ROI Analysis and Review of Foundations

- 1. Interpreting conditional assertions
- 2. Assertions of weak and strong program correctness are conditionals!
- 3. Where do ROI's come from?
- How to analyze the validity of a hypothesized ROI

- Interpretation of "A => B": "if A is true, then B must (in every case) be true."
- When is "A => B" <u>false</u> (i.e., "invalid")?

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Find a counterexample that PROVES

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is <u>false</u> (i.e., "invalid")?

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Let x=4. Then "x is even" is true, but "x>10" is false.

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• Interpretation of hypothesized ROI:

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• When is the ROI <u>false</u> (i.e., "invalid")?

If and only if there exists a case for which A and B and C are true, but D is false.

Find a counterexample that PROVES

$$(x=17)=>(x \text{ is odd})$$

 $x=17$

is "invalid" (i.e., false)?

Find a counterexample that PROVES

is "invalid" (i.e., false)?

Let x=4. Then "(x=17)=>(x is odd)" is true but "x=17" is false.

Find a counterexample that PROVES

is "invalid" (i.e., false)?

Let x=4. Then "(x=17)=>(x is odd)" is true but "x=17" is false.

Very important corollary: (A=>B) ≠> A

A really smokin' counter-example...

Consider the following assertion/ROI:

"People who wear red shirts do not smoke."

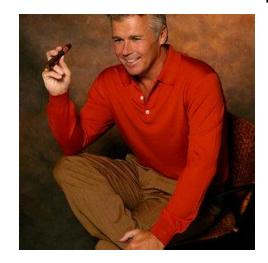
Wears red shirts(X) => Does not smoke(X)

Wears red shirts(X)

Does not smoke(X)

A really smokin' counter-example... (cont'd)

- Is the assertion/ROI valid (true)?
- No. Proof by counterexample:



 This person satisfies the antecedent, but not the consequent!

weak correctness: {P} S {Q}

Interpretation of {P} S {Q}: "if the input (initial state) satisfies pre-condition P and (if) program S executes and terminates, then the output (final state) must (in every case) satisfy post-condition Q."

weak correctness: {P} S {Q}

- Interpretation of {P} S {Q}: "if the input (initial state) satisfies pre-condition P and (if) program S executes and terminates, then the output (final state) must (in every case) satisfy post-condition Q."
- When is {P} S {Q} false (i.e., "invalid")?

weak correctness: {P} S {Q}

- Interpretation of {P} S {Q}: "if the input (initial state) satisfies pre-condition P and (if) program S executes and terminates, then the output (final state) must (in every case) satisfy post-condition Q."
- When is {P} S {Q} false (i.e., "invalid")?

If <u>and only if</u> there exists a case for which Q will be false when S terminates, given that P held before S executes.

weak correctness: {P} S {Q} (cont'd)

Example: true or false?

$$\{y=17\}$$
 if x<>0 do y := 0 $\{y=0\}$

weak correctness: {P} S {Q} (cont'd)

Example: true or false?

$$\{y=17\}$$
 if $x<>0$ do $y:=0$ $\{y=0\}$

The assertion is FALSE (invalid) since Q will not hold on termination when x is initially equal to 0 (which is not precluded by the given pre-condition).

strong correctness: {P} S {Q} strongly

Interpretation of {P} S {Q} strongly:
 "{P} S {Q} AND P => S terminates"

strong correctness: {P} S {Q} strongly

- Interpretation of {P} S {Q} strongly:
 "{P} S {Q} AND P => S terminates"
- When is "{P} S {Q} strongly" false (i.e., "invalid")?

strong correctness: {P} S {Q} strongly

- Interpretation of {P} S {Q} strongly:
 "{P} S {Q} AND P => S terminates"
- When is "{P} S {Q} strongly" false (i.e., "invalid")?

