



**Sunbeam Institute of Information Technology
Pune and Karad**

Data structures and Algorithms

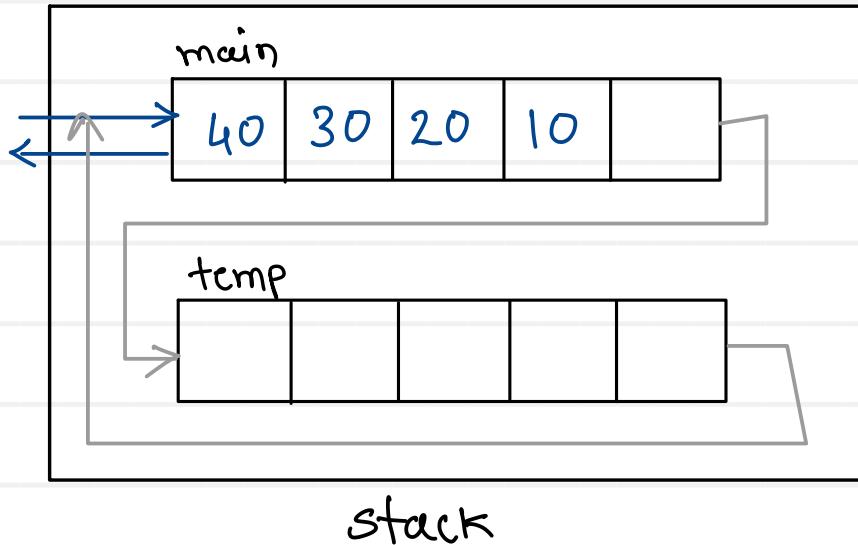
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Create stack using queue

LIFO

FIFO

Push order : 10, 20, 30, 40



Pop order : 40, 30, 20, 10

java.util

Queue - Interface

Queue<> q = new LinkedList<>();

q.offer(—) ← push op

q.poll() ← pop op

push() :

```
while (!mq.isEmpty())
    tq.offer(mq.poll());
mq.offer(value);
while (!tq.isEmpty())
    mq.offer(tq.poll());
```

pop() :

```
mq.poll()
```

peek() :

```
mq.peek()
```

push:

$$itr = n + n$$

$$T \propto 2n$$

$$T(n) = O(n)$$

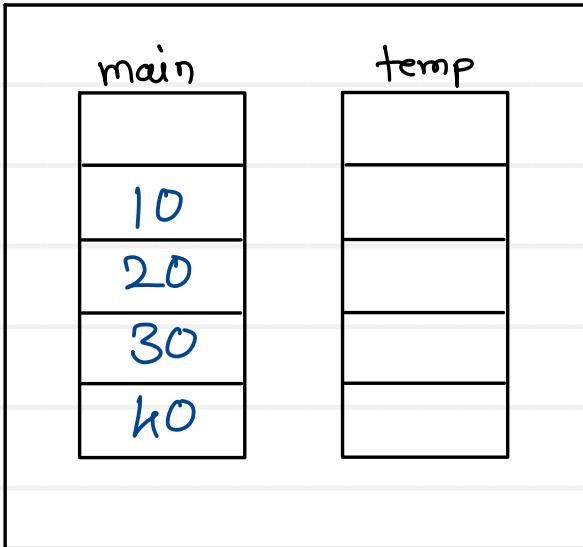
pop & peek:

