

Untyped Lambda Calculus

Structure and Evaluation, Currying, Church Encodings

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Dream IT

<https://dreamit.de>

Lambda Calculus

- Invented by Alonzo Church (1920s)
- Equally expressive to the Turing Machine(s)
- Formal Language
- Computational Model
 - Lisp (1950s)
 - ML
 - Haskell
- "Lambda Expressions" in almost every modern programming language

Why should I care?

- Simple Computational Model
 - to describe structure and behaviour (E.g. Operational Semantics, Type Systems)
 - to reason and prove

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- Explains why things in FP are like they are
 - pure functions
 - higher-order functions
 - currying
 - lazy evaluation

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- Explains why things in FP are like they are
 - pure functions
 - higher-order functions
 - currying
 - lazy evaluation
- Understand FP Compilers
 - Introduce FP stuff into other languages
 - Write your own compiler
 - GHC uses an enriched Lambda Calculus internally

Untyped Lambda Calculus

$t ::=$

x

$\lambda x. t$

$t \ t$

Terms:

Variable

Abstraction

Application

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Example - Identity

Lambda Calculus

$$\underbrace{\underbrace{\lambda x. x}_{\text{Abstraction}} \underbrace{y}_{\text{Variable}}}_{\text{Application}} \rightarrow y$$

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Example - Identity

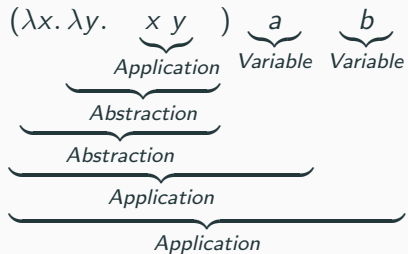
Lambda Calculus

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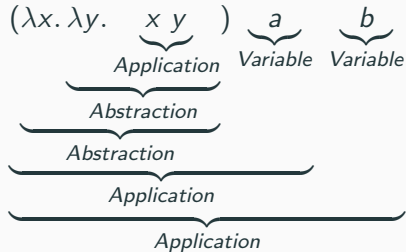
Javascript

$\underbrace{(\underbrace{\text{function } (x) \{ \text{return } x; \}}_{\text{Abstraction}}) (\underbrace{y}_{\text{Variable}})}_{\text{Application}}$

Evaluation / Reduction



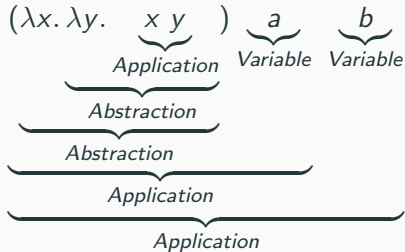
Evaluation / Reduction



$(\lambda x. \lambda y. x y) a b$

Parentheses are not part of the grammar? See next slide... :)

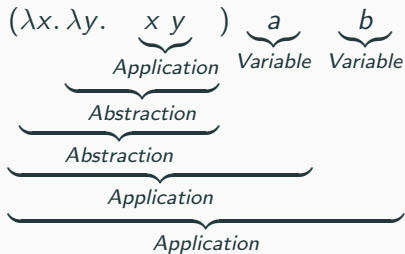
Evaluation / Reduction



$$\begin{aligned} & (\lambda x. \lambda y. x y) a b \\ \rightarrow & (\lambda y. a y) b \end{aligned}$$

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Evaluation / Reduction



$(\lambda x. \lambda y. x y) a b$
 $\rightarrow (\lambda y. a y) b$
 $\rightarrow a b$

Parentheses are not part of the grammar? See next slide... :)

Notational Conventions

- We use parentheses to clarify what's meant
- Applications associate to the left

$$s\ t\ u \equiv (s\ t)\ u$$

- Lambda Expressions expand as much to the right as possible

$$\lambda x. \lambda y. x\ y\ x \equiv \lambda x. (\lambda y. ((x\ y)\ x))$$

$\lambda x. \lambda y. x \ y \ z$

- *Bound and free*

λy y is *bound*, x and z are *free*

λx x and y are *bound*, z is *free*

$\lambda x, \lambda y$ *binder*

- A term with no free variables is "*closed*"
 - A "*combinator*"
 - $id \equiv \lambda x. x$

- We learned how to write down and talk about Lambda Calculus Terms
- How to evaluate them?
- Different Strategies
 - Interesting outcomes

Full Beta-Reduction

- RedEx
 - **R**educible **E**xpression
 - Always an Application

$$\begin{array}{c} (\lambda x.x) \ ((\lambda x.x) \ (\lambda z. \underbrace{(\lambda x.x) \ z}_{\text{RedEx}})) \\ \underbrace{\hspace{10em}}_{\text{RedEx}} \\ \underbrace{\hspace{10em}}_{\text{RedEx}} \end{array}$$

Full Beta-Reduction

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 - **R**educible **E**xpression
 - Always an Application

$$\underbrace{\underbrace{(\lambda x.x) \underbrace{((\lambda x.x) (\lambda z. \underbrace{(\lambda x.x) z}))}_{\text{RedEx}}}}_{\text{RedEx}}}_{\text{RedEx}}$$

- Full Beta-Reduction
 - Any RedEx, Any Time
 - Like in Arithmetics
 - Too vague for programming. . .
 - How to write a good test if the next step could be several expressions?

