# Assginment4

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**1.1**  $\neg \neg A \supset A$ 

$$\begin{array}{c} A \longrightarrow A \\ \hline \neg A, A \end{array} (\neg: left) \\ \hline \neg \neg A \longrightarrow A \\ \hline \neg \neg A \supset A \end{array} (\neg: left) \\ \hline \end{array}$$

**1.2**  $A\supset (A\vee B)$ 

$$\frac{A \rightarrow A, B}{A \rightarrow (A \vee B)} \stackrel{(\vee : right)}{(\supset : right)}$$
$$A \supset (A \vee B)$$

**1.3**  $A \lor (A \land B) \supset A$ 

$$\begin{array}{cccc} A & \rightarrow & A & \hline A, B & \rightarrow & A \\ \hline A & \rightarrow & A & A & (A \lor B) & (\lor : left) \\ \hline & A \lor (A \land B) & \rightarrow & A \\ \hline & A \lor (A \land B) \supset A & (\supset : right) \end{array}$$

#### 2 Question2

**2.1**  $A\supset (B\supset (A\land B))$ 

$$\frac{B,A \rightarrow A \qquad B,A \rightarrow B}{B,A \rightarrow (AB)} \stackrel{(\wedge : right)}{(\supset : right)} \\ \frac{A \rightarrow (B \supset (A \land B))}{A \supset (B \supset (A \land B))} \stackrel{(\supset : right)}{(\supset : right)}$$

**2.2**  $(A \supset B) \supset (\neg B \supset \neg A)$ 

**2.3**  $(A \supset C) \supset ((B \supset C) \supset ((A \lor B) \supset C))$ 

$$\frac{A, (B \supset C), (A \supset C) \rightarrow C \qquad B, (B \supset C), (A \supset C) \rightarrow C}{\underbrace{(A \lor B), (B \supset C), (A \supset C) \rightarrow C}_{(B \supset C), (A \supset C) \rightarrow (A \lor B) \supset C}}_{(C): right)} \xrightarrow{(\triangle : left)} \frac{(A \supset C) \rightarrow ((B \supset C) \supset ((A \lor B) \supset C))}{(A \supset C) \supset ((B \supset C) \supset ((A \lor B) \supset C))}} \xrightarrow{(\triangle : right)} \xrightarrow{(\triangle : right)}$$

#### 3.1 $P \equiv \neg P$

Based on  $:P \equiv \neg P = (P \supset \neg P) \land (\neg P \supset P)$ 

$$\frac{P, P \rightarrow P}{P, \rightarrow P} \xrightarrow{(\neg : right)} \frac{P, P}{\neg P \rightarrow P} \xrightarrow{(\neg : left)} \frac{P \rightarrow P}{\neg P \rightarrow P} \xrightarrow{(\neg : left)} \frac{P, P}{\neg P \rightarrow P} \xrightarrow{(\neg : right)} \frac{P, P}{\neg P \rightarrow P} \xrightarrow{(\neg : ri$$

=> The only axiom is  $P\supset P$ , which is true, so  $P\equiv \neg P$  is valid and satisfiable.

#### **3.2** $(P \wedge Q) \supset (P \vee Q)$

$$\frac{P,Q \rightarrow P,Q}{P,Q \rightarrow (P \lor Q)} \xrightarrow{(\lor : right)} \frac{(P \land Q) \rightarrow (P \lor Q)}{(P \land Q) \supset (P \lor Q)} \xrightarrow{(\supset : right)}$$

=> The only axiom is  $P,Q\supset P,Q$ , which is true, so  $(P\land Q)\supset (P\lor Q)$  is valid and satisfiable.

$$\textbf{3.3} \quad (Q \supset \neg P) \equiv (P \equiv Q)$$

Based on  $(Q \supset \neg P) \equiv (P \equiv Q) = ((Q \supset \neg P) \supset ((P \supset Q) \land (Q \supset P))) \land (((P \supset Q) \land (Q \supset P)) \supset (Q \supset \neg P))$ 

$$\frac{P \rightarrow Q, Q \qquad \frac{P \rightarrow P, Q}{\neg P, P \rightarrow Q} \qquad (\neg: left)}{P, (Q \supset \neg P) \supset Q} \qquad (\supset: left) \qquad Q \rightarrow Q, P \qquad \frac{Q \rightarrow P, P}{\neg P, Q \rightarrow P} \qquad (\supset: left)}{Q, (Q \supset \neg P) \supset P} \qquad (\supset: left) \qquad Q, (Q \supset \neg P) \supset P \qquad (\supset: right) \qquad (\supset: right) \qquad (\cap: left) \qquad$$