$$T(n) = O(1)$$

Create queue using stack

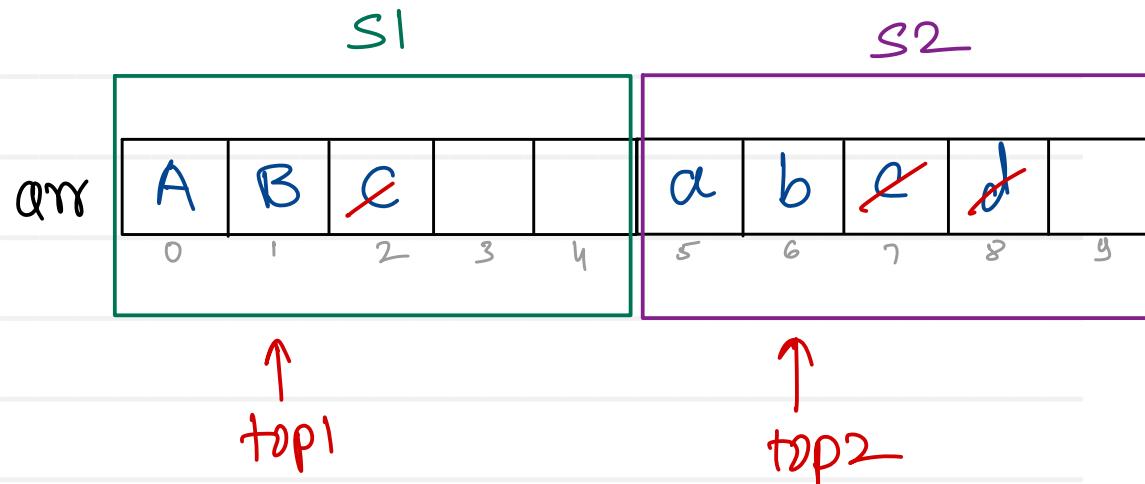
Push order : 10 , 20 , 30 , 40

Queue





How to Implement two stacks using array efficiently?



- two stacks with
fixed capacity

isEmpty : $\text{top1} == -1$

isFull : $\text{top1} == \text{length}/2$

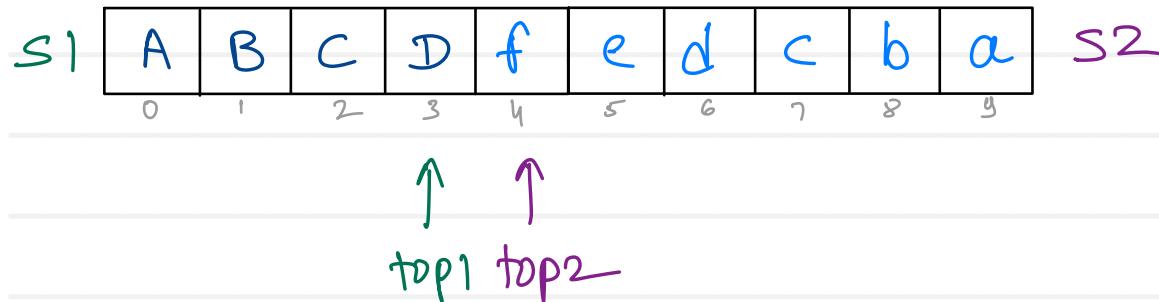
$\text{top2} == \text{length}/2$

$\text{top2} == \text{length}-1$





How to Implement two stacks using array efficiently?



- two stacks with variable capacity
- more memory efficient

Stack1
init : $\text{top1} = -1$
push : $\text{top1}++$
 $\text{arr}[\text{top1}] = \text{val}$
pop : $\text{top1}--$
isEmpty : $\text{top1} == -1$
isFull : $\text{top1} == \text{top2} - 1$

Stack2
init : $\text{top2} = \text{length}$
push : $\text{top2}--$
 $\text{arr}[\text{top2}] = \text{val}$
pop : $\text{top2}++$
isEmpty : $\text{top2} == \text{length}$
isFull : $\text{top1} == \text{top2} - 1$



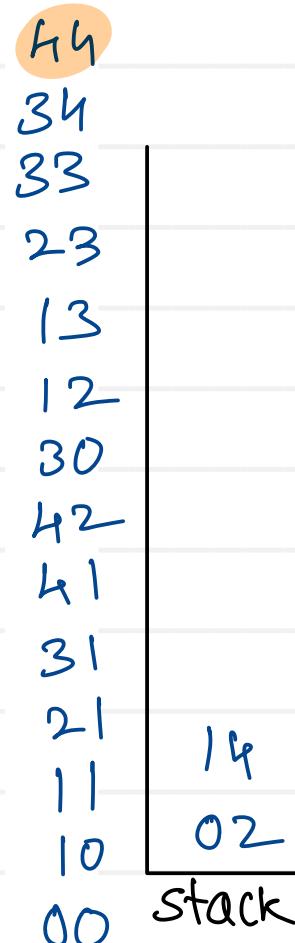


Rat in a maze problem

← backtracking

Algorithm:

