $a \vee z = 1$  and  $a \wedge z = 0$
- Suppose  $a$  has  $b$  as a unique predecessor
- Now,  $b \vee z = a \vee z = 1$  and  $b \wedge z = a \wedge z = 0$  as  $b$  is the immediate predecessor of  $a$
- So,  $b$  is also a complement of  $z$

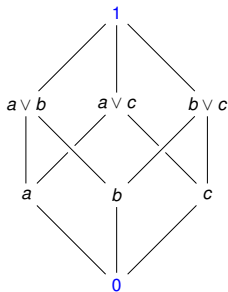


### Join irreducible elements in a Boolean lattice

- A lattice with an element having a non-zero join irreducible element as a predecessor will not have unique complements
- In a Boolean lattice all non-zero join irreducible elements are atoms



# Stone representation of Boolean lattices



- Atom of a Boolean lattice: Non-trivial minimal element of  $A \setminus \{0\}$
- $|A| = 2^n$  for some  $n$  for a Boolean lattice
- Its structure is that of the power set of the atomic elements

- Non-trivial atomic elements are present for  $|A| > 1$  directly above level 0, let those be  $S = \{a_1, \dots, a_n\}$ , akin to  $\{a_1\}, \{a_2\}, \dots, \{a_n\}$
- Join of pairs of elements  $Y_1, Y_2$  at level  $i$  ( $n > i > 1$ ) st  $|Y_1 - Y_2| = |Y_2 - Y_1| = 1$  at level  $i + 1$  is  $Y = Y_1 \cup Y_2$
- Meet of pairs of elements  $X_1, X_2$  at level  $i$  ( $n > i > 1$ ) st  $|X_1 - X_2| = |X_2 - X_1| = 1$  at level  $i - 1$  is  $Y = Y_1 \cap Y_2$
- There will be  $\binom{n}{i}$  such sets in level  $i$ , totaling to  $\sum_{i=0}^n \binom{n}{i} = 2^n$



# Boolean algebra from Boolean lattice

- For the Boolean lattice  $\langle \mathcal{A}, \preceq, 0, 1 \rangle$  consider the algebraic system  $\langle \mathcal{A}, +, \cdot, \bar{\phantom{x}}, 0, 1 \rangle$  where  $\vee \mapsto +$ ,  $\wedge \mapsto \cdot$  and  $\forall x \in \mathcal{A}$ ,  $\bar{x} \mapsto z \mid x + z = 1, x \cdot z = 0$
- This system satisfies the Huntington's postulates for a Boolean algebra

## B1: Commutative Laws

- 1  $x + y = y + x$
- 2  $x \cdot y = y \cdot x$

## B2: Distributive Laws

- 1  $x \cdot (y + z) = x \cdot y + x \cdot z$
- 2  $x + (y \cdot z) = (x + y) \cdot (x + z)$

## B3: Identity Laws

- 1  $x + 0 = x = 0 + x$
- 2  $x \cdot 1 = x = 1 \cdot x$

## B4: Complementation Laws

- 1  $x + \bar{x} = 1 = \bar{\bar{x}} + x$
- 2  $x \cdot \bar{x} = 0 = \bar{\bar{x}} \cdot x$



# Additional Boolean algebra properties

- These properties carry over from the Boolean lattice
- May be proven independently from the Huntington's postulates

## Idempotence:

- 1  $x + x = x$
- 2  $x \cdot x = x$

### Axiomatic proof

- $x + x = (x + x) \cdot 1$
- $= (x + x) \cdot (x + \bar{x})$
- $= x + (x \cdot \bar{x})$
- $= x + 0 = x$

## Absorption:

- 1  $x + xy = x$
- 2  $x \cdot (x + y) = x$
- 3  $x + \bar{x}y = x + y$
- 4  $x \cdot (\bar{x} + y) = xy$

### Axiomatic proof

- $x + xy = (x \cdot 1) + xy$
- $= x(1 + y) = x(y + 1)$
- $= x \cdot 1 = x$



# Boolean algebra (contd.)

## Boundedness/annihilation:

$$\textcircled{1} \quad x + 1 = 1$$

$$\textcircled{2} \quad x \cdot 0 = 0$$

## Axiomatic proof

$$\bullet \quad x + 1 = 1 \cdot (x + 1)$$

$$\bullet \quad = (x + \bar{x}) \cdot (x + 1)$$

$$\bullet \quad = x + (\bar{x} \cdot 1)$$

$$\bullet \quad = x + \bar{x} = 1$$

## Truth table for Boolean AND, OR, NOT:

$x$	$y$	$\bar{x}$	$x \cdot y$	$x + y$
0	0	1	0	0
0	1	1	0	1
1	0	0	0	1
1	1	0	1	1

## Associativity:

$$\textcircled{1} \quad (x + y) + z = x + (y + z)$$

$$\textcircled{2} \quad (x \cdot y) \cdot z = x \cdot (y \cdot z)$$



# Boolean algebra (contd.)

## Axiomatic proof of associativity of Boolean $+$

- Let  $x = a + (b + c)$  and  $y = (a + b) + c$
- $ax = aa + a(b + c) = a + a(b + c) = a$
- $bx = ba + b(b + c) = ba + (bb + bc) = ba + (b + bc) = ba + b = b$
- Similarly,  $cx = c$  and  $ay = a$ ,  $by = b$  and  $cy = c$
- $yx = ((a + b) + c)x = (a + b)x + cx = (ax + bx) + cx = (a + b) + c = y$
- $xy = (a + (b + c))y = ay + (b + c)y = ay + (by + cy) = a + (b + c) = x$
- Thus,  $x = xy = yx = y$



# Additional Boolean algebra properties (contd.)

## Uniqueness of Complement:

If  $(a + x) = 1$  and  $(a \cdot x) = 0$ , then  $x = \bar{a}$ .

## Involution:

$$\overline{(\bar{a})} = a$$

## Complements of extreme elements:

- $\bar{0} = 1$
- $\bar{1} = 0$

### Axiomatic proof

- $1 + 0 = 1$  [identity]
- $1 \cdot 0 = 0$  [boundedness]
- $\therefore 0$  is the complement of 1

## DeMorgan's laws:

- $\overline{(x + y)} = \bar{x} \cdot \bar{y}$
- $\overline{(x \cdot y)} = \bar{x} + \bar{y}$

- Let  $a \preceq b$  if  $a \cdot b = a$  or  $a + b = b$  then  $\cdot \mapsto \wedge$  and  $+$   $\mapsto \vee$
- Properties from axiomatic proofs allow Boolean algebras to be expressed as Boolean lattices – they are equivalent

