

# Software Testing and Verification

## Problem Set 6: Predicate Transforms

1. Consider the assertion of *weak* correctness:  $\{t=5 \wedge z<0\} \text{ s } \{y=z+1 \wedge t=z\}$ . Which of the following observations/facts would allow one to deduce that the assertion is FALSE and which would not? Consider the observations individually and briefly justify your answer for each.

- a. The  $\text{wp}(s, y=z)$  is  $z>-5$ .
- b. The  $\text{wlp}(s, y=z)$  is  $z>-5$ .
- c.  $(t=5 \wedge z<0) \not\Rightarrow \text{wlp}(s, y=z+1 \vee t=z)$  (Note the logical **OR**.)
- d.  $[\text{"sp}(s, t=5 \wedge z>-5) \text{ is undefined"}] \wedge [\text{sp}(s, t=5 \wedge z\leq-5) \Rightarrow (y=z+1 \wedge t=z)]$

2. Consider the program:

```
temp := x
x := y
y := temp
if temp > z then
  y := z
  z := temp
  if x > y then
    temp := x
    x := y
    y := temp
  end_if
end_if
```

Under what circumstances will the program result in  $\{x \leq y < z\}$ ? (Hint: determine the  $\text{wp}$  of the program w.r.t. the desired result.)

3. We have learned that  $(P \Rightarrow \text{wp}(s, Q))$  implies  $(\{P\} \text{ s } \{Q\})$ . However,  $(\{P\} \text{ s } \{Q\})$  does NOT imply  $(P \Rightarrow \text{wp}(s, Q))$ . Why?

4. a. Use the wlp rule for while-do statements given in Lecture Notes 20 to find the weakest liberal pre-condition of the following program with respect to the post-condition  $Z=XY$ .

```

while J<>Y do
  Z := Z+X;
  J := J+1
end_while

```

- b. Consider the invariant,  $I: Z=XJ$ , used with the while-loop ROI to prove the assertion given on slide 11 of Lecture Notes 18. How does it compare in "strength" to the weakest liberal pre-condition from part (a) above? (In particular, is one STRONGER than the other?) Briefly explain your answer.
5. Prove "finalization" of the weakest liberal pre-condition of a while statement. That is, prove the tautology:

$$(\text{wlp}(\text{while } b \text{ do } s, Q) \wedge \sim b) \Rightarrow Q.$$

6. a. Identify  $H_1, H_2, H_3$ , and  $H_k$  such that

$$\text{wp}(\text{Repeat } s \text{ Until } c, Q) = H_1 \vee H_2 \vee H_3 \vee \dots \vee H_k \vee \dots$$

where  $H_i$  represents the necessary and sufficient condition that "Repeat  $s$  Until  $c$ " terminates in state  $Q$  after  $i$  executions of  $s$ . For  $i > 1$ ,  $H_i$  should be expressed as a function of  $H_{i-1}$ .

- b. Use your formulations to determine the weakest pre-condition of:

```

Repeat
  x := x+1;
  y := y-1
Until y=0

```

with respect to the post-condition:  $x=17$ . (Show the values of  $H_1, H_2, H_3$ , and  $H_k$ , and then express the result in closed form.)