1. push src cell on stack & mark it.
2. while stack is not empty :
 - 2.1 pop a cell from stack (cur)
 - 2.2 if(cur cell == dest cell) return true;
 - 2.3 push all valid & unmarked neighbor cells on stack.
3. if destination is not reachable return false.



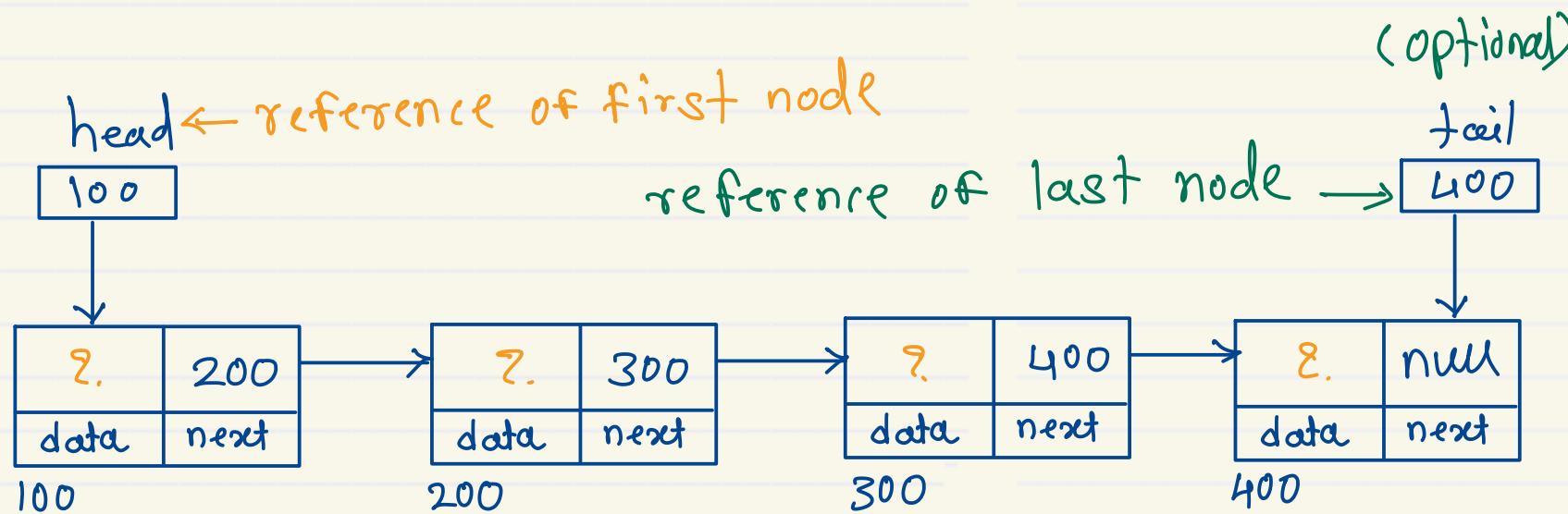
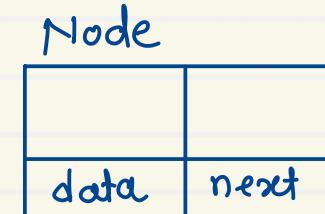
maze

	0	1	2	3	4
0	S 00	01	02	03	04
1	10	11	12	13	14
2	20	21	22	23	24
3	30	31	32	33	34
4	40	41	42	43	d 44

marked →
(true/false)

