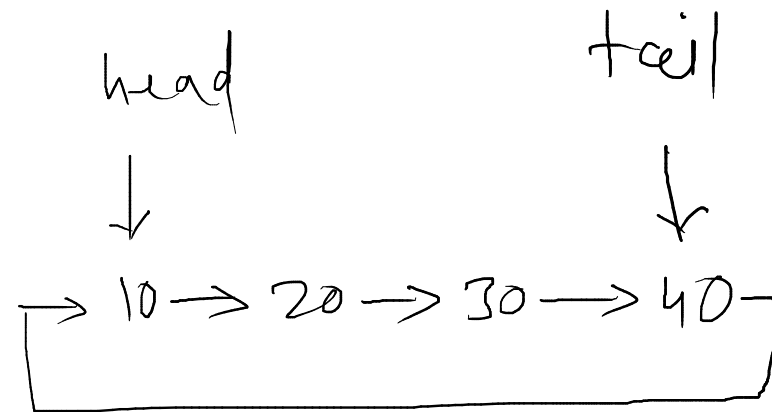
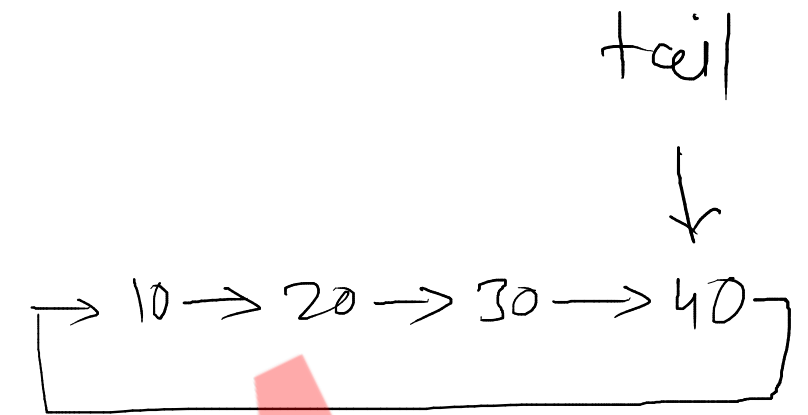


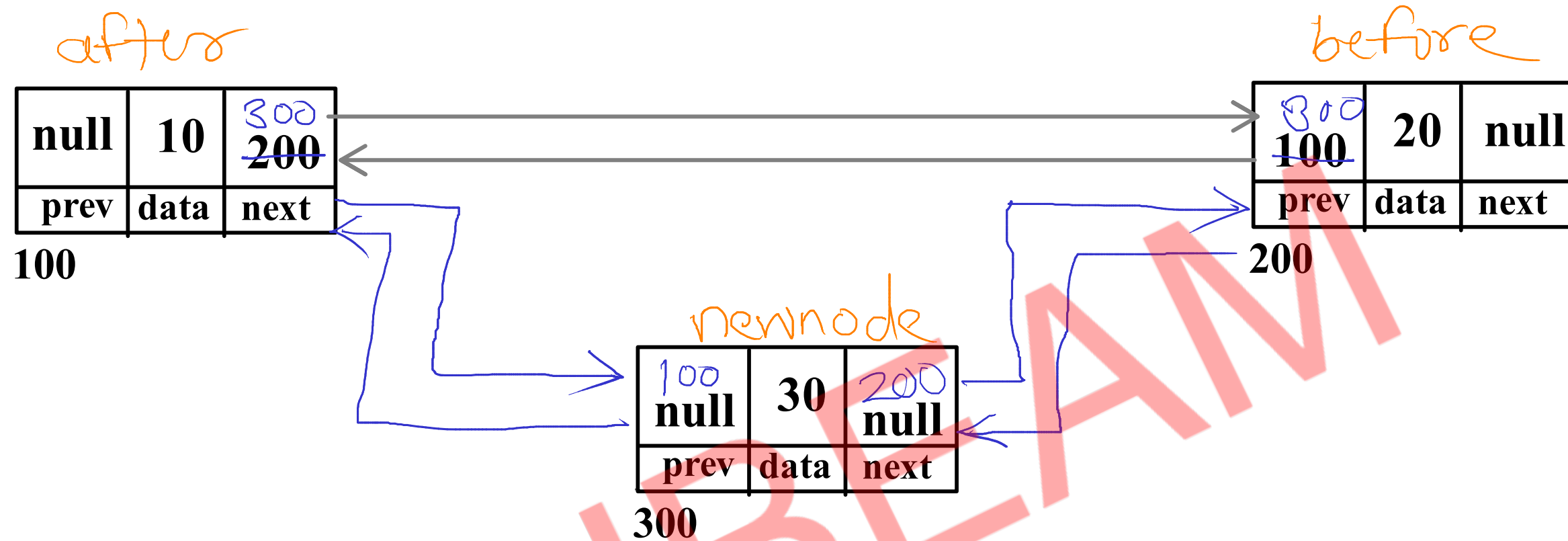
Add First $\rightarrow O(n)$
 Add Last $\rightarrow O(n)$
 Add Pos $\rightarrow O(n)$
 Delete First $\rightarrow O(n)$
 Delete Last $\rightarrow O(n)$
 Delete Pos $\rightarrow O(n)$
 Traversal $\rightarrow O(n)$