If and only if either {P} S {Q} is false OR
P #> S terminates

strong correctness: {P} S {Q} strongly (cont'd)

• Example: true or false?

```
\{y=17 \land |x|=1\} while y<>0 do y := y-x \{y=0\} strongly
```

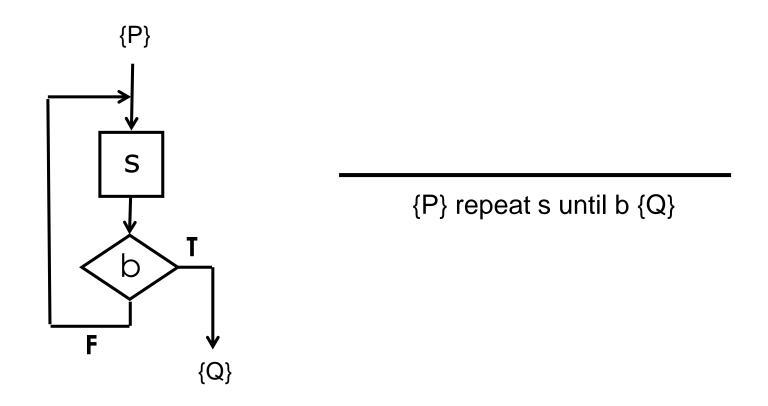
strong correctness: {P} S {Q} strongly (cont'd)

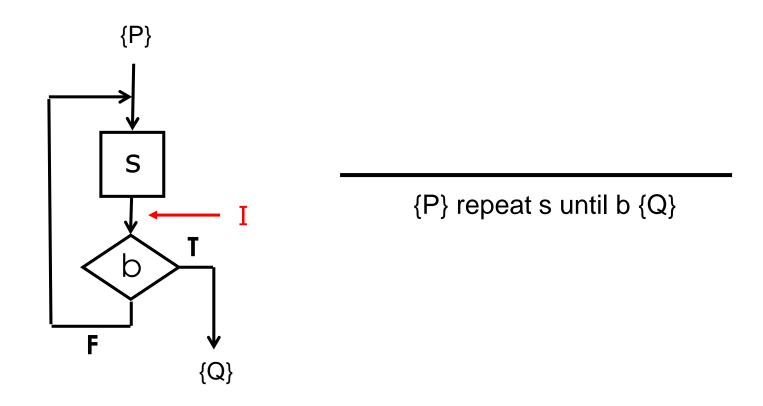
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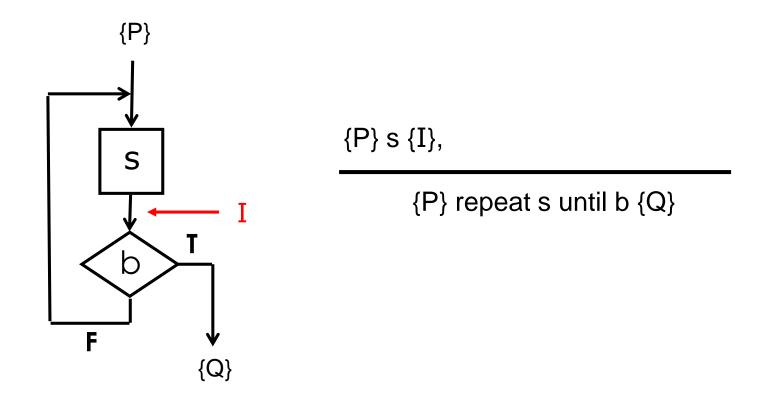
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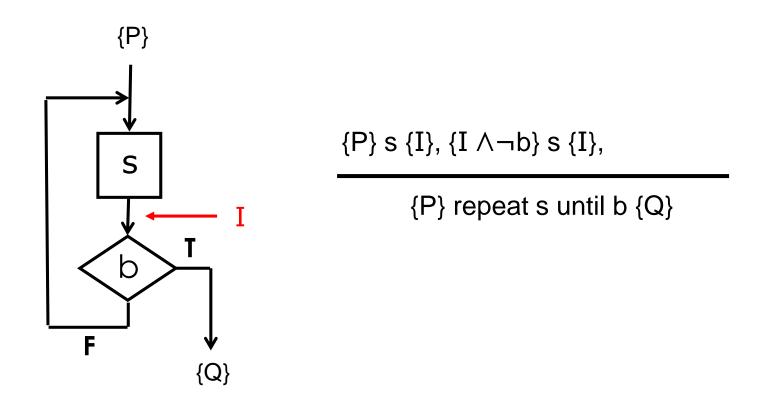
The assertion is FALSE (invalid) since $P \neq S$ terminates (x may equal -1).