	0	1	2	3	4
0	00	01	02	03	04
1	10	11	12	13	14
2	20	21	22	23	24
3	30	31	32	33	34
4	40	41	42	43	44

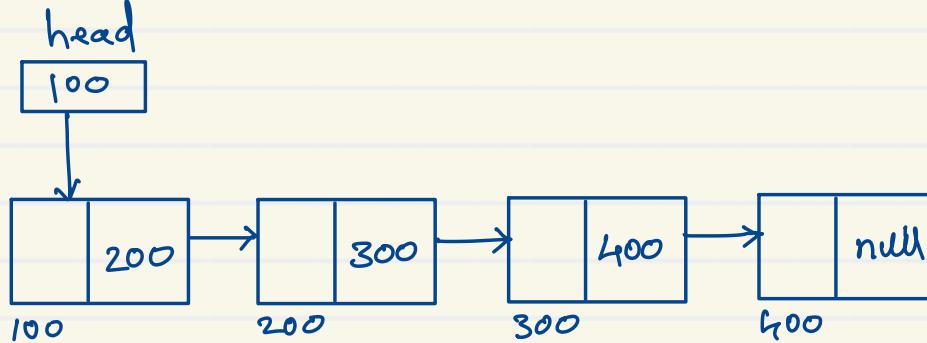
- Linked list is a linear data structure
- Linked list is list of items linked together
- Each item in linked list is called as Node
- Each node contains data and pointer / reference to the next node



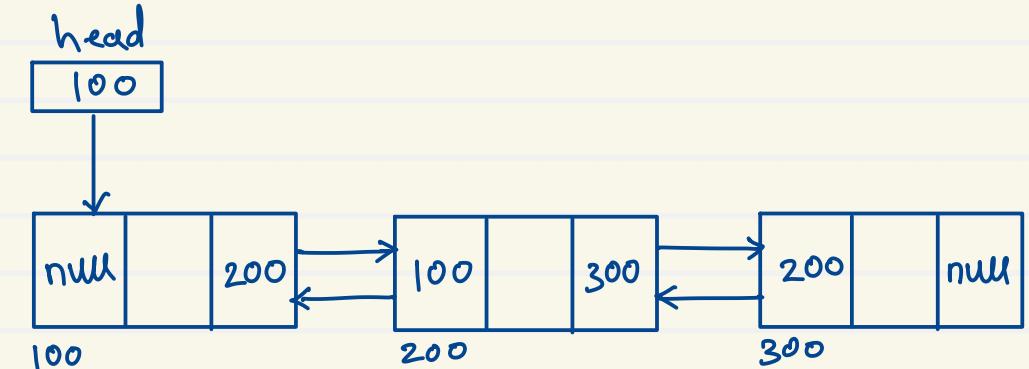
- Operations
1. Add first
 2. Add last
 3. Add position
 4. Delete first
 5. Delete last
 6. Delete position
 7. Delete all
 8. Display

