

Branch & Bound

- Used to find optimal solution to many optimization problems.
- Systematic enumerates all candidate solutions, by discarding in between candidate state by using upper and lower bounds.

→ Terminology:

1. Live Node: Node that is generated but whose child nodes are not yet generated completely.
2. E-Node: Is a live node that is currently being explored.
3. Dead Node: Is a generated node that is not to be expanded or explored any further.

Branch & Bound (BB):

All state space search methods in which all children of E-node are generated before any other live node can become the E-node.

→ General Method

- Both BFS and DFS generalized to BB strategies

- | | |
|---|---|
| 1) <u>BFS</u> - FIFO Search
In terms of live node
- List of live node is a <u>queue</u> | 2) <u>DFS</u> - LIFO Search
In terms of live node
- List of live node is a <u>stack</u> |
|---|---|

- Like Backtracking, BB uses bounding function to avoid generating subtrees that do not contain an answer node.

- FIFO - Branch & Bound : - Search the tree using BFS with ranking fn.
- LIFO - Branch & Bound : Search the tree as a BFS, but replace the FIFO queue with a stack, with ranking fn. - also known as D-Search.
- LC (Least Cost) Branch & Bound : Replace the FIFO queue with priority queue.
 - The priority of a node 'p' in the queue is based on an estimate of the likelihood that the answer node is in the subtree whose root is 'p'.
- LC Branch & Bound directs the search to parts of the space most likely to contain the answer.
- Ranking function :
 - In LC Branch & Bound, Ranking fn is the basis on which next E-node is selected.
 - In general for a node 'n', ranking function is,

$$C(n) = g(n) + f(n) \rightarrow \text{Estimate of additional effort required to reach an answer node.}$$

$g(n)$: cost of reaching n from root
 $f(n)$: estimate of additional effort required to reach an answer node.

15-Puzzle Problem

[I] Problem :

- Given a 4×4 board with 15 tiles, and every tile has one number from 1 to 15 and one empty space.
- Also, initial state and goal state is given, objective is to transform from initial state to goal state using legal moves.

1	3	4	15
2		5	12
7	6	11	14
8	9	10	13

1	2	3	4
5	6	7	8
9	10	11	12
13	14	15	

Initial state (IS)

Goal state (GS)

- Legal move involves moving a tile adjacent to the empty spot E to E.
- For this initial state, four possible moves are moving tiles 2, 3, 5, 6.
- A state is reachable from the initial state iff there is a sequence of legal moves for IS to GS.

[II] Check whether GS is reachable from IS :

- Number the frame positions from 1 to 16.
- P_i is the frame position containing tile numbered i .
- P_{16} is the position of empty spot.

$P_{IS} = [1, 5, 2, 3, 7, 10, 9, 13, 14, 15, 11, 8, 16, 12, 4, 6]$

$P_{GS} = [1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15, 16]$

→ For any state, let l_i be the number of files j such that

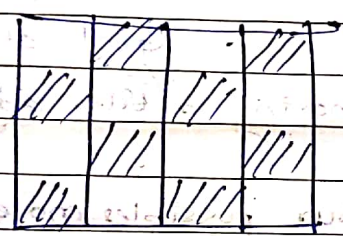
(*) $j < i$ and $p_j > p_i$

i.e. l_i gives the number of lower-numbered files that are to the right of file i in that state.

- Example,

$l_1 = 0, l_2 = 1, l_3 = 6, l_4 = 0$

→ let $n = 1$ if in the IS, the empty spot is in one of the following positions i.e. 2, 4, 5, 7, 10, 12, 13, 15. $n = 0$ otherwise.



(*) shaded, $n = 1$

Theorem

GS is reachable from IS iff

$\sum_{i=1}^{16} l_i + n$ is even.

$l_i = [0, 0, 1, 1, 0, 0, 1, 0, 0, 0, 3, 6, 0, 11, 14]$

$\sum_{i=1}^{16} l_i = 40, n = 0$

$40 + 0 = 40 \Rightarrow R$

$(R \% 2 = 0) \Rightarrow$ possible
 $f = 0 \Rightarrow$ Not possible

III] Organize state space as a tree

- child of each node 'x' represents the states that are reachable from 'x' in one legal move.
- Consider, the move as move of empty space rather than the tile.
- Empty space can have 4 legal moves:- up, down, left, right.
- No node should have a child state that is same as its parent.
- let us order the move of empty space as up, right, down, left (clockwise moves).