Add First $\rightarrow \underline{O(1)}$
 Add Last $\rightarrow \underline{O(1)}$
 Add Pos $\rightarrow O(n)$
 Delete First $\rightarrow \underline{O(1)}$
 Delete Last $\rightarrow O(n)$
 Delete Pos $\rightarrow O(n)$
 Traversal $\rightarrow O(n)$



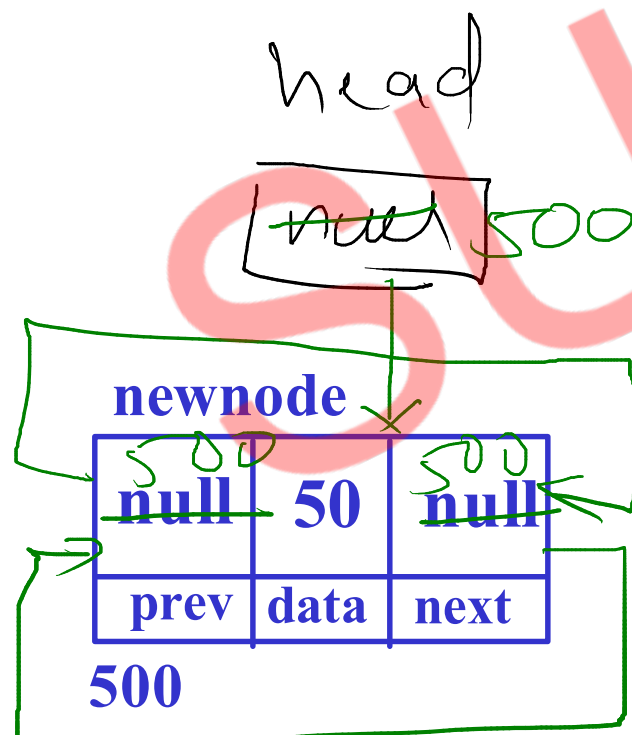
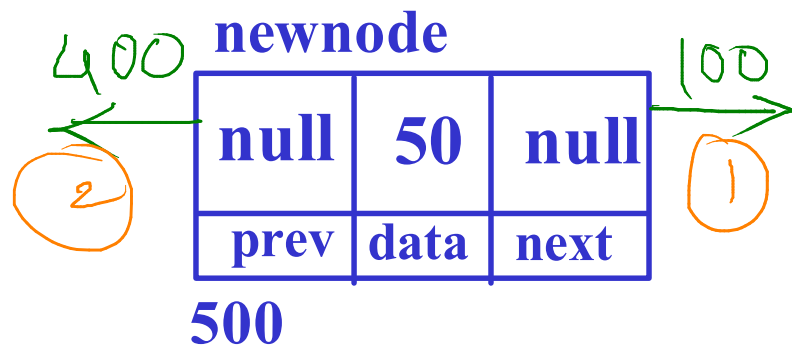
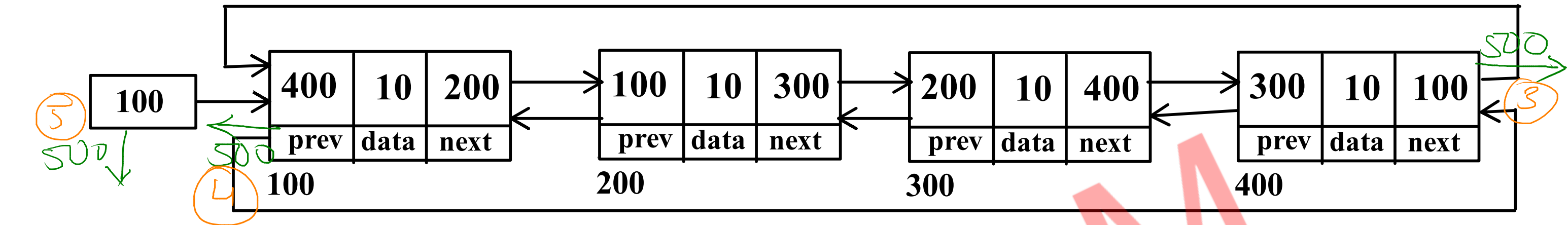
Add First $\rightarrow \underline{O(1)}$
 Add Last $\rightarrow \underline{O(1)}$
 Add Pos $\rightarrow O(n)$
 Delete First $\rightarrow \underline{O(1)}$
 Delete Last $\rightarrow O(n)$
 Delete Pos $\rightarrow O(n)$
 Traversal $\rightarrow O(n)$



