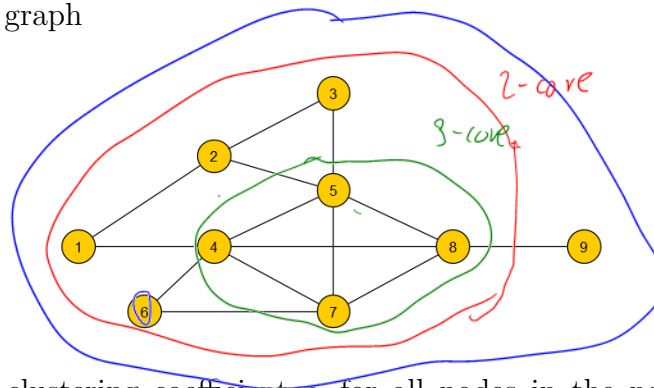


Pen & Paper Exercise 2

Social Networks

1 Local Clustering and Coreness

Consider the following graph



there cannot be
 a 4-core, because
 in the 3-core
 subgraph, no node
 has degree 4

- Compute the local clustering coefficient c_i for all nodes in the network. What is the average local clustering?
- Compute the coreness of each node in the network!
- Calculate the closeness centralities of nodes 1 and 5!

a) clustering coeff of node i :

$$c_i = \frac{2e_i}{k_i(k_i-1)}$$

$$c_1 = \frac{2 \cdot 0}{2 \cdot 1} = 0$$

$$c_2 = \frac{2 \cdot 1}{3 \cdot 2} = \frac{1}{3}$$

$$c_3 = \frac{2 \cdot 1}{2 \cdot 1} = 1$$

$c_i = \frac{\text{\# links between neighbors of } i}{k_i(k_i-1)}$
 $k_i = \text{degree of node } i$

$$c_4 = \frac{2 \cdot 4}{5 \cdot 4} = \frac{2}{5}$$

$$c_5 = \frac{2}{5}$$

$$c_6 = 1$$

$$c_7 = \frac{2}{3}$$

$$c_8 = \frac{1}{2}$$

$$c_9 = 0$$

Corenesses:

1	2	3	4	5	6	7	8	9
2	2	2	3	3	2	3	3	1

avg local clustering

$$c_{avg} = \frac{1}{9} \sum_{i=1}^9 c_i$$

$$= \frac{43}{90}$$

c) Closeness Centralities

Let $d_i = \frac{1}{N-1} \sum_{j \neq i} d_{ij}$ with d_{ij} is distance between nodes i and j

↳ avg distance

Then closeness centrality of node i
is defined as

$$C_i = \frac{1}{d_i}$$

$$d_1 = \frac{1}{8} \cdot (1+2+1+2+2+2+2+3) = \frac{15}{8}$$

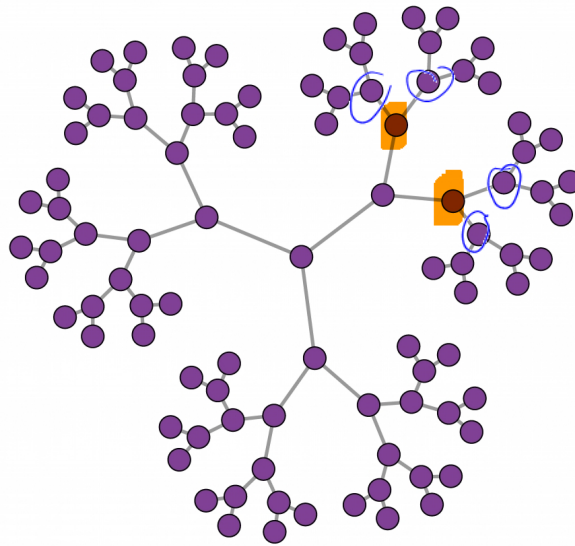
$$\hookrightarrow C_1 = \frac{1}{d_1} = \frac{8}{15}$$

$$d_5 = \frac{1}{8} (2+1+1+2+1+1+2) = \frac{11}{8}$$

$$\hookrightarrow C_5 = \frac{8}{11}$$

2 Small Worlds

A Cayley tree is a symmetric regular tree in which each node is connected to the same number of k others, until we get out to the leaves. A Cayley tree with $k = 3$ looks like this:



Show that the number of nodes which are reachable in exactly d steps from the central node is $k(k-1)^{d-1}$. Further, find an expression for the diameter of the network in terms of k and the number of nodes n . Does the Cayley tree display the small-world effect? Explain your answer.

let $r(d)$ be the number of nodes reachable from the center in d steps. We want to show that $r(d) = k(k-1)^{d-1}$

↳ we prove this claim via mathematical induction

Base Case: $d=1$

• the central node has k neighbors, so $r(1) = k = k(k-1)^{1-1}$ ✓

Induction step $d \rightarrow d+1$:

• by induction hypothesis, we have that $r(d) = k(k-1)^{d-1}$.

now every node that is reachable in d steps, has $k-1$ neighbors reachable

in the next step.

$$\hookrightarrow r(d+1) = r(d) \cdot (k-1) \stackrel{\text{Ind. Hyp.}}{=} k(k-1)^{d-1} \cdot (k-1) = k(k-1)^{d+1-1}$$

