

# ACM/ICPC Template Manaual

# QUST

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# 0 Include

```
1 #include <bits/stdc++.h>
 2 using namespace std;
 3 #define clr(a, x) memset(a, x, sizeof(a))
 4 #define mp(x, y) make_pair(x, y)
5 #define pb(x) push_back(x)
 6 #define X first
   #define Y second
 7
 8 #define fastin
        ios_base::sync_with_stdio(0); \
 9
         cin.tie(0);
10
typedef long long ll;
typedef long double ld;
13 typedef pair<int, int> PII;
14 typedef vector<int> VI;
15 const int INF = 0x3f3f3f3f;
16 const int mod = 1e9 + 7;
17 const double eps = 1e-6;
18
   int main()
19
20 {
21 #ifndef ONLINE_JUDGE
        freopen("test.in", "r", stdin);
freopen("test.out", "w", stdout);
22
23
    #endif
24
25
         return 0;
26
27 }
```

# 1 Math

# 1.1 Prime

#### 1.1.1 Eratosthenes Sieve

```
O(n \log \log n) maxn
   notprime[i] = 0/1 0  1
1 const int maxn = "Edit";
   bool notprime[maxn] = {1, 1};
                                     // 0 && 1
   void GetPrime()
3
   {
4
        for (int i = 2; i < maxn; i++)</pre>
5
            if (!notprime[i] && i <= maxn / i) // √n</pre>
6
                for (int j = i * i; j < maxn; j += i)
7
                    notprime[j] = 1;
8
9
   }
   1.1.2 Eular Sieve
   O(n)
           phi[] prime[]
                           tot
     \mathbf{n}
1 const int maxn = "Edit";
2 bool vis[maxn];
3 int tot, phi[maxn], prime[maxn];
4 void CalPhi(int n)
5
        clr(vis, 0);
6
7
        phi[1] = 1;
8
        tot = 0;
9
        for (int i = 2; i < n; i++)
10
            if (!vis[i])
11
                prime[tot++] = i, phi[i] = i - 1;
12
            for (int j = 0; j < tot; j++)
13
14
                if (i * prime[j] > n) break;
15
                vis[i * prime[j]] = 1;
16
                if (i % prime[j] == 0)
17
18
                     phi[i * prime[j]] = phi[i] * prime[j];
19
20
21
                }
22
                else
                     phi[i * prime[j]] = phi[i] * (prime[j] - 1);
```

#### 1.1.3 Prime Factorization

```
fact[i][0]^{fact[i][1]}
```

}

}

 $\frac{23}{24}$ 

25 26 }

```
ll fact[100][2];
   int getFactors(ll x)
2
3
        int cnt = 0;
4
        for (int i = 0; prime[i] <= x / prime[i]; i++)</pre>
5
6
            fact[cnt][1] = 0;
7
            if (x % prime[i] == 0)
8
9
                fact[cnt][0] = prime[i];
10
                while (x % prime[i] == 0) fact[cnt][1]++, x /= prime[i];
11
12
                cnt++;
            }
13
        }
14
        if (x != 1) fact[cnt][0] = x, fact[cnt++][1] = 1;
15
        return cnt;
16
17
   }
   1.1.4 Miller Rabin
              2^{63}
   O(s \log n)
   bool Miller_Rabin(ll n, int s)
2
   {
        if (n == 2) return 1;
3
        if (n < 2 | | !(n & 1)) return 0;
4
        int t = 0;
5
        ll x, y, u = n - 1;
6
        while ((u \& 1) == 0) t++, u >>= 1;
7
        for (int i = 0; i < s; i++)
8
9
10
            ll\ a = rand() \% (n - 1) + 1;
            11 x = Pow(a, u, n);
11
            for (int j = 0; j < t; j++)
12
13
                ll y = Mul(x, x, n);
14
                if (y == 1 \&\& x != 1 \&\& x != n - 1) return 0;
15
16
                x = y;
17
            if (x != 1) return 0;
18
19
20
        return 1;
21 }
   1.1.5 Segment Sieve
      [a,b)
   is_prime[i-a]=true i
   a < b \le 10^{12}, b - a \le 10^6
1 const int maxn = "Edit";
2 bool is_prime_small[maxn], is_prime[maxn];
3 int prime[maxn];
4 int segment_sieve(ll a, ll b)
5
   {
6
        int tot = 0;
```

```
for (ll i = 0; i * i < b; ++i)
7
            is_prime_small[i] = true;
8
       for (ll i = 0; i < b - a; ++i)
9
            is_prime[i] = true;
10
       for (ll i = 2; i * i < b; ++i)
11
            if (is_prime_small[i])
12
13
                for (ll j = 2 * i; j * j < b; j += i)
14
                    is_prime_small[j] = false;
15
                for (ll j = max(2LL, (a + i - 1) / i) * i; j < b; j += i)
16
                    is_prime[j - a] = false;
17
18
       for (ll i = 0; i < b - a; ++i)
19
20
           if (is_prime[i]) prime[tot++] = i + a;
21
       return tot;
   }
22
   1.2 Eular phi
   1.2.1 Eular
   ll Euler(ll n)
1
2
3
       ll rt = n;
       for (int i = 2; i * i <= n; i++)
4
           if (n \% i == 0)
5
6
7
                rt -= rt / i;
8
                while (n \% i == 0) n /= i;
9
       if (n > 1) rt -= rt / n;
10
       return rt;
11
12 }
   1.2.2 Sieve
1 const int N = "Edit";
   int phi[N] = \{0, 1\};
   void CalEuler()
3
   {
4
       for (int i = 2; i < N; i++)
5
            if (!phi[i])
6
                for (int j = i; j < N; j += i)
7
8
                    if (!phi[j]) phi[j] = j;
9
                    phi[j] = phi[j] / i * (i - 1);
10
                }
11
12 }
   1.3 Basic Number Theory
   1.3.1 Extended Euclidean
  ll exgcd(ll a, ll b, ll &x, ll &y)
1
2
   {
3
       if (b) d = exgcd(b, a \% b, y, x), y -= x * (a / b);
```

```
else x = 1, y = 0;
       return d;
7 }
   1.3.2 ax+by=c
      : X = x + k * dx, Y = y - k * dy
1 #define Mod(a, b) (((a) % (b) + (b)) % (b))
   bool solve(ll a, ll b, ll c, ll& x, ll& y, ll& dx, ll& dy)
3
       if (a == 0 \&\& b == 0) return 0;
4
5
       11 x0, y0;
6
       11 d = exgcd(a, b, x0, y0);
       if (c % d != 0) return 0;
7
       dx = b / d, dy = a / d;
8
       x = Mod(x0 * c / d, dx);
9
       y = (c - a * x) / b;
10
       // y = Mod(y0 * c / d, dy); x = (c - b * y) / a;
11
12
       return 1;
13 }
   1.3.3 Multiplicative Inverse Modulo
     exgcd a m , gcd(a, m) == 1.
1 ll inv(ll a, ll m)
2
   {
3
       11 x, y;
       ll d = exgcd(a, m, x, y);
       return d == 1 ? (x + m) % m : -1;
5
6 }
   a 
1 ll inv(ll a, ll p) { return Pow(a, p - 2, p); }
1 for (int i = 2; i < n; i++) inv[i] = inv[p % i] * (p - p / i) % p;
   1.4 Modulo Linear Equation
   1.4.1 Chinese Remainder Theory
   X = r_i(modm_i); m_i
       X = re + k * mo
1 void crt(ll r[], ll m[], ll n, ll &re, ll &mo)
2
       mo = 1, re = 0;
3
       for (int i = 0; i < n; i++) mo *= m[i];</pre>
4
       for (int i = 0; i < n; i++)
5
6
           ll x, y, tm = mo / m[i];
7
           ll d = exgcd(tm, m[i], x, y);
8
           re = (re + tm * x * r[i]) % mo;
```

```
}
10
        re = (re + mo) \% mo;
11
   }
12
   1.4.2 ExCRT
   X = r_i(modm_i); m_i
       X = re + k * mo;
   bool excrt(ll r[], ll m[], ll n, ll &re, ll &mo)
1
2
3
        11 x, y;
        mo = m[0], re = r[0];
4
        for (int i = 1; i < n; i++)
5
6
7
            ll d = exgcd(mo, m[i], x, y);
            if ((r[i] - re) % d != 0) return 0;
8
            x = (r[i] - re) / d * x % (m[i] / d);
9
            re += x * mo;
10
            mo = mo / d * m[i];
11
            re %= mo;
12
13
        re = (re + mo) \% mo;
14
15
        return 1;
16 }
         Combinatorics
   1.5.1 Combination
   0 \leq m \leq n \leq 1000
   const int maxn = 1010;
1
   11 C[maxn][maxn];
2
  void CalComb()
3
   {
4
        C[0][0] = 1;
5
6
        for (int i = 1; i < maxn; i++)
7
            C[i][0] = 1;
8
            for (int j = 1; j \leftarrow i; j++) C[i][j] = (C[i-1][j-1] + C[i-1][j]) % mod;
9
10
   }
11
   0 \le m \le n \le 10^5, p
   const int maxn = 100010;
  ll f[maxn];
  ll inv[maxn]; //
   void CalFact()
4
5
        f[0] = 1;
6
7
        for (int i = 1; i < maxn; i++) f[i] = (f[i - 1] * i) % p;
        inv[maxn - 1] = Pow(f[maxn - 1], p - 2, p);
8
        for (int i = maxn - 2; \sim i; i--) inv[i] = inv[i + 1] * (i + 1) % p;
9
10
   ll C(int n, int m) { return f[n] * inv[m] % p * inv[n - m] % p; }
```

#### 1.5.2 Lucas

```
1 \le n, m \le 1000000000, 1 
1 const int maxp = 100010;
2 11 f[maxn];
   ll inv[maxn]; //
3
   void CalFact()
5
6
        f[0] = 1;
        for (int i = 1; i < maxn; i++) f[i] = (f[i - 1] * i) % p;
7
        inv[maxn - 1] = Pow(f[maxn - 1], p - 2, p);
8
        for (int i = maxn - 2; \sim i; i--) inv[i] = inv[i + 1] * (i + 1) % p;
9
10 }
11 ll Lucas(ll n, ll m, ll p)
12 {
        ll ret = 1;
13
        while (n && m)
14
15
            ll a = n \% p, b = m \% p;
16
            if (a < b) return 0;
17
            ret = ret * f[a] % p * inv[b] % p * inv[a - b] % p;
18
19
            n \neq p, m \neq p;
20
21
        return ret;
22 }
   1.5.3 Big Combination
   0 \le n \le 10^9, 0 \le m \le 10^4, 1 \le k \le 10^9 + 7
1 vector<int> v;
   int dp[110];
3 ll Cal(int l, int r, int k, int dis)
   {
4
        ll res = 1;
5
        for (int i = 1; i <= r; i++)</pre>
6
7
8
            int t = i;
9
            for (int j = 0; j < v.size(); j++)</pre>
10
11
                int y = v[j];
12
                while (t % y == 0) dp[j] += dis, t /= y;
13
            res = res * (11)t % k;
14
15
16
        return res;
17
   11 Comb(int n, int m, int k)
19
   {
        clr(dp, 0);
20
        v.clear();
21
22
        int tmp = k;
        for (int i = 2; i * i <= tmp; i++)</pre>
23
            if (tmp \% i == 0)
24
25
            {
26
                int num = 0;
27
                while (tmp % i == 0) tmp /= i, num++;
```

```
v.pb(i);
28
29
         if (tmp != 1) v.pb(tmp);
30
         ll ans = Cal(n - m + 1, n, k, 1);
31
         for (int j = 0; j < v.size(); j++) ans = ans * Pow(v[j], dp[j], k) % k;
32
         ans = ans * inv(Cal(2, m, k, -1), k) % k;
33
34
         return ans;
35
   }
    1.5.4 Polya
                     gcd(i, n)
    N*N\atop,\frac{m^8+17m^4+6m^2}{24},\frac{c^{n^2+3}}{m^4+11m^2}+2c^{\frac{n^2+1}{2}}+2c^{n\frac{n+1}{2}}+2c^{\frac{n(n+1)}{2}}
1 // n c
   ll solve(int c, int n)
3
         if (n == 0) return 0;
4
         11 \text{ ans} = 0;
5
         for (int i = 1; i \le n; i++) ans += Pow(c, __gcd(i, n));
6
         if (n & 1) ans += n * Pow(c, n + 1 >> 1);
else ans += n / 2 * (1 + c) * Pow(c, n >> 1);
7
9
         return ans / n / 2;
10 }
    1.6 Fast Power
   ll Mul(ll a, ll b, ll mod)
1
2
         11 t = 0;
3
         for (; b; b >>= 1, a = (a << 1) % mod)
4
             if (b \& 1) t = (t + a) \% mod;
5
         return t;
6
7
8
   ll Pow(ll a, ll n, ll mod)
9
    {
10
         ll t = 1;
         for (; n; n >>= 1, a = (a * a % mod))
11
              if (n \& 1) t = (t * a % mod);
12
13
         return t;
14 }
         Mobius Inversion
    1.7.1 Mobius
    F(n) = \sum_{d|n} f(d) \Rightarrow f(n) = \sum_{d|n} \mu(d) F(\frac{n}{d})
    F(n) = \sum_{n|d} f(d) \Rightarrow f(n) = \sum_{n|d} \mu(\frac{d}{n}) F(d)
1 ll ans;
2 const int maxn = "Edit";
int n, x, prime[maxn], tot, mu[maxn];
4 bool check[maxn];
5 void calmu()
```

```
{
6
7
        mu[1] = 1;
        for (int i = 2; i < maxn; i++)
8
9
            if (!check[i]) prime[tot++] = i, mu[i] = -1;
10
            for (int j = 0; j < tot; j++)
11
12
                if (i * prime[j] >= maxn) break;
13
                check[i * prime[j]] = true;
14
                if (i % prime[j] == 0)
15
16
17
                    mu[i * prime[j]] = 0;
18
                    break;
19
                else mu[i * prime[j]] = -mu[i];
20
            }
21
22
        }
23
   }
```