- //1. add before node into next of newnode
- //2. add after node into prev of newnode
- //3. add newnode into next of after node
- //4. add newnode into prev of before node

`newnode.next = before;`
`newnode.prev = after;`
`after.next = newnode;`
`before.prev = newnode;`

Doubly Circular Linked List - Add First



//1. create newnode

//2. if list is empty

//a. add newnode into head

//b. make list circular

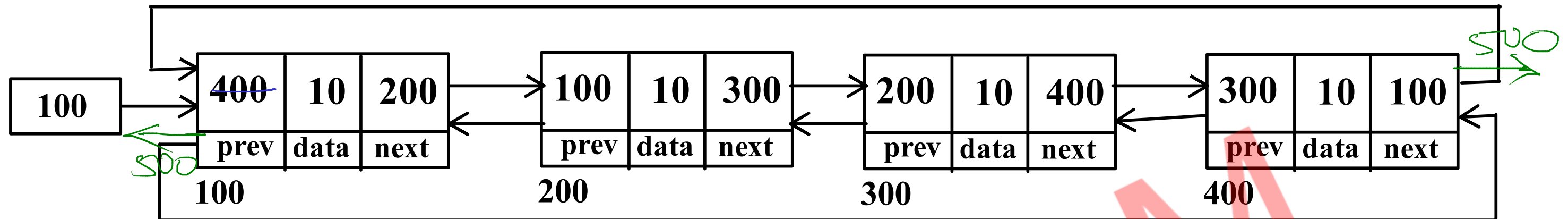
//3. if list is not empty

//a. call insert method

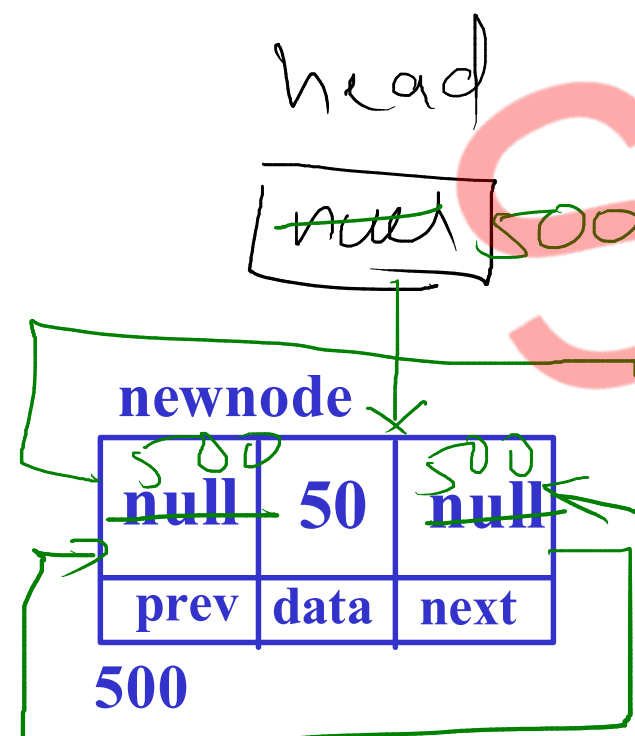
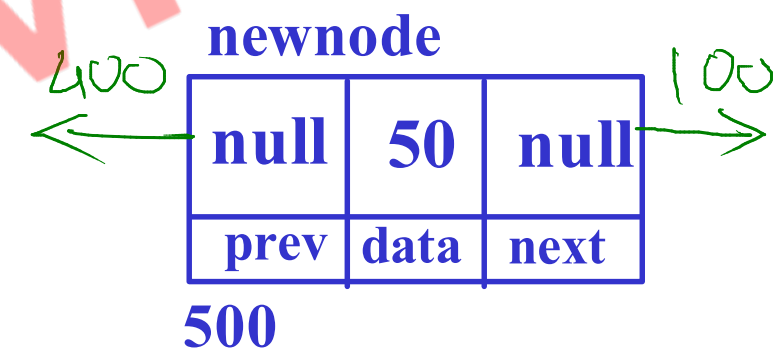
//b. move head on newnode

