Aggregate Methoditoso & sprand ad Illow nothership to

-> In the aggregate method, we analyze the total cost for a sequence of operations and then divide it by the no. of operations to get the average cost (amortized Cost) per operation que (1) 0 south

1) Inserting an element without resizing; for each invertion where prostructed to feed out to programme of the form of the table, the Cost is Constant. O(1)

2) Resizing: Every time the totale doubles in size, the cost is proportional. to the no. of elements being copied to the new table if the table size is k before doubling, copying only k elements takes O(k), times from

Total Cost of n Insertions; (1)0 efros ti, anaggod prister to ments; The cost of inserting n elements; O(n), for the linestions is too

the cost of resigning: The first doubling involves Copying lelement, the next involves Copying 2, then 4 and Go-on.

O(1+2+4+8+ - 5-12k) NO(n)

so, using aggregate method, the amortized time Complexity

Descenting Methodison bound by the Sound constant of points of account for the costs of tuture expensive operations.

Surparkith Redisherty 1) Assigning Credits;

Each insertion will be charged 3 creditetted stopping to

> 2 Credits for the insertion it self which pays constant the overage lost (amortierd lost) per inoitary que (1) 0, smit

=> To Credit to help pay for the Cost of future resizing operation.

2. Cost of insertion; reduce doubles in the prising is the sold of the series of insertion;

insertion: (1) o mand thomas of the pringer fullow the

but since it costs o(c) for copying k elements. but since we have I credit soved for each previous insertion,

=> when the table doubles the cost of copying elements doubles as well.

insection, for n insections, resulting in 3 n Credits with asset is

=> Each resizing is covered by the Soved Credits. The total no. of cesizing operations is propostionale to them no itational oblings of noitoesso was it its propostionale to them no itational buildings (about log in times) wings a subst to stee set in toward

final Amorttzed cost;

- Inserting n elements costs I credit each.
- · Total exedits collected 3n
- · Cost of each resizing operation is orbready covered by Sowed

-Amortized Cost per Insertion: $= \frac{3n}{n}$

$$=\frac{3n}{n}$$

.. O(1) is the amostized time complexity for inserting n elements using a courting method for a dynamic bable that doubles in size.

2 to a fix