## 1.7.2 Number of Coprime-pair

```
n \quad (n \le 100000), \quad n
   ll solve()
1
2
   {
        int b[100005];
3
        11 \text{ _max}, \text{ ans } = 0;
4
        clr(b, 0);
5
        for (int i = 0; i < n; i++)
6
7
8
             scanf("%d", &x);
9
             if (x > _max) _max = x;
             b[x]++;
10
11
        for (int i = 1; i <= _max; i++)
12
13
             int cnt = 0;
14
             for (ll j = i; j \le \max; j += i) cnt += b[j];
15
             ans += 1LL * mu[i] * cnt * cnt;
16
17
        return (ans - b[1]) / 2;
18
   }
19
```

#### 1.7.3 VisibleTrees

```
gcd(x,y) = 1 , x ≤ n,y ≤ m

1  ll solve(int n, int m)
2  {
3     if (n < m) swap(n, m);
4     ll ans = 0;
5     for (int i = 1; i <= m; ++i) ans += (ll)mu[i] * (n / i) * (m / i);
6     return ans;
7  }</pre>
```

### 1.8 Fast Transformation

#### 1.8.1 FFT

```
1 const double PI = acos(-1.0);
3
  struct Complex
4
        double x, y; //
5
                         x+yi
6
        Complex(double _x = 0.0, double _y = 0.0) { x = _x, y = _y; }
7
        Complex operator-(const Complex& b) const { return Complex(x - b.x, y - b.y); }
        Complex operator+(const Complex& b) const { return Complex(x + b.x, y + b.y); }
8
        Complex operator*(const Complex& b) const { return Complex(x * b.x - y * b.y, x * b
9
        .y + y * b.x); }
10 };
11
   * FFT IFFT
13 * i (i
  * len 2
14
   */
15
   void change(Complex y[], int len)
16
17
   {
        for (int i = 1, j = len / 2; i < len - 1; i++)
18
19
            if (i < j) swap(y[i], y[j]);</pre>
20
            //
21
                   ,i<j
            //i +1,j +1, ij
22
            int k = len / 2;
23
            while (j >= k) j -= k, k /= 2;
24
            if (j < k) j += k;
25
        }
26
   }
27
28
   * FFT
29
  * len 2^k
   * on==1 DFT.on==-1 IDFT
32
33 void fft(Complex y[], int len, int on)
34 {
        change(y, len);
35
        for (int h = 2; h <= len; h <<= 1)
36
37
            Complex wn(cos(-on * 2 * PI / h), sin(-on * 2 * PI / h));
38
            for (int j = 0; j < len; <math>j += h)
39
40
                Complex w(1, 0);
41
                for (int k = j; k < j + h / 2; k++)
42
43
                    Complex u = y[k];
44
45
                    Complex t = w * y[k + h / 2];
                    y[k] = u + t, y[k + h / 2] = u - t;
46
                    W = W * Wn;
47
                }
48
            }
49
50
        if (on == -1)
51
52
            for (int i = 0; i < len; i++) y[i].x /= len;
53 }
```

#### 1.8.2 NTT

```
.G P G^{\frac{P-1}{n}} w_n = e^{\frac{2i\pi}{n}}
                                       P G 1.11
   const int mod = 119 << 23 | 1;</pre>
   const int G = 3;
3
   int wn[20];
4 void getwn()
5
  { //
6
        for (int i = 0; i < 20; i++) wn[i] = Pow(G, (mod - 1) / (1 << i), mod);
7
  void change(int y[], int len)
8
9
10
        for (int i = 1, j = len / 2; i < len - 1; i++)
11
            if (i < j) swap(y[i], y[j]);</pre>
12
            int k = len / 2;
13
            while (j >= k) j -= k, k /= 2;
14
            if (j < k) j += k;
15
16
17
   }
18
   void ntt(int y[], int len, int on)
19
        change(y, len);
20
21
        for (int h = 2, id = 1; h <= len; h <<= 1, id++)
22
            for (int j = 0; j < len; <math>j += h)
23
24
25
                int w = 1;
                for (int k = j; k < j + h / 2; k++)
26
27
                     int u = y[k] \% mod;
28
                     int t = 1LL * w * (y[k + h / 2] % mod) % mod;
29
                     y[k] = (u + t) \% \mod, y[k + h / 2] = ((u - t) \% \mod + \mod) \% \mod;
30
                     w = 1LL * w * wn[id] % mod;
31
32
                }
            }
33
34
        if (on == -1)
35
36
37
            int inv = Pow(len, mod - 2, mod);
38
            for (int i = 1; i < len / 2; i++) swap(y[i], y[len - i]);
39
            for (int i = 0; i < len; i++) y[i] = 1LL * y[i] * inv % mod;
40
41
   }
42
   1.8.3 FWT
   void fwt(int f[], int m)
        int n = __builtin_ctz(m);
3
        for (int i = 0; i < n; ++i)
4
            for (int j = 0; j < m; ++j)
5
6
                if (j & (1 << i))
7
                 {
                     int l = f[j \land (1 << i)], r = f[j];
8
                     f[j \land (1 << i)] = l + r, f[j] = l - r;
9
```

```
// or: f[j] += f[j \land (1 << i)];
10
                      // and: f[j \land (1 << i)] += f[j];
11
12
13
   void ifwt(int f[], int m)
14
   {
15
        int n = __builtin_ctz(m);
16
        for (int i = 0; i < n; ++i)
17
             for (int j = 0; j < m; ++j)
18
                 if (j & (1 << i))
19
20
21
                      int l = f[j \land (1 << i)], r = f[j];
                      f[j \land (1 \lessdot i)] = (l + r) / 2, f[j] = (l - r) / 2;
22
23
                     // or: f[j] -= f[j \land (1 << i)];
24
                     // and: f[j \land (1 << i)] -= f[j];
25
                 }
26
27 }
    1.9 Numerical Integration
   1.9.1 Adaptive Simpson's Rule
    \int_{a}^{b} f(x)dx \approx \frac{b-a}{6} [f(a) + 4f(\frac{a+b}{2}) + f(b)]
    |\ddot{S}(a,c) + S(c,b) - S(a,b)|/15 < \epsilon
1 double F(double x) {}
   double simpson(double a, double b)
2
   { // Simpson
3
        double c = a + (b - a) / 2;
4
        return (F(a) + 4 * F(c) + F(b)) * (b - a) / 6;
5
6
   double asr(double a, double b, double eps, double A)
7
   { // Simpson ( )
                        [a,b] Simpson A
        double c = a + (b - a) / 2;
9
        double L = simpson(a, c), R = simpson(c, b); if (fabs(L + R - A) \ll 15 * eps) return L + R + (L + R - A) / 15.0;
10
11
        return asr(a, c, eps / 2, L) + asr(c, b, eps / 2, R);
12
13
   double asr(double a, double b, double eps) { return asr(a, b, eps, simpson(a, b)); }
   1.9.2 Berlekamp-Massey
   const int N = 1 \ll 14;
   11 res[N], base[N], _c[N], _md[N];
   vector<int> Md;
3
   void mul(ll* a, ll* b, int k)
4
   {
5
        for (int i = 0; i < k + k; i++) _{c}[i] = 0;
6
        for (int i = 0; i < k; i++)
7
8
             if (a[i])
                 for (int j = 0; j < k; j++) _{c[i + j]} = (_{c[i + j]} + a[i] * b[j]) % mod;
9
10
        for (int i = k + k - 1; i >= k; i--)
11
             if (_c[i])
                 for (int j = 0; j < Md.size(); j++) _c[i - k + Md[j]] = (_c[i - k + Md[j]]
12
        - _c[i] * _md[Md[j]]) % mod;
        for (int i = 0; i < k; i++) a[i] = _c[i];
13
14 }
```

```
int solve(ll n, VI a, VI b)
16
   {
17
        ll ans = 0, pnt = 0;
18
        int k = a.size();
        assert(a.size() == b.size());
19
        for (int i = 0; i < k; i++) _md[k - 1 - i] = -a[i];
20
21
        _{md[k]} = 1;
        Md.clear();
22
23
        for (int i = 0; i < k; i++)
            if (_md[i] != 0) Md.push_back(i);
24
25
        for (int i = 0; i < k; i++) res[i] = base[i] = 0;
26
        res[0] = 1;
        while ((1LL << pnt) <= n) pnt++;</pre>
27
        for (int p = pnt; p >= 0; p--)
28
29
            mul(res, res, k);
30
31
            if ((n >> p) & 1)
32
                for (int i = k - 1; i >= 0; i--) res[i + 1] = res[i];
33
                res[0] = 0;
34
                for (int j = 0; j < Md.size(); j++) res[Md[j]] = (res[Md[j]] - res[k] * _md
35
        [Md[j]]) % mod;
36
37
38
        for (int i = 0; i < k; i++) ans = (ans + res[i] * b[i]) % mod;
39
        if (ans < 0) ans += mod;
        return ans;
40
41
   VI BM(VI s)
42
43
   {
        VI C(1, 1), B(1, 1);
44
        int L = 0, m = 1, b = 1;
45
        for (int n = 0; n < s.size(); n++)
46
47
            11 d = 0;
48
            for (int i = 0; i \le L; i++) d = (d + (ll)C[i] * s[n - i]) % mod;
49
            if (d == 0)
50
51
                ++m;
            else if (2 * L <= n)
52
53
                VI T = C;
54
                11 c = mod - d * Pow(b, mod - 2) % mod;
55
                while (C.size() < B.size() + m) C.pb(0);</pre>
56
                for (int i = 0; i < B.size(); i++) C[i + m] = (C[i + m] + c * B[i]) % mod;
57
                L = n + 1 - L, B = T, b = d, m = 1;
58
            }
59
            else
60
            {
61
                11 c = mod - d * Pow(b, mod - 2) % mod;
62
63
                while (C.size() < B.size() + m) C.pb(0);
64
                for (int i = 0; i < B.size(); i++) C[i + m] = (C[i + m] + c * B[i]) % mod;
65
                ++m;
66
            }
67
        return C;
68
69
   int gao(VI a, ll n)
70
71
   {
72
        VI c = BM(a);
```

```
c.erase(c.begin());
73
           for (int i = 0; i < c.size(); i++) c[i] = (mod - c[i]) % mod;
74
           return solve(n, c, VI(a.begin(), a.begin() + c.size()));
75
76 }
     1.10 Others
    n , , m
    int josephus(int n, int m)
 2
 3
           int r = 0;
           for (int k = 1; k \le n; ++k) r = (r + m) \% k;
 4
           return r + 1;
 5
    }
 6
    n^n
    int leftmost(int n)
 1
 2
           double m = n * log10((double)n);
 3
           double g = m - (11)m;
 4
 5
           return (int)pow(10.0, g);
    }
 6
    n!
 1
    int count(ll n)
 2
           if (n == 1) return 1;
 3
           return (int)ceil(0.5 * log10(2 * M_PI * n) + n * log10(n) - n * log10(M_E));
 4
 5
     1.11 Formula
            : n = \prod_{i=1}^{k} p_i^{a_i},
        1.
             (a) f(n) = \prod_{i=1}^{k} (a_i + 1)
             (b) g(n) = \prod_{i=1}^{k} (\sum_{j=0}^{a_i} p_i^j)
        2. n
                     n\varphi(n)/2
        3. gcd(n, i) = 1, gcd(n, n - i) = 1(1 \le i \le n)
            D(n) = (n-1)(D(n-2) + D(n-1)) = \sum_{i=2}^{n} \frac{(-1)^{k} n!}{k!} = \left[\frac{n!}{e} + 0.5\right]
        5. : p \text{ is prime } \Rightarrow (p-1)! \equiv -1 \pmod{p}
        6. : gcd(a, n) = 1 \Rightarrow a^{\varphi(n)} \equiv 1 \pmod{n}
        7. : gcd(n,p) = 1 \Rightarrow a^n \equiv a^{n\%\varphi(p)} \pmod{p}
            : n \qquad \pi(n), \lim_{n \to \infty} \pi(n) = \frac{n}{\ln n}
        9. : x N = log 10(n) + 1
            n! \approx \sqrt{2\pi n} (\frac{n}{2})^n
       10.
       11. a > 1, m, n > 0, gcd(a^m - 1, a^n - 1) = a^{gcd(m,n)} - 1
       12. a > b, gcd(a, b) = 1, gcd(a^m - b^m, a^n - b^n) = a^{gcd(m, n)} - b^{gcd(m, n)}
                                      G = \gcd(C_n^1, C_n^2, ..., C_n^{n-1}) = \begin{cases} n, & n \text{ is prime} \\ 1, & n \text{ has multy prime factors} \\ p, & n \text{ has single prime factor } p \end{cases}
            gcd(Fib(m), Fib(n)) = Fib(gcd(m, n))
```

13. 
$$gcd(m,n) = 1$$
, :

(a) 
$$m*n-m-n$$

(b) 
$$N = \frac{(m-1)(n-1)}{2}$$

14. 
$$(n+1)lcm(C_n^0, C_n^1, ..., C_n^{n-1}, C_n^n) = lcm(1, 2, ..., n+1)$$

15. 
$$p$$
 ,  $(x+y+...+w)^p \equiv x^p + y^p + ... + w^p \pmod{p}$ 

16. :1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, 208012 
$$h(0) = h(1) = 1, h(n) = \frac{(4n-2)h(n-1)}{n+1} = \frac{C_{2n}^n}{n+1} = C_{2n}^n - C_{2n}^{n-1}$$

