

2. Patricia

- o Node : (NodeLabel:BitNumber:Data)
 - NodeLabel is used to show the link of pointer to node.
- 3. This problem can be solved using suffix tree. Let S and T be given two strings.

First, combine the two strings into one as follows: ST=S#T\$, where special characters # and \$ are appended to each string.

Second, construct suffix tree of concatenated string ST. This takes linear time, O(m+n)

While constructing suffix tree, we add information, length, substring to each internal node. The LCS is a substring with maximum length.

Example)
$$S = "abc", T = "bcd"$$

ST=abc#bcd\$

S1=abc#bcd\$

S2=bc#bcd\$

S3=c#bcd\$

S4=#bcd\$

S5=bcd\$

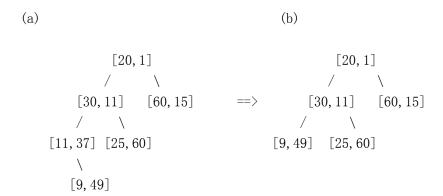
S6=cd\$

S7=d\$

S8=\$

LCS: "bc"

4.



5. 2D range tree

Assume we have N records with 2 keys, x and y.

A 2-D range tree is a binary tree using the range of x, while in each node we keep a sorted list based on y. The tree always splits the records in half, using a median x key value as the discriminator.

Preprocessing time P:

we need to build a sorted list on y, and another on x. Since we have the sorted list on x, we can find the discriminator easily. And since we have the sorted list on y, we can split the list in linear time on each level (there are logN levels overall). So the preprocessing time is:

$$P = NlogN + NlogN + N*logN = O(NlogN)$$

Space required S:

In each node we store an array of the records. At each level the number of records is N. So the space required is:

$$S = N*logN = O(NlogN)$$

Query time Q:

Upon each query, we start at the root, decide which branch to take by comparing the boundaries of x with the discriminator. Once we arrive a node, we use binary search on y in the array. Then

we report those records that fit the query. So the query time is:

$$Q = logN*logN + F = O((logN)^2 + F)$$

where F is the number of records reported.