$$(\lambda x.x) ((\lambda x.x) (\lambda z.(\lambda x.x) z))$$

- Normal Order Reduction
 - Left-most, Outer-most RedEx

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- Normal Order Reduction
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$$(\lambda x.x) ((\lambda x.x) (\lambda z.(\lambda x.x) z))$$

- Call-by-Name
 - like Normal Order Reduction, but NO reductions inside Abstractions
 - lazy, non-strict
 - Parameters are NOT evaluated before they are used
 - Lambdas are values
 - Optimization: Save result \rightarrow *Call-by-Need*

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Higher Order Functions

- Functions that take or return functions
 - Are there "by definition"

$$\underbrace{\underbrace{\lambda x.x}_{\text{Abstraction}} \underbrace{\lambda y.y}_{\text{Abstraction}}}_{\text{Application}} \rightarrow \underbrace{\lambda y.y}_{\text{Abstraction}}$$

$$(\lambda x. \lambda y. x \ y) \ z \rightarrow \lambda y. z \ y$$

- Example
 - (+1) Section in Haskell
 - $(\lambda x. \lambda y. + \ x \ y) \ 1 \rightarrow \lambda y. + \ 1 \ y$
- Partial Application is there "by definition"

- Everything (Term) is an Expression
 - No statements
- No "destructive" Assignments
 - The reason why FP Languages promote pure functions
 - But you could invent a built-in function to manipulate "state"...

Alpha conversion

$$\lambda x.x \rightarrow_{\alpha} \lambda y.y$$

Reductions and Conversions

Alpha conversion

$$\lambda x.x \rightarrow_{\alpha} \lambda y.y$$

Beta reduction

$$(\lambda x.x) y \rightarrow_{\beta} y$$

Reductions and Conversions

Alpha conversion

$$\lambda x.x \rightarrow_{\alpha} \lambda y.y$$

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$$(\lambda x.x) y \rightarrow_{\beta} y$$

Eta conversion

Iff (if and only if) x is not free in f :

$$\lambda x.f \ x \rightarrow_{\eta} f$$

$$(\lambda x.(\lambda y.y) \ x) \ a \rightarrow_{\eta} (\lambda y.y) \ a$$

If x is free in f , no η conversion possible:

$$\lambda x.(\lambda y.y \overset{\text{Bound}}{\downarrow} x) \ x \not\rightarrow_{\eta} (\lambda y.y \overset{\text{Free?!}}{\downarrow} x)$$

- Encode Data into the Lambda Calculus
- To simplify our formulas, let's say that we have declarations

$$id \equiv \lambda x.x$$

$$id\ y \rightarrow y$$

Booleans

true \equiv $\lambda t.\lambda f.t$

false \equiv $\lambda t.\lambda f.f$

Booleans

true \equiv $\lambda t.\lambda f.t$

false \equiv $\lambda t.\lambda f.f$

test \equiv $\lambda c.\lambda t.\lambda f.c\ t\ f$

Booleans

test true a b

true \equiv $\lambda t.\lambda f.t$

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$\rightarrow (\lambda t.\lambda f.true\ t\ f)\ a\ b$

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$\equiv (\lambda t. \lambda f. t) \ a \ b$

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$\rightarrow a$

And

$true \equiv \lambda t. \lambda f. t$

$false \equiv \lambda t. \lambda f. f$

And

$true \equiv \lambda t. \lambda f. t$

$false \equiv \lambda t. \lambda f. f$

$and \equiv \lambda p. \lambda q. p \ q \ p$

And

and true false

true \equiv $\lambda t.\lambda f.t$

false \equiv $\lambda t.\lambda f.f$

and \equiv $\lambda p.\lambda q.p\ q\ p$

And

and true false

true \equiv $\lambda t.\lambda f.t$

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$\rightarrow (\lambda q. true \ q \ true) \ false$

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$\equiv (\lambda t. \lambda f. t) \ false \ true$

$\rightarrow (\lambda f. false) \ true$

$\rightarrow false$

Or

$\lambda p. \lambda q. p \ p \ q$

$$\textit{pair} \equiv \lambda x.\lambda y.\lambda z.z \times y$$

pair $\equiv \lambda x.\lambda y.\lambda z.z\ x\ y$

first $\equiv (\lambda p.p)\ \lambda x.\lambda y.x$

second $\equiv (\lambda p.p)\ \lambda x.\lambda y.y$

Pairs

$pair_{AB} \equiv$

$pair\ a\ b$

$pair \equiv \lambda x.\lambda y.\lambda z.z\ x\ y$

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Pairs

$pair_{AB} \equiv$

pair $a\ b$

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Pairs

$pair \equiv \lambda x.\lambda y.\lambda z.z \ x \ y$

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$second \equiv (\lambda p.p) \ \lambda x.\lambda y.y$

$pair_{AB} \equiv \text{pair} \ a \ b$
 $\equiv (\lambda x.\lambda y.\lambda z.z \ x \ y) \ a \ b$

Pairs

$pair \equiv \lambda x.\lambda y.\lambda z.z \ x \ y$

$first \equiv (\lambda p.p) \ \lambda x.\lambda y.x$

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$pair_{AB} \equiv \text{pair } a \ b$
 $\equiv (\lambda x.\lambda y.\lambda z.z \ x \ y) \ a \ b$

Pairs

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$first \equiv (\lambda p.p) \ \lambda x.\lambda y.x$

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$pair_{AB} \equiv$ $pair$ $a \ b$
 \equiv $(\lambda x.\lambda y.\lambda z.z \ x \ y)$ $a \ b$
 \rightarrow $(\lambda y.\lambda z.z \ a \ y)b$