Linked List - Types

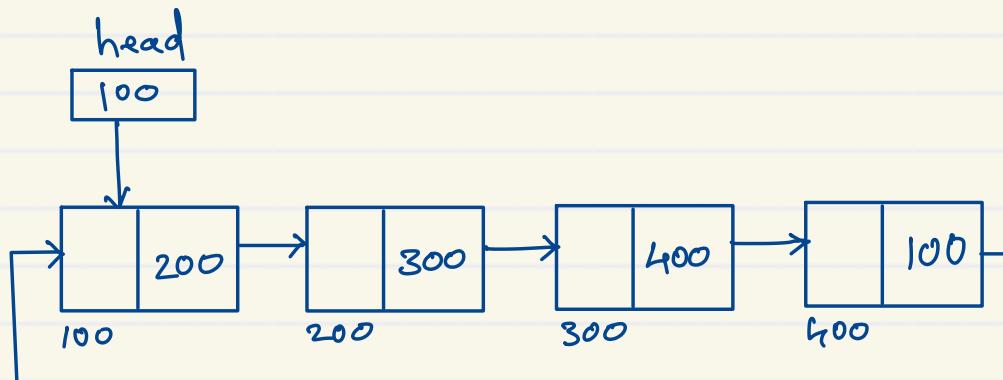
1. Singly linear linked list



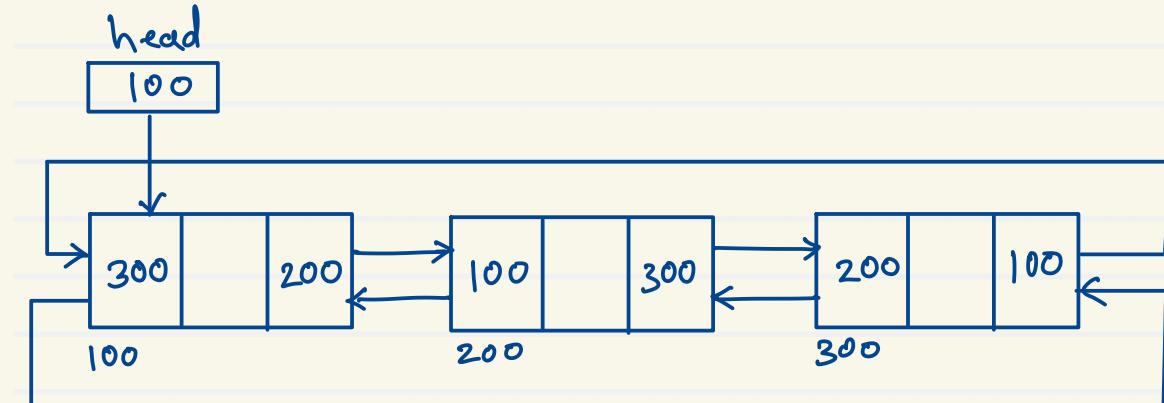
3. Doubly linear linked list



2. Singly circular linked list



4. Doubly circular linked list



Linked List

Node :

data - int, char, double, class

next - reference of next node

class Node { ← self referential class

int data;

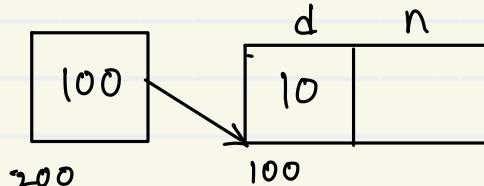
Node next;

}

- In java, Node class is inner class of LinkedList class

- In C++, Node class is friend class of LinkedList class.

Node n = new Node(10)



class LinkedList {

Node head;

Node tail;

int count;

public LinkedList() {} } }

public add() {} ... }

public delete() {} ... }

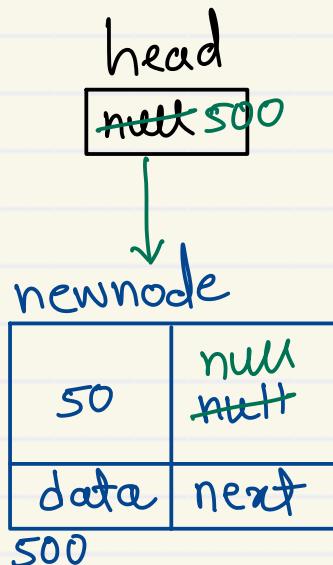
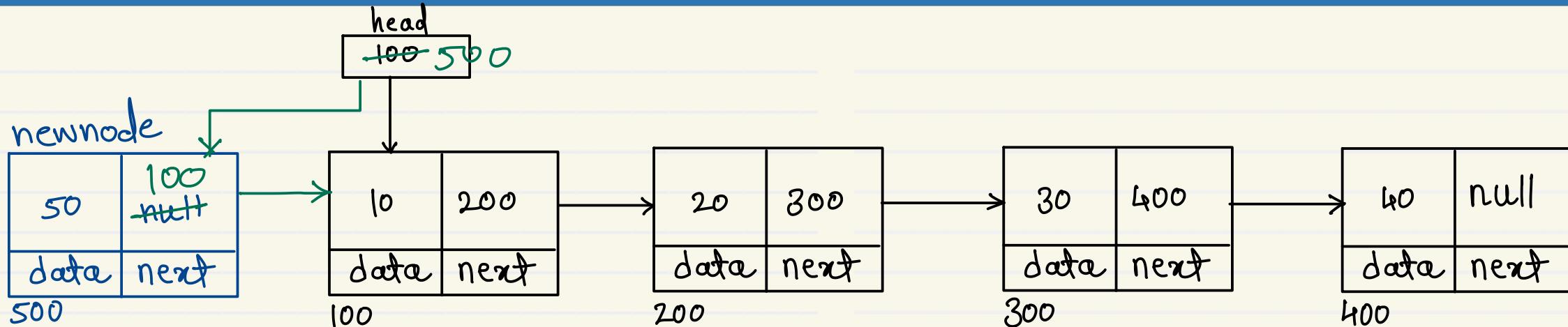
public display() {} ... }