17. 
$$:B_n = -\frac{1}{n+1} \sum_{i=0}^{n-1} C_{n+1}^i B_i$$

$$\sum_{i=1}^{n} i^{k} = \frac{1}{k+1} \sum_{i=1}^{k+1} C_{k+1}^{i} B_{k+1-i} (n+1)^{i}$$

## 18. FFT

ггі			
$r \ 2^k + 1$	r	k	g
3	1	1	$\frac{g}{2}$
5	1	2	2
17	1	4	3
97	3	5	5
193	3	6	5
257	1	8	3
7681	15	9	17
12289	3	12	11
40961	5	13	3
65537	1	16	3
786433	3	18	10
5767169	11	19	3
7340033	7	20	3
23068673	11	21	3
104857601	25	22	3
167772161	5	25	3
469762049	7	26	3
998244353	119	23	3
1004535809	479	21	3
2013265921	15	27	31
2281701377	17	27	3
3221225473	3	30	5
75161927681	35	31	3
77309411329	9	33	7
206158430209	3	36	22
2061584302081	15	37	7
2748779069441	5	39	3
6597069766657	3	41	5
39582418599937	9	42	5
79164837199873	9	43	5
263882790666241	15	44	7
1231453023109121	35	45	3
1337006139375617	19	46	3
3799912185593857	27	47	5
4222124650659841	15	48	19
7881299347898369	7	50	6
31525197391593473	7	52	3
180143985094819841	5	55	6
1945555039024054273	27	56	5
4179340454199820289	29	57	3

# 2 String Processing

## 2.1 KMP

```
1 // yx
  const int N = "Edit";
  int next[N];
3
   void initkmp(char x[], int m)
4
5
        int i = 0, j = next[0] = -1;
6
7
        while (i < m)
8
        {
            while (j != -1 \&\& x[i] != x[j]) j = next[j];
9
10
            next[++i] = ++j;
        }
11
12
int kmp(char x[], int m, char y[], int n)
   {
14
        int i, j, ans;
15
        i = j = ans = 0;
16
        initkmp(x, m);
17
       while (i < n)
18
19
            while (j != -1 \&\& y[i] != x[j]) j = next[j];
20
            i++, j++;
if (j >= m) ans++, j = next[j];
21
22
23
24
        return ans;
  }
25
```

## 2.2 ExtendKMP

```
1 //next[i]:x[i...m-1] x[0...m-1]
^2 //extend[i]:y[i...n-1] x[0...m-1]
3 const int N = "Edit"
4 int next[N], extend[N];
5 void pre_ekmp(char x[], int m)
6
7
       next[0] = m;
8
       int j = 0;
       while (j + 1 < m \&\& x[j] == x[j + 1]) j++;
9
       next[1] = j;
10
11
       int k = 1
       for (int i = 2; i < m; i++)
12
13
            int p = next[k] + k - 1;
14
            int L = next[i - k];
15
            if (i + L 
16
                next[i] = L;
17
            else
18
19
            {
                j = max(0, p - i + 1);
20
21
                while (i + j < m \&\& x[i + j] == x[j]) j++;
22
                next[i] = j;
                k = i;
23
           }
24
       }
25
26 }
```

```
void ekmp(char x[], int m, char y[], int n)
27
28
   {
       pre_ekmp(x, m, next);
29
30
       int j = 0;
       while (j < n \&\& j < m \&\& x[j] == y[j]) j++;
31
       extend[0] = j;
32
       int k = 0;
33
       for (int i = 1; i < n; i++)
34
35
            int p = extend[k] + k - 1;
36
37
            int L = next[i - k];
38
            if (i + L 
                extend[i] = L;
39
            else
40
            {
41
                j = max(0, p - i + 1);
42
                while (i + j < n \&\& j < m \&\& y[i + j] == x[j]) j++;
43
                extend[i] = j, k = i;
44
            }
45
       }
46
  }
47
   2.3 Manacher
   O(n)
1 const int N = "Edit";
2 char s[N], str[N << 1];</pre>
3 int p[N << 1];</pre>
   void Manacher(char s□, int& n)
4
5
       str[0] = '$', str[1] = '#';
6
       for (int i = 0; i < n; i++) str[(i << 1) + 2] = s[i], <math>str[(i << 1) + 3] = '\#';
7
       n = 2 * n + 2;
8
9
       str[n] = 0;
       int mx = 0, id;
10
       for (int i = 1; i < n; i++)
11
12
            p[i] = mx > i ? min(p[2 * id - i], mx - i) : 1;
13
            while (str[i - p[i]] == str[i + p[i]]) p[i]++;
14
15
            if (p[i] + i > mx) mx = p[i] + i, id = i;
       }
16
17
   }
  int solve(char s[])
18
19
   {
       int n = strlen(s);
20
21
       Manacher(s, n);
22
       return *max_elememt(p, p + n) - 1;
23
  }
   2.4 Aho-Corasick Automaton
1 const int maxn = "Edit";
   struct Trie
^{2}
3
       int ch[maxn][26], f[maxn], val[maxn];
```

int sz, rt;

5

```
int newnode() { clr(ch[sz], -1), val[sz] = 0; return sz++; }
6
        void init() { sz = 0, rt = newnode(); }
7
        inline int idx(char c) { return c - 'A'; };
8
        void insert(const char* s)
9
10
            int u = 0, n = strlen(s);
11
            for (int i = 0; i < n; i++)
12
            {
13
                int c = idx(s[i]);
14
                if (ch[u][c] == -1) ch[u][c] = newnode();
15
16
                u = ch[u][c];
17
            }
            val[u]++;
18
        }
19
        void build()
20
21
            queue<int> q;
22
23
            f[rt] = rt;
            for (int c = 0; c < 26; c++)
24
25
                if (~ch[rt][c])
26
27
                     f[ch[rt][c]] = rt, q.push(ch[rt][c]);
                else
28
29
                     ch[rt][c] = rt;
30
            while (!q.empty())
31
32
                int u = q.front();
33
                q.pop();
34
                // val[u] |= val[f[u]];
35
                for (int c = 0; c < 26; c++)
36
37
                     if (~ch[u][c])
38
                         f[ch[u][c]] = ch[f[u]][c], q.push(ch[u][c]);
39
                     else
40
                         ch[u][c] = ch[f[u]][c];
41
42
                }
43
            }
        }
44
45
        int query(const char* s)
46
47
            int u = rt, n = strlen(s);
48
49
            int res = 0;
            for (int i = 0; i < n; i++)
50
51
52
                int c = idx(s[i]);
                u = ch[u][c];
53
                int tmp = u;
54
55
                while (tmp != rt)
56
57
                     res += val[tmp];
                     val[tmp] = 0;
58
                     tmp = f[tmp];
59
60
61
62
            return res;
63
        }
64 };
```

## 2.5 Suffix Array

```
, O(nlogn)
   const int maxn = "Edit";
2
   char s[maxn];
4 int sa[maxn], t[maxn], t2[maxn], c[maxn], rank[maxn], height[maxn];
              0~m-1
5
  void build_sa(int m, int n)
6
   {
7
8
       int *x = t, *y = t2;
9
10
       for (int i = 0; i < m; i++) c[i] = 0;
11
12
       for (int i = 0; i < n; i++) c[x[i] = s[i]]++;
13
       for (int i = 1; i < m; i++) c[i] += c[i - 1];
       for (int i = n - 1; \sim i; i--) sa[--c[x[i]]] = i;
14
       for (int k = 1; k <= n; k <<= 1)
15
16
            // sa
17
            int p = 0;
18
            for (int i = n - k; i < n; i++) y[p++] = i;
19
            for (int i = 0; i < n; i++)
20
                if (sa[i] >= k) y[p++] = sa[i] - k;
21
22
            for (int i = 0; i < m; i++) c[i] = 0;
23
            for (int i = 0; i < n; i++) c[x[y[i]]]++;
24
25
            for (int i = 0; i < m; i++) c[i] += c[i - 1];
26
            for (int i = n - 1; \sim i; i--) sa[--c[x[y[i]]]] = y[i];
            // say
27
28
            swap(x, y);
29
            p = 1;
            x[sa[0]] = 0;
30
            for (int i = 1; i < n; i++)
31
                x[sa[i]] = y[sa[i - 1]] == y[sa[i]] && y[sa[i - 1] + k] == y[sa[i] + k] ? p
32
        -1:p++;
33
            if (p >= n) break; //
                                      ,sa
34
           m = p;
35
       }
       n--;
36
37
       int k = 0;
38
       for (int i = 0; i <= n; i++) rank[sa[i]] = i;</pre>
       for (int i = 0; i < n; i++)
39
40
            if (k) k--;
41
            int j = sa[rank[i] - 1];
42
43
            while (s[i + k] == s[j + k]) k++;
            height[rank[i]] = k;
44
       }
45
   }
46
47
   int dp[maxn][30];
   void initrmq(int n)
50
   {
        for (int i = 1; i <= n; i++)
51
            dp[i][0] = height[i];
52
       for (int j = 1; (1 << j) <= n; j++)
53
            for (int i = 1; i + (1 << j) - 1 <= n; i++)
54
                dp[i][j] = min(dp[i][j - 1], dp[i + (1 << (j - 1))][j - 1]);
55
56 }
```

```
57 int rmq(int l, int r)
58
   {
        int k = 31 - \_builtin\_clz(r - l + 1);
59
        return min(dp[l][k], dp[r - (1 << k) + 1][k]);
60
61
62 int lcp(int a, int b)
   { //
63
        a = rank[a], b = rank[b];
64
        if (a > b) swap(a, b);
65
66
        return rmq(a + 1, b);
67
  }
   2.6 Suffix Automation
1 const int maxn = "Edit";
   struct SAM
3
   {
        int len[maxn << 1], link[maxn << 1], ch[maxn << 1][26];</pre>
4
        int sz, rt, last;
5
        int newnode(int x = 0)
6
7
            len[sz] = x;
8
            link[sz] = -1;
9
            clr(ch[sz], -1);
10
            return sz++;
11
12
        void init() { sz = last = 0, rt = newnode(); }
13
        void extend(int c)
14
15
            int np = newnode(len[last] + 1);
17
            for (p = last; \sim p \& ch[p][c] == -1; p = link[p]) ch[p][c] = np;
18
            if (p == -1)
19
                link[np] = rt;
20
            else
21
            {
22
                int q = ch[p][c];
23
24
                if (len[p] + 1 == len[q])
25
                     link[np] = q;
                else
26
27
                {
                     int nq = newnode(len[p] + 1);
28
                    memcpy(ch[nq], ch[q], sizeof(ch[q]));
29
30
                    link[nq] = link[q], link[q] = link[np] = nq;
                     for (; \sim p \&\& ch[p][c] == q; p = link[p]) ch[p][c] = nq;
31
                }
32
            }
33
            last = np;
34
35
        int topcnt[maxn], topsam[maxn << 1];</pre>
36
37
        void sort()
38
        { //
            clr(topcnt, 0);
39
            for (int i = 0; i < sz; i++) topcnt[len[i]]++;</pre>
40
            for (int i = 0; i < maxn - 1; i++) topcnt[i + 1] += topcnt[i];
41
            for (int i = 0; i < sz; i++) topsam[--topcnt[len[i]]] = i;
42
43
        }
44 };
```