$$\frac{(Q\supset \neg P)\supset ((P\supset Q)\land (Q\supset P))) \qquad ((P\supset Q)\land (Q\supset P))\supset (Q\supset \neg P)}{(Q\supset \neg P)\supset ((P\supset Q)\land (Q\supset P)))\land ((P\supset Q)\land (Q\supset P))\supset (Q\supset \neg P)} \land :right)$$

This formula is satisfiable and invalid.

**4.1**  $(P \supset Q) \land (\neg P \supset Q)$ 

$$\frac{P \to Q}{P \supset Q} \text{ ($\supset: right$)} \qquad \frac{ \to P, Q}{\neg P \to Q} \text{ ($\neg: left$)}$$
$$\frac{P \supset Q}{(P \supset Q) \land (\neg P \supset Q)} \text{ ($\wedge: right$)}$$

This formula is unsatisfiable and invalid.

**4.2**  $((P \supset Q) \supset R) \supset S$ 

$$\frac{P \to Q, S}{(P \supset Q), S} \xrightarrow{(\supset : right)} R \to S \xrightarrow{(\supset : left)} \frac{((P \supset Q) \supset R) \to S}{((P \supset Q) \supset R) \supset S} \xrightarrow{(\supset : right)}$$

This formula is unsatisfiable and invalid.

**4.3**  $(p \wedge Q) \supset (P \supset Q)$ 

$$\frac{P,P,Q \rightarrow Q}{P,(PQ) \rightarrow Q} \stackrel{(\wedge : left)}{(> : right)} \frac{(p \land Q) \supset (P \rightarrow Q)S}{(p \land Q) \supset (P \supset Q)} \stackrel{(\supset : right)}{(> : right)}$$

This formula is satisfiable and valid.

#### 5 Question5

**5.1** 
$$(P \supset R) \supset ((Q \supset S) \supset ((P \land Q) \supset R))$$

$$\frac{P, (Q \supset S), (P \supset R) \rightarrow R \qquad Q, (Q \supset S), (P \supset R) \rightarrow R}{\frac{(P \lor Q), (Q \supset S), (P \supset R) \rightarrow R}{(Q \supset S), (P \supset R) \rightarrow (P \lor Q) \supset R}} \underset{(\supset : right)}{\overset{(\supset : right)}{(P \supset R) \rightarrow ((Q \supset S) \supset ((P \lor Q) \supset R))}} \underset{(\supset : right)}{\overset{(\supset : right)}{(P \supset R) \supset ((Q \supset S) \supset ((P \lor Q) \supset R))}}$$

The false axiom is  $S,Q \rightarrow P,R$ , so the conjunctive normal form is :  $\neg S \lor \neg Q \lor P \lor R$ 

**5.2**  $(P \supset Q) \supset ((Q \supset \neg R) \supset \neg P)$ 

$$\frac{P \rightarrow P, Q \quad Q, P \rightarrow Q}{P, (P \supset Q) \rightarrow Q \quad (\neg : right)} \xrightarrow{\begin{array}{c} P, (P \supset Q) \rightarrow R \\ \hline P, (P \supset Q) \rightarrow Q, \neg P \end{array}} (\neg : right) \xrightarrow{\begin{array}{c} P, \neg R, (P \supset Q) \rightarrow R \\ \hline P, \neg R, (P \supset Q) \rightarrow \end{array}} (\neg : right) \xrightarrow{\begin{array}{c} P, \neg R, (P \supset Q) \rightarrow R \\ \hline P, \neg R, (P \supset Q) \rightarrow \end{array}} (\neg : right) \xrightarrow{\begin{array}{c} P, \neg R, (P \supset Q) \rightarrow R \\ \hline P, \neg R, (P \supset Q) \rightarrow \end{array}} (\neg : right) \xrightarrow{\begin{array}{c} P, \neg R, (P \supset Q) \rightarrow R \\ \hline P, \neg R, (P \supset Q) \rightarrow \end{array}} (\neg : right) \xrightarrow{\begin{array}{c} P, \neg R, (P \supset Q) \rightarrow R \\ \hline P, \neg R, (P \supset Q) \rightarrow \end{array}} (\neg : right) \xrightarrow{\begin{array}{c} P, \neg R, (P \supset Q) \rightarrow R \\ \hline P, \neg R, (P \supset Q) \rightarrow \end{array}} (\neg : right) \xrightarrow{\begin{array}{c} P, \neg R, (P \supset Q) \rightarrow R \\ \hline P, \neg R, (P \supset Q) \rightarrow \neg P \\ \hline P, \neg R,$$