$$T(n) = O(1)$$

Doubly Circular Linked List - Add Last



after . nn before
last first



//1. create newnode

//2. if list is empty

//a. add newnode into head

//b. make list circular

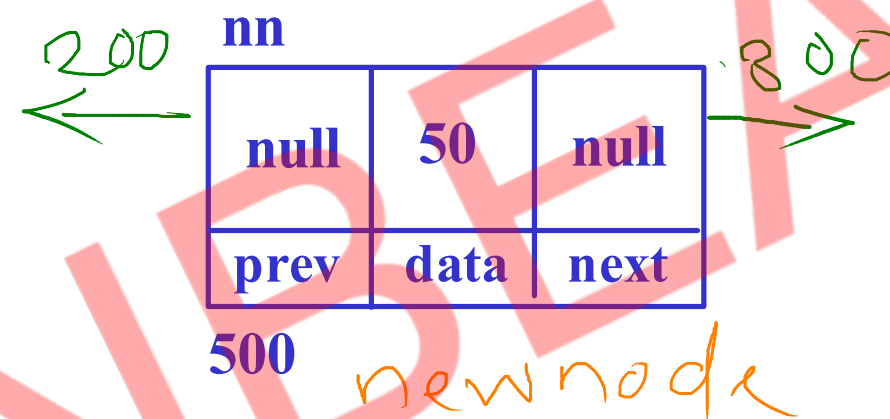
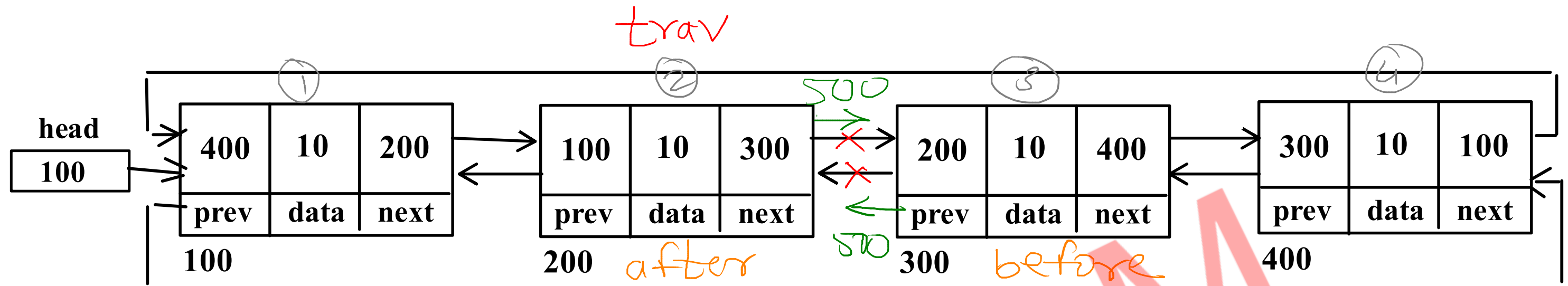
//3. if list is not empty