# 3 Data Structure

# 3.1 Binary Indexed Tree

```
O(\log n)
  //
   const int maxn = "Edit";
  int bit[N], n;
4 int sum(int x)
5 {
6
        int s = 0;
        for (int i = x; i; i -= i & -i) s += bit[i];
7
8
        return s;
9 }
10 void add(int x, int v)
11 {
        for (int i = x; i <= n; i += i & -i) bit[i] += v;
12
13 }
   3.2 Segment Tree
1 #define lson rt << 1</pre>
                                 //
2 #define rson rt << 1 | 1
                                 //
3 #define Lson l, m, lson
                                 //
4 #define Rson m + 1, r, rson //
5 void PushUp(int rt);
                               // lson rson rt
  void PushDown(int rt[, int m]);
                                                      // rt ,m ( )
                                                      // rt , [l, r]
// rt[l, r]
  void build(int l, int r, int rt);
void update([...,] int l, int r, int rt)
9 int query(int L, int R, int l, int r, int rt)
                                                      // rt[l, r] [L, R]
   3.2.1 Single-point Update
1 const int maxn = "Edit";
2 int sum[maxn << 2]; // sum[rt]</pre>
3 void PushUp(int rt) { sum[rt] = sum[lson] + sum[rson]; }
4 void build(int l, int r, int rt)
5
        if (l == r)
6
7
            scanf("%d", &sum[rt]); //
8
9
            return;
10
        int m = (l + r) >> 1;
11
        build(Lson);
12
        build(Rson);
13
        PushUp(rt);
14
  void update(int p, int add, int l, int r, int rt)
17
18
        if (l == r)
19
            sum[rt] += add;
20
21
            return;
22
23
        int m = (l + r) >> 1;
```

```
if (p \ll m)
24
25
            update(p, add, Lson);
26
        else
            update(p, add, Rson);
27
28
        PushUp(rt);
29
  int query(int L, int R, int l, int r, int rt)
30
31
        if (L <= 1 && r <= R) return sum[rt];</pre>
32
        int m = (l + r) >> 1, s = 0;
33
34
        if (L \le m) s += query(L, R, Lson);
35
        if (m < R) s += query(L, R, Rson);
36
        return s;
37 }
   3.2.2 Interval Update
   const int maxn = "Edit";
   int seg[maxn << 2], sum[maxn << 2]; // seg[rt] , PushDown</pre>
   void PushUp(int rt) { sum[rt] = sum[lson] + sum[rson]; }
   void PushDown(int rt, int m)
4
5
   {
6
        if (seg[rt] == 0) return;
7
        seg[lson] += seg[rt];
        seg[rson] += seg[rt];
sum[lson] += seg[rt] * (m - (m >> 1));
8
9
10
        sum[rson] += seg[rt] * (m >> 1);
11
        seg[rt] = 0;
12
   void build(int 1, int r, int rt)
13
14
15
        seg[rt] = 0;
16
        if (l == r)
17
        {
            scanf("%lld", &sum[rt]);
18
19
            return;
20
21
        int m = (l + r) >> 1;
22
        build(Lson);
23
        build(Rson);
24
        PushUp(rt);
25
   void update(int L, int R, int add, int l, int r, int rt)
27
   {
28
        if (L <= 1 && r <= R)
29
30
            seg[rt] += add;
            sum[rt] += add * (r - l + 1);
31
            return;
32
33
        PushDown(rt, r - l + 1);
34
        int m = (l + r) >> 1;
35
36
        if (L <= m) update(L, R, add, Lson);</pre>
        if (m < R) update(L, R, add, Rson);</pre>
37
        PushUp(rt);
38
39
   int query(int L, int R, int l, int r, int rt)
40
41 {
```

```
if (L <= 1 && r <= R) return sum[rt];</pre>
42
        PushDown(rt, r - l + 1);
43
        int m = (l + r) >> 1, ret = 0;
44
        if (L <= m) ret += query(L, R, Lson);</pre>
45
46
        if (m < R) ret += query(L, R, Rson);</pre>
        return ret;
47
48 }
   3.3 Splay Tree
   #define key_value ch[ch[root][1]][0]
   const int maxn = "Edit";
3
   struct Splay
4
   {
        int a[maxn];
5
        int sz[maxn], ch[maxn][2], fa[maxn];
int key[maxn], rev[maxn];
6
7
8
        int root, tot;
9
        int stk[maxn], top;
10
        void init(int n)
        {
11
            tot = 0, top = 0;
12
            root = newnode(0, -1);
13
            ch[root][1] = newnode(root, -1);
14
             for (int i = 0; i < n; i++) a[i] = i + 1;
15
16
            key_value = build(0, n - 1, ch[root][1]);
17
            pushup(ch[root][1]);
18
            pushup(root);
19
        int newnode(int p = 0, int k = 0)
20
21
22
            int x = top ? stk[top--] : ++tot;
23
            fa[x] = p;
24
            sz[x] = 1;
            ch[x][0] = ch[x][1] = 0;
25
            key[x] = k;
26
27
            rev[x] = 0;
28
            return x;
29
30
        void pushdown(int x)
31
32
            if (rev[x])
33
            {
                 swap(ch[x][0], ch[x][1]);
34
                 if (ch[x][0]) rev[ch[x][0]] ^= 1;
35
                 if (ch[x][1]) rev[ch[x][1]] ^= 1;
36
                 rev[x] = 0;
37
38
            }
39
        }
        void pushup(int x) { sz[x] = sz[ch[x][0]] + sz[ch[x][1]] + 1; }
40
41
        void rotate(int x, int d)
42
             int y = fa[x];
43
            pushdown(y), pushdown(x);
ch[y][d ^ 1] = ch[x][d];
44
45
            fa[ch[x][d]] = y;
46
            if (fa[y]) ch[fa[y]][ch[fa[y]][1] == y] = x;
47
            fa[x] = fa[y];
48
```

```
49
            ch[x][d] = y;
50
            fa[y] = x;
51
            pushup(y);
52
        void splay(int x, int goal = 0)
53
54
            pushdown(x);
55
            while (fa[x] != goal)
56
57
                 if (fa[fa[x]] == goal)
58
                     rotate(x, ch[fa[x]][0] == x);
59
60
                else
                 {
61
                     int y = fa[x];
62
                     int d = ch[fa[y]][0] == y;
63
                     ch[y][d] == x ? rotate(x, d \land 1) : rotate(y, d);
64
65
                     rotate(x, d);
                }
66
67
            }
            pushup(x);
68
            if (goal == 0) root = x;
69
70
        int kth(int r, int k)
71
72
73
            pushdown(r);
            int t = sz[ch[r][0]] + 1;
74
            if (t == k) return r;
75
            return t > k ? kth(ch[r][0], k) : kth(ch[r][1], k - t);
76
77
        int build(int 1, int r, int p)
78
79
            if (l > r) return 0;
80
            int mid = l + r \gg 1;
81
            int x = newnode(p, a[mid]);
82
            ch[x][0] = build(l, mid - 1, x);
83
            ch[x][1] = build(mid + 1, r, x);
84
85
            pushup(x);
86
            return x;
87
        }
        void select(int l, int r)
88
89
            splay(kth(root, 1), 0);
90
            splay(kth(ch[root][1], r - l + 2), root);
91
92
93
   };
94
```

## 3.4 Functional Segment Tree

k

```
const int maxn = "Edit";
int a[maxn], rt[maxn];
int cnt;
int lson[maxn << 5], rson[maxn << 5], sum[maxn << 5];
#define Lson l, m, lson[x], lson[y]
#define Rson m + 1, r, rson[x], rson[y]</pre>
```

```
void update(int p, int l, int r, int& x, int y)
8
   {
       lson[++cnt] = lson[y], rson[cnt] = rson[y], sum[cnt] = sum[y] + 1, x = cnt;
9
10
       if (l == r) return;
       int m = (l + r) >> 1;
11
       if (p <= m) update(p, Lson);</pre>
12
       else update(p, Rson);
13
14
   }
  int query(int 1, int r, int x, int y, int k)
15
16
       if (l == r) return l;
17
18
       int m = (l + r) >> 1;
       int s = sum[lson[y]] - sum[lson[x]];
19
       if (s >= k) return query(Lson, k);
20
       else return query(Rson, k - s);
21
22
   3.5 Sparse Table
   const int maxn = "Edit";
   int mmax[maxn][30], mmin[maxn][30];
   int a[maxn], n, k;
   void init()
   {
5
        for (int i = 1; i \le n; i++) mmax[i][0] = mmin[i][0] = a[i];
6
       for (int j = 1; (1 << j) <= n; j++)
7
            for (int i = 1; i + (1 << j) - 1 <= n; i++)
8
9
                mmax[i][j] = max(mmax[i][j - 1], mmax[i + (1 << (j - 1))][j - 1]);
10
                mmin[i][j] = min(mmin[i][j - 1], mmin[i + (1 << (j - 1))][j - 1]);
11
            }
12
   }
13
   // op=0/1 [l,r] /
14
   int rmq(int 1, int r, int op)
15
16
        int k = 31 - \_builtin\_clz(r - l + 1);
17
       if (op == 0)
18
            return max(mmax[l][k], mmax[r - (1 << k) + 1][k]);
19
        return min(mmin[l][k], mmin[r - (1 << k) + 1][k]);</pre>
20
21 }
     RMQ
   void init()
2
   {
3
        for (int i = 0; (1 << i) <= n; i++)
            for (int j = 0; (1 << j) <= m; j++)
4
5
                if (i == 0 \&\& j == 0) continue;
6
                for (int row = 1; row + (1 << i) - 1 <= n; row++)
7
                    for (int col = 1; col + (1 << j) - 1 <= m; col++)
8
                        if (i)
9
                            dp[row][col][i][j] = max(dp[row][col][i - 1][j],
10
                                                 dp[row + (1 << (i - 1))][col][i - 1][j]);
11
                        else
12
                            dp[row][col][i][j] = max(dp[row][col][i][j - 1],
13
                                                  dp[row][col + (1 << (j - 1))][i][j - 1]);
14
15
           }
16
  }
```

```
int rmq(int x1, int y1, int x2, int y2)
18
   {
       int kx = 31 - \_builtin\_clz(x2 - x1 + 1);
19
       int ky = 31 - \_builtin_clz(y2 - y1 + 1);
20
       int m1 = dp[x1][y1][kx][ky];
21
       int m2 = dp[x2 - (1 << kx) + 1][y1][kx][ky];
22
       int m3 = dp[x1][y2 - (1 << ky) + 1][kx][ky];
23
       int m4 = dp[x2 - (1 << kx) + 1][y2 - (1 << ky) + 1][kx][ky];
24
       return max(max(m1, m2), max(m3, m4));
25
26
  }
   3.6 Heavy-Light Decomposition
   const int maxn = "Edit";
1
   struct HLD
2
3
   {
       int n, dfs_clock;
4
       int sz[maxn], top[maxn], son[maxn], dep[maxn], fa[maxn], id[maxn];
5
       vector<int> G[maxn];
6
       void init(int n)
7
       {
8
            this->n = n, clr(son, -1), dfs_clock = 0;
9
            for (int i = 0; i < n; i++) G[i].clear();
10
11
       void add_edge(int u, int v) { G[u].pb(v), G[v].pb(u); }
12
13
       void dfs(int u, int p, int d)
14
            dep[u] = d, fa[u] = p, sz[u] = 1;
15
16
            for (auto& v : G[u])
17
            {
                if (v == p) continue;
18
                dfs(v, u, d + 1);
19
                sz[u] += sz[v];
20
                if (son[u] == -1 \mid | sz[v] > sz[son[u]]) son[u] = v;
21
            }
22
23
       void link(int u, int t)
24
25
26
            top[u] = t, id[u] = ++dfs\_clock;
            if (son[u] == -1) return;
27
            link(son[u], t);
28
            for (auto& v : G[u])
29
                if (v != son[u] \&\& v != fa[u]) link(v, v);
30
       }
31
       //
32
       int query_path(int u, int v)
33
34
            int ret = 0;
35
            while (top[u] != top[v])
36
37
                if (dep[top[u]] < dep[top[v]]) swap(u, v);</pre>
38
39
                ret += query(id[top[u]], id[u]);
                u = fa[top[u]];
40
41
            if (dep[u] > dep[v]) swap(u, v);
42
43
            ret += query(id[u], id[v]);
44
       }
45 };
```

#### 3.7 Link-Cut Tree

```
const int maxn = "Edit";
   struct LCT
3
   {
        int val[maxn], sum[maxn]; //
4
       int rev[maxn], ch[maxn][2], fa[maxn];
5
       int stk[maxn];
6
       inline void init(int n)
7
        { //
8
            for (int i = 1; i <= n; i++) scanf("%d", val + i);</pre>
9
10
       inline bool isroot(int x) { return ch[fa[x]][0] != x && ch[fa[x]][1] != x; }
11
       inline bool get(int x) { return ch[fa[x]][1] == x; }
12
       void pushdown(int x)
13
14
        {
15
            if (!rev[x]) return;
            swap(ch[x][0], ch[x][1]);
16
            if (ch[x][0]) rev[ch[x][0]] ^= 1;
17
            if (ch[x][1]) rev[ch[x][1]] ^= 1;
18
            rev[x] \sim 1;
19
20
       void pushup(int x) { sum[x] = val[x] + sum[ch[x][0]] + sum[ch[x][1]]; }
21
       void rotate(int x)
22
23
        {
            int y = fa[x], z = fa[fa[x]], d = get(x);
24
25
            if (!isroot(y)) ch[z][get(y)] = x;
            fa[x] = z;
26
            ch[y][d] = ch[x][d \wedge 1], fa[ch[y][d]] = y;
27
28
            ch[x][d \land 1] = y, fa[y] = x;
29
            pushup(y), pushup(x);
30
       void splay(int x)
31
32
            int top = 0;
33
            stk[++top] = x;
34
            for (int i = x; !isroot(i); i = fa[i]) stk[++top] = fa[i];
35
            for (int i = top; i; i--) pushdown(stk[i]);
36
            for (int f; !isroot(x); rotate(x))
37
                if (!isroot(f = fa[x])) rotate(get(x) == get(f) ? f : x);
38
39
       void access(int x)
40
41
        {
            for (int y = 0; x; y = x, x = fa[x]) splay(x), ch[x][1] = y, pushup(x);
42
43
       int find(int x) { access(x), splay(x); while (ch[x][0]) x = ch[x][0]; return x; }
44
       void makeroot(int x) { access(x), splay(x), rev[x] ^= 1; }
45
       void link(int x, int y) { makeroot(x), fa[x] = y, splay(x); }
46
       void cut(int x, int y) { makeroot(x), access(y), splay(y), fa[x] = ch[y][0] = 0; }
47
       void update(int x, int v) { val[x] = v, access(x), splay(x); }
48
        int query(int x, int y) { makeroot(y), access(x), splay(x); return sum[x]; }
49
50
  };
```

