

NAME (from your UF ID): \_\_\_\_\_ UF ID#: \_\_\_\_\_  
(Please **PRINT**)

----- CEN 4072/6070 Software Testing & Verification -----

Quiz 3 -- Spring 2017

You have 30 minutes to work on this quiz. It is a "closed-book/closed-notes" test. Pay attention to point values, since you may not have time to work all 3 problems.

PRINT your name and UF ID# above NOW and sign the pledge at the bottom of this page, if appropriate, when you are finished.

PLEASE PRINT – **do NOT write *cursively*** – ANSWERS IN THE SPACE PROVIDED ONLY – **NOT IN THE MARGINS** – PREFERABLY USING A BALLPOINT PEN TO INCREASE LEGIBILITY. Good luck!

On my honor, I have neither given nor received unauthorized aid on this exam and I pledge not to divulge information regarding its contents to those who have not yet taken it.

\_\_\_\_\_  
SIGNATURE

1. a. (3 pts.) Give the antecedents ("initialization, preservation, and finalization") to complete the **Repeat\_Until** Rule of Inference (ROI) given in the Solution Notes of Problem Set 5:

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**{P} repeat S until b {Q}**

- b. (3 pts.) Consider the following assertion of weak correctness. (You may assume that N represents the number of elements in integer array List[1:N] and that the value of integer variable Index is always  $\leq N$ .)

$\{N \geq 1\}$

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Index := 0
repeat
    Index := Index+1;
    Found := (Key=List[Index])
until (Found OR Index=N)
  
```

$\{(Found \wedge Key=List[Index]) \vee (\neg Found \wedge \forall (1 \leq i \leq N) \bullet Key \neq List[i])\}$

Which one of the following expressions could be used as a Q-adequate Invariant to prove this assertion using the **Repeat\_Until** ROI given in the Problem Set 5 Solution Notes? (Circle ONE only.)

- i.  $(Found \wedge Key=List[Index]) \vee (\neg Found \wedge \forall (Index < i \leq N) \bullet Key \neq List[i])$
- ii.  $[(Found \wedge Key=List[Index]) \vee (\neg Found \wedge \forall (i \in [1, Index) \cup (Index, N]) \bullet Key \neq List[i])] \wedge iorder$  (where  $iorder = \forall 1 \leq i < N \bullet List[i] \geq List[i+1]$ )
- iii.  $(Found \wedge Key=List[Index]) \vee (\neg Found \wedge \forall (1 \leq i \leq Index) \bullet Key \neq List[i])$
- iv.  $Key=List[Index] \vee (\neg Found)$
- v.  $(Found \wedge Key=List[Index] \wedge \forall (1 \leq i < Index) \bullet Key \neq List[i])$
- vi.  $(Found \wedge Key=List[Index] \wedge \forall (1 \leq i < Index) \bullet Key \neq List[i]) \vee (\neg Found)$
- vii. true

- c. (12 pts.) Use the Q-adequate Invariant expression you identified in (b) above to prove that the "Initialization" and "Preservation" antecedents (only) of the **Repeat\_Until** ROI for the given assertion hold. **Simplify logical expressions where appropriate and show all steps, including the effect of each individual assignment statement.**

(cont'd)

i. Proof that “Initialization” holds:

ii. Proof that “Preservation” holds:

2. (10 pts.) Consider the assertion of **weak** correctness:  $\{z < 0\} S \{y = z + 1\}$  where  $z, y$  are integers. Which of the following observations/facts **would** allow one to deduce that the assertion is false, and which **would not**? Circle either "would" or "would not" as appropriate, considering the observations individually. To compensate for random guessing, your score in points will be 2 times the number of [correct minus incorrect] answers, or 0 – whichever is greater. Therefore, if you are not more than 50% sure of your answer, consider skipping the problem.

- |   |       |           |
|---|-------|-----------|
| a. When the initial value of $z$ is $-1$ , $S$ terminates and the final value of $y$ is NOT equal to $0$ .                        | would | would not |
| b. When the initial value of $z$ is even, $S$ terminates and the final value of $y$ is even.                                      | would | would not |
| c. When the initial value of $z$ is $\geq -5$ , $S$ terminates and the final value of $y$ is less than the initial value of $z$ . | would | would not |
| d. When the initial value of $z$ is $-17$ , $S$ terminates and the final value of $y$ is twice the final value of $z$ .           | would | would not |
| e. The program $S$ is:<br>$z := z + 1;$<br>$y := z + 1$<br>end_while  | would | would not |

3. (10 pts.) Circle either "valid" or "invalid" for each of the following *hypothesized* Rules of Inference. To compensate for random guessing, your score in points will be 2 times the number of [correct minus incorrect] answers, or 0 – whichever is greater. Therefore, if you are not more than 50% sure of your answer, consider skipping the problem.

- |  |       |         |
|--|-------|---------|
| a. $\frac{P \Rightarrow (b \wedge Q)}{\{P\} \text{ repeat } S \text{ until } b \{Q\}} \quad ?$   | valid | invalid |
| b. $\frac{I \Rightarrow P, \{I\} S \{I\}, P \Rightarrow Q}{\{P\} \text{ repeat } S \text{ until } b \{Q\}} \quad ?$                                | valid | invalid |
| c. $\frac{(\neg b) \Rightarrow (P \wedge Q)}{\{P\} \text{ while } b \text{ do } S \{Q\}} \quad ?$  | valid | invalid |
| d. $\frac{\{true\} S \{I\} \text{ strongly}, (I \wedge b) \Rightarrow Q}{\{P\} \text{ repeat } S \text{ until } b \{Q\} \text{ strongly}} \quad ?$ | valid | invalid |
| e. $\frac{\{true\} \text{ if } b \text{ then } S \{\neg b \wedge Q\} \text{ strongly}}{\{P\} \text{ while } b \text{ do } S \{Q\}} \quad ?$        | valid | invalid |