public search() {} ... }

public deleteAll() {} ... }

}

Singly linear Linked List - Add first



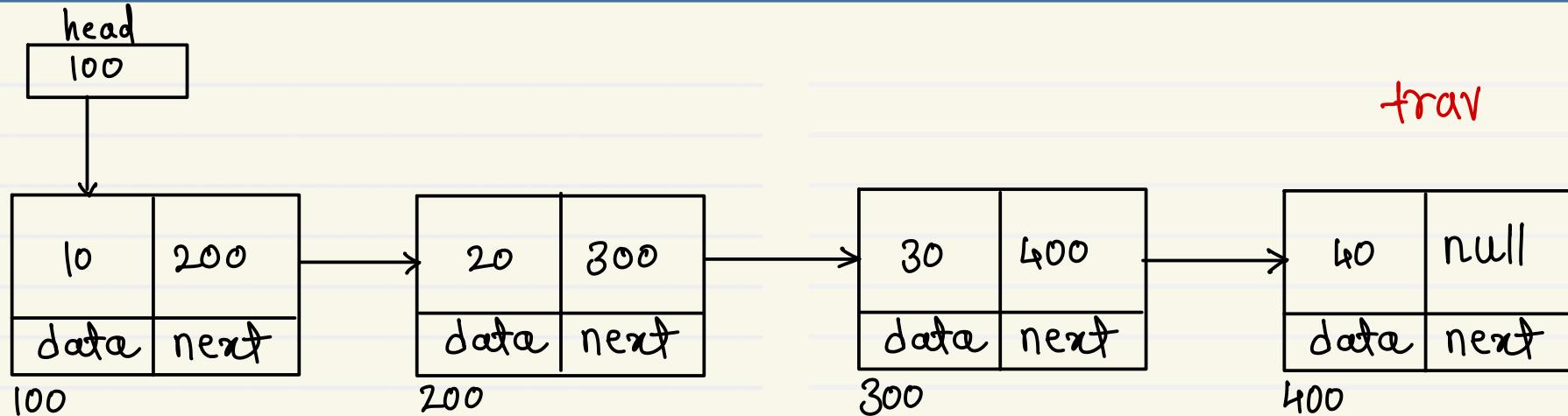
newnode.next = head
head = newnode

1. create a newnode with given data
2. add first node into next of newnode
3. update head by newnode