# Graph Theory

#### 4.1 Union-Find Set

1 // //

3 const int maxn = "Edit";

```
1 const int maxn = "Edit";
  int n, fa[maxn], ra[maxn];
  void init()
   {
4
       clr(ra, 0);
5
       iota(fa, fa + n, 0);
6
   }
7
   int find(int x) { return fa[x] != x ? fa[x] = find(fa[x]) : x; }
   void unite(int x, int y)
9
10
       x = find(x), y = find(y);
11
12
       if (x == y) return;
13
       if (ra[x] < ra[y])
            fa[x] = y;
14
15
       else
16
       {
17
            fa[y] = x;
            if (ra[x] == ra[y]) ra[x]++;
18
       }
19
20
   bool same(int x, int y) { return find(x) == find(y); }
   4.2
        Minimal Spanning Tree
   4.2.1 Kruskal
1 typedef pair<int, PII> Edge;
   vector<Edge> G;
   void add_edge(int u, int v, int d) { G.pb(mp(d, mp(u, v))); }
4 int Kruskal(int n)
5
       init(n); //
6
       sort(G.begin(), G.end());
7
8
       int num = 0, ret = 0;
9
       for (auto& e : G)
10
            int x = e.Y.X, y = e.Y.Y;
11
            int d = e.X;
12
            if (!same(x, y))
13
15
                unite(x, y);
16
                num++;
17
                ret += d;
18
            if (num == n - 1) break;
19
20
21
       return ret;
22
  }
   4.2.2 Prim
        cost[][], 0,0~n-1
```

```
bool vis[maxn];
   int lowc[maxn];
   int Prim(int cost[][maxn], int n)
6
7
   {
8
       int ans = 0;
       clr(vis, 0);
9
       vis[0] = 1;
10
       for (int i = 1; i < n; i++)
11
            lowc[i] = cost[0][i];
12
       for (int i = 1; i < n; i++)
13
14
15
            int minc = INF;
16
            int p = -1;
            for (int j = 0; j < n; j++)
17
                if (!vis[j] && minc > lowc[j])
18
19
                    minc = lowc[j];
20
                    p = j;
21
22
            if (minc == INF) return -1;
23
            vis[p] = 1;
24
25
            ans += minc;
            for (int j = 0; j < n; j++)
26
27
                if (!vis[j] && lowc[j] > cost[p][j])
28
                    lowc[j] = cost[p][j];
29
30
       return ans;
  }
31
   4.3 Shortest Path
   struct Edge
1
2
       int from, to, dist;
3
       Edge() {}
4
       Edge(int u, int v, int d) : from(u), to(v), dist(d) {}
5
6
  };
   4.3.1 Dijkstra
   struct HeapNode
2
   {
3
       bool operator<(const HeapNode& rhs) const
4
5
        {
            return d > rhs.d;
6
7
   };
8
   const int maxn = "Edit";
   struct Dijkstra
   {
11
12
       int n, m;
                              //
       vector<Edge> edges;
                             //
13
                                      (0)
       vector<int> G[maxn]; //
14
       bool done[maxn];
                              //
15
       int d[maxn];
                              // s
16
                              //
17
       int p[maxn];
```

```
void init(int n)
18
19
            this->n = n;
20
            for (int i = 0; i < n; i++) G[i].clear(); //</pre>
21
22
            edges.clear();
23
        void AddEdge(int from, int to, int dist)
24
                      AddEdge
25
        { //
26
            edges.emplace_back(from, to, dist);
            m = edges.size();
27
28
            G[from].push_back(m - 1);
29
        }
        void dijkstra(int s)
30
31
            priority_queue<HeapNode> q;
32
            for (int i = 0; i < n; i++) d[i] = INF;</pre>
33
            d[s] = 0;
34
            memset(done, 0, sizeof(done));
35
            q.push({0, s});
36
            while (!q.empty())
37
            {
38
                 HeapNode x = q.top();
39
                 q.pop();
40
                 int u = x.u;
41
42
                 if (done[u]) continue;
                 done[u] = true;
43
                 for (auto& id : G[u])
44
45
                     Edge& e = edges[id];
46
                     if (d[e.to] > d[u] + e.dist)
47
48
                     {
                         d[e.to] = d[u] + e.dist;
49
50
                         p[e.to] = id;
                         q.push({d[e.to], e.to});
51
                     }
52
                 }
53
54
            }
55
        }
   };
56
   4.3.2 Bellman-Ford
   const int maxn = "Edit";
   struct BellmanFord
2
3
   {
        int n, m;
4
        vector<Edge> edges;
5
        vector<int> G[maxn];
6
        bool inq[maxn]; //
7
        int d[maxn];
                         // s
8
        int p[maxn];
                         //
9
        int cnt[maxn];
10
        void init(int n)
11
12
13
            this -> n = n;
            for (int i = 0; i < n; i++) G[i].clear();</pre>
14
15
            edges.clear();
        }
16
```

```
void AddEdge(int from, int to, int dist)
17
18
            edges.emplace_back(from, to, dist);
19
            m = edges.size();
20
            G[from].push_back(m - 1);
21
22
        bool bellmanford(int s)
23
24
25
            queue<int> q;
            memset(inq, 0, sizeof(inq));
26
27
            memset(cnt, 0, sizeof(cnt));
28
            for (int i = 0; i < n; i++) d[i] = INF;
            d[s] = 0;
29
            inq[s] = true;
30
            q.push(s);
31
            while (!q.empty())
32
33
34
                int u = q.front();
35
                q.pop();
                inq[u] = false;
36
                for (auto& id : G[u])
37
38
                     Edge& e = edges[id];
39
                     if (d[u] < INF \&\& d[e.to] > d[u] + e.dist)
40
41
                     {
                         d[e.to] = d[u] + e.dist;
42
                         p[e.to] = id;
43
                         if (!inq[e.to])
44
45
                              q.push(e.to);
46
47
                              inq[e.to] = true;
                              if (++cnt[e.to] > n) return false;
48
49
                         }
                     }
50
                }
51
52
            }
53
            return true;
54
        }
55
   };
```

### 4.4 Topo Sort

```
G, deg, map
   Ans
1 const int maxn = "Edit";
2 int Ans[maxn];
3 vector<int> G[maxn];
4 int deg[maxn];
5 map<PII, bool> S;
   void init(int n)
6
7
8
       S.clear();
       for (int i = 0; i < n; i++) G[i].clear();</pre>
9
       clr(deg, 0), clr(Ans, 0);
10
11
   void add_edge(int u, int v)
```

```
13
   {
        if (S[mp(u, v)]) return;
14
        G[u].pb(v), S[mp(u, v)] = 1, deg[v]++;
15
16
   bool Toposort(int n)
17
18
   {
        int tot = 0;
19
        queue<int> q;
20
        for (int i = 0; i < n; ++i)
21
            if (deg[i] == 0) q.push(i);
22
23
        while (!q.empty())
24
        {
            int u = q.front();
25
            que.pop();
26
            Ans[tot++] = u;
27
            for (auto& v : G[u])
28
                if (--deg[v] == 0) q.push(t);
29
30
        if (tot < n - 1) return false;
31
32
        return true;
  }
33
   4.5 LCA
   4.5.1 Tarjan
   Tarjan
      O(n+q)
 1 const int maxn = "Edit";
   int par[maxn];
3 int ans[maxn];
                              //
4 vector<int> G[maxn];
                              //
5 vector<PII> query[maxn]; //
6 bool vis[maxn];
                              //
7 inline void init(int n)
   {
8
        for (int i = 1; i <= n; i++)
9
10
11
            G[i].clear(), query[i].clear();
12
            par[i] = i, vis[i] = 0;
        }
13
14
   inline void add_edge(int u, int v) { G[u].pb(v); }
   inline void add_query(int id, int u, int v)
17
   {
        query[u].pb(mp(v, id));
18
19
        query[v].pb(mp(u, id));
20
   void tarjan(int u)
21
22
   {
        vis[u] = 1;
23
        for (auto& v : G[u])
24
25
            if (vis[v]) continue;
26
27
            tarjan(v);
            unite(u, v);
28
29
        for (auto& q : query[u])
30
```

```
{
31
            int &v = q.X, &id = q.Y;
32
            if (!vis[v]) continue;
33
34
            ans[id] = find(v);
35
   }
36
   4.5.2 DFS+ST
   DFS+ST
      O(nlogn + q)
   const int maxn = "Edit";
  vector<int> G[maxn], sp;
3 int dep[maxn], dfn[maxn];
4 PII dp[21][maxn << 1];
5
   void init(int n)
6
7
       for (int i = 0; i < n; i++) G[i].clear();</pre>
8
       sp.clear();
   }
9
10 void dfs(int u, int fa)
11
   {
12
       dep[u] = dep[fa] + 1;
       dfn[u] = sp.size();
13
       sp.push_back(u);
14
       for (auto& v : G[u])
16
            if (v == fa) continue;
17
            dfs(v, u);
18
            sp.push_back(u);
19
20
21
   void initrmq()
22
23
   {
       int n = sp.size();
24
       for (int i = 0; i < n; i++) dp[0][i] = {dfn[sp[i]], sp[i]};
25
       for (int i = 1; (1 << i) <= n; i++)
26
27
            for (int j = 0; j + (1 << i) - 1 < n; j++)
                dp[i][j] = min(dp[i - 1][j], dp[i - 1][j + (1 << (i - 1))]);
28
  int lca(int u, int v)
30
31
       int l = dfn[u], r = dfn[v];
32
       if (l > r) swap(l, r);
33
       int k = 31 - \_builtin\_clz(r - l + 1);
34
       return min(dp[k][l], dp[k][r - (1 << k) + 1]).Y;
35
36
  }
       Depth-First Traversal
   4.6.1 Biconnected-Component
1 // bccno
2 const int maxn = "Edit";
3 int pre[maxn], iscut[maxn], bccno[maxn], dfs_clock, bcc_cnt;
4 vector<int> G[maxn], bcc[maxn];
5 stack<PII> s;
```

```
void init(int n)
7
   {
        for (int i = 0; i < n; i++) G[i].clear();</pre>
8
   }
9
   inline void add_edge(int u, int v) { G[u].pb(v), G[v].pb(u); }
10
   int dfs(int u, int fa)
11
12
   {
        int lowu = pre[u] = ++dfs_clock;
13
        int child = 0;
14
        for (auto& v : G[u])
15
16
17
            PII e = mp(u, v);
18
            if (!pre[v])
            {
19
                // v
20
                s.push(e);
21
22
                child++;
23
                int lowv = dfs(v, u);
                lowu = min(lowu, lowv); // low
24
25
                if (lowv >= pre[u])
                {
26
                     iscut[u] = true;
27
                     bcc_cnt++;
28
29
                     bcc[bcc_cnt].clear(); // !bcc 1
30
                     for (;;)
31
                         PII x = s.top();
32
33
                         s.pop();
                         if (bccno[x.X] != bcc_cnt)
34
                             bcc[bcc\_cnt].pb(x.X), bcc[x.X] = bcc\_cnt;
35
                         if (bccno[x.Y] != bcc_cnt)
36
37
                             bcc[bcc\_cnt].pb(x.Y), bcc[x.Y] = bcc\_cnt;
                         if (x.X == u \&\& x.Y == v) break;
38
                     }
39
                }
40
41
42
            else if (pre[v] < pre[u] && v != fa)</pre>
43
                s.push(e);
44
                lowu = min(lowu, pre[v]); //
45
            }
46
47
        if (fa < 0 && child == 1) iscut[u] = 0;</pre>
48
49
        return lowu;
50 }
   void find_bcc(int n)
51
52
   {
53
        clr(pre, 0), clr(iscut, 0), clr(bccno, 0);
54
        dfs_clock = bcc_cnt = 0;
56
        for (int i = 0; i < n; i++)
57
            if (!pre[i]) dfs(i, -1);
58 }
   4.6.2 Strongly Connected Component
1 const int maxn = "Edit";
2 vector<int> G[maxn];
```

```
int pre[maxn], lowlink[maxn], sccno[maxn], dfs_clock, scc_cnt;
   stack<int> S;
   inline void init(int n)
5
6
   {
        for (int i = 0; i < n; i++) G[i].clear();</pre>
7
8
   inline void add_edge(int u, int v) { G[u].pb(v); }
9
   void dfs(int u)
10
   {
11
        pre[u] = lowlink[u] = ++dfs_clock;
12
13
        S.push(u);
14
        for (auto& v : G[u])
15
            if (!pre[v])
16
17
                 dfs(v);
18
                 lowlink[u] = min(lowlink[u], lowlink[v]);
19
20
            else if (!sccno[v])
21
                 lowlink[u] = min(lowlink[u], pre[v]);
22
23
24
        if (lowlink[u] == pre[u])
25
26
            scc_cnt++;
27
            for (;;)
28
                 int x = S.top();
29
                 S.pop();
30
                 sccno[x] = scc_cnt;
31
32
                 if (x == u) break;
33
            }
34
        }
35
   void find_scc(int n)
36
37
        dfs_clock = 0, scc_cnt = 0;
38
39
        clr(sccno, 0), clr(pre, 0);
40
        for (int i = 0; i < n; i++)
            if (!pre[i]) dfs(i);
41
  }
42
   4.6.3 2-SAT
   struct TwoSAT
2
   {
3
        int n;
        vector<int> G[maxn << 1];</pre>
4
        bool mark[maxn << 1];</pre>
5
        int S[maxn << 1], c;</pre>
6
        void init(int n)
7
8
9
10
            for (int i = 0; i < (n << 1); i++) G[i].clear();</pre>
            clr(mark, 0);
11
12
13
        bool dfs(int x)
14
            if (mark[x ^ 1]) return false;
15
```

```
if (mark[x]) return true;
16
            mark[x] = true;
17
            S[c++] = x;
18
            for (auto\& y : G[x])
19
                if (!dfs(y)) return false;
20
21
            return true;
22
23
        //x = xval or y = yval
        void add_clause(int x, int xval, int y, int yval)
24
25
26
            x = (x << 1) + xval;
27
            y = (y << 1) + yval;
            G[x \wedge 1].pb(y);
28
            G[y \wedge 1].pb(x);
29
30
        bool solve()
31
32
            for (int i = 0; i < (n << 1); i += 2)
33
                if (!mark[i] && !mark[i + 1])
34
35
                     c = 0;
36
                     if (!dfs(i))
37
38
39
                         while (c > 0) mark[S[--c]] = false;
40
                         if (!dfs(i + 1)) return false;
                     }
41
                }
42
43
            return true;
        }
44
   };
45
   4.7 Eular Path
                           )
                    ( ,
                              ),
           G
           - G
           – G
                (
                          0 2.
                      )
           G
           - G
           - G
           G
           - G
                   _{1,v}
                         1,
                                 (u ,v )
           – u
           G
           - G
```