//a. call insert method

$$T(n) = O(1)$$

Doubly Circular Linked List - Add pos

pos=3



//1. create newnode

//2. if list is empty

//a. add newnode into head

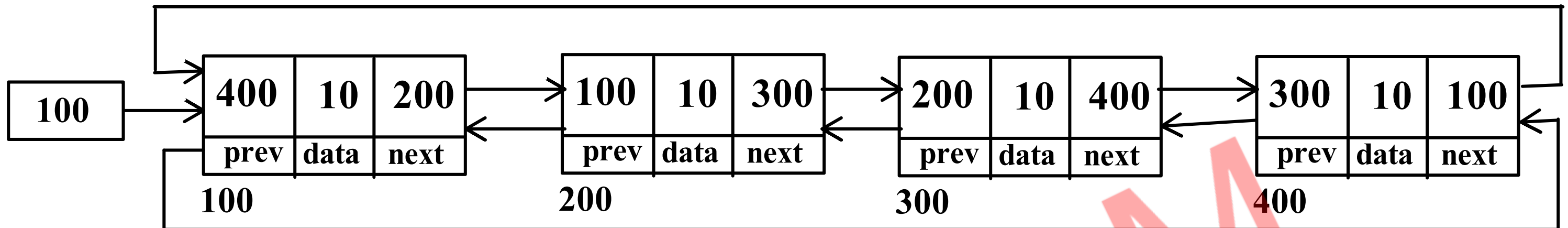
//b. make list circular

//3. if list is not empty

//a. traverse till pos-1 node

//b. call insert method

Doubly Circular Linked List - Display



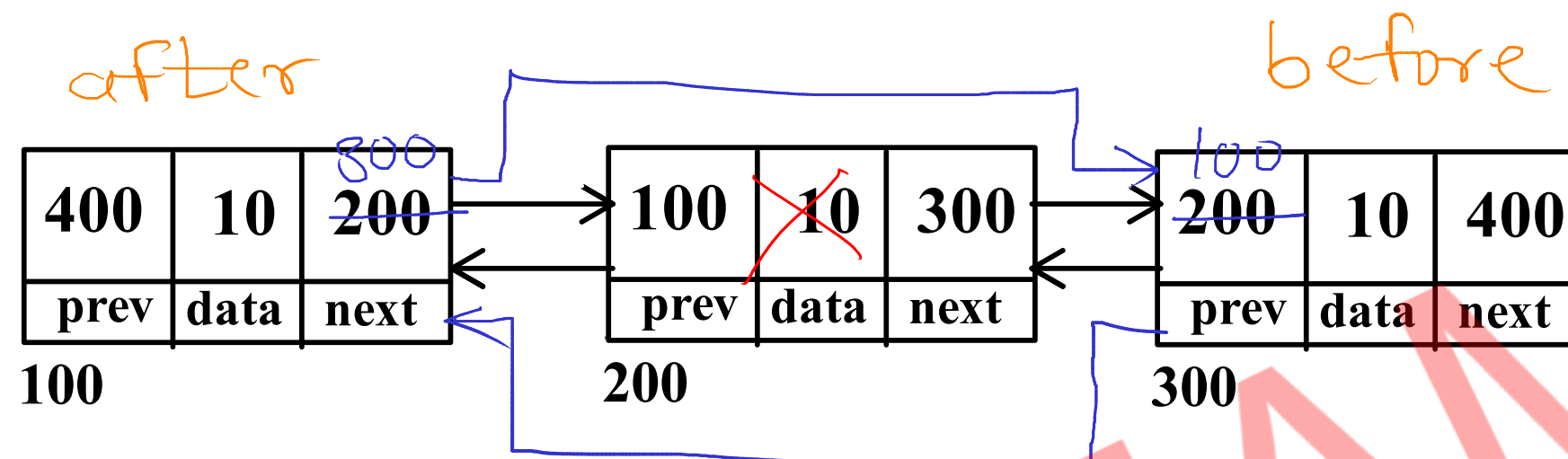
Forward Display

- //1. start at first node
- //2. print current node (trav.data)
- //3. go on next node (trav.next)
- //4. repeat step 2 and 3 till last node

Reverse Display

- //1. start at last node
- //2. print current node (trav.data)
- //3. go on prev node (trav.prev)
- //4. repeat step 2 and 3 till first node

