Instructor: Dr. Sartaj Sahni Spring, 2005

Advanced Data Structures (COP5536) **Exam 01 - Make-up**

CLOSED BOOK 50 Minutes

Name:

PLEASE READ THE FOLLOWING INSTRUCTIONS CAREFULLY

- 1. For all problems, use only the algorithms discussed in class/text.
- 2. Write your name at the top of every exam sheet.
- 3. Write your answers directly on the exam question sheet. You may use scrap paper for work, but these will not be graded.
- 4. All answers will be graded on correctness, efficiency, clarity, elegance and other normal criteria that determine quality.
- 5. The points assigned to each question are provided in parentheses.
- 6. You may use only a pen or a pencil. No calculators allowed.
- 7. Do not write on the reverse side of the exam sheet.
- 8. Do not write close to the margins since those areas do not always make it through when faxed.

Name:

1. (10) Consider two operations put(x) and get(n) on an initially empty jar. put(x) pushes a ball x into the jar and get(n) pulls out n balls from the jar (If the number of balls b in the jar is less than n, just pull out b balls). A put(x) operation takes O(n) time and a get(n) operation takes O(n) time.

Show that the amortized complexity of the put and get operations is O(1).

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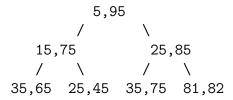
2. (10) You are given 8 runs with 100, 200, 300, 400, 500, 600, 700, and 800 equal-length records. The block size is 100 records. The runs are to be merged using either an optimal 4-way or 8-way merge scheme. Assume that each merge is done using a loser tree.

Determine the number of comparisons and the number of disk I/Os for both merge schemes. Which scheme do you recommend when all input, output, and CPU processing are sequential?

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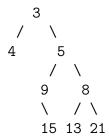
3. (10) For the interval heap,



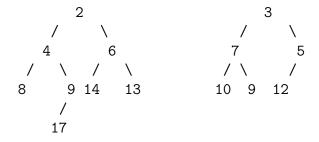
- (a) (4) Insert 90 into the interval tree, showing steps (Use the algorithm discussed in class).
- (b) (6) Perform *RemoveMin* from the original interval heap above, showing each step (Use the algorithm discussed in class).

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- 4. (10) Consider a height-biased min leftist tree:
 - (a) (4) Convert the following min tree to a height-biased min leftist tree and label each node x with its shortest(x) value. Do this by swapping left and right subtrees as needed.



(b) (6) Draw the min leftist tree that results from when the *combine* operation is performed on the two min leftist trees. Show each step.



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- 5. (10) For the min bionomial heap,
 - (a) (4) Insert the keys in sequence: 6, 2, 5, 13, 10, 8, 3, 9, 1, 12, and 4 into an initially empty min bionomial heap (Use the algorithm discussed in class). Show the resulting min bionomial heap.
 - (b) (6) Perform *RemoveMin* on the tree of (a) (Use the algorithm discussed in class), showing each step.

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