- G

```
4.7.1 Fleury
```

15

```
1 const int maxn = "Edit";
2 int G[maxn][maxn];
3 int deg[maxn][maxn];
4 vector<int> Ans;
5 inline void init() { clr(G, 0), clr(deg, 0); }
  inline void AddEdge(int u, int v) { deg[u]++, deg[v]++, G[u][v]++, G[v][u]++; }
   void Fleury(int s)
7
   {
8
        for (int i = 0; i < n; i++)
9
10
            if (G[s][i])
11
12
                G[s][i]--, G[i][s]--;
13
                Fleury(i);
14
15
        Ans.pb(s);
16 }
   4.8 Bipartite Graph Matching
      1.
      2.
            =|G|-
          N \times N
                                );
         (a)
         (b)
                p_1, p_2, \dots, p_k, p_1, p_k, p_1, p_2, \dots, p_k
                    G
                   =|G|-;
      3.
             =
         :
   4.8.1 Hungry(Matrix)
      :O(VE).
1 const int maxn = "Edit";
2 int uN, vN; //uN ,vN
3 int g[maxn][maxn]; // g[i][j] i->j
4 int linker[maxn];
   bool used[maxn];
5
   bool dfs(int u)
6
7
   {
        for (int V = 0; V < VN; V++)
8
9
            if (g[u][v] && !used[v])
10
11
                used[v] = true;
                if (linker[v] == -1 || dfs(linker[v]))
12
13
                {
                     linker[v] = u;
14
```

return true;

```
}
16
            }
17
        return false;
18
   }
19
   int hungary()
20
21
   {
        int res = 0;
22
23
        clr(linker, -1);
24
        for (int u = 0; u < uN; u++)
25
26
            clr(used, 0);
27
            if (dfs(u)) res++;
        }
28
29
        return res;
30 }
   4.8.2 Hungry(List)
      init()
       addedge(u,v)
1 const int maxn = "Edit";
2 int n;
   vector<int> G[maxn];
   int linker[maxn];
   bool used[maxn];
5
   inline void init(int n)
6
7
   {
        for (int i = 0; i < n; i++) G[i].clear();</pre>
8
9
inline void addedge(int u, int v) { G[u].pb(v); }
11 bool dfs(int u)
12
   {
        for (auto& v : G[u])
13
14
            if (!used[v])
15
16
                used[v] = true;
17
                if (linker[v] == -1 || dfs(linker[v]))
18
19
20
                     linker[v] = u;
21
                     return true;
22
                }
23
            }
        }
24
        return false;
25
   }
26
27
   int hungary()
28
29
        int ans = 0;
        clr(linker, -1);
30
        for (int u = 0; u < n; v++)
31
32
        {
33
            clr(used, 0);
            if (dfs(u)) ans++;
34
35
36
        return ans;
37 }
```

# 4.8.3 Hopcroft-Carp

```
O(\sqrt{n}*E)
   uN
         , \qquad (\quad 0 \quad )
1 const int maxn = "Edit";
2
   vector<int> G[maxn];
3
  int uN;
  int Mx[maxn], My[maxn];
  int dx[maxn], dy[maxn];
  int dis;
  bool used[maxn];
7
8 inline void init(int n)
9
        for (int i = 0; i < n; i++) G[i].clear();</pre>
10
   }
11
   inline void addedge(int u, int v) { G[u].pb(v); }
12
   bool bfs()
13
14
   {
        queue<int> q;
15
16
        dis = INF;
        clr(dx, -1), clr(dy, -1);
17
        for (int i = 0; i < uN; i++)
18
            if (Mx[i] == -1)
19
                 q.push(i), dx[i] = 0;
20
21
        while (!q.empty())
22
23
            int u = q.front();
24
            q.pop();
            if (dx[u] > dis) break;
25
26
            for (auto& v : G[u])
27
                 if (dy[v] == -1)
28
29
30
                     dy[v] = dx[u] + 1;
                     if (My[v] == -1)
31
                         dis = dy[v];
32
33
                     else
34
                     {
35
                         dx[My[v]] = dy[v] + 1;
36
                         q.push(My[v]);
37
                     }
                 }
38
            }
39
40
41
        return dis != INF;
42
   bool dfs(int u)
43
44
   {
        for (auto& v : G[u])
45
46
            if (!used[v] && dy[v] == dx[u] + 1)
47
48
            {
                 used[v] = true;
49
                 if (My[v] != -1 \&\& dy[v] == dis) continue;
50
                 if (My[v] == -1 \mid | dfs(My[v]))
51
52
                 {
                     My[v] = u, Mx[u] = v;
53
                     return true;
54
```

```
}
55
            }
56
57
        return false;
58
59
  int MaxMatch()
60
61
   {
        int res = 0;
62
        clr(Mx, -1), clr(My, -1);
63
        while (bfs())
64
65
66
            clr(used, false);
            for (int i = 0; i < uN; i++)
67
                if (Mx[i] == -1 \&\& dfs(i)) res++;
68
69
        return res;
70
71
  }
   4.8.4 Hungry(Multiple)
1 const int maxn = "Edit";
  const int maxm = "Edit";
3 int uN, vN;
                       //u,v
  int g[maxn][maxm]; //
  int linker[maxm][maxn];
6 bool used[maxm];
   int num[maxm]; //
7
   bool dfs(int u)
8
9
   {
        for (int v = 0; v < vN; v++)
10
            if (g[u][v] && !used[v])
11
            {
12
                used[v] = true;
13
                if (linker[v][0] < num[v])</pre>
14
15
                     linker[v][++linker[v][0]] = u;
16
                     return true;
17
18
                for (int i = 1; i <= num[0]; i++)
19
20
                     if (dfs(linker[v][i]))
21
22
                         linker[v][i] = u;
23
                         return true;
24
25
            }
26
        return false;
   }
27
28
   int hungary()
29
30
        int res = 0;
        for (int i = 0; i < vN; i++) linker[i][0] = 0;
31
        for (int u = 0; u < uN; u++)
32
33
        {
            clr(used, 0);
34
35
            if (dfs(u)) res++;
36
37
        return res;
38 }
```

### 4.8.5 Kuhn-Munkres

```
const int maxn = "Edit";
                                            //
2
  int nx, ny;
3 int g[maxn][maxn];
4 int linker[maxn], lx[maxn], ly[maxn]; //y
                                                   ,x,y
5 int slack[N];
6 bool visx[N], visy[N];
7 bool dfs(int x)
   {
8
9
       visx[x] = true;
10
       for (int y = 0; y < ny; y++)
11
12
            if (visy[y]) continue;
            int tmp = lx[x] + ly[y] - g[x][y];
13
            if (tmp == 0)
14
15
                visy[y] = true;
16
                if (linker[y] == -1 || dfs(linker[y]))
17
18
19
                    linker[y] = x;
20
                    return true;
                }
21
22
            else if (slack[y] > tmp)
23
24
                slack[y] = tmp;
25
       }
26
       return false;
27
   }
  int KM()
28
29
   {
30
       clr(linker, -1), clr(ly, 0);
31
       for (int i = 0; i < nx; i++)
32
            lx[i] = -INF;
33
            for (int j = 0; j < ny; j++)
34
                if (g[i][j] > lx[i]) lx[i] = g[i][j];
35
36
       for (int x = 0; x < nx; x++)
37
38
39
            clr(slack, 0x3f);
40
            for (;;)
41
                clr(visx, 0), clr(visy, 0);
42
                if (dfs(x)) break;
43
                int d = INF;
44
                for (int i = 0; i < ny; i++)
45
                    if (!visy[i] && d > slack[i]) d = slack[i];
46
                for (int i = 0; i < nx; i++)
47
                    if (visx[i]) lx[i] -= d;
48
                for (int i = 0; i < ny; i++)
49
50
                    if (visy[i])
                         ly[i] += d;
51
52
                    else
                        slack[i] -= d;
53
           }
54
55
56
       int res = 0;
       for (int i = 0; i < ny; i++)
57
```

```
if (~linker[i]) res += g[linker[i]][i];
58
59
        return res;
60 }
   4.9 Network Flow
1 struct Edge
2
        int from, to, cap, flow;
3
        Edge(int u, int v, int c, int f)
4
           : from(u), to(v), cap(c), flow(f) {}
5
6 };
   struct Edge
2
   {
        int from, to, cap, flow, cost;
3
        Edge(int u, int v, int c, int f, int w)
4
            : from(u), to(v), cap(c), flow(f), cost(w) {}
6 };
              S T, S X , Y T , X Y
                                                       \{p_1, p_2\}
           m
                          ,
                                                                 , 1,
                , , T
                              m
                                                O(\log m)
        k
                                 k
                    , \qquad w \qquad [u,v) \quad u {\rightarrow} v, \quad 1, \quad -w \qquad \qquad i {\rightarrow} i+1, \quad k, \quad 0 \quad ,
             G( ),
           s t, s
                      , ;
                                           S - \{s\}
   4.9.1 EdmondKarp
   const int maxn = "Edit";
   struct EdmonsKarp // O(v*E*E)
3
    {
        int n, m;
4
        vector<Edge> edges; //
5
        vector<int> G[maxn]; // ,G[i][j] i j e
6
        int a[maxn];
                            // i
7
        int p[maxn];
                              // p
8
9
        void init(int n)
10
            for (int i = 0; i < n; i++) G[i].clear();</pre>
11
            edges.clear();
12
13
        void AddEdge(int from, int to, int cap)
14
15
```

```
edges.pb(Edge(from, to, cap, 0));
16
            edges.pb(Edge(to, from, 0, 0)); //
17
            m = edges.size();
18
            G[from].pb(m - 2);
19
20
            G[to].pb(m - 1);
21
        int Maxflow(int s, int t)
22
23
            int flow = 0;
24
            for (;;)
25
26
            {
27
                 clr(a, 0);
                queue<int> q;
28
                q.push(s);
29
                a[s] = INF;
30
                while (!q.empty())
31
32
                     int x = q.front();
33
34
                     q.pop();
                     for (int i = 0; i < G[x].size(); i++)</pre>
35
36
                         Edge& e = edges[G[x][i]];
37
                         if (!a[e.to] && e.cap > e.flow)
38
39
40
                              p[e.to] = G[x][i];
                             a[e.to] = min(a[x], e.cap - e.flow);
41
                              q.push(e.to);
42
43
44
                     if (a[t]) break;
45
46
                if (!a[t]) break;
47
                for (int u = t; u != s; u = edges[p[u]].from)
48
49
                     edges[p[u]].flow += a[t];
50
                     edges[p[u] ^1].flow -= a[t];
51
52
53
                flow += a[t];
            }
54
            return flow;
55
        }
56
   };
57
   4.9.2 Dinic
   const int maxn = "Edit";
1
   struct Dinic
2
3
   {
                              // , ( ),
4
        int n, m, s, t;
        vector<Edge> edges; // edge[e] edge[e^1]
5
        vector<int> G[maxn]; // ,G[i][j] i j e
6
        bool vis[maxn];
                              //BFS
7
8
        int d[maxn];
                               // i
        int cur[maxn];
                               //
9
        void init(int n)
10
11
        {
12
            this -> n = n;
            for (int i = 0; i < n; i++) G[i].clear();</pre>
13
```

```
edges.clear();
14
15
        void AddEdge(int from, int to, int cap)
16
17
            edges.pb(Edge(from, to, cap, 0));
18
            edges.pb(Edge(to, from, 0, 0));
19
            m = edges.size();
20
            G[from].pb(m - 2);
21
22
            G[to].pb(m - 1);
23
        }
24
        bool BFS()
25
            clr(vis, 0);
26
            clr(d, 0);
27
28
            queue<int> q;
29
            q.push(s);
30
            d[s] = 0;
            vis[s] = 1;
31
            while (!q.empty())
32
33
                 int x = q.front();
34
35
                 q.pop();
                 for (int i = 0; i < G[x].size(); i++)</pre>
36
37
                     Edge& e = edges[G[x][i]];
38
                     if (!vis[e.to] && e.cap > e.flow)
39
40
                          vis[e.to] = 1;
41
                          d[e.to] = d[x] + 1;
42
43
                          q.push(e.to);
44
                     }
                 }
45
46
            }
            return vis[t];
47
48
        int DFS(int x, int a)
49
50
51
            if (x == t | | a == 0) return a;
            int flow = 0, f;
52
            for (int& i = cur[x]; i < G[x].size(); i++)</pre>
53
54
55
                 Edge& e = edges[G[x][i]];
56
                 if (d[x] + 1 == d[e.to] && (f = DFS(e.to, min(a, e.cap - e.flow))) > 0)
57
58
                     e.flow += f;
59
                     edges[G[x][i] \land 1].flow -= f;
60
                     flow += f;
61
                     a -= f;
62
63
                     if (a == 0) break;
64
                 }
65
            }
            return flow;
66
67
        int Maxflow(int s, int t)
68
69
70
            this -> s = s;
71
            this->t = t;
            int flow = 0;
72
```

```
while (BFS())
73
74
                clr(cur, 0);
75
                flow += DFS(s, INF);
76
77
            return flow;
78
79
       }
80
   };
   4.9.3 ISAP
   const int maxn = "Edit";
2
   struct ISAP
3
   {
       4
5
       vector<int> G[maxn]; // ,G[i][j] i j e
6
                             //BFS
       bool vis[maxn];
7
8
       int d[maxn];
                             // i
9
       int cur[maxn];
                             //
       int p[maxn];
                             //
10
       int num[maxn];
                             //
11
       void init(int n)
12
13
       {
            this->n = n;
14
15
            for (int i = 0; i < n; i++) G[i].clear();</pre>
16
            edges.clear();
17
       }
       void AddEdge(int from, int to, int cap)
18
19
20
            edges.pb(Edge(from, to, cap, 0));
21
            edges.pb(Edge(to, from, 0, 0));
22
            int m = edges.size();
23
           G[from].pb(m - 2);
           G[to].pb(m - 1);
24
25
26
       int Augument()
27
28
            int x = t, a = INF;
29
           while (x != s)
30
                Edge& e = edges[p[x]];
31
32
                a = min(a, e.cap - e.flow);
                x = edges[p[x]].from;
33
34
            }
35
           x = t;
           while (x != s)
36
37
            {
                edges[p[x]].flow += a;
38
                edges[p[x] ^ 1].flow -= a;
39
                x = edges[p[x]].from;
40
            }
41
42
            return a;
43
       }
       void BFS()
44
45
            clr(vis, 0);
46
            clr(d, 0);
47
```

```
48
             queue<int> q;
49
             q.push(t);
             d[t] = 0;
50
             vis[t] = 1;
51
52
             while (!q.empty())
             {
53
                  int x = q.front();
54
55
                  q.pop();
                  int len = G[x].size();
56
                  for (int i = 0; i < len; i++)
57
58
59
                      Edge& e = edges[G[x][i]];
                      if (!vis[e.from] && e.cap > e.flow)
60
61
                           vis[e.from] = 1;
62
                           d[e.from] = d[x] + 1;
63
64
                           q.push(e.from);
                      }
65
                  }
66
             }
67
68
         int Maxflow(int s, int t)
69
70
71
             this -> s = s;
72
             this->t = t;
             int flow = 0;
73
             BFS();
74
             clr(num, 0);
75
             for (int i = 0; i < n; i++)
76
                  if (d[i] < INF) num[d[i]]++;</pre>
77
78
             int x = s;
79
             clr(cur, 0);
80
             while (d[s] < n)
             {
81
                  if(x == t)
82
83
                      flow += Augumemt();
84
85
                      X = S;
86
                  int ok = 0;
87
                  for (int i = cur[x]; i < G[x].size(); i++)</pre>
88
89
                      Edge& e = edges[G[x][i]];
90
91
                      if (e.cap > e.flow && d[x] == d[e.to] + 1)
92
                           ok = 1;
93
                           p[e.to] = G[x][i];
94
                           cur[x] = i;
95
96
                           x = e.to;
97
                           break;
98
                      }
99
                  if (!ok) //Retreat
100
101
                      int m = n - 1;
102
                      for (int i = 0; i < G[x].size(); i++)</pre>
103
104
                           Edge& e = edges[G[x][i]];
105
                           if (e.cap > e.flow) m = min(m, d[e.to]);
106
```

```
107
                      if (--num[d[x]] == 0) break; //gap
108
                      num[d[x] = m + 1] ++;
109
110
                      cur[x] = 0;
111
                      if (x != s) x = edges[p[x]].from;
112
             }
113
114
             return flow;
         }
115
116
    };
    4.9.4 MinCost MaxFlow
    const int maxn = "Edit";
    struct MCMF
 2
    {
 3
         int n, m;
 4
 5
         vector<Edge> edges;
 6
         vector<int> G[maxn];
 7
         int inq[maxn]; //
         int d[maxn];
                        //bellmanford
 8
 9
         int p[maxn];
                         //
                        //
         int a[maxn];
10
         void init(int n)
11
12
         {
13
             this->n = n;
14
             for (int i = 0; i < n; i++) G[i].clear();</pre>
15
             edges.clear();
16
         void AddEdge(int from, int to, int cap, int cost)
17
18
             edges.pb(Edge(from, to, cap, 0, cost));
19
20
             edges.pb(Edge(to, from, 0, 0, -cost));
21
             m = edges.size();
             G[from].pb(m - 2);
22
23
             G[to].pb(m - 1);
24
25
         bool BellmanFord(int s, int t, int& flow, ll& cost)
26
27
             for (int i = 0; i < n; i++) d[i] = INF;
             clr(inq, 0);
28
             d[s] = 0;
29
             inq[s] = 1;
30
31
             p[s] = 0;
             a[s] = INF;
32
             queue<int> q;
33
             q.push(s);
34
35
             while (!q.empty())
36
                 int u = q.front();
37
38
                 q.pop();
                 inq[u] = 0;
39
                 for (int i = 0; i < G[u].size(); i++)</pre>
40
41
                      Edge& e = edges[G[u][i]];
42
                      if (e.cap > e.flow && d[e.to] > d[u] + e.cost)
43
44
                          d[e.to] = d[u] + e.cost;
45
```

```
p[e.to] = G[u][i];
a[e.to] = min(a[u], e.cap - e.flow);
46
47
                          if (!inq[e.to])
48
49
                              q.push(e.to);
50
                              inq[e.to] = 1;
51
                          }
52
53
                     }
                 }
54
            }
55
             if (d[t] == INF) return false; //
56
             flow += a[t];
57
             cost += (ll)d[t] * (ll)a[t];
58
             for (int u = t; u != s; u = edges[p[u]].from)
59
60
                 edges[p[u]].flow += a[t];
61
                 edges[p[u] \land 1].flow -= a[t];
62
63
64
             return true;
65
66
        int MincostMaxflow(int s, int t, ll& cost)
67
68
            int flow = 0;
             cost = 0;
69
            while (BellmanFord(s, t, flow, cost));
70
71
             return flow;
72
        }
73 };
```