$$T(n) = O(1)$$

Singly linear Linked List - Display

traverse - visiting each node at least once

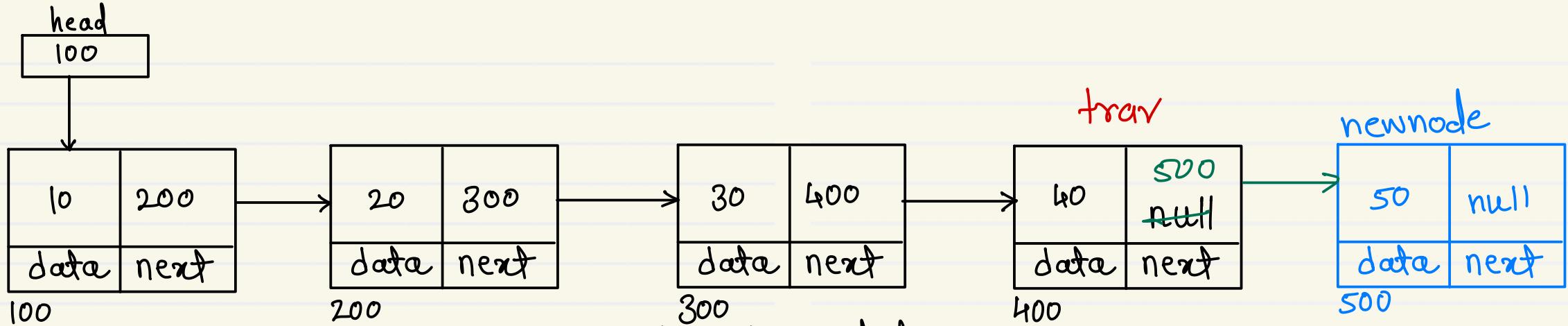


trav	trav.data	trav.next
100	10	200
200	20	300
300	30	400
400	40	null
null		

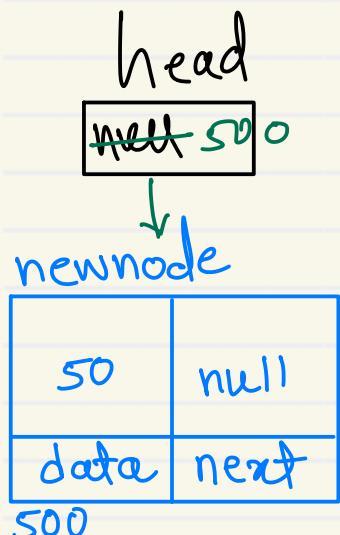
1. Create reference p start at first node
2. visit/print current node
3. go on next node
4. repeat step 2 & 3 till last node

$$T(n) = O(n)$$

Singly linear Linked List - Add last



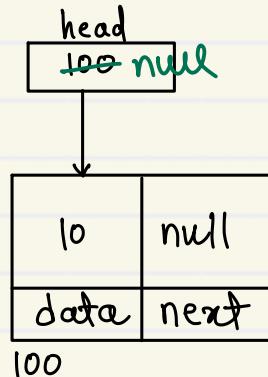
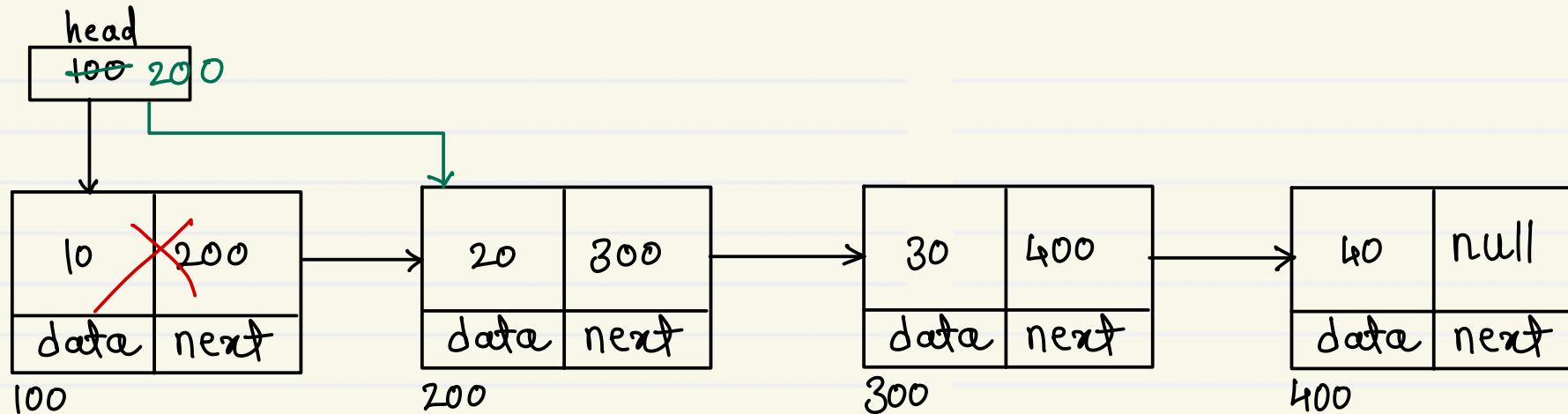
0. create newnode with data
1. if list is empty,
add newnode into head itself
2. if list is not empty,
 - a. traverse till last node
 - b. add newnode into next of last node



