

CHEAT SHEET

HACK PROGRAMMING

1 Memory Elements : $\underbrace{D, A}_{\text{on-chip}}, M$ RAM[A] for HACK asm

2 For any operation (ALU), data needs to be in CPU
⇒ otherwise, copy from RAM → CPU
M D

3 Always start with a Pseudocode comprising D, A, M
⇒ Use R0, ..., R15, SCREEN, KBD
for accessing RAM[0], ..., RAM[15], RAM[16,384], RAM[24576]
⇒ Use (LABEL) if code has any Branching

example: if ($D < 0$)
 // jump to TRUE
 ==

else

// FALSE

==

// Remaining code

==

// END with oo loop

0 @TRUE
1 D ; JLT
2 // FALSE
3 ==
4 @REMAINING
5 0 ; JMP
(TRUE)
6 ==
7 ==
8 (REMAINING)
9 ==
(END)
10 @END
11 0 ; JMP

Note: (LABEL) declaration is not assigned any address on ROM (\equiv comment)

\Rightarrow LABEL is stored in Symbol Table with value as "rom" address of HACK asm instruction after declaration
 $\text{@REMAINING} \parallel @8$ i.e. $\boxed{8}$ A

\Rightarrow example: $\text{@TRUE} \parallel @6$ i.e. $\boxed{6}$
refer above $\text{@END} \parallel @10$ i.e. $\boxed{10}$ A

\Rightarrow User-defined variables are mapped based on order of first appearance in asm code

example: sum=0;

@sum
 $M=0$

$i=2;$

@i
 ~~$M=2$~~ sin -1, 0, 1

valid constants as per HACK "comp"

@2
 $D=A$
 @i
 $M=D$

sum=i
in HACK, any operation involving RAM registers has to go through CPU

@i
 $D=M$
 @sum
 $M=D$

$\left. \begin{matrix} \text{CPU} \leftarrow \text{RAM} \\ \text{RAM} \leftarrow \end{matrix} \right\}$

Note: User defined variables are stored in data memory starting RAM[16] onwards

example: Refer above

$\text{@sum} \quad // @16 \text{ i.e. } \boxed{16} \equiv 16 \boxed{\text{sum}}$

$\text{@i} \quad // @17 \text{ i.e. } \boxed{17} \equiv 17 \boxed{\text{i}}$

everytime @sum @i are now used in rest of the code, A-register on CPU will store $\boxed{16}$ $\boxed{17}$ respectively

4 { Keywords (in-built symbols)
Labels
User-defined Variables

are for Programmer's convenience to write HACK asm code

However,

what gets stored in ROM (instruction memory) is the corresponding "numeric values" they are mapped to in Symbol Table

Example:
refer above

@RD
 @sum
 @END

$\boxed{@0}$
 $\boxed{@16}$
 $\boxed{@10}$

ROM view

5

A - register can have 3 interpretations based on the HACK asm statement which immediately follows

