

Do something today  
that your future self  
will thank you for.



Good

Morning

## Today's content

01. Introduction to Queues

02. Implementation of Queues

→ Using Arrays

→ Using Linkedlist

→ Using stacks

03. Introduction of Deque

04. Max of every window

Sliding Window Maximum

Queue → Linear data structure where elements are inserted from one end & removed from the other end

⇒ First In First Out



### Some real world Examples

- 01. Ticket counter
- 02. Message queue
- 03. Token entry
- 04. Call center
- 05. playlist Queue

### \* Common operations in Queue

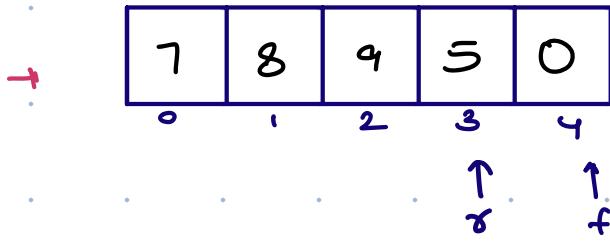
- 01. Add( $x$ ) / Enqueue( $x$ ) → Add  $x$  inside queue from last
- 02. Remove() / Deque() → Remove first ele present in queue
- 03. top() / peek() → return first ele present in queue
- 04. size() → no. of elements inside queue

# Implementation of Queue

## 01. Array

$f = -1 \Rightarrow$  idx of last ele which was removed

$r = -1 \Rightarrow$  idx of last ele which got inserted



if ( $f == r$ )  $\rightarrow$  q is empty  
↳ q is full

Add(5)	remove() $\rightarrow$ 5	add(6)	remove() $\rightarrow$ 6
Add(4)	remove() $\rightarrow$ 4	add(7)	top() $\rightarrow$ 7
Add(3)	remove() $\rightarrow$ 3	add(8)	
Add(2)	remove() $\rightarrow$ 2	add(9)	
		add(5)	

## class Queue {

```
int sz=0;  
  
int [ ] q;  
  
int f=-1, r=-1;  
  
Queue () {  
    q = new int [5];  
}
```

```
int size () {  
    return sz;  
}  
  
int top () {  
    if (sz==0) return -1;  
  
    int idx = (f+1)%N  
    return A [idx];  
}
```

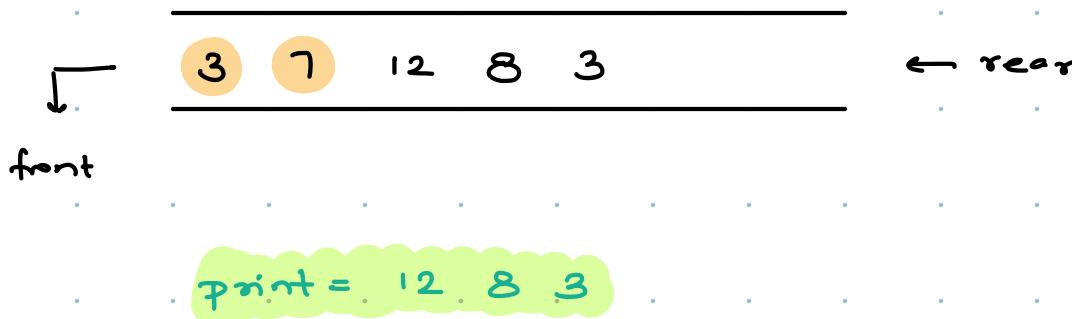
```
void add ( val )
```

```
if (sz == n) return;  
r = (r + 1) % n  
A[r] = val;  
sz++;
```