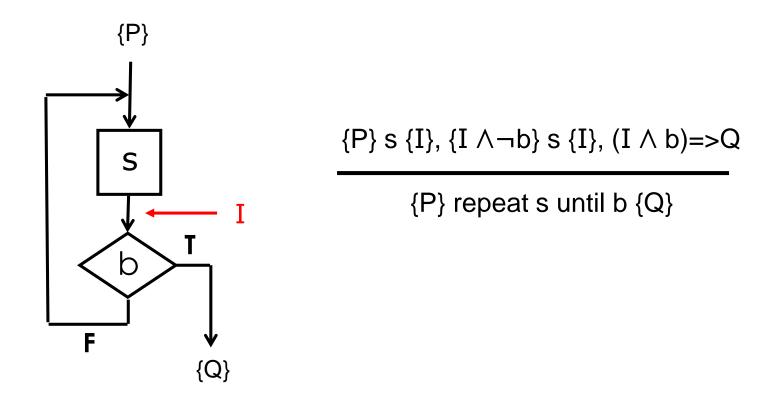
From Problem Set 5: Axiomatic Verification "Hints and Notes"

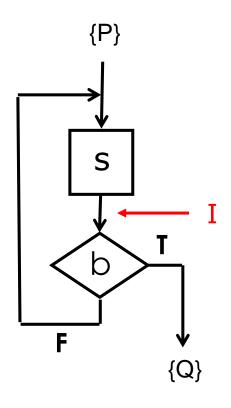












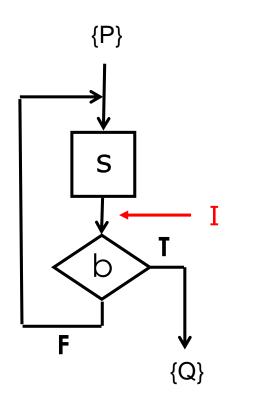
{P} s {I}, {I
$$\land \neg b$$
} s {I}, (I $\land b$)=>Q

{P} repeat s until b {Q}

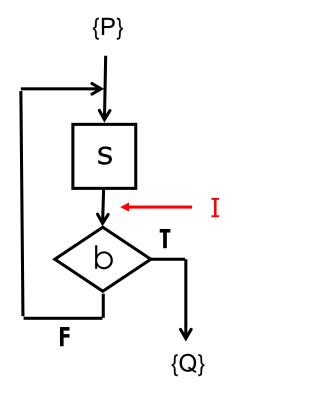
case 1 (s executes just once): ?

case 2 (s executes just twice): ?

case n>2: ?



case n>2: ?



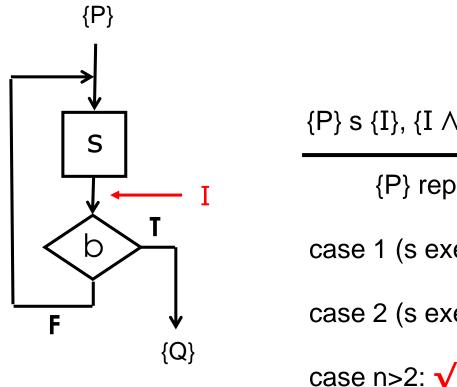
{P} s {I}, {I
$$\land \neg b$$
} s {I}, (I $\land b$)=>Q

{P} repeat s until b {Q}

case 1 (s executes just once): \checkmark

case 2 (s executes just twice): \checkmark

case n>2: ?