The conjunctive form is  $\neg Q \lor \neg P \lor R$ 

#### 6 Question6

#### 6.1

Based on Question 5.1, the disjunctive normal form is  $(\neg S \land \neg P) \lor Q \lor R$ 

#### 6.2

Based on Question 5.2, the disjunctive normal form is  $(\neg P \lor Q) \lor (P \land \neg R) \lor ((Q \land P) \land \neg R)$ 

### 7 Question7

**7.1**  $\forall x A \supset A[t/x]$ 

$$\frac{A[t/x] \rightarrow A[t/x]}{\forall xA \rightarrow A[t/x]} (\forall : left)$$
$$\frac{\forall xA \rightarrow A[t/x]}{\forall xA \supset A[t/x]} (\supset : right)$$

**7.2**  $A[t/x] \supset \exists x A$ 

$$\frac{A[t/x] \rightarrow A[t/x]}{A[t/x] \rightarrow \exists xA \atop A[t/x] \supset \exists xA} (\exists : left)$$

### 8 Question9

Let  $L = \{CS, FS, PS\}$ 

1.  $CS = \emptyset$ 

2.  $FS = \emptyset$ 

3.  $PS = \{\leq\} \text{ with } r(\leq) = 2$ 

Let  $\Gamma$  be the following set of axioms:

1.  $\forall x \neg (x \leq x)$ ,

2.  $\forall x \forall y ((x < y) \supset \neg (y < x))$ 

3.  $\forall x \forall y \forall z (((x < y) \land (y < z)) \supset (x < z))$ 

Then  $T = (L, \Gamma)$  is a first-order theory of the weak partial orders.

### 9 Question11

The first-order language of (directed) graphs is L = r, where r is a binary relation symbol.

The only terms are the variables x.

Atomic formulas look like  $(x \approx y)$  or (rxy).

Example:

The subformulas of  $\forall x((rxy) \rightarrow \exists y(ryx))$ :

 $\forall x((rxy) \rightarrow \exists y(ryx))$ 

 $rxy(rxy) \rightarrow \exists y(ryx)$ 

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\exists y(ryx) \\ ryx.
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define a first-order language LK for K-vector spaces (or left K-modules) as follows. LK has all the usual logical symbols from first-order logic as well as a constant symbol 0 for the zero vector, a binary function symbol + for vector addition, a unary function symbol for additive inverse of vectors, and for each  $\lambda \in K$  a unary function symbol  $\lambda$  for scalar mulitiplication by  $\lambda$ . In this language we axiomatise the notion of K-vector space of positive dimension. The axioms are,

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Axioms for an abelian group:
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 \forall x \forall y \forall z (x + (y + z) = (x + y) + z) \forall x \forall y \forall z (x + (y + z) = (x + y) + z)   \forall x \forall y (x + y = y + x) \forall x \forall y (x + y = y + x)   \forall x (x + 0 = x) \forall x (x + 0 = x)   \forall x (x + (x) = 0) \forall x (x + (-x) = 0)   Axioms for scalar multiplication:   \forall x (s\lambda(s\mu(x)) = s\lambda(x)) \forall x (s\lambda(s\mu(x)) = s\lambda\mu(x)), for all \lambda, \mu \in K\lambda, \mu \in K   \forall x (s\lambda(x) + s\mu(x) = s(\lambda + \mu)(x)) \forall x (s\lambda(x) + s\mu(x) = s(\lambda + \mu)(x)), for all \lambda, \mu \in K\lambda, \mu \in K   \forall x (s\lambda(x + y) = s\lambda(x) + s\lambda(x)) \forall x (s\lambda(x + y) = s\lambda(x) + s\lambda(x)), for all \lambda \in K\lambda \in K   \forall x (s1(x) = x)
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