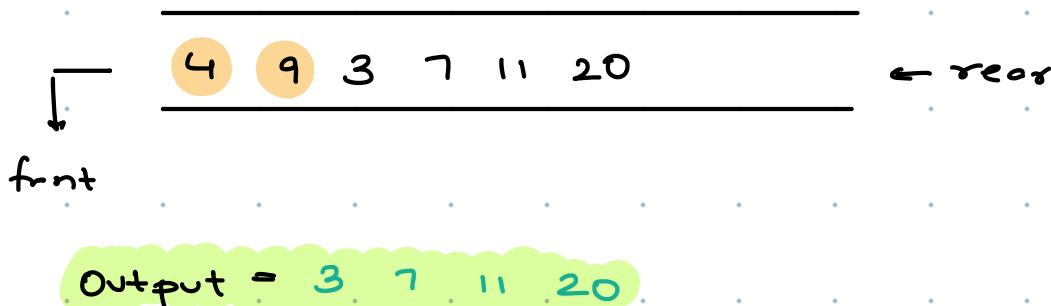
```
int remove ( )
```

```
if (sz == 0) return -1;  
f = (f + 1) % n  
int val = A[f];  
A[f] = 0;  
sz = sz - 1;  
return val;
```

What will be the state of the queue after these operations enqueue(3), enqueue(7), enqueue(12),  
dqueue(), dqueue(), enqueue(8), enqueue(3)



What will be the state of the queue after these operations enqueue(4), dqueue(), enqueue(9),  
enqueue(3), enqueue(7), enqueue(11), enqueue(20), dqueue()

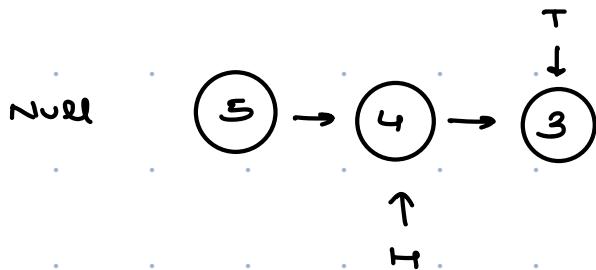


## \* Using LinkedList

Insert at tail  
remove at head } TC: O(1)

Add(s)      Add(z)

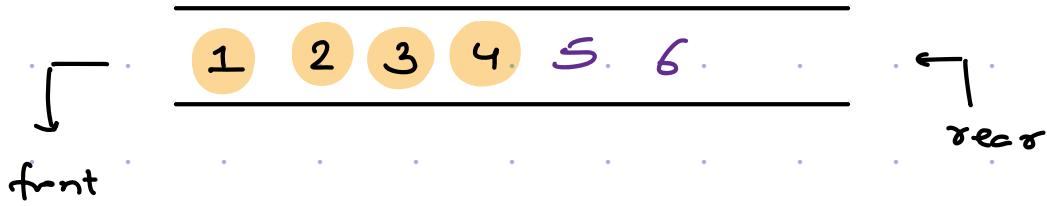
Add(4)      remove()



## \* Implement queue using Stacks

Stacks

→ push      → peek  
→ pop      → size



Add(1)

Add(2)

Add(3)

Add(4)

remove()

remove()

Add(5)

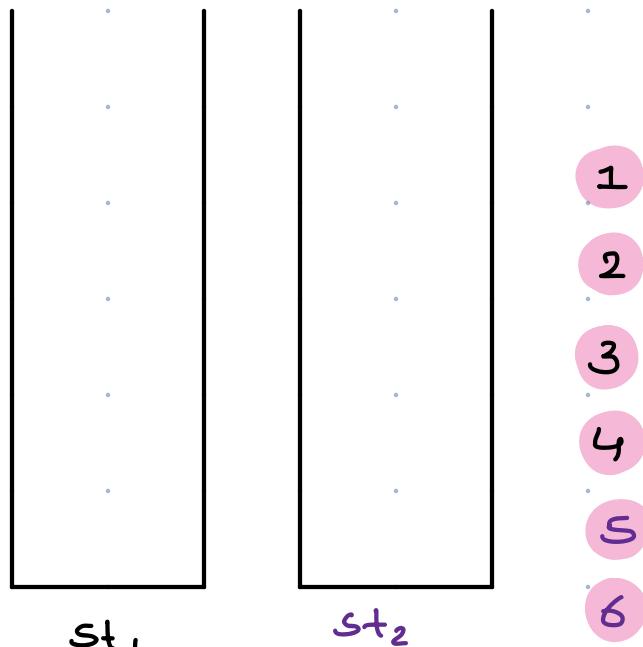
Add(6)

remove()

remove()

remove()

remove()



