

# Stack Usage

VE281 - Data Structures and Algorithms, Xiaofeng Gao, Autumn 2019

## Hanoi Tower

### 1. Pseudocodes

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**Algorithm 1:** Hanoi( $n, a, b, c$ )

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**Input:**  $a, b$  and  $c$ .  $n$  disks

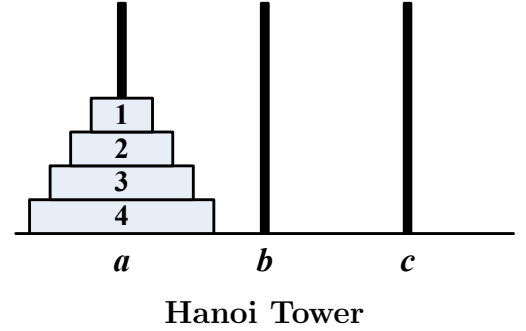
**Output:** Moving strategy from  $a$  to  $c$ .

```

1 if  $n = 1$  then
2   | Move( $a, 1, c$ );
3 else
4   | Hanoi( $n - 1, a, c, b$ );
5   | Move( $a, n, c$ );
6   | Hanoi( $n - 1, b, a, c$ );

```

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### 2. Function Operations

Hanoi(4,  $a, b, c$ )

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⊢ Hanoi(3,  $a, c, b$ )
|   ⊢ Hanoi(2,  $a, b, c$ )
|   |   ⊢ Hanoi(1,  $a, c, b$ )
|   |   |   ⊢ Move( $a, 1, b$ )
|   |   |   ⊢ Move( $a, 2, c$ )
|   |   |   ⊢ Hanoi(1,  $b, a, c$ )
|   |   |   |   ⊢ Move( $b, 1, c$ )
|   |   ⊢ Move( $a, 3, b$ )
|   |   ⊢ Hanoi(2,  $c, a, b$ )
|   |   |   ⊢ Hanoi(1,  $c, b, a$ )
|   |   |   |   ⊢ Move( $c, 1, a$ )
|   |   |   |   ⊢ Move( $c, 2, b$ )
|   |   |   |   ⊢ Hanoi(1,  $a, c, b$ )
|   |   |   |   |   ⊢ Move( $a, 1, b$ )
|   |   ⊢ Move( $a, 4, c$ )
|   ⊢ Hanoi(3,  $b, a, c$ )
|   |   ⊢ Hanoi(2,  $b, c, a$ )
|   |   |   ⊢ Hanoi(1,  $b, a, c$ )
|   |   |   |   ⊢ Move( $b, 1, c$ )
|   |   |   |   ⊢ Move( $b, 2, a$ )
|   |   |   |   ⊢ Hanoi(1,  $c, b, a$ )
|   |   |   |   |   ⊢ Move( $c, 1, a$ )
|   |   ⊢ Move( $b, 3, c$ )
|   ⊢ Hanoi(2,  $a, b, c$ )
|   |   ⊢ Hanoi(1,  $a, c, b$ )
|   |   |   ⊢ Move( $a, 1, b$ )
|   |   |   ⊢ Move( $a, 2, c$ )
|   |   |   ⊢ Hanoi(1,  $b, a, c$ )
|   |   |   |   ⊢ Move( $b, 1, c$ )

```

**Stack Usage:**