Pairs

$pair \equiv \lambda x.\lambda y.\lambda z.z \ x \ y$

$first \equiv (\lambda p.p) \ \lambda x.\lambda y.x$

$second \equiv (\lambda p.p) \ \lambda x.\lambda y.y$

$pair_{AB} \equiv$ $pair \ a \ b$
 \equiv $(\lambda x.\lambda y.\lambda z.z \ x \ y) \ a \ b$
 \rightarrow $(\lambda y.\lambda z.z \ a \ y)b$

Pairs

$pair \equiv \lambda x.\lambda y.\lambda z.z \ x \ y$

$first \equiv (\lambda p.p) \ \lambda x.\lambda y.x$

$second \equiv (\lambda p.p) \ \lambda x.\lambda y.y$

$pair_{AB} \equiv$ $pair \ a \ b$
 \equiv $(\lambda x.\lambda y.\lambda z.z \ x \ y) \ a \ b$
 \rightarrow $(\lambda y.\lambda z.z \ a \ y)b$
 \rightarrow $\lambda z.z \ a \ b$

Pairs

$$pair \equiv \lambda x. \lambda y. \lambda z. z \ x \ y$$

$$first \equiv (\lambda p. p) \ \lambda x. \lambda y. x$$

$$second \equiv (\lambda p. p) \ \lambda x. \lambda y. y$$

$$\begin{aligned} pair_{AB} &\equiv pair \ a \ b \\ &\equiv (\lambda x. \lambda y. \lambda z. z \ x \ y) \ a \ b \\ &\rightarrow (\lambda y. \lambda z. z \ a \ y) b \\ &\rightarrow \lambda z. z \ a \ b \\ &\equiv pair'_{ab} \end{aligned}$$

Pairs (continued)

$$pair \equiv \lambda x. \lambda y. \lambda z. z \ x \ y$$

$$first \equiv (\lambda p. p) \ \lambda x. \lambda y. x$$

$$pair'_{ab} \equiv \lambda z. z \ a \ b$$

Pairs (continued)

first *pair'*_{ab}

pair $\equiv \lambda x. \lambda y. \lambda z. z \ x \ y$

first $\equiv (\lambda p. p) \ \lambda x. \lambda y. x$

*pair'*_{ab} $\equiv \lambda z. z \ a \ b$

Pairs (continued)

first *pair'*_{ab}

pair $\equiv \lambda x. \lambda y. \lambda z. z \ x \ y$

first $\equiv (\lambda p. p) \ \lambda x. \lambda y. x$

*pair'*_{ab} $\equiv \lambda z. z \ a \ b$

Pairs (continued)

$$\begin{aligned} pair &\equiv \lambda x. \lambda y. \lambda z. z \ x \ y \\ first &\equiv (\lambda p. p) \ \lambda x. \lambda y. x \\ pair'_{ab} &\equiv \lambda z. z \ a \ b \end{aligned} \qquad \begin{aligned} &first \ pair'_{ab} \\ \equiv &(\lambda p. p) \ (\lambda x. \lambda y. x) \ pair'_{ab} \end{aligned}$$

Pairs (continued)

$pair \equiv \lambda x.\lambda y.\lambda z.z\ x\ y$

$first \equiv (\lambda p.p)\ \lambda x.\lambda y.x$

$pair'_{ab} \equiv \lambda z.z\ a\ b$

$\equiv \text{first}\ pair'_{ab}$
 $\equiv (\lambda p.p)\ (\lambda x.\lambda y.x)\ pair'_{ab}$

Pairs (continued)

$$\begin{aligned} pair &\equiv \lambda x. \lambda y. \lambda z. z \ x \ y \\ first &\equiv (\lambda p. p) \ \lambda x. \lambda y. x \\ pair'_{ab} &\equiv \lambda z. z \ a \ b \end{aligned}$$
$$\begin{aligned} &first \ pair'_{ab} \\ \equiv & \ (\lambda p. p) \ (\lambda x. \lambda y. x) \ pair'_{ab} \\ \rightarrow & \ pair'_{ab} \ (\lambda x. \lambda y. x) \end{aligned}$$

Pairs (continued)

$$\begin{aligned} \text{pair} &\equiv \lambda x. \lambda y. \lambda z. z \ x \ y \\ \text{first} &\equiv (\lambda p. p) \ \lambda x. \lambda y. x \\ \text{pair}'_{ab} &\equiv \lambda z. z \ a \ b \end{aligned}$$
$$\begin{aligned} &\text{first } \text{pair}'_{ab} \\ &\equiv (\lambda p. p) \ (\lambda x. \lambda y. x) \ \text{pair}'_{ab} \\ &\rightarrow \text{pair}'_{ab} \ (\lambda x. \lambda y. x) \end{aligned}$$

Pairs (continued)

$$\begin{aligned} pair &\equiv \lambda x. \lambda y. \lambda z. z \ x \ y \\ first &\equiv (\lambda p. p) \ \lambda x. \lambda y. x \\ pair'_{ab} &\equiv \lambda z. z \ a \ b \end{aligned} \qquad \begin{aligned} &first \ pair'_{ab} \\ \equiv &(\lambda p. p) \ (\lambda x. \lambda y. x) \ pair'_{ab} \\ \rightarrow &pair'_{ab} \ (\lambda x. \lambda y. x) \\ \equiv &(\lambda z. z \ a \ b) \ (\lambda x. \lambda y. x) \end{aligned}$$