6. Consider the following HYPOTHESIZED rules of inference for the "while" construct:

a.
$$\frac{P \Rightarrow (\neg b \land Q)}{\{P\} \text{ while b do s } \{Q\}}?$$
b.
$$\frac{\{P \land b\} \text{ s } \{I\}, \{I \land b\} \text{ s } \{I\}, (I \land \neg b) \Rightarrow Q}{\{P\} \text{ while b do s } \{Q\}}$$

...Clearly indicate whether or not the rule is **valid**. If valid, provide an assertion of the form {P} while b do S {Q} for which it *could* be used. If not valid, prove this by providing a counterexample.

6. Consider the following HYPOTHESIZED rules of inference for the "while" construct:

a.
$$\frac{P \Rightarrow (\neg b \land Q)}{P}$$
?
$$\frac{P}{P} \text{ while b do s } \{Q\}$$

b.
$$\frac{\{P \land b\} \ s \ \{I\}, \{I \land b\} \ s \ \{I\}, (I \land \neg b) => Q}{\{P\} \ while \ b \ do \ s \ \{Q\}}$$

...Clearly indicate whether or not the rule is **valid**. If valid, provide an assertion of the form {P} while b do S {Q} for which it *could* be used. If not valid, prove this by providing a counterexample.

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$$\frac{P \Rightarrow (\neg b \land Q)}{\{P\} \text{ while b do s } \{Q\}}$$
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Is the rule **valid**?

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 case 0: Does the ROI antecedent => that Q will hold if s executes 0 times given that P holds initially?

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$$\frac{P}{P} \text{ while b do s } \{Q\}$$

 case 0: Does the ROI antecedent => that Q will hold if s executes 0 times given that P holds initially? YES

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$$\frac{P \Rightarrow (\neg b \land Q)}{P}$$
?
$$\frac{P}{P} \text{ while b do s } \{Q\}$$

- case 0: Does the ROI antecedent => that Q will hold if s executes 0 times given that P holds initially? YES
- case N>0: Does the ROI antecedent => that Q will hold if s executes 1 or more times given that P holds initially?

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 N/A (since the antecedent implies that when P holds initially, the loop must terminate without s being executed)

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 N/A (since the antecedent implies that when P holds initially, the loop must terminate without s being executed)

The rule could be employed, for example, to prove: $\{x=17\}$ while x<0 do x:=0 $\{x>0\}$

b.
$$\frac{\{P \land b\} s \{I\}, \{I \land b\} s \{I\}, (I \land \neg b) \Rightarrow Q}{\{P\} \text{ while b do s } \{Q\}}$$
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...Clearly indicate whether or not the rule is **valid**. If valid, provide an assertion of the form {P} while b do S {Q} for which it *could* be used. If not valid, prove this by providing a counterexample.

b.
$$\frac{\{P \land b\} \ s \ \{I\}, \{I \land b\} \ s \ \{I\}, (I \land \neg b) \Rightarrow Q}{\{P\} \ while \ b \ do \ s \ \{Q\}}$$

Is the rule **valid**?

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$$\frac{\{P \land b\} s \{I\}, \{I \land b\} s \{I\}, (I \land \neg b) \Rightarrow Q}{\{P\} \text{ while b do s } \{Q\}}$$
?

Is the rule **valid**?

case 0: Do the ROI antecedents => that Q will hold if b is initially false and the loop terminates with s not executing, given that P holds initially?

b.
$$\frac{\{P \land b\} s \{I\}, \{I \land b\} s \{I\}, (I \land \neg b) \Rightarrow Q}{\{P\} \text{ while b do s } \{Q\}}$$
?

Is the rule **valid**?

case 0: Do the ROI antecedents => that Q will hold if b is initially false and the loop terminates with s not executing, given that P holds initially?

NO, since neither of the first two antecedents implies that I will hold when the loop terminates for this case. (Therefore, the third antecedent cannot be used to deduce that Q will hold.)

b.
$$\frac{\{P \land b\} \ s \ \{I\}, \ \{I \land b\} \ s \ \{I\}, \ (I \land \neg b) \Rightarrow Q}{\{P\} \ while \ b \ do \ s \ \{Q\}}?$$

Question: How can this be proven using a counterexample?

b.
$$\frac{\{P \land b\} s \{I\}, \{I \land b\} s \{I\}, (I \land \neg b) \Rightarrow Q}{\{P\} \text{ while b do s } \{Q\}}$$
?