```

⊢ H(4,  $a, b, c$ )
⊢ H(3,  $a, c, b$ ), H(4,  $a, b, c$ )
⊢ H(2,  $a, b, c$ ), H(3,  $a, c, b$ ), H(4,  $a, b, c$ )
⊢ H(2,  $a, b, c$ ), H(3,  $a, c, b$ ), H(4,  $a, b, c$ )
⊢ H(3,  $a, c, b$ ), H(4,  $a, b, c$ )
⊢ H(2,  $a, b, c$ ), H(3,  $a, c, b$ ), H(4,  $a, b, c$ )
⊢ H(2,  $a, b, c$ ), H(3,  $a, c, b$ ), H(4,  $a, b, c$ )
⊢ H(4,  $a, b, c$ )
⊢ H(3,  $a, c, b$ ), H(4,  $a, b, c$ )
⊢ H(2,  $c, a, b$ ), H(3,  $a, c, b$ ), H(4,  $a, b, c$ )
⊢ H(2,  $c, a, b$ ), H(3,  $a, c, b$ ), H(4,  $a, b, c$ )
⊢ H(3,  $a, c, b$ ), H(4,  $a, b, c$ )
⊢ H(2,  $c, a, b$ ), H(3,  $a, c, b$ ), H(4,  $a, b, c$ )
⊢ H(2,  $c, a, b$ ), H(3,  $a, c, b$ ), H(4,  $a, b, c$ )
⊢ H(2,  $c, a, b$ ), H(3,  $a, c, b$ ), H(4,  $a, b, c$ )
⊢ H(2,  $c, a, b$ ), H(3,  $a, c, b$ ), H(4,  $a, b, c$ )
⊢ H(2,  $c, a, b$ ), H(3,  $a, c, b$ ), H(4,  $a, b, c$ )
⊢ H(4,  $a, b, c$ )
⊢ H(3,  $b, a, c$ ), H(4,  $a, b, c$ )
⊢ H(2,  $b, c, a$ ), H(3,  $b, a, c$ ), H(4,  $a, b, c$ )
⊢ H(2,  $b, c, a$ ), H(3,  $b, a, c$ ), H(4,  $a, b, c$ )
⊢ H(3,  $b, a, c$ ), H(4,  $a, b, c$ )
⊢ H(2,  $b, c, a$ ), H(3,  $b, a, c$ ), H(4,  $a, b, c$ )
⊢ H(2,  $b, c, a$ ), H(3,  $b, a, c$ ), H(4,  $a, b, c$ )
⊢ H(2,  $b, c, a$ ), H(3,  $b, a, c$ ), H(4,  $a, b, c$ )
⊢ H(4,  $a, b, c$ )
⊢ H(3,  $b, a, c$ ), H(4,  $a, b, c$ )
⊢ H(2,  $a, b, c$ ), H(3,  $b, a, c$ ), H(4,  $a, b, c$ )
⊢ H(2,  $a, b, c$ ), H(3,  $b, a, c$ ), H(4,  $a, b, c$ )
⊢ H(3,  $b, a, c$ ), H(4,  $a, b, c$ )
⊢ H(2,  $a, b, c$ ), H(3,  $b, a, c$ ), H(4,  $a, b, c$ )
⊢ H(2,  $a, b, c$ ), H(3,  $b, a, c$ ), H(4,  $a, b, c$ )
⊢ H(2,  $a, b, c$ ), H(3,  $b, a, c$ ), H(4,  $a, b, c$ )

```

# Parenthesis Matching

## 1. Introduction

Stack can be used to deal with the **Parenthesis Matching** problem. Actually, we can implement a formula calculator with stack.

When the calculator is doing the operation, we must compare the priority order of the operators. The relative orders are shown in table on the right side, where  $\theta_1$  and  $\theta_2$  represent the input operator and “>, <, =” represent the relative order between them.

Every time the calculator reads an input operator  $\theta_2$ , it should always be compared with  $\theta_1$  that is on the top of “StackR”. If  $\theta_1 < \theta_2$ , then  $\theta_2$  should be pushed into “StackR”. Otherwise,  $\theta_2$  should be popped out of stack. Then the calculator should fetch two numbers stored in “StackD” and carry out the operation.

**Relative order between operators.**

$\theta_1 \backslash \theta_2$	+	-	*	/	(	)	#
+	>	>	<	<	<	>	>
-	>	>	<	<	<	>	>
*	>	>	>	>	<	>	>
/	>	>	>	>	<	>	>
(	<	<	<	<	<	=	
)	>	>	>	>		>	>
#	<	<	<	<	<		=

## 2. Example

Here we provide an example with the input “#(9 - 3) \* 2/(5 + 1)#”. Note that the operator “#” indicates the start and end points of the formula.

No	StackR	StackD	Inputs	Operations	Explanations
0			#(9 - 3) * 2/(5 + 1)#	push(StackR, #)	Push starting # into stack.
1	#		(9 - 3) * 2/(5 + 1)#	push(StackR, (	StackR top, “#” < “(”. Push “(”.
2	#(		9 - 3) * 2/(5 + 1)#	push(StackD, 9)	Number 9 in StackD.
3	#(	9	- 3) * 2/(5 + 1)#	push(StackR, -)	“(” < “-”. Push “-”.
4	#(-	9	3) * 2/(5 + 1)#	push(StackD, 3)	Number 3 in StackD.
5	#(-	9 3	) * 2/(5 + 1)#	operate(9 - 3)	“-” > “)”. Execute “9 - 3”
6	#(	6	) * 2/(5 + 1)#	pop(StackR)	Eliminate a pair of parentheses.
7	#	6	* 2/(5 + 1)#	push(StackR, *)	“#” < “*”. Push “*”.
8	#*	6	2/(5 + 1)#	push(StackD, 2)	Number 2 in StackD.
9	#*	6 2	/(5 + 1)#	operate(6 * 2)	“*” > “/”. Execute “6 * 2”
10	#	12	/(5 + 1)#	push(StackR, /)	“#” < “/”. Push “/”.
11	#/	12	(5 + 1)#	push(StackR, (	“#” < “(”. Push “(”.
12	#/(	12	5 + 1)#	push(StackD, 5)	Number 5 in StackD.
13	#/(	12 5	+ 1)#	push(StackR, +)	“(” < “+”. Push “+”.
14	#/(+	12 5	1)#	push(StackD, 1)	Number 1 in StackD.
15	#/(+	12 5 1	)#	operate(5 + 1)	“+” > “)”. Execute “5 + 1”
16	#/(	12 6	)#	pop(StackR)	Eliminate a pair of parentheses.
17	#/	12 6	#	operate(12/6)	“/” > “#”. Execute “12/6”
18	#	2	#	pop(StackD)	Read “#”, end.