Pairs (continued)

$$\begin{aligned} pair &\equiv \lambda x. \lambda y. \lambda z. z \ x \ y \\ first &\equiv (\lambda p. p) \ \lambda x. \lambda y. x \\ pair'_{ab} &\equiv \lambda z. z \ a \ b \end{aligned} \qquad \begin{aligned} &first \ pair'_{ab} \\ &\equiv (\lambda p. p) \ (\lambda x. \lambda y. x) \ pair'_{ab} \\ &\rightarrow pair'_{ab} \ (\lambda x. \lambda y. x) \\ &\equiv (\lambda z. z \ a \ b) \ (\lambda x. \lambda y. x) \end{aligned}$$

Pairs (continued)

$$\begin{aligned} pair &\equiv \lambda x. \lambda y. \lambda z. z \ x \ y \\ first &\equiv (\lambda p. p) \ \lambda x. \lambda y. x \\ pair'_{ab} &\equiv \lambda z. z \ a \ b \end{aligned} \qquad \begin{aligned} &first \ pair'_{ab} \\ \equiv &(\lambda p. p) \ (\lambda x. \lambda y. x) \ pair'_{ab} \\ \rightarrow &pair'_{ab} \ (\lambda x. \lambda y. x) \\ \equiv &(\lambda z. z \ a \ b) \ (\lambda x. \lambda y. x) \\ \rightarrow &(\lambda x. \lambda y. x) \ a \ b \end{aligned}$$

Pairs (continued)

$$\begin{aligned} pair &\equiv \lambda x. \lambda y. \lambda z. z \ x \ y \\ first &\equiv (\lambda p. p) \ \lambda x. \lambda y. x \\ pair'_{ab} &\equiv \lambda z. z \ a \ b \end{aligned} \qquad \begin{aligned} &first \ pair'_{ab} \\ \equiv &(\lambda p. p) \ (\lambda x. \lambda y. x) \ pair'_{ab} \\ \rightarrow &pair'_{ab} \ (\lambda x. \lambda y. x) \\ \equiv &(\lambda z. z \ a \ b) \ (\lambda x. \lambda y. x) \\ \rightarrow &(\lambda x. \lambda y. x) \ a \ b \end{aligned}$$

Pairs (continued)

$$\begin{aligned} pair &\equiv \lambda x. \lambda y. \lambda z. z \ x \ y \\ first &\equiv (\lambda p. p) \ \lambda x. \lambda y. x \\ pair'_{ab} &\equiv \lambda z. z \ a \ b \end{aligned} \qquad \begin{aligned} &first \ pair'_{ab} \\ \equiv &(\lambda p. p) \ (\lambda x. \lambda y. x) \ pair'_{ab} \\ \rightarrow &pair'_{ab} \ (\lambda x. \lambda y. x) \\ \equiv &(\lambda z. z \ a \ b) \ (\lambda x. \lambda y. x) \\ \rightarrow &(\lambda x. \lambda y. x) \ a \ b \\ \rightarrow &(\lambda y. a) \ b \end{aligned}$$

Pairs (continued)

$$\begin{aligned} pair &\equiv \lambda x. \lambda y. \lambda z. z \ x \ y \\ first &\equiv (\lambda p. p) \ \lambda x. \lambda y. x \\ pair'_{ab} &\equiv \lambda z. z \ a \ b \end{aligned} \qquad \begin{aligned} &first \ pair'_{ab} \\ \equiv & (\lambda p. p) \ (\lambda x. \lambda y. x) \ pair'_{ab} \\ \rightarrow & pair'_{ab} \ (\lambda x. \lambda y. x) \\ \equiv & (\lambda z. z \ a \ b) \ (\lambda x. \lambda y. x) \\ \rightarrow & (\lambda x. \lambda y. x) \ a \ b \\ \rightarrow & (\lambda y. a) \ b \end{aligned}$$

Pairs (continued)

$$\begin{aligned} pair &\equiv \lambda x. \lambda y. \lambda z. z \ x \ y \\ first &\equiv (\lambda p. p) \ \lambda x. \lambda y. x \\ pair'_{ab} &\equiv \lambda z. z \ a \ b \end{aligned} \qquad \begin{aligned} &first \ pair'_{ab} \\ \equiv &(\lambda p. p) \ (\lambda x. \lambda y. x) \ pair'_{ab} \\ \rightarrow &pair'_{ab} \ (\lambda x. \lambda y. x) \\ \equiv &(\lambda z. z \ a \ b) \ (\lambda x. \lambda y. x) \\ \rightarrow &(\lambda x. \lambda y. x) \ a \ b \\ \rightarrow &(\lambda y. a) \ b \\ \rightarrow &a \end{aligned}$$