$\Delta d(x) \rightarrow st_1.push(x) \rightarrow Tc:O(1)$

$remove() \rightarrow \text{if } (st_2.size() == 0)$

Yes      No

$\rightarrow$  move all ele

$st_2.pop();$

from  $st_1$  to  $st_2$

$\rightarrow Tc:O(1)$

& then pop

$\rightarrow Tc:O(1)$

$size() \rightarrow st_1.size() + st_2.size();$

$peek() \rightarrow$  similar to remove  
function without popping

Add(1)

Add(2)

Add(3)

Add(4)

remove(1)

remove(1)

Add(5)

remove(1)

remove(1)

1<sup>st</sup> rem = Move all elements from st<sub>1</sub>,

to st<sub>2</sub> & then st<sub>2</sub>.pop();

$$8 \text{ ops} + 1 \text{ op} = 9 \text{ ops}$$

2<sup>nd</sup> rem = st<sub>2</sub>.pop() → 1 op

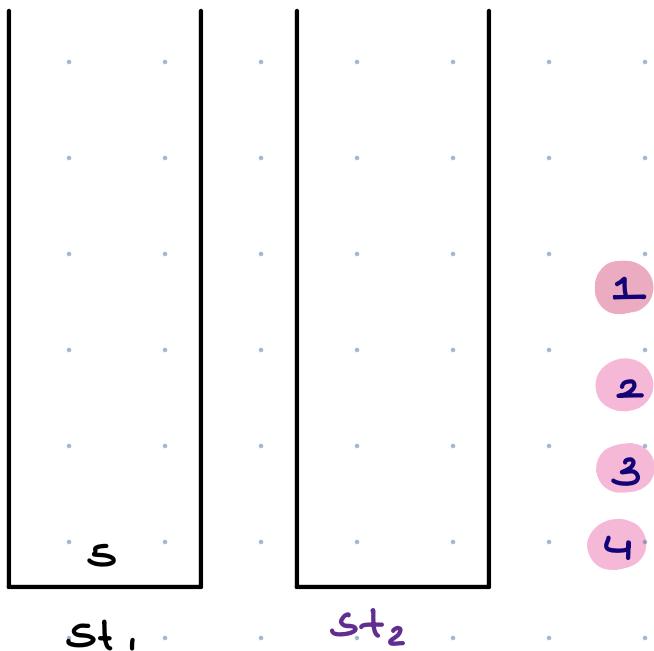
3<sup>rd</sup> rem = st<sub>2</sub>.pop() → 1 op

4<sup>th</sup> rem = st<sub>2</sub>.pop() → 1 op

4 removals = 12 ops

1 removal =  $\frac{12}{4}$  ops ≈ 3 ops

≈ O( $k$ ) ⇒ O(1)



## Doubly Ended Queue → Doubly Linked List

→ Allows insertion & removal from both the ends

- 01. AddFirst()
- 02. AddLast()
- 03. getFirst()
- 04. removeFirst()
- 05. removeLast()
- 06. getLast()

TC: O(1)