example : @20

$D = \underline{A}$ // used as constant

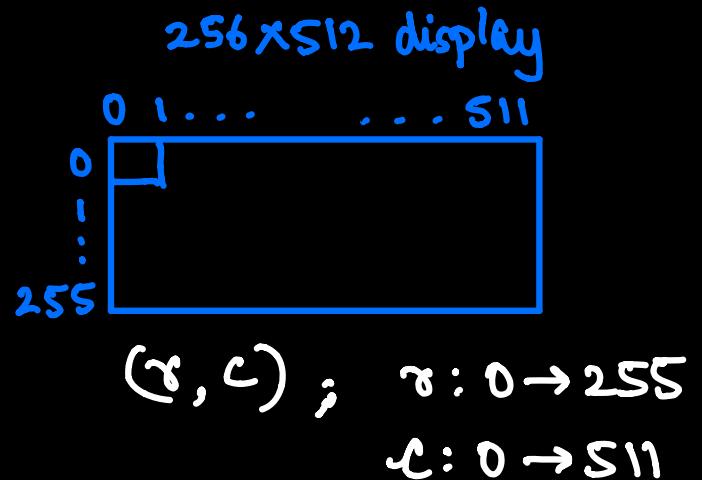
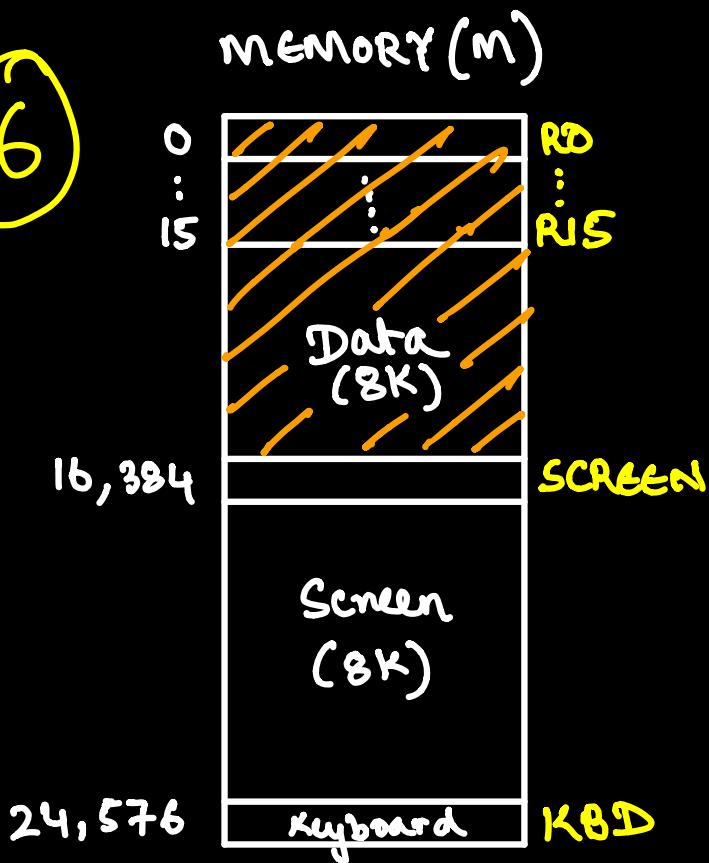
example : @20

$D = \underline{M}$ // used as $\text{RAM}[20]$

example : @20

$D; \underline{\text{JLT}}$ // used as $\text{ROM}[20]$
any jump

6



$\Rightarrow 256 \times 512$ pixels

$$\equiv 2^8 \times 2^9 = 2^{17} \text{ pixels}$$

$$\equiv 2^{17} \text{ bits} \quad \left(\begin{array}{l} 1 \text{ pixel} \equiv \text{B/W} \\ \equiv 2 \text{ values} \\ \equiv 1 \text{ bit} \end{array} \right)$$

$$\equiv \frac{2^{17}}{16} \text{ Registers} \quad (16 \text{ bits} \equiv 1 \text{ Register})$$

$$\equiv 2^{13} \text{ Registers}$$

$$\equiv 2^3 \times 2^{10} \text{ Registers}$$

$$\equiv 8K \text{ Registers} \quad (1K = 2^{10})$$

\equiv RAM 8K chip to implement
Screen Memory

\Rightarrow Each row in Display,

$$\equiv 512 \text{ pixels}$$

$$\equiv 512 \text{ bits}$$

$$\equiv \frac{512}{16} = \frac{2^9}{2^4} = 32 \text{ Registers}$$

i.e. 256 rows each requires 32 Registers

$$\equiv 256 \times 32 \text{ Registers}$$

$$\equiv 2^8 \times 2^5$$

$$\equiv 2^{13} = 8K \text{ Registers} \text{ (again verified)}$$

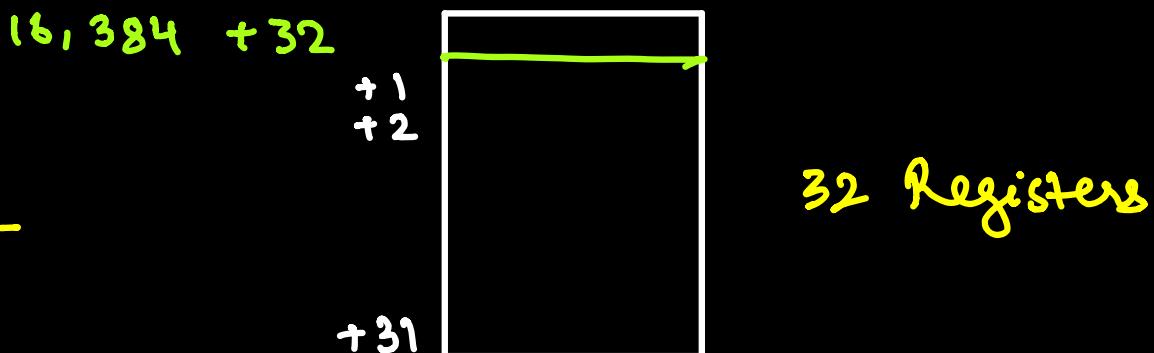
$\Rightarrow (0,0) \rightarrow (0,511) : \text{RAM}[16,384] \text{ or screen onwards}$