Numerals

Peano axioms

Every natural number can be defined with 0 and a successor function

$$0 \equiv \lambda f. \lambda x. x$$

$$1 \equiv \lambda f. \lambda x. f \ x$$

$$2 \equiv \lambda f. \lambda x. f \ (f \ x)$$

$$3 \equiv \lambda f. \lambda x. f \ (f \ (f \ x))$$

Meaning

0 f is evaluated 0 times

1 f is evaluated once

x can be every lambda term

Numerals Example - Successor

$$0 \equiv \lambda f. \lambda x. x$$

$$1 \equiv \lambda f. \lambda x. f \ x$$

Numerals Example - Successor

$$0 \equiv \lambda f. \lambda x. x$$

$$1 \equiv \lambda f. \lambda x. f \ x$$

$$\textit{successor} \equiv \lambda n. \lambda f. \lambda x. f \ (n \ f \ x)$$

Numerals Example - Successor

successor 1

0 \equiv $\lambda f.\lambda x.x$

1 \equiv $\lambda f.\lambda x.f\ x$

successor \equiv $\lambda n.\lambda f.\lambda x.f\ (n\ f\ x)$

Numerals Example - Successor

successor 1

$0 \equiv \lambda f. \lambda x. x$

$1 \equiv \lambda f. \lambda x. f \ x$

$\text{successor} \equiv \lambda n. \lambda f. \lambda x. f \ (n \ f \ x)$

Numerals Example - Successor

$$\begin{aligned} 0 &\equiv \lambda f. \lambda x. x && \equiv && \text{successor } 1 \\ 1 &\equiv \lambda f. \lambda x. f \ x && && (\lambda n. \lambda f. \lambda x. f \ (n \ f \ x)) \ 1 \end{aligned}$$

$$\text{successor} \equiv \lambda n. \lambda f. \lambda x. f \ (n \ f \ x)$$

Numerals Example - Successor

$$\begin{aligned} 0 &\equiv \lambda f. \lambda x. x && \equiv \text{successor } 1 \\ 1 &\equiv \lambda f. \lambda x. f \ x && \equiv (\lambda n. \lambda f. \lambda x. f \ (n \ f \ x)) \ 1 \end{aligned}$$

$$\text{successor} \equiv \lambda n. \lambda f. \lambda x. f \ (n \ f \ x)$$

Numerals Example - Successor

$$\begin{array}{llll} 0 \equiv & \lambda f. \lambda x. x & \equiv & \text{successor } 1 \\ 1 \equiv & \lambda f. \lambda x. f \ x & \rightarrow & (\lambda n. \lambda f. \lambda x. f \ (n \ f \ x)) \ 1 \\ & & & \lambda f. \lambda x. f \ (1 \ f \ x) \end{array}$$

$$\text{successor} \equiv \lambda n. \lambda f. \lambda x. f \ (n \ f \ x)$$

Numerals Example - Successor

$$\begin{array}{llll} & & & \text{successor } 1 \\ 0 \equiv & \lambda f. \lambda x. x & \equiv & (\lambda n. \lambda f. \lambda x. f (n f x)) \text{ } 1 \\ 1 \equiv & \lambda f. \lambda x. f x & \rightarrow & \lambda f. \lambda x. f (\text{ } 1 f x) \end{array}$$

$$\text{successor} \equiv \lambda n. \lambda f. \lambda x. f (n f x)$$

Note

We use *Normal Order Reduction* to reduce inside abstractions!

Numerals Example - Successor

$$\begin{aligned} 0 &\equiv \lambda f. \lambda x. x && \equiv \text{successor } 1 \\ 1 &\equiv \lambda f. \lambda x. f \ x && \rightarrow \lambda f. \lambda x. f \ (\textcolor{brown}{1} \ f \ x) \\ &&& \equiv \lambda f. \lambda x. f \ ((\lambda f. \lambda x. f \ x) \ f \ x) \end{aligned}$$

$$\text{successor} \equiv \lambda n. \lambda f. \lambda x. f \ (n \ f \ x)$$

Note

We use *Normal Order Reduction* to reduce inside abstractions!

Numerals Example - Successor

$$\begin{aligned} 0 &\equiv \lambda f. \lambda x. x && \equiv \text{successor } 1 \\ 1 &\equiv \lambda f. \lambda x. f \ x && \rightarrow (\lambda n. \lambda f. \lambda x. f \ (n \ f \ x)) \ 1 \\ &&& \equiv \lambda f. \lambda x. f \ ((\lambda f. \lambda x. f \ x) \ f \ x) \end{aligned}$$

$$\text{successor} \equiv \lambda n. \lambda f. \lambda x. f \ (n \ f \ x)$$

Note

We use *Normal Order Reduction* to reduce inside abstractions!

Numerals Example - Successor

$$\begin{aligned} 0 &\equiv \lambda f. \lambda x. x && \equiv \text{successor } 1 \\ 1 &\equiv \lambda f. \lambda x. f \ x && \rightarrow (\lambda n. \lambda f. \lambda x. f \ (n \ f \ x)) \ 1 \\ &&& \rightarrow \lambda f. \lambda x. f \ (1 \ f \ x) \\ &&& \equiv \lambda f. \lambda x. f \ ((\lambda f. \lambda x. f \ x) \ f \ x) \\ &&& \rightarrow \lambda f. \lambda x. f \ ((\lambda x. f \ x) \ x) \\ \text{successor} &\equiv \lambda n. \lambda f. \lambda x. f \ (n \ f \ x) \end{aligned}$$

Note

We use *Normal Order Reduction* to reduce inside abstractions!

Numerals Example - Successor

$$\begin{aligned} 0 &\equiv \lambda f. \lambda x. x && \equiv \text{successor } 1 \\ 1 &\equiv \lambda f. \lambda x. f \ x && \rightarrow (\lambda n. \lambda f. \lambda x. f \ (n \ f \ x)) \ 1 \\ &&& \rightarrow \lambda f. \lambda x. f \ (1 \ f \ x) \\ &&& \equiv \lambda f. \lambda x. f \ ((\lambda f. \lambda x. f \ x) \ f \ x) \\ &&& \rightarrow \lambda f. \lambda x. f \ ((\lambda x. f \ x) \ x) \\ \text{successor} &\equiv \lambda n. \lambda f. \lambda x. f \ (n \ f \ x) \end{aligned}$$

Note

We use *Normal Order Reduction* to reduce inside abstractions!