$\text{trav} = \text{head}$
 $\text{while}(\text{trav}. \text{next} \neq \text{null})$
 $\quad \text{trav} = \text{trav}. \text{next};$

$$T(n) = O(n)$$

Singly linear Linked List - Delete first



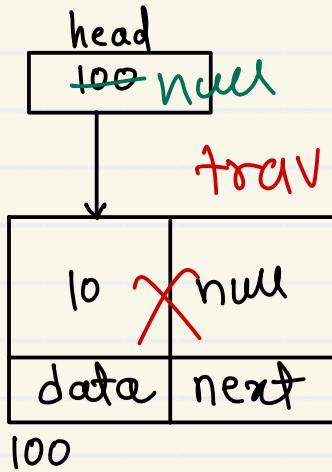
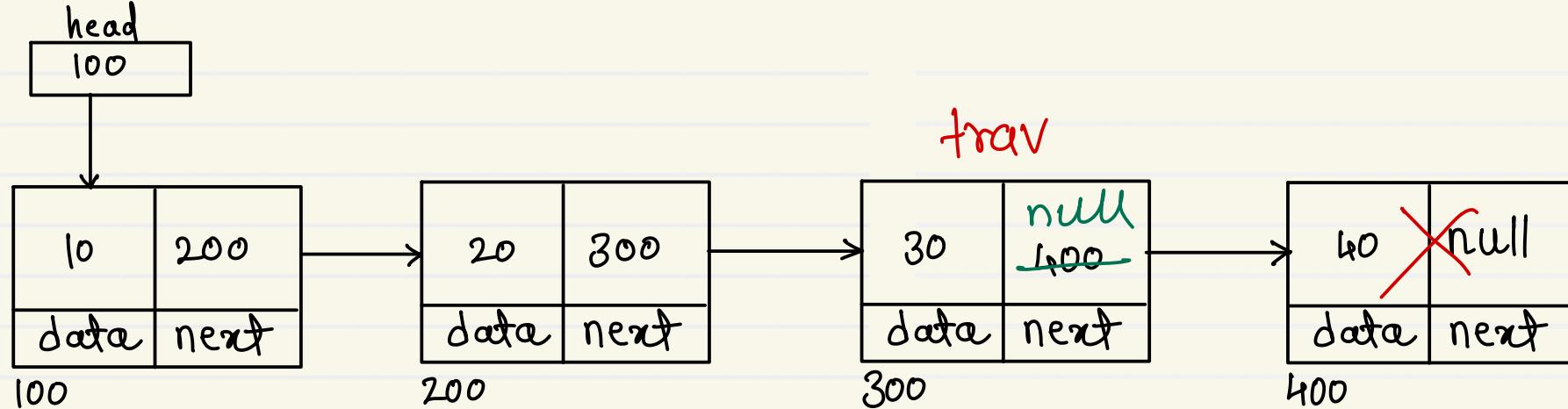
1. if list is empty , return
2. if list is not empty
 - a. move head on second node

$$T(n) = O(1)$$

head
null

head = head.next

Singly linear Linked List - Delete last



1. if list is empty , return
2. if list has single node, delete it
3. if list has multiple nodes
 - a. traverse till second last node
 - b. make next of second last node null

Node trav = head;
 while(trav.next != null)
 trav = trav.next;

$T(n) = O(n)$