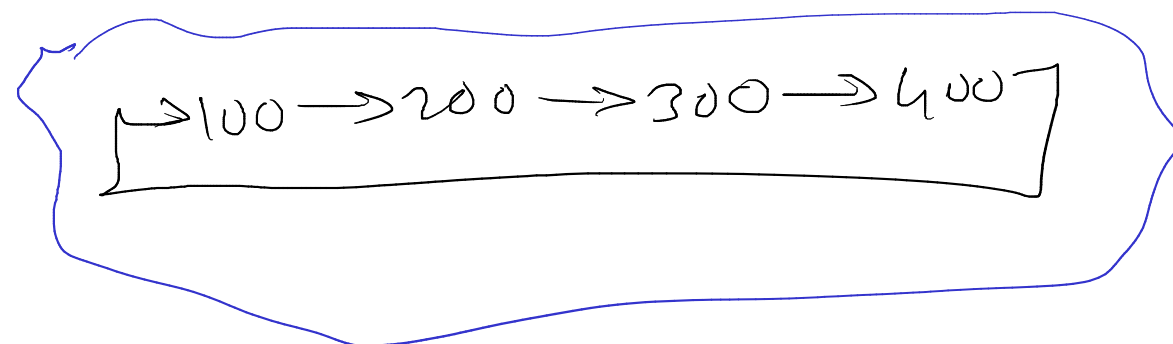
$$T(n) = O(n)$$



- //1. add before node into next of after node
- //2. add after node into prev of before node

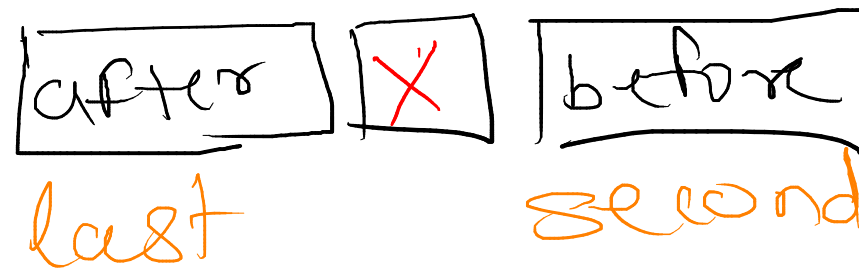
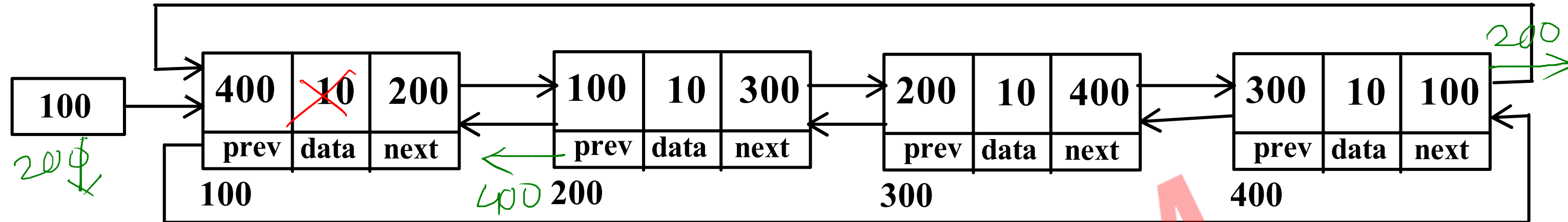
after.next = before
before.prev = after

head
↓



Island of Isolation

Doubly Circular Linked List - Delete First



//1. if list is empty
return;

//2. if list has single node
head = null;

