# PPC3 Report 110062120 高小榛

## 1. Code explanation

## (1) testpreempt.c

Different with the previous checkpoint, I add a function SemaphoreCreate. This function is to create a semaphore with initial value n.

First, I create global variable Cur\_ID, Mutex, Full, Empty, Word, Head, Tail and Buffer. The meaning of these variable is define in this picture.

```
// Global variables for shared buffer and threading
__data __at (0x25) ThreadID Cur_ID; // Current thread ID
__data __at (0x36) char Mutex; // Mutex semaphore for mutual exclusion
__data __at (0x37) char Full; // Semaphore indicating buffer is full
__data __at (0x38) char Empty; // Semaphore indicating buffer is empty
__data __at (0x39) char Word; // Character to be produced
__data __at (0x3A) char Head; // Head index for buffer
__data __at (0x3B) char Tail; // Tail index for buffer
__data __at (0x3D) char Buffer[3] = {' ', ' ', ' '}; // Circular buffer
```

In the Producer function, I assign first word produce with 'A', then in the while loop, I use SemaphoreWait and SemaphoreSignal to design. First wait Empty until buffer has empty space, then we can enter critical section with wait for Mutex. In the critical section, first add new character to buffer, update tail index then update character. Finally leave the critical section with signal mutex, then signal buffer is full.

In Consumer function, first I initialize Tx for polling using the information in class., change the TMOD to TMOD  $\mid$ = 0x20. Then we should wait for new data from producer, write data to serial port Tx, poll for Tx to finish writing (TI), then finally clear the flag. So first we should wait for data in buffer with wait Full, then wait for mutex to enter critical section. In the critical section, we can do thing that mention above. Wait for Tx to be ready, send character from buffer, clear Tx interrupt flag then circular increment of head. Finally exit critical section then signal buffer is empty.

In main function, I create 3 semaphore Mutex, Full and Empty. Then initialize head index and tail index. Use ThreadCreate function to create producer thread and store with cur\_id. Finally, use assembly to set up stack pointer, call consumer function.

Additionally, I add timer0\_ISR to let ISR call my routine timer.

### (2) preemptive.c

First, I create the static global SP0 ~ SP3, Map, Cur\_ID, SP\_old, SP\_new and New\_ID. Then I define 2 macro save state and restore state in a similar way. In savestate we should saving

the current thread context. So first push ACC, B register, Data pointer registers (DPL, DPH), PSW with assembly. Then save SP into the saved Stack Pointers array with if–else block to check which place we should store the SP into it. In restorestate, I first assign SP to the saved SP from the saved stack pointer array, then pop the registers PSW, data pointer registers, B reg, and ACC.

In bootstrap function, I clear thread bitmap indicating no threads are running and initialize stack pointer. I initialize timer and interrupts for preemption with set TMOD = 0, IE = 0x82 to enable timer 0 interrupt, and set TR0 = 1 to start running timer0. Then I create a thread for main and store with  $Cur_ID$  and restore the content to run main function.

In myTimer0Handler function, first I use EA=0 to disable interrupts, then save the current thread state. Next I cycle through thread IDs to find the next valid thread and break out of the loop if a valid thread is found. Finally restore the state, re-enable interrupts then return.

In ThreadCreate function, first disable interrupt, check to see we have not reached the max #threads and return –1 if no available thread ID. Then initialize a New\_ID, run a if-else block to check each thread ID (0 to 3) to find an available one. For each thread ID, perform bitwise AND with Map and check if the result is 0. If 0, it means the thread ID is available for use. If Thread I is available: first set New\_ID to 'i', then update Map to indicate Thread I is now in use, finally set SP\_new to the starting stack location for Thread i. Then save the current SP in SP\_old, set SP to the new thread's starting stack location. Then push the return address (fp) onto the stack for the new thread. Then initialize the registers to 0 for the new thread by set A to 0 then push into each register. Then set up the PSW for the new thread and save its SP, and PSW setup depends on the thread ID to use the correct register bank. Finally restore the original SP from SP\_old, re-enable interrupt, and return new\_ID.

In ThreadYield function, I use Round–Robin to find the next thread to run. First cycle through thread IDs and select the next valid thread, then break out of the loop if a valid thread is found, finally restore state.

In ThreadExit function, I clear the current thread from the thread bitmap and set up for context switch, then modify the thread bitmap to indicate the current thread is no longer valid. After dealing with Map, also cycle through thread IDs and select the next valid thread, then break out of the loop if a valid thread is found, finally restore state.

#### (3) preemptive.h

In preemptive.h, I add some definition according to the document. First, CNAME and LABELNAME is to concatenating symbols, it can add \_ in front and \$ in tail. Then I define SemaphoreSignal is to signal a semaphore. SemaphoreWaitBody do lots of things. First check semaphore value, decoded wait if it is 0, finally decrement semaphore. Last I define SemaphoreWait, it will use SemaphoreWaitBody to wait a semaphore.

## 2. Typescript

(1) Makefile

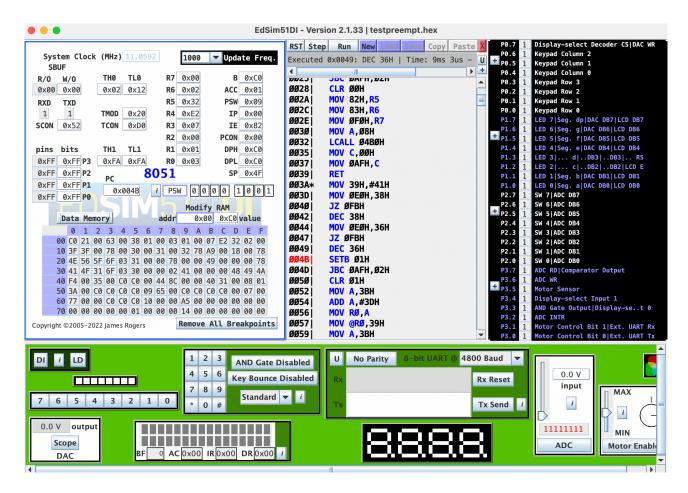
### (2) Screenshot

```
sunny@xiaozhendeAir 110062120 % cd 110062120-ppc3
sunny@xiaozhendeAir 110062120-ppc3 % make
sdcc -c testpreempt.c
sdcc -c preemptive.c
preemptive.c:281: warning 85: in function ThreadCreate unreferenced function arg
ument : 'fp'
sdcc -o testpreempt.hex testpreempt.rel preemptive.rel
sunny@xiaozhendeAir 110062120-ppc3 % make clean
rm *.hex *.ihx *.lnk *.lst *.map *.mem *.rel *.rst *.sym
rm: *.ihx: No such file or directory
rm: *.lnk: No such file or directory
make: *** [clean] Error 1
```

## 2. Screenshots and explanation

(1) Producer running

Producer is at 003A according to map.



### (2) Consumer running

Consumer is at 007F according to map.