Numerals Example - Successor

$$\begin{aligned} 0 &\equiv \lambda f. \lambda x. x && \equiv \text{successor } 1 \\ 1 &\equiv \lambda f. \lambda x. f \ x && \rightarrow (\lambda n. \lambda f. \lambda x. f \ (n \ f \ x)) \ 1 \\ &&& \rightarrow \lambda f. \lambda x. f \ (1 \ f \ x) \\ &&& \equiv \lambda f. \lambda x. f \ ((\lambda f. \lambda x. f \ x) \ f \ x) \\ &&& \rightarrow \lambda f. \lambda x. f \ ((\lambda x. f \ x) \ x) \\ \text{successor} &\equiv \lambda n. \lambda f. \lambda x. f \ (n \ f \ x) && \rightarrow \lambda f. \lambda x. f \ (f \ x) \end{aligned}$$

Note

We use *Normal Order Reduction* to reduce inside abstractions!

Numerals Example - Successor

$$\begin{aligned} 0 &\equiv \lambda f. \lambda x. x && \equiv \text{successor } 1 \\ 1 &\equiv \lambda f. \lambda x. f \ x && \rightarrow (\lambda n. \lambda f. \lambda x. f \ (n \ f \ x)) \ 1 \\ &&& \rightarrow \lambda f. \lambda x. f \ (\textcolor{brown}{1} \ f \ x) \\ &&& \equiv \lambda f. \lambda x. f \ ((\textcolor{brown}{\lambda f. \lambda x. f \ x}) \ f \ x) \\ &&& \rightarrow \lambda f. \lambda x. f \ ((\textcolor{brown}{\lambda x. f \ x}) \ x) \\ \text{successor} &\equiv \lambda n. \lambda f. \lambda x. f \ (n \ f \ x) && \rightarrow \lambda f. \lambda x. f \ (f \ x) \\ &&& \equiv 2 \end{aligned}$$

Note

We use *Normal Order Reduction* to reduce inside abstractions!

Numerals Example - $0 + 0$

$0 \equiv$

$\lambda f.\lambda x.x$

Numerals Example - $0 + 0$

$$0 \equiv \lambda f. \lambda x. x$$

$$plus \equiv \lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)$$

Numerals Example - $0 + 0$

plus 0 0

$0 \equiv \lambda f. \lambda x. x$

plus $\equiv \lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)$

Numerals Example - $0 + 0$

plus 0 0

$0 \equiv \lambda f. \lambda x. x$

$plus \equiv \lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)$

Numerals Example - $0 + 0$

plus 0 0

$\equiv (\lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)) \ 0 \ 0$

$0 \equiv \lambda f. \lambda x. x$

$plus \equiv \lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)$

Numerals Example - $0 + 0$

plus 0 0

$\equiv (\lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)) \ 0 \ 0$

$0 \equiv \lambda f. \lambda x. x$

plus $\equiv \lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)$

Numerals Example - $0 + 0$

plus 0 0

$\equiv (\lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)) \ 0 \ 0$

$\rightarrow (\lambda n. \lambda f. \lambda x. 0 \ f \ (n \ f \ x)) \ 0$

$0 \equiv \lambda f. \lambda x. x$

plus $\equiv \lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)$

Numerals Example - $0 + 0$

plus 0 0

$\equiv (\lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)) \ 0 \ 0$

$\rightarrow (\lambda n. \lambda f. \lambda x. 0 \ f \ (n \ f \ x)) \ 0$

$0 \equiv \lambda f. \lambda x. x$

plus $\equiv \lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)$

Numerals Example - $0 + 0$

plus 0 0

$\equiv (\lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)) \ 0 \ 0$

$\rightarrow (\lambda n. \lambda f. \lambda x. 0 \ f \ (n \ f \ x)) \ 0$

$\rightarrow \lambda f. \lambda x. 0 \ f \ (0 \ f \ x)$

$0 \equiv \lambda f. \lambda x. x$

plus $\equiv \lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)$

Numerals Example - $0 + 0$

plus 0 0

$\equiv (\lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)) \ 0 \ 0$

$\rightarrow (\lambda n. \lambda f. \lambda x. 0 \ f \ (n \ f \ x)) \ 0$

$\rightarrow \lambda f. \lambda x. 0 \ f \ (0 \ f \ x)$

$0 \equiv \lambda f. \lambda x. x$

plus $\equiv \lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)$

Numerals Example - $0 + 0$

$$\begin{aligned} 0 &\equiv \lambda f.\lambda x.x \\ &\equiv \text{plus } 0 \ 0 \\ &\equiv (\lambda m.\lambda n.\lambda f.\lambda x.m \ f \ (n \ f \ x)) \ 0 \ 0 \\ &\rightarrow (\lambda n.\lambda f.\lambda x.0 \ f \ (n \ f \ x)) \ 0 \\ &\rightarrow \lambda f.\lambda x.0 \ f \ (0 \ f \ x) \\ 0 &\equiv \lambda f.\lambda x.x \quad \equiv \lambda f.\lambda x.(\lambda f.\lambda x.x) \ f \ (0 \ f \ x) \end{aligned}$$

$$\text{plus} \equiv \lambda m.\lambda n.\lambda f.\lambda x.m \ f \ (n \ f \ x)$$