## \* Sliding Window Maximum

Given an integer [] A & a window of size k  
Find the maximum of all the windows

$A[] = \{10, 1, 9, 3, 7, 6, 5, 11, 8\}$     $k = 4$

Ans = {10, 9, 9, 7, 11, 11}

$A[] = \{1, 4, 3, 2, 5\}$     $k = 3$

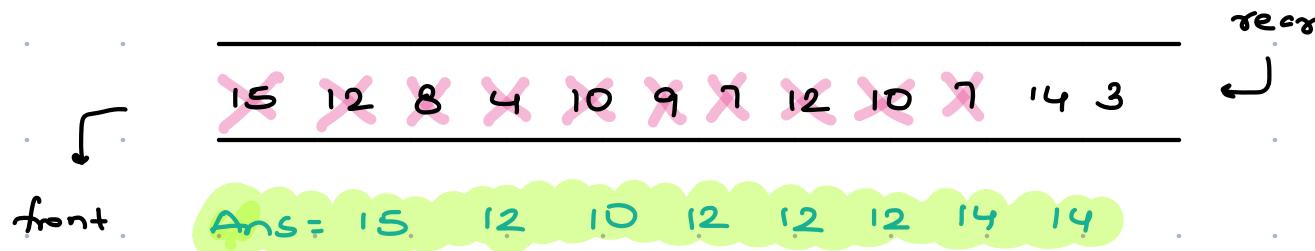
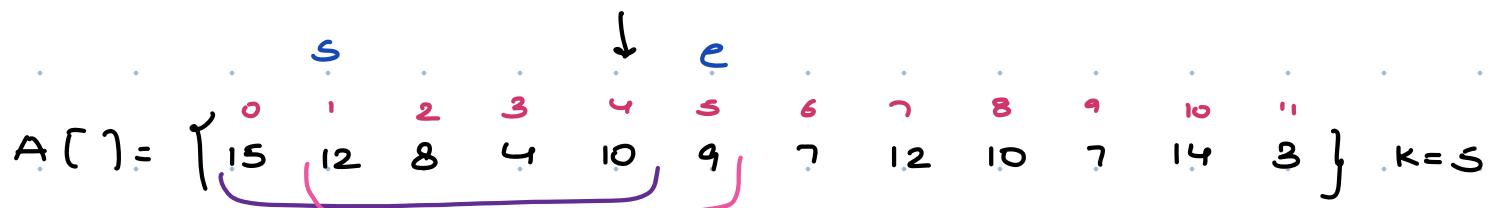
Ans = 4 4 5

Brute force  $\rightarrow$  Generate all subarrays of size  $k$ ,  
 iterate on subarr & get maximum

$$TC: O(N - k + 1) * k \Rightarrow O(N^2)$$

$$k = n/2$$

Optimal



$\rightarrow$  Insert at last & delete from last

$\rightarrow$  `deletefirst()`

$\rightarrow$  `accessfirst()`

```
int [] maxwindow ( int [] A, int k )
```

```
int idx = 0
```

```
int [] ans = new int [ N - k + 1 ]
```

```
Deque < I > dq = new ArrayDeque < > ( );
```

```

for ( i=0; i<k; i++ ) { . . . . . // insert first window

    while ( dq.size ()>0 && A[i] > dq.getLast() )
        dq.removeLast();

    dq.addLast ( A[i] );
}

ans [idx ++] = dq.getFirst();

s=1, e=k.

while ( e < n )

    if ( A[s-1] == dq.getFirst() )
        dq.removeFirst(); } removal

    while ( dq.size ()>0 && A[e] > dq.getLast() )
        dq.removeLast(); } Insertion

    dq.addLast ( A[e] );
    ans [idx ++] = dq.getFirst();
    s++, e++;

return ans;
}

```

## Scenario:

At TechTrade Inc., a trading team focuses on day-trading technology stocks. They employ a strategy that involves selling stocks at short-term peaks to maximize profits. The team uses an algorithm to determine the best time to sell stocks based on minute-to-minute price data.

### Stock Prices Array A

Consider the minute-by-minute stock prices of a tech company, TechCorp, over a 10-minute interval:

A=[220, 215, 230, 225, 240, 235, 230, 245, 250, 240]  
  ◦ 1 2 3

Queue <I> q = new LinkedList<>();