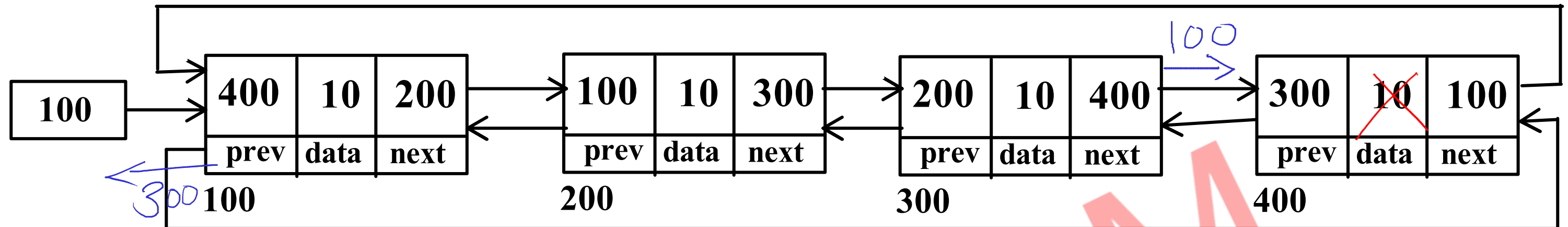
//3. if list has multiple nodes

//a. call delete method

//b. move head on second node

$$T(n) = O(1)$$

Doubly Circular Linked List - Delete Last



after ~~X~~ before
second last first

$$T(n) = O(1)$$

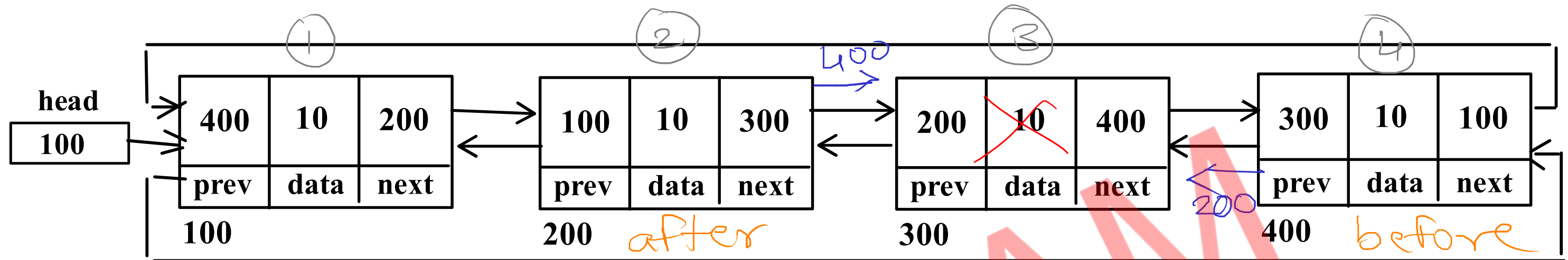
//1. if list is empty
return;

//2. if list has single node
head = null;

//3. if list has multiple nodes
//a. call delete method

Doubly Circular Linked List - Del Pos

pos = 3



//1. if list is empty
return;

//2. if list has single node
head = null;

//3. if list has multiple nodes

//a. traverse till pos node

//a. call delete method

$T(n) = O(n)$

Linked List Applications

- dynamic data structure - grow / shrink at runtime
- due to this dynamic nature, it is used to implement other data structures
 - ✓ 1. Stack
 - ✓ 2. Queue
 - ✓ 3. Hash Table (Seperate chaining)
 - ✓ 4. Graph (Adjacency list)
- Operating system - job queue, ready queue, waiting queues
(Doubly circular linked list)
- Ring topology in Networks - (Doubly circular linked list)

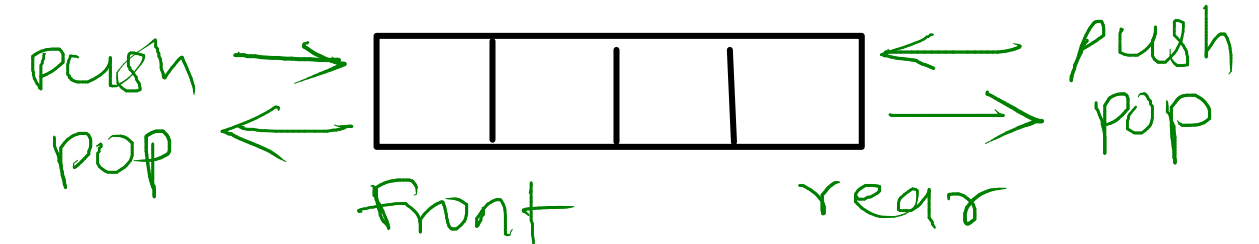
Stack(LIFO)

1. Add First
Delete First
2. Add Last
Delete Last

Queue(FIFO)

1. Add First
Delete Last
2. Add Last
Delete First

Deque(Double Ended Queue)



push front
add first

pop front
delete first

push rear
add last

pop rear
delete last

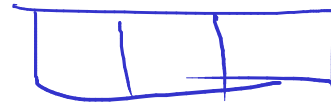
Types of deque

- 1) input restricted deque
- 2) output restricted deque

Array Vs Linked List

Array

1. Array space in memory is contiguous



2. Array can not grow or shrink at runtime

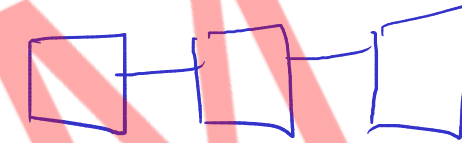
3. Random access of elements is allowed

4. Insert or Delete, needs shifting of array elements

5. Array needs less space

Linked List

1. Linked list space in memory is not contiguous



2. Linked list can grow or shrink at runtime

3. Random access of elements is not allowed(sequential)

4. Insert or Delete, do not need shifting of nodes

5. Linked lists need more space