# 5 Computational Geometry

### 5.1 Basic Function

```
#define zero(x) ((fabs(x) < eps ? 1 : 0))
   #define sqn(x) (fabs(x) < eps ? 0 : ((x) < 0 ? -1 : 1))
4 struct point
5
       double x, y;
6
       point(double a = 0, double b = 0) { x = a, y = b; }
7
       point operator-(const point& b) const { return point(x - b.x, y - b.y); }
8
       point operator+(const point& b) const { return point(x + b.x, y + b.y); }
9
10
       bool operator==(point& b) { return zero(x - b.x) && zero(y - b.y); }
11
12
       // ( )
       double operator*(const point& b) const { return x * b.x + y * b.y; }
13
       // ( )
14
       double operator^(const point& b) const { return x * b.y - y * b.x; }
15
       point rotate(point b, double a)
17
18
           double dx, dy;
19
           (*this - b).split(dx, dy);
20
           double tx = dx * cos(a) - dy * sin(a);
21
           double ty = dx * sin(a) + dy * cos(a);
22
23
           return point(tx, ty) + b;
24
       //
25
              a b
26
       void split(double& a, double& b) { a = x, b = y; }
27
   };
28 struct line
29 = \{
       point s, e;
30
31
       line() {}
       line(point ss, point ee) { s = ss, e = ee; }
32
   };
33
   5.2 Position
   5.2.1 Point-Point
double dist(point a, point b) { return sqrt((a - b) * (a - b)); }
   5.2.2 Line-Line
1 // <0, *> ; <1, *> ; <2, P>
  pair<int, point> spoint(line l1, line l2)
2
3
       point res = l1.s;
4
       if (sgn((11.s - 11.e) \wedge (12.s - 12.e)) == 0)
5
           return mp(sqn((l1.s - l2.e) ^ (l2.s - l2.e)) != 0, res);
6
       double t = ((11.s - 12.s) \wedge (12.s - 12.e)) / ((11.s - 11.e) \wedge (12.s - 12.e));
7
       res.x += (l1.e.x - l1.s.x) * t;
8
       res.y += (l1.e.y - l1.s.y) * t;
9
       return mp(2, res);
10
11 }
```

```
5.2.3 Segment-Segment
```

```
1 bool segxseg(line l1, line l2)
2
   {
3
       return
4
           max(11.s.x, 11.e.x) >= min(12.s.x, 12.e.x) &&
5
            max(12.s.x, 12.e.x) >= min(11.s.x, 11.e.x) &&
            max(11.s.y, 11.e.y) >= min(12.s.y, 12.e.y) &&
6
            max(12.s.y, 12.e.y) >= min(11.s.y, 11.e.y) &&
7
            sgn((l2.s - l1.e) \land (l1.s - l1.e)) * sgn((l2.e-l1.e) \land (l1.s - l1.e)) <= 0 &&
8
            sgn((11.s - 12.e) \wedge (12.s - 12.e)) * sgn((11.e-12.e) \wedge (12.s - 12.e)) <= 0;
9
10 }
   5.2.4 Line-Segment
1 //11 ,12
2 bool segxline(line l1, line l2)
3
       return sgn((l2.s - l1.e) ^ (l1.s - l1.e)) * sgn((l2.e - l1.e) ^ (l1.s - l1.e)) <=
4
       0;
5 }
   5.2.5 Point-Line
1 double pointtoline(point p, line l)
2
       point res;
3
       double t = ((p - l.s) * (l.e - l.s)) / ((l.e - l.s) * (l.e - l.s));
4
       res.x = 1.s.x + (1.e.x - 1.s.x) * t, res.y = 1.s.y + (1.e.y - 1.s.y) * t;
5
       return dist(p, res);
6
7
  }
   5.2.6 Point-Segment
   double pointtosegment(point p, line l)
2
3
       point res:
       double t = ((p - l.s) * (l.e - l.s)) / ((l.e - l.s) * (l.e - l.s));
4
       if (t >= 0 && t <= 1)
5
            res.x = l.s.x + (l.e.x - l.s.x) * t, res.y = l.s.y + (l.e.y - l.s.y) * t;
6
7
       else
            res = dist(p, l.s) < dist(p, l.e) ? l.s : l.e;
8
9
       return dist(p, res);
10 }
   5.2.7 Point on Segment
   bool PointOnSeg(point p, line l)
1
2
3
       return
            sgn((1.s - p) \wedge (1.e-p)) == 0 \&\&
4
5
            sgn((p.x - l.s.x) * (p.x - l.e.x)) <= 0 &&
6
            sgn((p.y - l.s.y) * (p.y - l.e.y)) <= 0;
7 }
```

#### 5.3 Polygon 5.3.1 Area 1 double area(point p[], int n) 2 3 double res = 0; for (int i = 0; i < n; i++) res $+= (p[i] \land p[(i + 1) \% n]) / 2;$ 4 return fabs(res); 6 } 5.3.2 Point in Convex < 0 > 0)// : [0,n) 3 // -1: 4 // 0 : 5 // 1 6 int PointInConvex(point a, point p∏, int n) 7 { for (int i = 0; i < n; i++) 8 if $(sgn((p[i] - a) \land (p[(i + 1) \% n] - a)) < 0)$ 9 10 return -1; else if (PointOnSeg(a, line(p[i], p[(i + 1) % n]))) 11 return 0; 1213 return 1; 14 } 5.3.3 Point in Polygon 3, 0~n-1 1 // ,poly[] 2 // -1: 3 // 0 : 4 // 1 5 int PointInPoly(point p, point poly[], int n) { 6 int cnt; 7 line ray, side; 8 9 cnt = 0;10 ray.s = p;11 ray.e.y = p.y; ray.e.x = -100000000000.0; // -INF, 12 for (int i = 0; i < n; i++) 13 14 side.s = poly[i], side.e = poly[(i + 1) % n]; 15 if (PointOnSeg(p, side)) return 0; 16 17 if (sgn(side.s.y - side.e.y) == 0)18

19

20 21

22

23

2425

26 27

28 }

continue;

cnt++;

return cnt % 2 == 1 ? 1 : -1;

if (PointOnSeg(sid e.s, r ay))

else if (segxseg(ray, side))

else if (PointOnSeg(side.e, ray))

cnt += (sgn(side.s.y - side.e.y) > 0);

cnt += (sgn(side.e.y - side.s.y) > 0);

```
5.3.4 Judge Convex
1 //
2 // 1~n-1
3 bool isconvex(point poly[], int n)
4
       bool s[3];
5
       clr(s, 0);
6
       for (int i = 0; i < n; i++)
7
8
           s[sgn((poly[(i + 1) % n] - poly[i]) ^ (poly[(i + 2) % n] - poly[i])) + 1] = 1;
9
           if (s[0] && s[2]) return 0;
10
11
12
       return 1;
13 }
   5.4 Integer Points
   5.4.1 On Segment
int OnSegment(line l) { return __gcd(fabs(l.s.x - l.e.x), fabs(l.s.y - l.e.y)) + 1; }
   5.4.2 On Polygon Edge
1 int OnEdge(point p□, int n)
2
       int i, ret = 0;
3
       for (i = 0; i < n; i++)
4
           ret += \__gcd(fabs(p[i].x - p[(i + 1) % n].x), fabs(p[i].y - p[(i + 1) % n].y));
5
       return ret;
6
7
   }
   5.4.3 Inside Polygon
1 int InSide(point p□, int n)
2
   {
3
       int i, area = 0;
4
       for (i = 0; i < n; i++)
           area += p[(i + 1) % n].y * (p[i].x - p[(i + 2) % n].x);
5
       return (fabs(area) - OnEdge(n, p)) / 2 + 1;
6
   }
7
   5.5 Circle
   5.5.1 Circumcenter
   point waixin(point a, point b, point c)
2
       double a1 = b.x - a.x, b1 = b.y - a.y, c1 = (a1 * a1 + b1 * b1) / 2;
3
       double a2 = c.x - a.x, b2 = c.y - a.y, c2 = (a2 * a2 + b2 * b2) / 2;
4
       double d = a1 * b2 - a2 * b1;
5
       return point(a.x + (c1 * b2 - c2 * b1) / d, a.y + (a1 * c2 - a2 * c1) / d);
6
7 }
```

# 6 Dynamic Programming

## 6.1 Subsequence

```
6.1.1 Max Sum
```

```
1 // a n,
2 int MaxSeqSum(int a[], int n)
3 {
4    int rt = 0, cur = 0;
5    for (int i = 0; i < n; i++)
6         cur += a[i], rt = max(cur, rt), cur = max(0, cur);
7    return rt;
8 }</pre>
```