* DFS

- takes us away from goal GS
- Search of state space tree is blind.
- taking the leftmost path from the root regardless of IS.

* BFS

- Find a GS nearest to the root.
- It is still blind because no matter what IS, the algo. attempts the same sequence of moves.

* Intelligent fn : (LC - Branch & Bound)

- Also known as ranking fn / cost fn.
- In search using ranking fn, the next E-node generated is the one with least cost.

- For 15 puzzle, we assume that moving one tile in any direction will have 1 unit cost.

Cost fn of node(n) $\Rightarrow c(n) = g(n) + h(n)$

where, $g(n)$ = length of path from root to n
(No. of moves so far)

$h(n)$ = Number of non-blank tiles not in their goal position
(No. of misplaced tiles)

Initial state

1	2	3	4
5	6		8
9	10	7	11
13	14	15	12

Goal state

1	2	3	4
5	6	7	8
9	10	11	12
13	14	15	

①

1	2	3	4
5	6		8
9	10	7	11
13	14	15	12

up

right

down

left

$c=1+4$ ②

$c=1+4$ ③

$c=1+2$ ④

$c=1+4$ ⑤

1	2		4
5	6	3	8
9	10	7	11
13	14	15	12

1	2	3	4
5	6	8	
9	10	7	11
13	14	15	12

1	2	3	4
5	6	7	8
9	10		11
13	14	15	12

1	2	3	4
5		6	8
9	10	7	11
13	14	15	12

(blank & (down) right)

⑥ $c=2+2$

$c=2+3$ ⑦

$c=2+3$ ⑧

1	2	3	4
5	6	7	8
9	10	11	
13	14	15	12

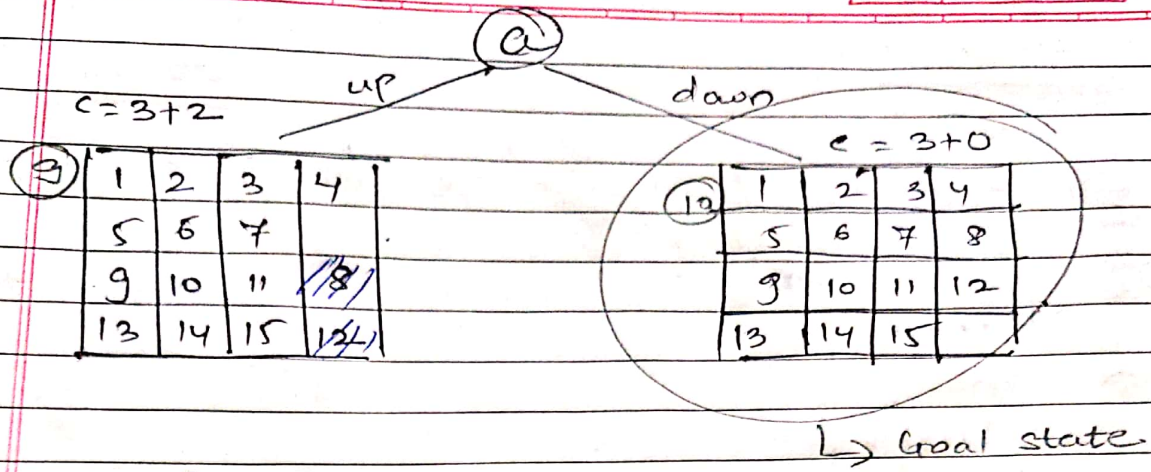
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1	2	3	4
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Algorithm

Iteration	Given Nodes (min priority Queue)	E-node
0.	$C(1)$	Node 1
1.	$C(2)=5, C(3)=5, C(4)=3, C(5)=5$	Node 4
2.	$C(2)=5, C(3)=5, C(5)=5, C(6)=3, C(7)=5, C(8)=5$	Node 6
3.	$C(2)=5, C(3)=5, C(5)=5, C(7)=5, C(8)=5, C(9)=5, C(10)=3 (GS)$	Goal state