Numerals Example - $0 + 0$

$$\begin{aligned} 0 &\equiv \text{plus } 0 \ 0 \\ &\equiv (\lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)) \ 0 \ 0 \\ &\rightarrow (\lambda n. \lambda f. \lambda x. 0 \ f \ (n \ f \ x)) \ 0 \\ &\rightarrow \lambda f. \lambda x. 0 \ f \ (0 \ f \ x) \\ 0 &\equiv \lambda f. \lambda x. x \equiv \lambda f. \lambda x. (\lambda f. \lambda x. x) \ f \ (0 \ f \ x) \end{aligned}$$

$$\text{plus} \equiv \lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)$$

Numerals Example - $0 + 0$

$$\begin{aligned} 0 &\equiv \lambda f.\lambda x.x \\ plus &\equiv \lambda m.\lambda n.\lambda f.\lambda x.m\ f\ (n\ f\ x) \end{aligned}$$
$$\begin{aligned} &\equiv plus\ 0\ 0 \\ &\equiv (\lambda m.\lambda n.\lambda f.\lambda x.m\ f\ (n\ f\ x))\ 0\ 0 \\ &\rightarrow (\lambda n.\lambda f.\lambda x.0\ f\ (n\ f\ x))\ 0 \\ &\rightarrow \lambda f.\lambda x.0\ f\ (0\ f\ x) \\ &\equiv \lambda f.\lambda x.(\lambda f.\lambda x.x)\ f\ (0\ f\ x) \\ &\rightarrow \lambda f.\lambda x.(\lambda x.x)\ (0\ f\ x) \end{aligned}$$

Numerals Example - $0 + 0$

$$\begin{aligned} 0 &\equiv \lambda f.\lambda x.x \\ plus &\equiv \lambda m.\lambda n.\lambda f.\lambda x.m\ f\ (n\ f\ x) \end{aligned}$$
$$\begin{aligned} &\equiv plus\ 0\ 0 \\ &\equiv (\lambda m.\lambda n.\lambda f.\lambda x.m\ f\ (n\ f\ x))\ 0\ 0 \\ &\rightarrow (\lambda n.\lambda f.\lambda x.0\ f\ (n\ f\ x))\ 0 \\ &\rightarrow \lambda f.\lambda x.0\ f\ (0\ f\ x) \\ &\equiv \lambda f.\lambda x.(\lambda f.\lambda x.x)\ f\ (0\ f\ x) \\ &\rightarrow \lambda f.\lambda x.(\lambda x.x)\ (0\ f\ x) \end{aligned}$$

Numerals Example - $0 + 0$

$$\begin{aligned} & \text{plus } 0 \ 0 \\ & \equiv (\lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)) \ 0 \ 0 \\ & \rightarrow (\lambda n. \lambda f. \lambda x. 0 \ f \ (n \ f \ x)) \ 0 \\ & \rightarrow \lambda f. \lambda x. 0 \ f \ (0 \ f \ x) \\ 0 & \equiv \lambda f. \lambda x. x \\ & \equiv \lambda f. \lambda x. (\lambda f. \lambda x. x) \ f \ (0 \ f \ x) \\ & \rightarrow \lambda f. \lambda x. (\lambda x. x) \ (0 \ f \ x) \\ \text{plus} & \equiv \lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x) \rightarrow \lambda f. \lambda x. 0 \ f \ x \end{aligned}$$

Numerals Example - $0 + 0$

$$\begin{aligned} & \text{plus } 0 \ 0 \\ & \equiv (\lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)) \ 0 \ 0 \\ & \rightarrow (\lambda n. \lambda f. \lambda x. 0 \ f \ (n \ f \ x)) \ 0 \\ & \rightarrow \lambda f. \lambda x. 0 \ f \ (0 \ f \ x) \\ 0 & \equiv \lambda f. \lambda x. x \\ & \equiv \lambda f. \lambda x. (\lambda f. \lambda x. x) \ f \ (0 \ f \ x) \\ & \rightarrow \lambda f. \lambda x. (\lambda x. x) \ (0 \ f \ x) \\ \text{plus} & \equiv \lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x) \rightarrow \lambda f. \lambda x. 0 \ f \ x \end{aligned}$$

Numerals Example - $0 + 0$

$$\begin{aligned} 0 &\equiv \lambda f.\lambda x.x \\ plus &\equiv \lambda m.\lambda n.\lambda f.\lambda x.m\ f\ (n\ f\ x) \\ 0 + 0 &\equiv plus\ 0\ 0 \\ &\equiv (\lambda m.\lambda n.\lambda f.\lambda x.m\ f\ (n\ f\ x))\ 0\ 0 \\ &\rightarrow (\lambda n.\lambda f.\lambda x.0\ f\ (n\ f\ x))\ 0 \\ &\rightarrow \lambda f.\lambda x.0\ f\ (0\ f\ x) \\ &\equiv \lambda f.\lambda x.(\lambda f.\lambda x.x)\ f\ (0\ f\ x) \\ &\rightarrow \lambda f.\lambda x.(\lambda x.x)\ (0\ f\ x) \\ &\rightarrow \lambda f.\lambda x.0\ f\ x \\ &\equiv \lambda f.\lambda x.(\lambda f.\lambda x.x)\ f\ x \end{aligned}$$

