

Assginment3

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1 Question1

1.1 Addition

Based on the observation:

$$\text{add}(n,0)=n+0=n$$

$$\text{add}(n,m+1)=n+(m+1)=(n+m)+1=\text{succ}(\text{add}(n,m))$$

we should find two primitive recursive functions

$f:N \Rightarrow N$ and $g:N^3 \Rightarrow N$ such that:

when $m=0$, $\text{add}(n,m)=f(n)$

when $m>0$, $\text{add}(n,m)=g(n,m-1,\text{add}(n,m-1))$

since $\text{add}(n,0)=n \Rightarrow f(n)=n$

since $\text{add}(n,m+1)=\text{succ}(\text{add}(n,m))$, g needs to be a function which maps the ordered triple $(n,m,\text{add}(n,m))$

$g:N^3 \Rightarrow N$ where $g(n1,n2,n3)=\text{succ}(n3)$

because then $g(n,m,\text{add}(n,m))=\text{succ}(\text{add}(n,m))$

To show g is primitive recursive:

$g(n1,n2,n3)=\text{succ}(\text{Pr}_3^3(n1,n2,n3))$, which is achieved by using the projection function

g is obtained by substitution from the basic primitive recursive functions succ and Pr_3^3 and so is primitive recursive

In conclusion, Addition is primitive recursive

1.2 Multiplication

Based on the observation:

$$\text{mult}(n,0)=n*0=0$$

$$\text{mult}(n,m+1)=n*(m+1)=nm+n=\text{add}(n,\text{mult}(n,m))$$

we should find two primitive recursive functions

$f:N \Rightarrow N$ and $g:N^3 \Rightarrow N$ such that:

when $m=0$, $\text{mult}(n,m)=f(n)$

when $m>0$, $\text{mult}(n,m)=g(n,m-1,\text{mult}(n,m-1))$

since $\text{mult}(n,0)=0 \Rightarrow f(n)=0$, which is the basic primitive recursive function: zero function

since $\text{mult}(n,m+1)=\text{add}(n,\text{mult}(n,m))$, g needs to be a function which maps the ordered triple $(n,m,\text{mult}(n,m))$

$g:N^3 \Rightarrow N$ where $g(n1,n2,n3)=\text{add}((\text{Pr}_3^1),(\text{Pr}_3^3))$

g is obtained by substitution from the primitive recursive functions Addition and Pr_3^3 , Pr_3^1 and so is primitive recursive

In conclusion, Multiplication is primitive recursive

1.3 Exponentiation

Based on the observation:

$$\exp(n, 0) = n^0 = 1$$

$$\exp(n, m+1) = n * n^m = \text{mult}(n, \exp(n, m))$$

we should find two primitive recursive functions

$f: N \Rightarrow N$ and $g: N^3 \Rightarrow N$ such that:

when $m=0$, $\exp(n, m) = f(n)$

when $m>0$, $\exp(n, m) = g(n, m-1, \exp(n, m-1))$

since $\exp(n, 0) = 1 \Rightarrow f(n) = 1$, which is the basic primitive recursive function: constant function

since $\exp(n, m+1) = \text{mult}(n, \exp(n, m))$, g needs to be a function which maps the ordered triple $(n, m, \exp(n, m))$

$g: N^3 \Rightarrow N$ where $g(n_1, n_2, n_3) = \exp((Pr_3^1), (Pr_3^3))$

g is obtained by substitution from the primitive recursive functions Multiplication and Pr_3^3 , Pr_3^1 and so is primitive recursive

In conclusion, Multiplication is primitive recursive

2 Question2

We can define the Ackerman function as follows:

$$A(n, m) = \text{if}(m=0, 2n, \text{if}(n=0, 0, \text{if}(n=1, 2, A(m-1, A(m, n-1)))))$$

We should show this is well-founded recursive.

For that consider the strict lexicographical order between pairs:

$(n_1, m_1) < (n_2, m_2)$ iff $(n_1 < n_2)$ or $(n_1 = n_2 \text{ and } m_1 < m_2)$

Then we have that for base cases the function is well-defined. For the general case we have that: $(m-1, A(m, n-1)) < (m, n)$

3 Question3

Tree is an example. A tree has three nodes. Let A be the root node, B is A's child and C is B's child. This tree is well-founded, because A is the minimum element. The tree is not a partial order, because it is not transitive: A is not C's parent.