Question: How can this be proven using a counterexample?

Answer: (1) Identify a specific, concrete program of the form while b do s together with pre- and post-conditions such that $\{P\}$ while b do s $\{Q\}$ does NOT hold. (2) Identify an invariant I such that all three antecedents of the rule DO hold. This would prove that the rule is not valid for the same reason that x=4 serves as a counterexample proving the rule: ["x is even" \Rightarrow x>10] is not valid.

b.
$$\frac{\{P \land b\} s \{I\}, \{I \land b\} s \{I\}, (I \land \neg b) \Rightarrow Q}{\{P\} \text{ while b do s } \{Q\}}$$
?

Proof:

```
\{y \neq 17\} I: y=17
while x>0 do
y := 17
x := x-1
end_while
\{y=17\}
```

The three antecedents hold for the invariant y=17 but the consequent does not since the initial value of x may be ≤ 0 initially, in which case Q would not hold on termination.

```
\{\text{true}\}\ S\ \{I\}\ \text{strongly, } (I\ \land\ b)\Rightarrow Q
\{P\}\ \text{repeat S until b } \{Q\}\ \text{strongly}
```

7. Consider the following HYPOTHESIZED rule of inference for the "repeat until" construct:

```
\{\text{true}\}\ S\ \{I\}\ \text{strongly, } (I\ \Lambda\ b)\Rightarrow Q
\{P\}\ \text{repeat S until b } \{Q\}\ \text{strongly}
```

- case 1: Do the ROI antecedents => that Q will hold if S executes just once given that P holds initially?

7. Consider the following HYPOTHESIZED rule of inference for the "repeat until" construct:

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\{ true \} S \{ I \} strongly, (I \land b) \Rightarrow Q
= \{ P \} repeat S until b \{ Q \} strongly
```

- case 1: Do the ROI antecedents => that Q will hold if S executes just once given that P holds initially? YES
- case 2: ...that Q will hold if S executes exactly twice...?

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- case 1: Do the ROI antecedents => that Q will hold if S executes just once given that P holds initially? YES
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```

- case 1: Do the ROI antecedents => that Q will hold if S executes just once given that P holds initially? YES
- case 2: ...that Q will hold if S executes exactly twice...? YES
- case N: ...that Q hold if S executes exactly N times...?

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```

- case 1: Do the ROI antecedents => that Q will hold if S executes just once given that P holds initially? YES
- case 2: ...that Q will hold if S executes exactly twice...? YES
- case N: ...that Q hold if S executes exactly N times...? YES
- Do the antecedents => that (<u>repeat</u> S <u>until b</u>) will **terminate** given that P holds initially?

```
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{P} repeat S until b {Q} strongly
```

- case 1: Do the ROI antecedents => that Q will hold if S executes just once given that P holds initially? YES
- case 2: ...that Q will hold if S executes exactly twice...? YES
- case N: ...that Q hold if S executes exactly N times...? YES
- Do the antecedents => that (<u>repeat</u> S <u>until b</u>) will **terminate** given that P holds initially? NO, so the ROI is invalid.

$$P => (b \land Q)$$

$$P => (b \land Q)$$

$$P => (b \land Q)$$

$$Q = P$$

8. Consider the following HYPOTHESIZED rule of inference for the "repeat until" construct:

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$$P => (b \land Q)$$

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NO, since P and ($b \land Q$) holding initially do NOT imply that either P or ($b \land Q$) will hold after S executes.

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$$P => (b \land Q)$$

– case 1: Does the ROI antecedent => that Q will hold if S executes just once given that P holds initially?

NO, since P and $(b \land Q)$ holding initially do NOT imply that either P or $(b \land Q)$ will hold after S executes.

Exercise: PROVE the ROI is invalid using a suitable counterexample.

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