AI506: Data Mining and Search (Spring 2020)

Homework 1: Locality Sensitive Hashing

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1 Min-Hashing

1.1 Shingling

1. get_shingles

```
start = time.time()
shingles = get_shingles(documents)
end = time.time()

# Check whether your implementation is correct [5pt]
if len(shingles) == 1766049:
    pass_test1_1_1 = True
    print('Test1 passed')

# Check whether your implementation is efficient enough [5pt]
    # With 4-lines of my implementations, it took 4.8 seconds with i7-8700 cpu
    if (end - start) < 20:
        pass_test1_1_2 = True
        print('Test2 passed')</pre>
Test1 passed
Test2 passed
```

2. build_doc_to_shingle_dictionary

>>> Test Result

```
doc_to_shingles = build_doc_to_shingle_dictionary(documents, shingles)

# Check whether your implementation is correct [5pt]
if len(doc_to_shingles) == 10882 and len(doc_to_shingles[0]) == 84:
    pass_test1_2 = True
    print('Test passed')
Test passed
```

1.2. Min-Hashing

1.2.1. Computing MinHash signatures

1. Calculate minhash signature

$$\begin{aligned} \min_{x \in S_1} h_1(x) &= \min\{h_1(2), h_1(5)\} = \min\{5, 5\} = 5 \\ \min_{x \in S_1} h_2(x) &= \min\{h_2(2), h_2(5)\} = \min\{2, 5\} = 2 \\ \min_{x \in S_1} h_3(x) &= \min\{h_3(2), h_3(5)\} = \min\{0, 3\} = 0 \\ \text{Similarly,} \end{aligned}$$

$$\begin{aligned} & [\min_{x \in S_2} h_1, \min_{x \in S_2} h_2, \min_{x \in S_2} h_3] = [1, 2, 1] \\ & [\min_{x \in S_3} h_1, \min_{x \in S_3} h_2, \min_{x \in S_3} h_3] = [1, 2, 4] \\ & [\min_{x \in S_4} h_1, \min_{x \in S_4} h_2, \min_{x \in S_4} h_3] = [1, 2, 0] \end{aligned}$$

S_1	S_2	S_3	S_4
5	1	1	1
2	2	2	2
0	1	4	0

Signature:

2. Calculate true and estimated Jaccard similarities.

	(S_1, S_2)	(S_1, S_3)	(S_1, S_4)	(S_2, S_3)	(S_2, S_4)	(S_3, S_4)
true	0	0	1/4	0	1/4	1/4
estimated	1/3	1/3	2/3	2/3	2/3	2/3

1.2.2. Implementation

1. jaccard_similarity

```
s1 = {1, 3, 4}
s2 = {3, 4, 6}

if (jaccard_similarity(s1, s2) - 0.5) < 1e-3:
    pass_test2_1 = True
    print('Test passed')

Test passed</pre>
```

2. min hash

I modified Hash class in order to parallelize the hash function calculation.

```
class Hash():
    def __init__(self, M, N):
        self.M = M
        self.N = N
        self.p = generate_prime_numbers(M, N)

        self.a = np.random.choice(9999, M)
        self.b = np.random.choice(9999, M)

    def __call__(self, x):
        return np.mod(np.mod((self.a * x + self.b), self.p), self.N))

    def __len__(self):
        return M

#primes = generate_prime_numbers(M, N)
hash_functions = Hash(M, N)
```

```
start = time.time()
signatures = min_hash(doc_to_shingles, hash_functions)
end = time.time()

diff_list = compare(signatures, doc_to_shingles)

# Check whether your implementation is correct [20pt]
# Average difference of document's jaccard similarity between the base of the b
```

2 Locality Sensitive Hashing (LSH)

2.1 Locality Sensitive Hashing

1. lsh

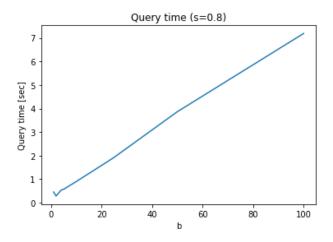
```
def lsh(signatures, b, r):
    M = signatures.shape[0] # The number of min-hash functions
C = signatures.shape[1] # The number of documents
    candidatePairs = set()
    for num_b in range(b):
        bucket = {}
bands = signatures[num_b*r:(num_b+1)*r]
        for col in range(C):
             if tuple(bands[:,col]) in bucket.keys():
                 bucket[tuple(bands[:,col])].append(col)
                 bucket[tuple(bands[:,col])] = [col]
        for value in bucket.values():
             if len(value) >= 2:
                 combi = combinations(value, 2)
                 candidatePairs.update(list(combi))
    return candidatePairs
```

2.2 Analysis

1. query_analysis

2. query time graph

The query time linearly increases with b. This is because the larger b is, the shorter each band is. Then this means each band is more likely to be in the same bucket, which is being a candidate pair. Increased candidate pairs cause linearly increased query time. Therefore, query time has time complexity of O(b).



3. precision, recall, and f1-score graph

Precision decreases with b, while recall increases with b. This is obvious because as b increases, the number of candidate pairs increase. Therefore, positive pairs that are assumed to be similar increase. For precision, there is FP in denominator, and FP increases as there are more positive pairs. So precision decreases. For recall, there is FN in denominator, and FN decreases as positive pairs increase. Therefore, recall increases.

In precision measure, the best b value is 1, but in recall measure, the best b value is 20~100. By the way, by looking at the f1-score, we can be between that trade-off. The best b, in terms of f1-score measure, is 10.