4 Question7

4.1 $P \equiv \neg P$

P	$\neg P$	$P \equiv \neg P$
T	F	F
F	T	F

4.2 $((P \wedge Q) \implies (P \vee Q))$

P	Q	$P \wedge Q$	$P \vee Q$	$((P \wedge Q) \implies (P \vee Q))$
T	T	T	T	T
T	F	F	T	T
F	T	F	T	T
F	F	F	F	T

4.3 $((Q \implies \neg P) \equiv (P \equiv Q))$

Q	$\neg P$	P	$Q \implies \neg P$	$P \equiv Q$	$((Q \implies \neg P) \equiv (P \equiv Q))$
T	T	F	T	F	F
F	T	F	T	T	T
T	F	T	F	T	F
F	F	T	T	F	F

4.4 $((P \implies Q) \wedge (\neg P \implies Q))$

P	Q	$\neg P$	$P \implies Q$	$\neg P \implies Q$	$((P \implies Q) \wedge (\neg P \implies Q))$
T	T	F	T	T	T
T	F	F	F	T	F
F	T	T	T	T	T
F	F	T	T	F	F

4.5 $((P \Rightarrow Q) \Rightarrow R) \Rightarrow S$

P	Q	R	R	$P \Rightarrow Q$	$(P \Rightarrow Q) \Rightarrow R$	$((P \Rightarrow Q) \Rightarrow R) \Rightarrow S$
T	T	T	T	T	T	T
F	T	T	T	T	T	T
T	F	T	T	F	T	T
T	T	F	T	T	F	T
T	T	T	F	T	T	F
F	F	T	T	T	T	T
T	F	F	T	F	T	T
T	T	F	F	T	F	T
F	T	T	F	T	T	F
F	T	F	T	T	F	T
T	F	T	F	F	T	F
F	F	F	T	T	F	T
F	T	F	F	T	F	T
F	F	T	F	T	T	F
T	F	F	F	F	T	F
F	F	F	F	T	F	T

4.6 $(P \wedge Q) \Rightarrow (P \Rightarrow Q)$

P	Q	$P \wedge Q$	$P \Rightarrow Q$	$(P \wedge Q) \Rightarrow (P \Rightarrow Q)$
T	T	T	T	T
T	F	F	F	T
F	T	F	T	T
F	F	F	T	T

5 Question8

5.1 Valid

2.2,2.6

5.2 Invalid

2.1,2.3,2.4,2.5

5.3 Satisfiable

2.2,2.3,2.4,2.5

5.4 Unsatisfiable

2.1

6 Question9

Assume B is a knave

\Rightarrow "A is a knave" is false

\Rightarrow A is knight

\Rightarrow "The two of us are both knights" is true, but we B is a knave

\Rightarrow The assumption is wrong

Assume B is a knight

\Rightarrow "A is a knave" is true

\Rightarrow A lies

\Rightarrow "The two of us are both knights" is false

\Rightarrow The assumption is right

In conclusion, B is a knight, A is a knave

7 Question10

A	B	
T	T	A:knight,B:knight
F	T	A:knave,B:knight
T	F	A:knight,B:knave
F	F	A:knave,B:knave

8 Question11

Let \uparrow denotes *nand* and Assume $(p \uparrow q) \uparrow r \Rightarrow p \uparrow (q \uparrow r)$

Assume $(p \wedge \neg r)$ is true

$\Rightarrow p, \neg r$

$\Rightarrow (\neg q) \vee (\neg r)$ (addition rule)

$\Rightarrow \neg(q \wedge r) = q \uparrow r$

$\Rightarrow p \wedge (q \uparrow r)$ (conjunction rule)

$\Rightarrow \neg \neg(p \wedge (q \uparrow r))$ (double negation introduction)

$\Rightarrow \neg(p \uparrow (q \uparrow r))$ (nand rule)

$\neg(p \uparrow q) \vee (\neg r)$ (addition)

$\Rightarrow \neg((p \uparrow q) \wedge r)$

$\Rightarrow (p \uparrow q) \uparrow r$ (nand rule)

$\Rightarrow \neg \neg((p \uparrow q) \uparrow r)$ (double negation introduction)

$(\neg \neg((p \uparrow q) \uparrow r)) \wedge (\neg(p \uparrow (q \uparrow r)))$ (conjunction rule)

$\Rightarrow \neg((\neg \neg((p \uparrow q) \uparrow r)) \vee (\neg(p \uparrow (q \uparrow r))))$

$\Rightarrow \neg((p \uparrow q) \uparrow r) \Rightarrow (p \uparrow (q \uparrow r))$

The assumption is wrong, nand is not associative.

9 Question12

$$p \uparrow r \equiv \neg(p \wedge r)$$

$$\implies p \uparrow r \equiv \neg p \vee \neg r$$

Since $\{\neg, \vee\}$ is a complete set, nand is a complete set.

10 Question13

10.1 conjunctive normal form

$$\neg(P \wedge Q \wedge R) \wedge \neg(P \wedge Q \wedge \neg R) \wedge \neg(P \wedge \neg Q \wedge R) \wedge \neg(P \wedge Q \wedge \neg R) \wedge \neg(\neg P \wedge \neg Q \wedge R)$$

10.2 disjunctive normal form

$$(P \wedge \neg Q \wedge \neg R) \vee (\neg P \wedge Q \wedge R) \vee (\neg P \wedge \neg Q \wedge \neg R)$$