Numerals Example - $0 + 0$

$$\begin{aligned} 0 &\equiv \lambda f.\lambda x.x \\ plus &\equiv \lambda m.\lambda n.\lambda f.\lambda x.m\ f\ (n\ f\ x) \\ plus\ 0\ 0 &\equiv (\lambda m.\lambda n.\lambda f.\lambda x.m\ f\ (n\ f\ x))\ 0\ 0 \\ &\rightarrow (\lambda n.\lambda f.\lambda x.0\ f\ (n\ f\ x))\ 0 \\ &\rightarrow \lambda f.\lambda x.0\ f\ (0\ f\ x) \\ &\equiv \lambda f.\lambda x.(\lambda f.\lambda x.x)\ f\ (0\ f\ x) \\ &\rightarrow \lambda f.\lambda x.(\lambda x.x)\ (0\ f\ x) \\ &\rightarrow \lambda f.\lambda x.0\ f\ x \\ &\equiv \lambda f.\lambda x.(\lambda f.\lambda x.x)\ f\ x \end{aligned}$$

Numerals Example - $0 + 0$

$$\begin{aligned} & \text{plus } 0 \ 0 \\ & \equiv (\lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)) \ 0 \ 0 \\ & \rightarrow (\lambda n. \lambda f. \lambda x. 0 \ f \ (n \ f \ x)) \ 0 \\ & \rightarrow \lambda f. \lambda x. 0 \ f \ (0 \ f \ x) \\ 0 & \equiv \lambda f. \lambda x. x \\ & \equiv \lambda f. \lambda x. (\lambda f. \lambda x. x) \ f \ (0 \ f \ x) \\ & \rightarrow \lambda f. \lambda x. (\lambda x. x) \ (0 \ f \ x) \\ \text{plus} & \equiv \lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x) \rightarrow \lambda f. \lambda x. 0 \ f \ x \\ & \equiv \lambda f. \lambda x. (\lambda f. \lambda x. x) \ f \ x \\ & \rightarrow \lambda f. \lambda x. (\lambda x. x) \ x \end{aligned}$$

Numerals Example - $0 + 0$

$$\begin{aligned} & \text{plus } 0 \ 0 \\ & \equiv (\lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)) \ 0 \ 0 \\ & \rightarrow (\lambda n. \lambda f. \lambda x. 0 \ f \ (n \ f \ x)) \ 0 \\ & \rightarrow \lambda f. \lambda x. 0 \ f \ (0 \ f \ x) \\ 0 & \equiv \lambda f. \lambda x. x \\ & \equiv \lambda f. \lambda x. (\lambda f. \lambda x. x) \ f \ (0 \ f \ x) \\ & \rightarrow \lambda f. \lambda x. (\lambda x. x) \ (0 \ f \ x) \\ \text{plus} & \equiv \lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x) \rightarrow \lambda f. \lambda x. 0 \ f \ x \\ & \equiv \lambda f. \lambda x. (\lambda f. \lambda x. x) \ f \ x \\ & \rightarrow \lambda f. \lambda x. (\lambda x. x) \ x \end{aligned}$$

Numerals Example - $0 + 0$

$$\begin{aligned}
 & \text{plus } 0 \ 0 \\
 & \equiv (\lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)) \ 0 \ 0 \\
 & \rightarrow (\lambda n. \lambda f. \lambda x. 0 \ f \ (n \ f \ x)) \ 0 \\
 & \rightarrow \lambda f. \lambda x. 0 \ f \ (0 \ f \ x) \\
 0 & \equiv \lambda f. \lambda x. x \equiv \lambda f. \lambda x. (\lambda f. \lambda x. x) \ f \ (0 \ f \ x) \\
 & \rightarrow \lambda f. \lambda x. (\lambda x. x) \ (0 \ f \ x) \\
 \text{plus} & \equiv \lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x) \rightarrow \lambda f. \lambda x. 0 \ f \ x \\
 & \equiv \lambda f. \lambda x. (\lambda f. \lambda x. x) \ f \ x \\
 & \rightarrow \lambda f. \lambda x. (\lambda x. x) \ x \\
 & \rightarrow \lambda f. \lambda x. x
 \end{aligned}$$

Numerals Example - $0 + 0$

$$\begin{aligned}
 & \text{plus } 0 \ 0 \\
 & \equiv (\lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x)) \ 0 \ 0 \\
 & \rightarrow (\lambda n. \lambda f. \lambda x. 0 \ f \ (n \ f \ x)) \ 0 \\
 & \rightarrow \lambda f. \lambda x. 0 \ f \ (0 \ f \ x) \\
 0 & \equiv \lambda f. \lambda x. x \\
 & \equiv \lambda f. \lambda x. (\lambda f. \lambda x. x) \ f \ (0 \ f \ x) \\
 & \rightarrow \lambda f. \lambda x. (\lambda x. x) \ (0 \ f \ x) \\
 \text{plus} & \equiv \lambda m. \lambda n. \lambda f. \lambda x. m \ f \ (n \ f \ x) \rightarrow \lambda f. \lambda x. 0 \ f \ x \\
 & \equiv \lambda f. \lambda x. (\lambda f. \lambda x. x) \ f \ x \\
 & \rightarrow \lambda f. \lambda x. (\lambda x. x) \ x \\
 & \rightarrow \lambda f. \lambda x. x \\
 & \equiv 0
 \end{aligned}$$

The implementation of programming languages Type Systems

Thanks

- Hope you enjoyed this talk and learned something new.
- Hope it wasn't too much math and dusty formulas ... :)