#### 6.1.2 Longest Increase

```
// 1 ,LIS() , lis[] const int N = "Edit";
1 //
  int len, a[N], b[N], f[N];
  int Find(int p, int l, int r)
   {
5
6
        while (l \ll r)
7
8
            int mid = (l + r) \gg 1;
9
            if (a[p] > b[mid])
                l = mid + 1;
10
            else
11
                r = mid - 1;
12
13
        return f[p] = 1;
14
15
16 int LIS(int lis[], int n)
17
   {
        int len = 1;
18
        f[1] = 1, b[1] = a[1];
19
        for (int i = 2; i <= n; i++)
20
21
            if (a[i] > b[len])
22
                b[++len] = a[i], f[i] = len;
23
24
            else
                b[Find(i, 1, len)] = a[i];
25
26
        for (int i = n, t = len; i >= 1 && t >= 1; i--)
27
28
            if (f[i] == t) lis[--t] = a[i];
29
        return len;
30 }
31
32 //
        (0,
  int dp[N];
  int LIS(int a[], int n)
35  {
        clr(dp, 0x3f);
36
        for (int i = 0; i < n; i++) *lower_bound(dp, dp + n, a[i]) = a[i];
37
        return lower_bound(dp, dp + n, INF) - dp;
38
39 }
```

### 6.1.3 Longest Common Increase

```
1
2 int LCIS(int a□, int b□, int n, int m)
3 {
       clr(dp, 0);
4
       for (int i = 1; i <= n; i++)
5
6
7
            int ma = 0;
            for (int j = 1; j <= m; j++)
8
9
                dp[i][j] = dp[i - 1][j];
10
                if (a[i] > b[j]) ma = max(ma, dp[i - 1][j]);
11
                if (a[i] == b[j]) dp[i][j] = ma + 1;
12
13
14
       return *max_element(dp[n] + 1, dp[n] + 1 + m);
15
16 }
   6.2 Digit Statistics
   int a[20];
   11 dp[20][state];
   ll dfs(int pos, /*state */, bool lead /* */, bool limit /* */)
3
4
       // , 0, pos==-1
5
       if (pos == -1) return 1;
6
7
           1,
8
            pos,
       if (!limit && !lead && dp[pos][state] != -1) return dp[pos][state];
9
10
       int up = limit ? a[pos] : 9; // limit
11
12
       11 \text{ ans} = 0;
       for (int i = 0; i \le up; i++) //,
13
14
            if () ...
15
           else if () ...
16
           ans += dfs(pos - 1, /* */, lead && i == 0, limit && i == a[pos])
17
18
           /*
19
                 , state i */
20
       }
21
22
       if (!limit && !lead) dp[pos][state] = ans;
23
24
             lead, lead
25
26
       return ans;
27
28 ll solve(ll x)
29
   {
       int pos = 0;
30
       do //
31
32
           a[pos++] = x \% 10;
33
       while (x \neq 10);
                                 */, /* */, true, true);
       return dfs(pos - 1 /*
34
35
36 }
```

# Others

### 7.1 Matrix

```
7.1.1 Matrix FastPow
```

```
typedef vector<ll> vec;
   typedef vector<vec> mat;
3
   mat mul(mat& A, mat& B)
4
        mat C(A.size(), vec(B[0].size()));
5
6
        for (int i = 0; i < A.size(); i++)</pre>
            for (int k = 0; k < B.size(); k++)
7
                 if (A[i][k]) //
8
9
                     for (int j = 0; j < B[0].size(); j++)
                         C[i][j] = (C[i][j] + A[i][k] * B[k][j]) % mod;
10
11
        return C;
12 }
13
   mat Pow(mat A, ll n)
14
        mat B(A.size(), vec(A.size()));
15
        for (int i = 0; i < A.size(); i++) B[i][i] = 1;
for (; n; n >>= 1, A = mul(A, A))
16
17
18
            if (n \& 1) B = mul(B, A);
        return B;
19
  }
20
   7.1.2 Gauss Elimination
1
   void gauss()
2
   {
        int now = 1, to;
3
        double t;
4
        for (int i = 1; i <= n; i++, now++)
5
6
            /*for (to = now; !a[to][i] && to <= n; to++);
7
8
            if (to != now)
9
                 for (int j = 1; j <= n + 1; j++)
10
                     swap(a[to][j], a[now][j]);*/
11
            t = a[now][i];
12
            for (int j = 1; j \le n + 1; j++) a[now][j] /= t;
13
            for (int j = 1; j <= n; j++)
14
                 if (j != now)
15
16
                 {
                     t = a[j][i];
17
                     for (int k = 1; k \le n + 1; k++) a[j][k] -= t * a[now][k];
18
19
20
        }
21
   }
   7.2
        Tricks
```

#### 7.2.1 Stack-Overflow

```
//
2 #pragma comment(linker, "/STACK:1024000000,1024000000")
```

```
7.2.2 Fast-Scanner
```

```
1 //
2 template <class T>
3 inline bool scan_d(T &ret)
4
5
       char c:
6
       int sqn;
       if (c = getchar(), c == EOF) return 0; //EOF
7
       while (c != '-' && (c < '0' || c > '9')) c = getchar();
8
       sgn = (c == '-') ? -1 : 1;
9
       ret = (c == '-') ? 0 : (c - '0');
10
       while (c = getchar(), c >= '0' \&\& c <= '9') ret = ret * 10 + (c - '0');
11
12
       ret *= san;
13
       return 1;
14 }
15 inline void out(int x)
16 {
       if (x > 9) out(x / 10);
17
       putchar(x % 10 + '0');
18
19 }
   7.2.3 Strok-Sscanf
1 //
gets(buf);
3 int v;
4 char *p = strtok(buf, " ");
5 while (p)
6
   {
       sscanf(p, "%d", &v);
7
       p = strtok(NULL," ");
8
   }
   7.3 Mo Algorithm
                 \sqrt{x},
   struct query { int L, R, id; };
   void solve(query node[], int m)
3
   {
4
       tmp = 0;
5
       clr(num, 0);
6
       clr(ans, 0);
       sort(node, node + m, [](query a, query b) { return a.l / unit < b.l / unit || a.l /
7
        unit == b.l / unit && a.r < b.r; });
       int L = 1, R = 0;
8
9
       for (int i = 0; i < m; i++)
10
           while (node[i].L < L) add(a[--L]);
11
           while (node[i].L > L) del(a[L++]);
12
           while (node[i].R < R) del(a[R--]);
13
           while (node[i].R > R) add(a[++R]);
14
           ans[node[i].id] = tmp;
15
16
       }
17 }
```

# 7.4 BigNum

#### 7.4.1 High-precision

```
1 //
2 struct bint
3
   {
        int 1;
4
        short int w[100];
5
        bint(int x = 0)
6
7
            l = x == 0, clr(w, 0);
8
            while (x) w[l++] = x \% 10, x /= 10;
9
10
        bool operator<(const bint& x) const</pre>
11
12
            if (l != x.l) return l < x.l;
13
            int i = l - 1;
14
            while (i >= 0 \&\& w[i] == x.w[i]) i--;
15
            return (i >= 0 \& w[i] < x.w[i]);
16
17
        bint operator+(const bint& x) const
18
19
20
            bint ans;
            ans.1 = 1 > x.1 ? 1 : x.1;
21
            for (int i = 0; i < ans.l; i++)
22
23
24
                ans.w[i] += w[i] + x.w[i];
                ans.w[i + 1] += ans.w[i] / 10;
25
                ans.w[i] = ans.w[i] % 10;
26
27
            if (ans.w[ans.l] != 0) ans.l++;
28
            return ans;
29
30
        bint operator*(const bint& x) const
31
32
33
            bint res;
            int up, tmp;
34
            for (int i = 0; i < 1; i++)
35
36
37
                up = 0;
                for (int j = 0; j < x.1; j++)
38
39
                     tmp = w[i] * x.w[j] + res.w[i + j] + up;
40
                     res.w[i + j] = tmp % 10;
41
                     up = tmp / 10;
42
43
                if (up != 0) res.w[i + x.l] = up;
44
45
46
            res.l = l + x.l;
            while (res.w[res.l - 1] == 0 && res.l > 1) res.l--;
47
48
            return res;
49
        void print()
50
51
            for (int i = l - 1; ~i; i--) printf("%d", w[i]);
52
            puts("");
53
        }
54
55 };
```

### 7.4.2 Complete High-precision

```
#define N 10000
   class bint
2
3
4
   private:
        int a[N]; // N
5
        int len; //
6
   public:
7
8
        bint() { len = 1, clr(a, 0); }
9
10
        // int -> bint
        bint(int n)
11
12
13
            len = 0;
            clr(a, 0);
14
            int d = n;
15
16
            while (n)
                d = n / 10 * 10, a[len++] = n - d, n = d / 10;
17
18
        }
19
        // char[] -> int
20
        bint(const char s[])
21
        {
            clr(a, 0);
22
23
            len = 0;
24
            int l = strlen(s);
25
            for (int i = l - 1; ~i; i--) a[len++] = s[i];
26
        }
        //
27
28
        bint(const bint& b)
29
30
            clr(a, 0);
31
            len = b.len;
            for (int i = 0; i < len; i++) a[i] = b.a[i];</pre>
32
33
        //
34
               bint = bint
        bint& operator=(const bint& n)
35
36
37
            len = n.len;
38
            for (int i = 0; i < len; i++) a[i] = n.a[i];
            return *this;
39
        }
40
        //
               bint + bint
41
        bint operator+(const bint& b) const
42
43
            bint t(*this);
44
            int res = b.len > len ? b.len : len;
45
            for (int i = 0; i < res; i++)
46
47
            {
                t.a[i] += b.a[i];
48
                if (t.a[i] >= 10) t.a[i + 1]++, t.a[i] -= 10;
49
50
            t.len = res + a[res] == 0;
51
52
            return t;
53
        }
        //
               bint - bint
54
        bint operator-(const bint& b) const
55
56
            bool f = *this > b;
57
```

```
bint t1 = f ? *this : b;
58
             bint t2 = f ? b : *this;
59
             int res = t1.len, j;
60
             for (int i = 0; i < res; i++)</pre>
61
                  if (t1.a[i] < t2.a[i])</pre>
62
63
                      j = i + 1;
64
                      while (t1.a[j] == 0) j++;
65
                      t1.a[j--]--;
66
                      while (j > i) t1.a[j--] += 9;
67
                      t1.a[i] += 10 - t1.a[i];
68
69
                  }
                  else
70
                      t1.a[i] -= t2.a[i];
71
             t1.len = res;
72
             while (t1.a[len - 1] == 0 && t1.len > 1) t1.len--, res--;
if (f) t1.a[res - 1] = 0 - t1.a[res - 1];
73
74
75
             return t1;
         }
76
         //
                bint * bint
77
         bint operator*(const bint& b) const
78
79
             bint t;
80
81
             int i, j, up, tmp, tmp1;
82
             for (i = 0; i < len; i++)
83
                  up = 0;
84
                  for (j = 0; j < b.len; j++)
85
86
                      tmp = a[i] * b.a[j] + t.a[i + j] + up;
87
88
                      if (tmp > 9)
                           tmp1 = tmp - tmp / 10 * 10, up = tmp / 10, t.a[i + j] = tmp1;
89
90
                      else
                          up = 0, t.a[i + j] = tmp;
91
92
                  if (up) t.a[i + j] = up;
93
94
             }
95
             t.len = i + j;
             while (t.a[t.len - 1] == 0 && t.len > 1) t.len--;
96
             return t;
97
         }
98
         //
                bint / int
99
         bint operator/(const int& b) const
100
101
             bint t;
102
             int down = 0;
103
             for (int i = len - 1; ~i; i--)
104
                  t.a[i] = (a[i] + down * 10) / b, down = a[i] + down * 10 - t.a[i] * b;
105
             t.len = len;
106
107
             while (t.a[t.len - 1] == 0 \&\& t.len > 1) t.len--;
108
             return t;
109
         }
         //
                bint ^ n (n ,
110
                                 n )
         bint operator^(const int n) const
111
112
             bint t(*this), rt(1);
113
114
             if (n == 0) return 1;
             if (n == 1) return *this;
115
             int m = n;
116
```

```
for (; m; m >>= 1, t = t * t)
117
                  if (m & 1) rt = rt * t;
118
             return rt;
119
         }
120
         //
121
                bint > bint
         bool operator>(const bint& b) const
122
123
             int p;
124
             if (len > b.len) return 1;
125
             if (len == b.len)
126
127
128
                  p = len - 1;
                  while (a[p] == b.a[p] \&\& p >= 0) p--;
129
                  return p >= 0 && a[p] > b.a[p];
130
             }
131
             return 0;
132
         }
133
         //
                bint > int
134
         bool operator>(const int& n) const { return *this > bint(n); }
135
136
         void out()
137
         {
138
             for (int i = len - 1; ~i; i--) printf("%d", a[i]);
139
140
             puts("");
141
         }
142 };
    7.5 VIM
 1 syntax on
    set cindent
 3
    set nu
 4 set tabstop=4
 5 set shiftwidth=4
 6 set background=dark
 7 set mouse=a
 8
 9 map<C-A> gg\veeG"+\vee
10 map<F5> :call Run()<CR>
11
    func! Run()
12
         exec "w"
13
         exec "!g++ -std=c++11 -02 % -o %<"
14
         exec "!time ./%<"
15
    endfunc
16
17
    autocmd BufNewFile *.cpp Or ~/include.cpp
18
    autocmd BufNewFile *.cpp normal G
19
20
21 inoremap (()<Esc>i
22 inoremap [ []<Esc>i
23 inoremap { {<CR>}<Esc>0
24 inoremap ' ''<Esc>i
25 inoremap " ""<Esc>i
26
   inoremap ) <c-r>=ClosePair(')')<CR>
inoremap ] <c-r>=ClosePair(']')<CR>
27
28
29
```

```
func ClosePair(char)
func
```