\rightarrow the claim follows by the principle of mathematical induction

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- Express number of nodes N in terms of d and k , consider the diameter is $2 \cdot d$

$$N = 1 + \sum_{l=1}^d r(l) = 1 + \sum_{l=1}^d k(k-1)^{l-1}$$

$\underbrace{\quad}_{\text{central node}}$
 $\underbrace{\quad}_{\text{\# nodes reachable in } l=1, \dots, d \text{ steps}}$

$$= 1 + k \sum_{l=1}^d (k-1)^{l-1}$$

reindex $c=l-1$

$$= 1 + k \sum_{c=0}^{d-1} (k-1)^c$$

$\underbrace{\quad}_{\text{geometric sum}}$

$$= 1 + k \cdot \frac{(k-1)^{d-1+1} - 1}{(k-1) - 1}$$

geometric sum

$$\sum_{l=0}^n a^l = \frac{a^{n+1} - 1}{a - 1}$$

$$a = k-1$$

$$n = d-1$$

$$\hookrightarrow N = 1 + k \cdot \frac{(k-1)^d - 1}{k-2} \quad (*)$$

\hookrightarrow we want to express diameter in terms of N and k

\hookrightarrow express d in terms of N, k

$(*)$ is equivalent to
$$d = \frac{\log(N(k-2)+2) - \log(k)}{\log(k-1)}$$

\hookrightarrow so in essence we have that $d \in O(\log(N))$

\hookrightarrow since diameter is $2d$, diameter is in $O(\log(N))$ as well

\hookrightarrow the network displays the small world property

3 Betweenness Centrality in a Ring Graph

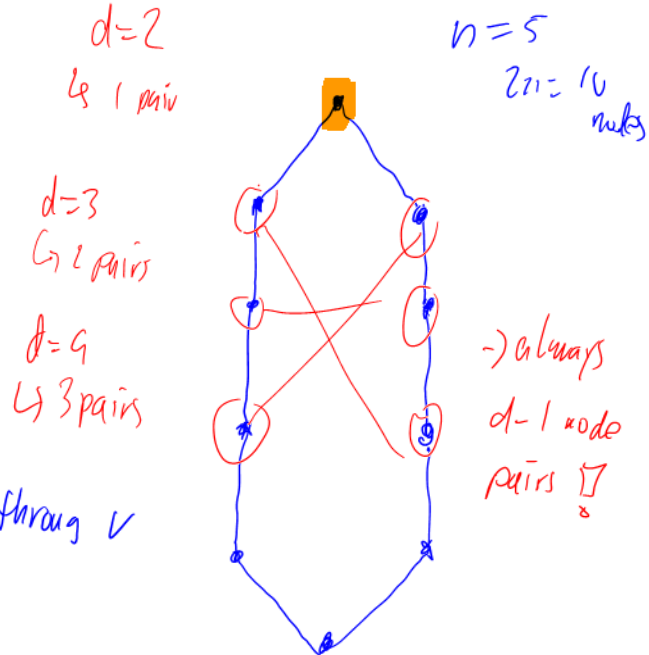
Consider a network of $2n$ nodes arranged in a ring. Each node has two neighbors. In such a network all nodes have the same betweenness centrality. Derive the formula for this value as a function of n .

Betweenness centrality of node v :

$$b(v) = \sum_{u \neq v \in V} \frac{\sigma_{uv}(v)}{\sigma_{uv}}$$

σ_{uw} = # number of shortest paths between u and w

$\sigma_{uv}(v) = \text{---} \text{---} \text{---}$ which go through v



- Case 1: nodes which are opposite in the graph
 - ↳ have maximum possible distance n
 - ↳ each of these opposite node pairs has 2 shortest paths, of which only 1 passes through node v → these contribute $\frac{1}{2}$ to $b(v)$
 - ↳ there are $n-1$ such pairs in total
- Case 2: node pairs with distance d between 2 and $n \rightarrow 2 \leq d \leq n-1$
 - ↳ How many node pairs with distance d are there which pass through v ? $d-1 \nabla_0$
 - ↳ each of these pairs has exactly one shortest path, which goes through v , so for these pairs u, w we have $\frac{\sigma_{uw}(v)}{\sigma_{uw}} = 1$

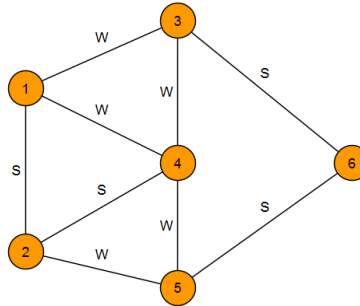
↳ If we sum up the contributions of all node pairs of cases 1+2, we get

$$b(v) = \underbrace{\frac{1}{2} \cdot (n-1)}_{\text{Case 1}} + \underbrace{\sum_{d=2}^{n-1} (d-1) \cdot 1}_{\text{Case 2}} = \frac{n-1}{2} + \sum_{d=1}^{n-2} d = \frac{n-1}{2} + \frac{(n-1)(n-2)}{2} = \frac{(n-1)^2}{2}$$

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4 Weak Ties and Strong Triadic Closure

Consider the following labeled graph, in which the are labeled as weak (W) or strong (S):



Do all nodes in this graph fulfill the strong triadic closure property? Provide an explanation to your answer!

Triadic Closure: is given iff for any 3 nodes A, B, C , from A having strong ties with B and C it follows that there is an edge between nodes B and C

↳ not given, because 6 has strong ties to 3+5, which are not connected