$r=0$



$\Rightarrow (1,0) \rightarrow (1,511) : \text{RAM}[16,384+32] \text{ onwards}$

$r=1$



$\Rightarrow (r,0) \rightarrow (r,511) \equiv \text{RAM}[16,384+(32*r)]$
as starting register

⇒ For any 'c' in a given 'r'

$\left\lfloor \frac{c}{16} \right\rfloor$ refers to the register index within the 32 Registers assigned to given 'r'

Example: $(0, \overset{c}{0}) \rightarrow (0, \overset{c}{15})$ pixels

map to register index 0
in $\text{RAM}[16, 384] + 0, \dots, \text{RAM}[16, 384 + 31]$

Example: $(0, \overset{c}{16}) \rightarrow (0, \overset{c}{31})$ pixels

map to register index 1 in
 $\text{RAM}[16, 384], \text{RAM}[16, 384 + 1], \dots, \text{RAM}[16, 384 + 31]$

Example: $(1, \overset{c}{16}) \rightarrow (1, \overset{c}{31})$ pixels map
to register index 1 in

$\text{RAM}[16, 384 + 32 * 1], \text{RAM}[16, 384 + 32 * 1 + 1], \dots, \text{RAM}[16, 384 + 32 * 1 + 31]$

So, for : (r, c) px mapped to

$\text{RAM}[16, 384 + (32 * r) + \left\lfloor \frac{c}{16} \right\rfloor]$

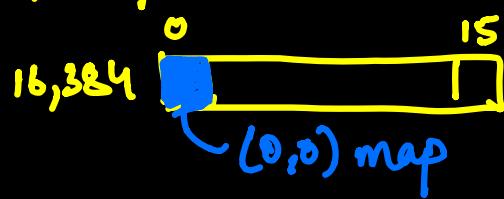
\Rightarrow Remainder of $\frac{C}{16}$ i.e. $(C \% 16)$
 specifies which cell position
(left to right) corresponds to
 (r, c) pixel

Example: $(0,0) \rightarrow (0,15)$

$$\rightarrow \gamma=0 \quad \equiv \quad \text{RAM} [16, 384 + 32*0] \\ \equiv \text{RAM} [16, 384]$$

$$\begin{aligned} \rightarrow c=0 &= \text{RAM}\left[16, 384 + \left\lfloor \frac{0}{16} \right\rfloor\right] \\ &\in \text{RAM}\left[16, 384\right] \end{aligned}$$

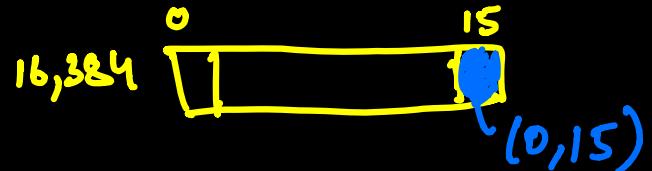
$$\rightarrow C \cdot 10\% = D \cdot 10\% = D$$



$$\rightarrow \mathcal{L} = 15 \equiv \text{RAM}\left[16, 384 + \left\lfloor \frac{15}{16} \right\rfloor \right]$$

| $\equiv \text{RAM}[16, 384]$

$$\rightarrow 15\% \cdot 16 = 15\%$$



7 Summary of (r, c) px mapping onto Screen Memory

(r, c)

$$\equiv \text{RAM} \left[16384 + (32 * r) + \left\lfloor \frac{c}{16} \right\rfloor \right]$$

with cell position $(C_{15} \dots C_0)$ from
Left \rightarrow Right
 (0) (15)

8 Pixel Coloring

Black or White at single pixel
 (1) (0)

Example

$$M = 0 \equiv [0|0|0| \dots |0]$$

16-bit
all corresponding 16 pixels
on screen display colored
WHITE

example:

$$M = -1 \equiv [1|1|1| \dots |1]$$

in 2's complement
16-bit
all corresponding 16 pixels
on screen display colored
BLACK

example: $\begin{array}{l} @3 \\ D = A \\ @SCREEN \\ M = D \end{array}$ } = storing 3 into RAM[16,384]

RAM[16,384]
0 0 0 0 1 1

will color (0,14) and (0,15)
pixels as BLACK
 $(0,0) \rightarrow (0,13)$ will be WHITE